

Lohmeyer

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draft proposed
American National Standard
for information systems -

SMALL COMPUTER SYSTEM INTERFACE - 2
(SCSI-2)

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Abstract: This standard defines mechanical, electrical, and functional requirements for attaching physically small computers with each other and with intelligent peripherals such as rigid disks, flexible disks, magnetic tape devices, printers, optical disks, and scanners. The resulting interface facilitates the interconnection of physically small computers and intelligent peripherals and thus provides a common interface specification for both systems integrators and suppliers of intelligent peripherals.

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1. Scope

This American National Standard defines an input/output bus for interconnecting computers and peripheral devices. The standard defines extensions to the Small Computer System Interface (X3.131-1986), referred to herein as SCSI-1. It also provides more complete standardization of the previously defined command sets. The document includes the necessary specification of the mechanical, electrical, and functional characteristics of the interface to allow inter-operability of devices meeting the standard. This standard is referred to herein as SCSI-2. The term SCSI is used wherever it is not necessary to distinguish between the two versions of SCSI.

SCSI is a local I/O bus that can be operated over a wide range of data rates. The primary objective of the interface is to provide host computers with device independence within a class of devices. Thus, different disk drives, tape drives, printers, optical media drives, and other devices can be added to the host computers without requiring modifications to generic system hardware or software. Provision is made for the addition of special features and functions through the use of vendor unique fields and codes. Reserved fields and codes are provided for future standardization.

A second key objective of SCSI-2 is to provide compatibility with those SCSI-1 devices that support bus parity and that meet conformance level 2 of SCSI-1. While some previously vendor unique commands and parameters have been defined by the SCSI-2 standard, devices meeting SCSI-1 and SCSI-2 can co-exist on the same bus. It is intended that those operating systems providing support for both command sets be able to operate in environments mixing SCSI-1 and SCSI-2 devices. Properly conforming SCSI-1 devices, both initiators and targets, should respond in an acceptable manner to reject SCSI-2 protocol extensions. All SCSI-2 protocol extensions are designed to be permissive of such rejections and to allow the SCSI-1 device to continue operation without requiring the use of the extension.

A third key objective of SCSI-2 is to move device-dependent intelligence out to the SCSI-2 devices. This requires the definition of a command set that allows a sophisticated operating system to obtain all required initialization information from the attached SCSI-2 devices. The formalized sequence of requests identify the type of attached SCSI-2 device, the characteristics of the device, and all the changeable parameters supported by the device. Further requests can determine the readiness of the device to operate, the types of media supported by the device, and all other pertinent system information. Those parameters not required by the operating system for operation, initialization, or system tuning are not exposed to the SCSI-2 interface, but are managed by the SCSI-2 device itself.

The interface uses logical rather than physical addressing for all data blocks. For direct-access devices, each logical unit may be interrogated to determine how many blocks it contains. A logical unit may coincide with all or part of a peripheral device.

The interface protocol includes provision for the connection of multiple initiators (SCSI devices capable of initiating an operation) and multiple targets (SCSI devices capable of responding to a request to perform an operation). Distributed arbitration (i.e., bus-contention logic) is built into the architecture of SCSI. A priority system awards interface control to the highest priority SCSI device that is contending for use of the bus. The time to complete arbitration is independent of the number of devices that are contending and can be completed in less than 10 microseconds.

The physical characteristics are described in Section 4. There are two electrical alternatives: single-ended and differential. Single-ended and differential devices are electrically different and should not be mixed on the same bus.

Provision is made for cable lengths up to 25 meters using differential drivers and receivers. A single-ended driver and receiver configuration is defined for cable lengths of up to 6 meters and is primarily intended for applications within a cabinet.

Section 5 describes the logical characteristics of the interface. Arbitration is defined to permit multiple initiators and to permit concurrent I/O operations. All SCSI devices are required to be capable of operating with the defined asynchronous transfer protocol. In addition, an optional synchronous transfer protocol is defined. Section 5 also specifies a message protocol for control of the interface. In most cases, messages are not directly apparent to the host computer software.

The SCSI command and status structure is specified in Section 6. Commands are classified as mandatory (M), optional (O), or vendor unique (V). SCSI devices are required to implement all mandatory commands defined for the appropriate device type and may implement other commands as well. SCSI devices contain commands that facilitate the writing of self-configuring software drivers that can "discover" all necessary attributes without prior knowledge of specific peripheral characteristics (such as storage capacity).

Many commands also implement a very large logical block address space (2³² blocks), although some commands implement a somewhat smaller logical block address space (2²¹ blocks).

Section 7 specifies those commands that have a consistent meaning for all device types.

Sections 8 through 17 contain commands for direct-access (e.g., magnetic disk), sequential-access (e.g., magnetic tape), printer, processor, write-once (e.g., optical disk), CD-ROM devices, scanner devices, optical memory devices, medium changer devices, and communications devices, respectively. The commands in each of these sections are unique to the device type, or they have interpretations, fields, or features that are specific for the device type. Thus, for example, although the WRITE command is used for several device types, it has a somewhat different form for each type, with different parameters and meanings. Therefore, it is specified separately for each device type.

Appendixes A through C provide examples of SCSI signal sequences, timing, and phase sequences. Appendix D contains information on other standards related to medium types and density codes for flexible disks and magnetic tapes. Appendix E describes data integrity in command queuing environments. Appendix F describes normal procedures following a power-on condition. Appendix G describes the worst case skew times for a fast SCSI implementation. Appendix H contains information on other SCSI standardization activities. Appendix I contains the additional sense codes and operation codes in numerical order. Appendix J contains the vendor identification codes as of the date of this document. However, the appendixes are not a required part of this standard.

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2. Referenced Standards and Organizations

The requirements of ISO 8482:1987, "Information Processing - twisted pair multi-point interconnection data communication" apply to the use of differential drivers and receivers.

American National Standard Small Computer System Interface, X3.131-1986, may also be useful to achieve compatibility with devices that conform to version 1 of SCSI.

The medium catalog numbers in the CD-ROM section (13) are controlled by the Uniform Product Code Council, 8163 Old Yankee Road, Suite J, Dayton, Ohio 45459, U.S.A. and the European Article Number Council, Rue des Colonies, 54-BTE8, 1000 Brussels, Belgium.

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3. Glossary and Conventions

3.1. Glossary

This section contains a glossary of special terms used in this standard. Also see the glossaries for each device type in sections 8.4, 9.4, ..., 17.4.

active I/O process. An I/O process that is presently in execution (not queued).

AEN. Asynchronous event notification (see 6.5.5).

AWG. American Wire Gauge.

byte. In this standard, this term indicates an 8-bit (octet) construct.

CA. Contingent Allegiance (see 6.6).

command descriptor block (CDB). The structure used to communicate commands from an initiator to a target.

command queue. This term refers to the queue used to store the tagged I/O processes (see 6.8.2).

connect. The initiator function that selects a target to establish a nexus and to start an I/O process. The connection that results is an initial connection.

connection. An initial connection or reconnection. A connection can only occur between one initiator and one target.

contact. The electrically-conductive portion of a connector associated with a single conductor in a cable.

current I/O process. The I/O process that is presently connected on the SCSI bus.

disconnect. The action that occurs when an SCSI device releases control of the SCSI bus, allowing it to go to the BUS FREE phase.

ECA. Extended Contingent Allegiance (see 6.7).

field. A group of one or more contiguous bits. Fields containing only one bit are usually referred to as the xx bit instead of the xx field.

host adapter. A device which connects between a host system and the SCSI bus. The device usually performs the lower layers of the SCSI protocol and normally operates in the initiator role. This function may be integrated into the host system.

initial connection. An initial connection is the result of a connect and it exists from the assertion of the BSY signal in a SELECTION phase until the next BUS FREE phase occurs.

initiator. An SCSI device (usually a host system) that requests an I/O process to be performed by another SCSI device (a target).

invalid. An illegal (reserved) or unsupported bit, field, or code value.

I/O process. An I/O process consists of one initial connection and zero or more reconnections, all pertaining to a single command or a group of linked commands. More specifically, the connection(s) pertain to a nexus as defined below in which zero or more command descriptor blocks are transferred. An I/O process begins with the establishment of a nexus. An I/O process normally ends with the BUS FREE phase following successful transfer of a COMMAND COMPLETE or a RELEASE RECOVERY message. An I/O process also ends with the BUS FREE phase following an ABORT, ABORT TAG, BUS DEVICE RESET, or CLEAR QUEUE message or when a hard RESET condition or an unexpected disconnect occurs.

I_T nexus. A nexus which exists between an initiator and a target.

I_T_L nexus. A nexus which exists between an initiator, a target, and a logical unit. This relationship replaces the prior I_T nexus.

I_T_R nexus. A nexus which exists between an initiator, a target, and a target routine. This relationship replaces the prior I_T nexus.

I_T_x nexus. A nexus which is either an I_T_L or I_T_R nexus.

I_T_L_Q nexus. A nexus between an initiator, a target, a logical unit, and a queue tag following the successful receipt of one of the queue tag messages. This relationship replaces the prior I_T_L nexus.

I_T_x_y nexus. A nexus which is either an I_T_x or I_T_L_Q.

logical block. A unit of data supplied or requested by an initiator.

logical unit. A physical or virtual peripheral device addressable through a target.

logical unit number. An encoded three-bit identifier for the logical unit.

LSB. Least significant bit.

LUN. Logical unit number.

mm. Millimeter.

ms. Millisecond.

MSB. Most significant bit.

nexus. A relationship that begins with the establishment of an initial connection and ends with the completion of the I/O process. The relationship may be restricted to specify a single logical unit or target routine by the successful transfer of an IDENTIFY message. The relationship may be further restricted by the successful transfer of a queue tag message.

ns. Nanosecond.

one. A true signal value or a true condition of a variable.

page. Several commands use regular parameter structures that are referred to as pages. These pages are identified with a value known as a page code.

parameter. A structure containing one or more fields.

peripheral device. A physical peripheral device that can be attached to an SCSI device, which connects to the SCSI bus. The peripheral device and the SCSI device (peripheral controller) may be physically packaged together. Often there is a one-to-one mapping between peripheral devices and logical units, but this is not required. Examples of peripheral devices are: magnetic disks, printers, optical disks, and magnetic tapes.

queue tag. The value associated with an I/O process that uniquely identifies it from other queued I/O processes on the logical unit for the same initiator.

queued I/O process. An I/O process that is in the command queue and has not begun execution.

reconnect. The act of reviving a nexus to continue an I/O process. A target reconnects to an initiator by using the RESELECTION and MESSAGE IN phases after winning arbitration. An initiator reconnects to a target by using the SELECTION and MESSAGE OUT phases after winning arbitration (see 6.5.2).

reconnection. A reconnection is the result of a reconnect and it exists from the assertion of the BSY signal in a SELECTION or RESELECTION phase until the next BUS FREE phase occurs.

reserved. The term used for bits, fields, and code values that are set aside for future standardization.

SCSI. Either SCSI-1 or SCSI-2.

SCSI-1. The Small Computer System Interface (X3.131-1986).

SCSI-2. The Small Computer System Interface - 2 (X3.131-198X).

SCSI address. The octal representation of the unique address (0-7) assigned to an SCSI device. This address would normally be assigned and set in the SCSI device during system installation.

SCSI ID. The bit-significant representation of the SCSI address referring to one of the signal lines DB(7-0).

SCSI device. A host adapter or a target controller that can be attached to the SCSI bus.

signal assertion. The act of driving a signal to the true state.

signal negation. The act of driving a signal to the false state or allowing the cable terminators to bias the signal to the false state (by placing the driver in the high impedance condition).

signal release. The act of allowing the cable terminators to bias the signal to the false state (by placing the driver in the high impedance condition).

status. One byte of information sent from a target to an initiator upon completion of each command.

target. An SCSI device that performs an operation requested by an initiator.

target routine. A target routine is an I/O process directed to a target, and not to a logical unit.

third-party. When used in reference to COPY commands, third-party means a COPY command issued to one device to perform a copy operation between two other devices. When used in reference to RESERVE, or RELEASE commands, third-party means a reservation made on behalf of another device (e.g., A processor device requests that a direct-access device reserve itself for exclusive use by a sequential-access device).

unexpected disconnect. A disconnection that occurs as a result of a protocol error (see 5.1.1).

us. Microsecond.

vendor specific (VS). Something (e.g., a bit, field, code value, etc.) that is not defined by this standard and may be used differently in various implementations.

xx. Numbers that are not immediately followed by lower-case "b" or "h" are decimal values. Large Numbers are not separated by commas or spaces (e.g., 12345; not 12,345 or 12 345).

xxb. Numbers immediately followed by lower-case "b" are binary values.

xxh. Numbers immediately followed by lower-case "h" are hexadecimal values.

zero. A false signal value or a false condition of a variable.

3.2. Editorial Conventions

Certain words and terms used in this standard have a specific meaning beyond the normal English meaning. These words and terms are defined either in the glossary (see 3.1 and 8.4, 9.4, ..., 17.4) or in the text where they first appear. Names of signals, phases, messages, commands, statuses, sense keys, additional sense codes, and additional sense code qualifiers are in all upper-case (e.g., REQUEST SENSE). Lower-case is used for words having the normal English meaning.

4. Physical Characteristics

This section contains the physical definition of SCSI-2. The connectors, cables, signals, terminators, and bus timing values needed to implement the interface are covered.

4.1. Physical Description

SCSI devices are daisy-chained together using a common 50-conductor A cable and, optionally, a 68-conductor B cable. Both ends of each cable are terminated. All signals are common between all SCSI devices on the A cable. In systems that employ the wide SCSI option, wide SCSI devices additionally connect to the B cable. Various width SCSI devices may be mixed.

Two driver/receiver alternatives are specified:

- (1) Single-ended drivers and receivers, which allow a maximum cable length of six meters (primarily for connection within a cabinet).
- (2) Differential drivers and receivers, which allow a maximum cable length of 25 meters (primarily for connection outside of a cabinet).

The single-ended and differential alternatives are mutually exclusive within a system.

IMPLEMENTORS NOTE: Use of single-ended drivers and receivers with the fast synchronous data transfer option is not recommended.

4.2. Cable Requirements

An ideal impedance match with cable terminators implies a cable characteristic impedance of 132 ohms (single-ended option) or 122 ohms (differential option). The cable characteristic impedance shall be no less than 90 ohms and no greater than 140 ohms. It is recommended that the cable characteristic impedance be greater than 100 ohms for all cable types. See 4.2.3 for cable requirements when implementing the fast synchronous data transfer option.

A minimum conductor size of 0.08042 square mm (28 AWG) should be used to minimize noise effects and ensure proper distribution of terminator power.

IMPLEMENTORS NOTES:

- (1) To minimize discontinuities and signal reflections, cables of different impedances should not be used in the same bus. Implementations may require trade-offs in shielding effectiveness, cable length, the number of loads, transfer rates, and cost to achieve satisfactory system operation.
- (2) To minimize discontinuities due to local impedance variation, a flat cable should be spaced at least 1.27 mm (0.050 in) from other cables, any other conductor, or the cable itself when the cable is folded.
- (3) Regulatory agencies may require use of larger wire size.

4.2.1. Single-Ended Cable

A 50-conductor flat cable or 25-signal twisted-pair cable shall be used for the A cable. A 68-conductor flat cable or 34-signal twisted-pair cable shall be used for the B cable if the wide SCSI option is implemented. The maximum cumulative cable length shall be 6.0 meters. If twisted-pair cables are used, then twisted pairs in the cable shall be wired to physically opposing contacts in the connector.

A stub length of no more than 0.1 meters is allowed off the mainline interconnection within any connected equipment or from any connected point.

IMPLEMENTORS NOTE: Stub clustering should be avoided. Stubs should be spaced at least 0.3 meters apart.

SCSI bus termination shall be at each end of the cable and may be internal to the SCSI devices that are at the ends of the cable.

4.2.2. Differential Cable

A 50-conductor flat cable or 25-signal twisted-pair cable shall be used for the A cable. A 68-conductor flat cable or 34-signal twisted-pair cable shall be used for the B cable if the wide SCSI option is implemented. The maximum cumulative cable length shall be 25 meters. If twisted-pair cables are used, then twisted pairs in the cable shall be wired to physically opposing contacts in the connector.

A stub length of no more than 0.2 meters is allowed off the mainline interconnection within any connected equipment or from any connected point.

SCSI bus termination shall be at each end of the cable and may be internal to the SCSI devices that are at the ends of the cable.

IMPLEMENTORS NOTE: The use of twisted pair cable (either twisted-flat or discrete wire twisted pairs) is strongly recommended. Without twisted pairs, even at slow data rates and very short distances, crosstalk between adjacent signals causes spurious pulses with differential signals.

4.2.3. Cable Requirements for Fast Synchronous Data Transfer

In systems which use the fast synchronous data transfer option (see 4.8), the A and B cables should meet the conductor size recommendation in 4.2. The cable should have an overall shield suitable for termination in a shielded connector.

In such systems, the cables shall have the following electrical characteristics:

Characteristic Impedance: 90 to 132 ohms
 Signal Attenuation: 0.095 dB maximum per meter at 5 MHz
 Pair-to-Pair Propagation Delay Delta: 0.20 ns maximum per meter
 DC Resistance: 0.230 ohms maximum per meter at 20 degrees C

4.3. Connector Requirements

Two types of connectors are defined: nonshielded and shielded. The nonshielded connectors are typically used for in-cabinet applications. Shielded connectors are typically used for external applications where electromagnetic compatibility (EMC) and electrostatic discharge (ESD) protection may be required. Either type of connector may be used with the single-ended or differential drivers.

4.3.1. Nonshielded Connector Requirements

Two nonshielded connector alternatives are specified for the A cable and one nonshielded connector is specified for the B cable.

4.3.1.1. Nonshielded Connector Alternative 1 - A Cable

The alternative 1 nonshielded high-density SCSI device connector for the A cable (Figure 4-1) shall be a 50-conductor connector consisting of two rows of 25 female contacts with adjacent contacts 1.27 mm (0.05 in) apart. The nonmating portion of the connector is shown for reference only.

The alternative 1 nonshielded high-density cable connector for the A cable (Figure 4-2) shall be a 50-conductor connector consisting of two rows of 25 male contacts with adjacent contacts 1.27 mm (0.05 in) apart. The nonmating portion of the connector is shown for reference only.

4.3.1.2. Nonshielded Connector Alternative 2 - A Cable

The alternative 2 nonshielded low-density SCSI device connector for the A cable (Figure 4-3) shall be a 50-conductor connector consisting of two rows of 25 male pins with adjacent pins 2.54 mm (0.1 in) apart. A shroud and header body should be used. The nonmating portion of the connector is shown for reference only.

The alternative 2 nonshielded low-density cable connector for the A cable (Figure 4-4) shall be a 50-conductor connector consisting of two rows of 25 female contacts with adjacent contacts 2.54 mm (0.1 in) apart. It is recommended that keyed connectors be used.

4.3.1.3. Nonshielded Connector - B Cable

The nonshielded high-density SCSI device connector for the B cable (Figure 4-1) shall be a 68-conductor connector consisting of two rows of 34 female contacts with adjacent contacts 1.27 mm (0.05 in) apart. The nonmating portion of the connector is shown for reference only.

The nonshielded high-density cable connector for the B cable (Figure 4-2) shall be a 68-conductor connector consisting of two rows of 34 male contacts with adjacent contacts 1.27 mm (0.05 in) apart. The nonmating portion of the connector is shown for reference only.

4.3.2. Shielded Connector Requirements

Two shielded connector alternatives are specified for the A cable and one shielded connector is specified for the B cable. The connector shielding system should provide a dc resistance of less than 10 milliohms from the cable shield at its termination point to the SCSI device enclosure.

In order to support daisy-chain connections, SCSI devices that use shielded connectors should provide two shielded device connectors on the device enclosure. These two connectors may be wired "one-to-one" with a stub to the SCSI device's drivers and receivers provided the maximum stub length is not violated. Alternatively, two cables may be run from the two shielded connectors to the drivers and receivers so that the maximum stub length is not violated. The length of the cable within the device enclosure is included when calculating the total cable length of the SCSI bus.

IMPLEMENTORS NOTE: SCSI-1 defined three shielded connector systems in an appendix. The alternative 1 shielded connector of SCSI-1 has been replaced by a high-density connector in this standard. The alternative 2 shielded connector remains unchanged. The EUROCARD Boxes shielded connector system of SCSI-1 has been deleted in this standard.

4.3.2.1. Shielded Connector Alternative 1 - A Cable

The shielded high-density SCSI device connector for the A cable (Figure 4-5) is a 50-conductor connector consisting of two rows of 25 female contacts with adjacent contacts 1.27 mm (0.05 in) apart. The nonmating portion of the connector is shown for reference only.

The shielded high-density cable connector for the A cable (Figure 4-6) is a 50-conductor connector consisting of two rows of 25 male contacts with adjacent contacts 1.27 mm (0.05 in) apart. The nonmating portion of the connector is shown for reference only.

4.3.2.2. Shielded Connector Alternative 2 - A Cable

The shielded low-density device connector for the A cable (Figure 4-7) is a 50-conductor connector consisting of two rows of ribbon contacts spaced 2.16 mm (0.085 in) apart. The nonmating portion of the connector is shown for reference only.

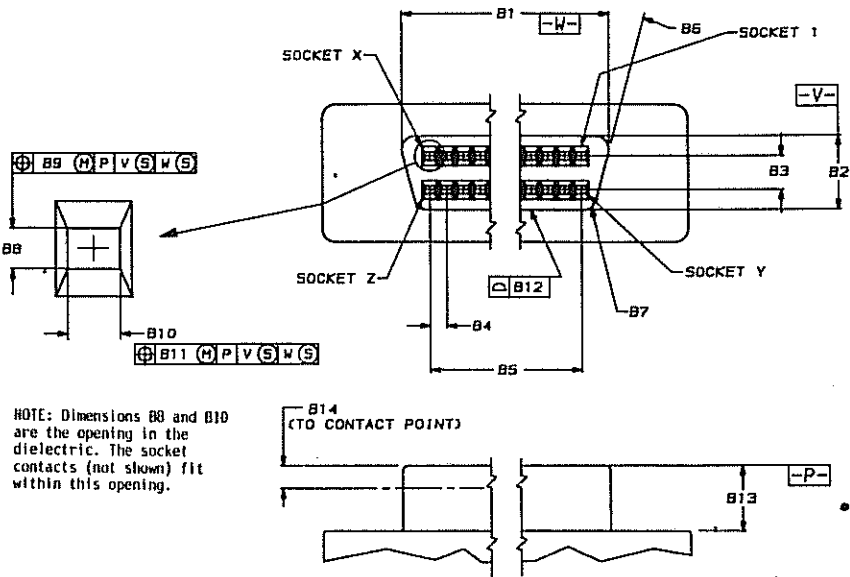
The shielded low-density cable connector for the A cable (Figure 4-8) is a 50-conductor connector consisting of two rows of ribbon contacts spaced 2.16 mm (0.085 in) apart. The nonmating portion of the connector is shown for reference only.

4.3.2.3. Shielded Connector - B Cable

The shielded high-density SCSI device connector for the B cable (Figure 4-5) is a 68-conductor connector consisting of two rows of 34 female contacts with adjacent contacts 1.27 mm (0.05 in) apart. The nonmating portion of the connector is shown for reference only.

The shielded high-density cable connector for the B cable (Figure 4-6) is a 68-conductor connector consisting of two rows of 34 male contacts with adjacent contacts 1.27 mm (0.05 in) apart. The nonmating portion of the connector is shown for reference only.

UNSHIELDED RECEPTACLE
(BOARD VERSION)



DIMENSIONS	50 POSITION		68 POSITION	
	MILLIMETERS	INCHES	MILLIMETERS	INCHES
B1	34.70	1.366	46.13	1.816
B2	5.54	.218	5.54	.218
B3	2.54	.100	2.54	.100
B4	1.27	.050	1.27	.050
B5	30.48	1.200	41.91	1.650
B6	15*	.15*	15*	.15*
B7	1.00 R	.039 R	1.00 R	.039 R
B8	0.61±.05	.024±.002	0.61±.05	.024±.002
B9	0.15	.006	0.15	.006
B10	0.86±.10	.034±.004	0.86±.10	.034±.004
B11	0.15	.006	0.15	.006
B12	0.05	.002	0.05	.002
B13	5.00±.13	.197±.005	5.00±.13	.197±.005
B14	1.75 MAX	.069 MAX	1.75 MAX	.069 MAX
SOCKET X	25		34	
SOCKET Y	26		35	
SOCKET Z	50		68	

Figure 4-1: 50/68-Contact Nonshielded High-Density SCSI Device Connector (A Cable/B Cable)

UNSHIELDED PLUG (CABLE VERSION)

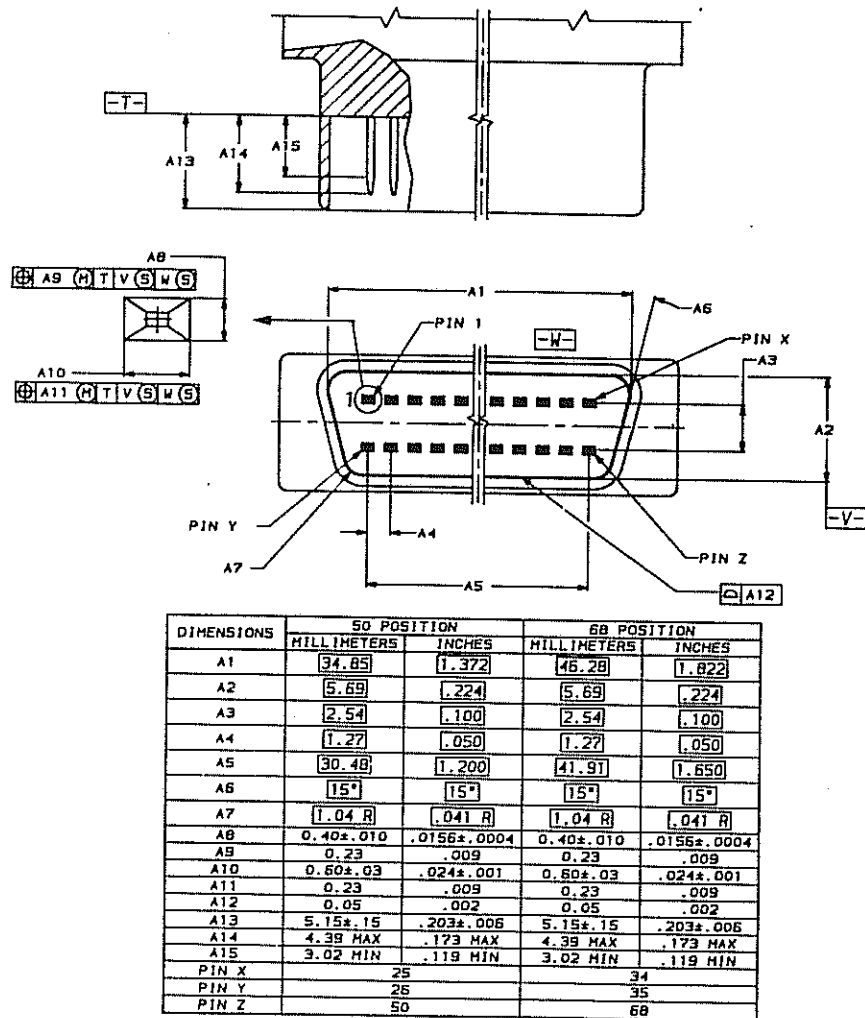
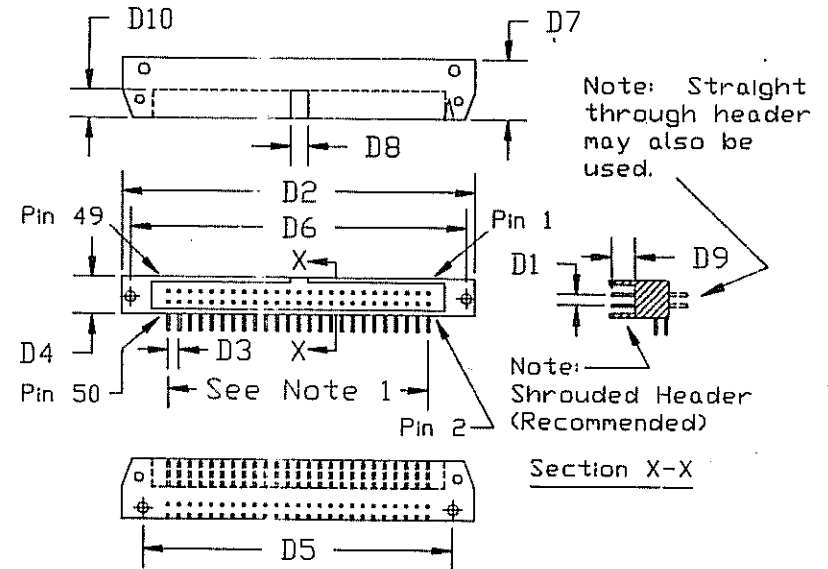


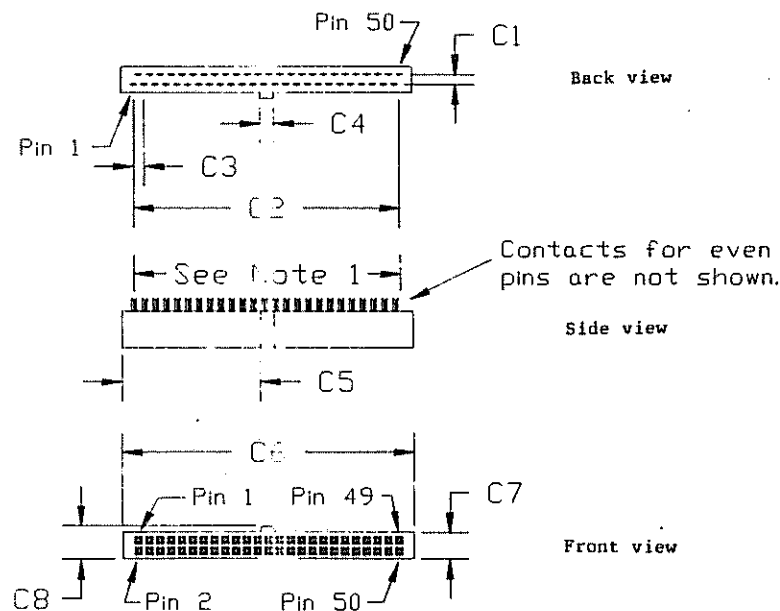
Figure 4-2: 50/68-Contact Nonshielded High-Density Cable Connector
(A Cable/B Cable)



NOTES:

- (1) Two rows of twenty five contacts on 2.54-mm (0.100-inch) spacing = 60.96 mm (2.400 inch).
- (2) Tolerances ±0.127 mm (0.005 inch) noncumulative, unless specified otherwise.

Figure 4-3: 50-Contact Nonshielded Low-Density SCSI Device Connector
(A Cable)

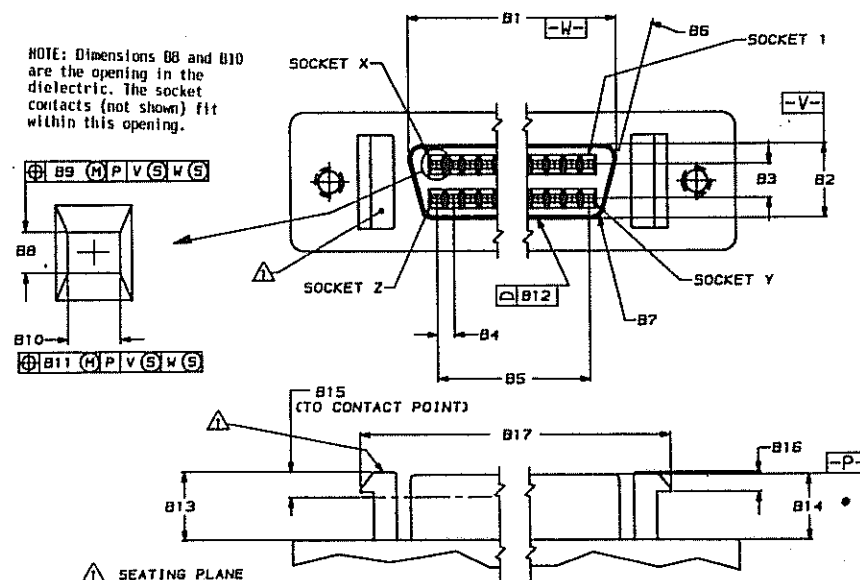


Dimensions	Millimeters	Inches	Comments
C1	2.540	0.100	
C2	60.960	2.400	
C3	2.540	0.100	
C4	3.302	0.130	
C5	32.385	1.275	
C6	68.072	2.680	
C7	6.096	0.240	
C8	7.620	0.300	Maximum

NOTES:

- (1) Fifty contacts on 1.27-mm (0.05-inch) staggered spacing = 62.23 mm (2.450 inch) [reference only].
- (2) Tolerances ± 0.127 mm (0.005) noncumulative, unless specified otherwise.
- (3) Connector cover and strain relief are optional.

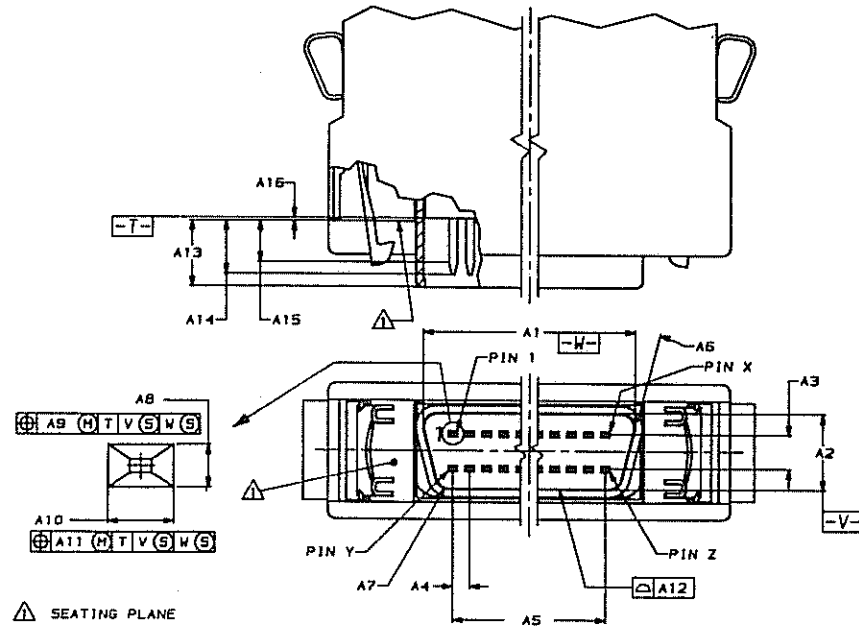
Figure 4-4: 50-Contact Nonshielded Low-Density Cable Connector (A Cable)

SHIELDED RECEPTACLE
(BOARD VERSION)

DIMENSIONS	50 POSITION		68 POSITION	
	MILLIMETERS	INCHES	MILLIMETERS	INCHES
B1	34.70	1.366	46.13	1.816
B2	5.54	.218	5.54	.218
B3	2.54	.100	2.54	.100
B4	1.27	.050	1.27	.050
B5	30.48	1.200	41.91	1.650
B6	15°	15°	15°	15°
B7	1.00 R	.039 R	1.00 R	.039 R
B8	0.61±.05	.024±.002	0.61±.05	.024±.002
B9	0.15	.006	0.15	.006
B10	0.86±.10	.034±.004	0.86±.10	.034±.004
B11	0.15	.006	0.15	.006
B12	0.05	.002	0.05	.002
B13	5.10±.05	.201±.002	5.10±.05	.201±.002
B14	5.00±.13	.197±.005	5.00±.13	.197±.005
B15	1.85 MAX	.073 MAX	1.85 MAX	.073 MAX
B16	1.50±.03	.059±.001	1.50±.03	.059±.001
B17	42.29±.10	1.665±.004	53.72±.10	2.115±.004
SOCKET X	25		34	
SOCKET Y	26		35	
SOCKET Z	50		68	

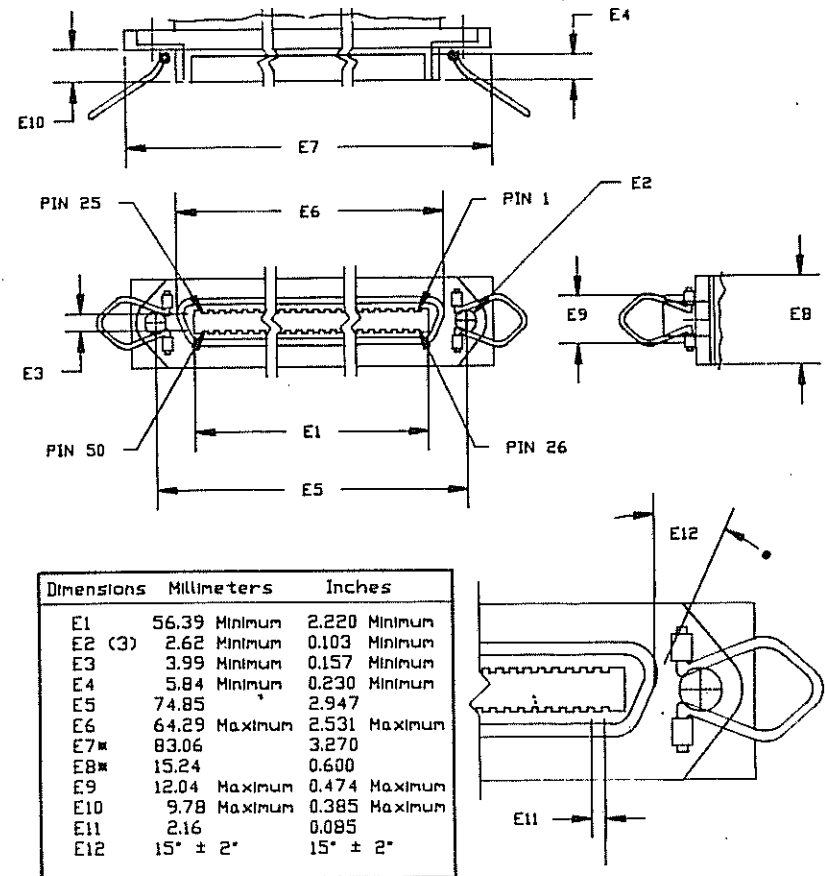
Figure 4-5: 50/68-Contact Shielded High-density SCSI Device Connector (A Cable/B Cable)

SHIELDED PLUG (CABLE VERSION)



DIMENSIONS	50 POSITION		68 POSITION	
	MILLIMETERS	INCHES	MILLIMETERS	INCHES
A1	34.85	1.372	46.28	1.822
A2	5.69	.224	5.69	.224
A3	2.54	.100	2.54	.100
A4	1.27	.050	1.27	.050
A5	30.48	1.200	41.91	1.650
A6	15*	.15*	15*	.15*
A7	1.04 R	.041 R	1.04 R	.041 R
A8	0.40±.010	.0156±.0004	0.40±.010	.0156±.0004
A9	0.23	.009	0.23	.009
A10	0.60±.03	.024±.001	0.60±.03	.024±.001
A11	0.23	.009	0.23	.009
A12	0.05	.002	0.05	.002
A13	4.90±.10	.193±.004	4.90±.10	.193±.004
A14	4.27 MAX	.168 MAX	4.27 MAX	.168 MAX
A15	2.64 MIN	.104 MIN	2.64 MIN	.104 MIN
A16	0.25±.13	.010±.005	0.25±.13	.010±.005
PIN X	25		34	
PIN Y	26		35	
PIN Z	50		68	

Figure 4-6: 50/68-Contact Shielded High-density Cable Connector
(A Cable/B Cable)

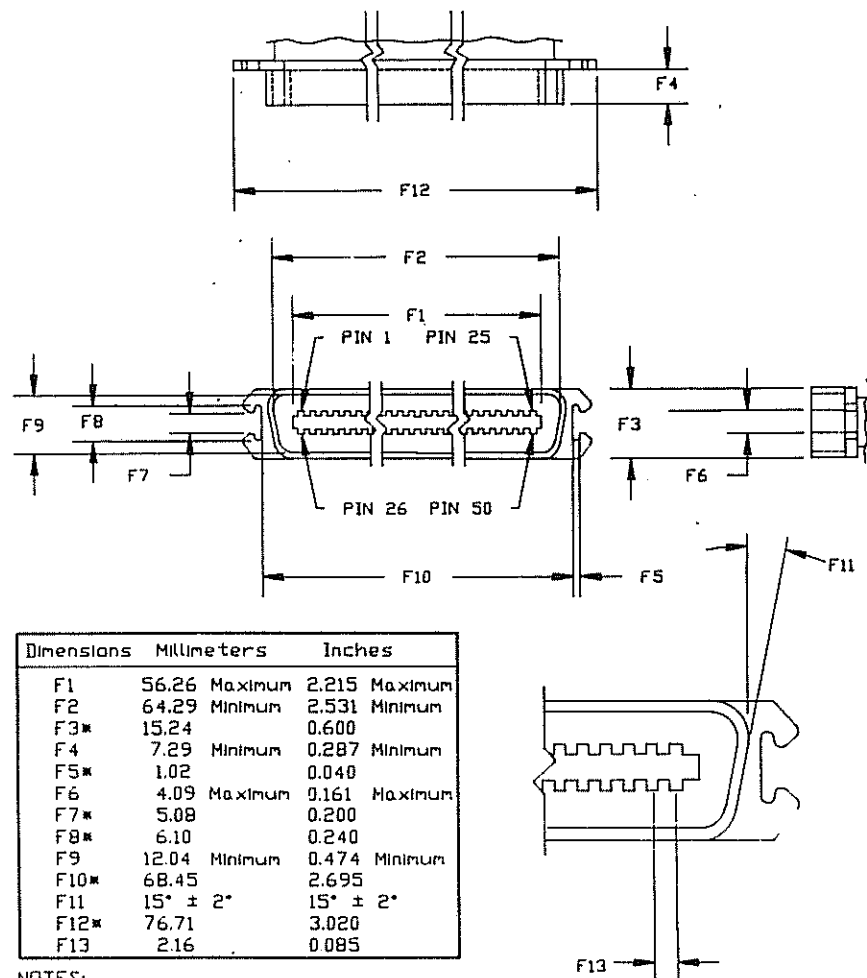


Dimensions	Millimeters	Inches
E1	56.39 Minimum	2.220 Minimum
E2 (3)	2.62 Minimum	0.103 Minimum
E3	3.99 Minimum	0.157 Minimum
E4	5.84 Minimum	0.230 Minimum
E5	74.85	2.947
E6	64.29 Maximum	2.531 Maximum
E7*	83.06	3.270
E8*	15.24	0.600
E9	12.04 Maximum	0.474 Maximum
E10	9.78 Maximum	0.385 Maximum
E11	2.16	0.085
E12	15" ± 2"	15" ± 2"

NOTES:

- (1) Tolerances ± 0.127 mm (0.005 inch) noncumulative, unless specified otherwise.
- (2) Dimensions listed with asterisks are shown for reference only.
- (3) Dimension E2 to accommodate 4-40 or 6-32 threaded screws.

Figure 4-7: 50-Contact Shielded Low-Density SCSI Device Connector



NOTES:

- (1) Tolerances ± 0.127 mm (0.005 inch) noncumulative, unless specified otherwise.
- (2) Dimensions listed with asterisks are shown for reference only.

Figure 4-8: 50-Contact Shielded Low-Density Cable Connector

4.3.3. Connector Contact Assignments

The connector contact assignments are defined in Tables 4-1 through 4-5. Table 4-1 defines which of the other four tables to use and which set of contact assignments to use.

Table 4-1: Cross-Reference to Connector Contact Assignments

Connector Type	Driver/Receiver Type	Cable	Connector Figure	Contact Assignment Table	Contact Set
Nonshielded Alternative 1	Single-Ended	A	4-1 & 4-2	4-2	2
Nonshielded Alternative 1	Single-Ended	B	4-1 & 4-2	4-3	
Nonshielded Alternative 1	Differential	A	4-1 & 4-2	4-4	2
Nonshielded Alternative 1	Differential	B	4-1 & 4-2	4-5	
Nonshielded Alternative 2	Single-Ended	A	4-3 & 4-4	4-2	1
Nonshielded Alternative 2	Differential	A	4-3 & 4-4	4-4	1
Shielded Alternative 1	Single-Ended	A	4-5 & 4-6	4-2	2
Shielded Alternative 1	Single-Ended	B	4-5 & 4-6	4-3	
Shielded Alternative 1	Differential	A	4-5 & 4-6	4-4	2
Shielded Alternative 1	Differential	B	4-5 & 4-6	4-5	
Shielded Alternative 2	Single-Ended	A	4-7 & 4-8	4-2	2
Shielded Alternative 2	Differential	A	4-7 & 4-8	4-4	2

Table 4-2: Single-Ended Contact Assignments - A Cable

Signal Name	Connector Contact Set 2	Connector Contact Set 1	Cable Conductor Number	Connector Contact Set 1	Connector Contact Set 2	Signal Name
GROUND	1	1	1	2	26	-DB(0)
GROUND	2	3	3	4	27	-DB(1)
GROUND	3	5	5	6	28	-DB(2)
GROUND	4	7	7	8	29	-DB(3)
GROUND	5	9	9	10	30	-DB(4)
GROUND	6	11	11	12	31	-DB(5)
GROUND	7	13	13	14	32	-DB(6)
GROUND	8	15	15	16	33	-DB(7)
GROUND	9	17	17	18	34	-DB(P)
GROUND	10	19	19	20	35	GROUND
GROUND	11	21	21	22	36	GROUND
RESERVED	12	23	23	24	37	RESERVED
OPEN	13	25	25	26	38	TERMPWR
RESERVED	14	27	27	28	39	RESERVED
GROUND	15	29	29	30	40	GROUND
GROUND	16	31	31	32	41	-ATN
GROUND	17	33	33	34	42	GROUND
GROUND	18	35	35	36	43	-BSY
GROUND	19	37	37	38	44	-ACK
GROUND	20	39	39	40	45	-RST
GROUND	21	41	41	42	46	-MSG
GROUND	22	43	43	44	47	-SEL
GROUND	23	45	45	46	48	-C/D
GROUND	24	47	47	48	49	-REQ
GROUND	25	49	49	50	50	-I/O

NOTES:

- (1) The minus sign next to a signal indicates active low.
- (2) The conductor number refers to the conductor position when using 0.050-inch centerline flat ribbon cable with a low-density connector or when using 0.025-inch centerline flat ribbon cable with a high-density connector. Other cable types may be used to implement equivalent contact assignments.
- (3) Two sets of contact assignments are shown. Refer to Table 4-1 to determine which set of contacts applies to each connector.
- (4) See 4.4.4 for a definition of the RESERVED lines.

Table 4-3: Single-Ended Contact Assignments - B Cable

Signal Name	Connector Contact Number	Cable Conductor Number	Connector Contact Number	Signal Name
GROUND	1	1	2	GROUND
GROUND	2	3	4	-DB(8)
GROUND	3	5	6	-DB(9)
GROUND	4	7	8	-DB(10)
GROUND	5	9	10	-DB(11)
GROUND	6	11	12	-DB(12)
GROUND	7	13	14	-DB(13)
GROUND	8	15	16	-DB(14)
GROUND	9	17	18	-DB(15)
GROUND	10	19	20	-DB(P1)
GROUND	11	21	22	-ACKB
GROUND	12	23	24	GROUND
GROUND	13	25	26	-REQB
GROUND	14	27	28	-DB(16)
GROUND	15	29	30	-DB(17)
GROUND	16	31	32	-DB(18)
TERMPWRB	17	33	34	TERMPWRB
TERMPWRB	18	35	36	TERMPWRB
GROUND	19	37	38	-DB(19)
GROUND	20	39	40	-DB(20)
GROUND	21	41	42	-DB(21)
GROUND	22	43	44	-DB(22)
GROUND	23	45	46	-DB(23)
GROUND	24	47	48	-DB(P2)
GROUND	25	49	50	-DB(24)
GROUND	26	51	52	-DB(25)
GROUND	27	53	54	-DB(26)
GROUND	28	55	56	-DB(27)
GROUND	29	57	58	-DB(28)
GROUND	30	59	60	-DB(29)
GROUND	31	61	62	-DB(30)
GROUND	32	63	64	-DB(31)
GROUND	33	65	66	-DB(P3)
GROUND	34	67	68	GROUND

NOTES:

- (1) The minus sign next to a signal indicates active low.
- (2) The conductor number refers to the conductor position when using 0.025-inch centerline flat ribbon cable. Other cable types may be used to implement equivalent contact assignments.

Table 4-4: Differential Contact Assignments - A Cable

Signal Name	Connector Contact Number		Cable Conductor Number	Connector Contact Number		Signal Name
	Set 2	Set 1		Set 1	Set 2	
GROUND	1	1	1	2	26	GROUND
+DB(0)	2	3	3	4	27	-DB(0)
+DB(1)	3	5	5	6	28	-DB(1)
+DB(2)	4	7	7	8	29	-DB(2)
+DB(3)	5	9	9	10	30	-DB(3)
+DB(4)	6	11	11	12	31	-DB(4)
+DB(5)	7	13	13	14	32	-DB(5)
+DB(6)	8	15	15	16	33	-DB(6)
+DB(7)	9	17	17	18	34	-DB(7)
+DB(P)	10	19	19	20	35	-DB(P)
DIFFSENS	11	21	21	22	36	GROUND
RESERVED	12	23	23	24	37	RESERVED
TERMPWR	13	25	25	26	38	TERMPWR
RESERVED	14	27	27	28	39	RESERVED
+ATN	15	29	29	30	40	-ATN
GROUND	16	31	31	32	41	GROUND
+BSY	17	33	33	34	42	-BSY
+ACK	18	35	35	36	43	-ACK
+RST	19	37	37	38	44	-RST
+MSG	20	39	39	40	45	-MSG
+SEL	21	41	41	42	46	-SEL
+C/D	22	43	43	44	47	-C/D
+REQ	23	45	45	46	48	-REQ
+I/O	24	47	47	48	49	-I/O
GROUND	25	49	49	50	50	GROUND

NOTES:

(1) The conductor number refers to the conductor position when using 0.050-inch centerline flat ribbon cable with a low-density connector or when using 0.025-inch centerline flat ribbon cable with a high-density connector. Other cable types may be used to implement equivalent contact assignments.

(2) Two sets of contact assignments are shown. Refer to Table 4-1 to determine which set of contacts applies to each connector.

(3) See 4.4.4 for a definition of the RESERVED lines.

Table 4-5: Differential Contact Assignments - B Cable

Signal Name	Connector Contact Number	Cable Conductor Number	Connector Contact Number	Signal Name
GROUND	1	1	2	GROUND
+DB(8)	2	3	4	-DB(8)
+DB(9)	3	5	6	-DB(9)
+DB(10)	4	7	8	-DB(10)
+DB(11)	5	9	10	-DB(11)
+DB(12)	6	11	12	-DB(12)
+DB(13)	7	13	14	-DB(13)
+DB(14)	8	15	16	-DB(14)
+DB(15)	9	17	18	-DB(15)
+DB(P1)	10	19	20	-DB(P1)
+ACKB	11	21	22	-ACKB
GROUND	12	23	24	DIFFSENS
+REQB	13	25	26	-REQB
+DB(16)	14	27	28	-DB(16)
+DB(17)	15	29	30	-DB(17)
+DB(18)	16	31	32	-DB(18)
TERMPWRB	17	33	34	TERMPWRB
TERMPWRB	18	35	36	TERMPWRB
+DB(19)	19	37	38	-DB(19)
+DB(20)	20	39	40	-DB(20)
+DB(21)	21	41	42	-DB(21)
+DB(22)	22	43	44	-DB(22)
+DB(23)	23	45	46	-DB(23)
+DB(P2)	24	47	48	-DB(P2)
+DB(24)	25	49	50	-DB(24)
+DB(25)	26	51	52	-DB(25)
+DB(26)	27	53	54	-DB(26)
+DB(27)	28	55	56	-DB(27)
+DB(28)	29	57	58	-DB(28)
+DB(29)	30	59	60	-DB(29)
+DB(30)	31	61	62	-DB(30)
+DB(31)	32	63	64	-DB(31)
+DB(P3)	33	65	66	-DB(P3)
GROUND	34	67	68	GROUND

NOTE: The conductor number refers to the conductor position when using 0.025-inch centerline flat ribbon cable. Other cable types may be used to implement equivalent contact assignments.

4.4. Electrical Description

For the measurements in this section, SCSI bus termination is assumed to be external to the SCSI device. See 4.4.1 for the terminating requirements for the RESERVED lines. SCSI devices may have the provision for allowing optional internal termination.

4.4.1. Single-Ended Alternative

All signals not defined as RESERVED, GROUND, or TERMPWR shall be terminated at both ends of the cable. The implementor may choose one of the following two methods to terminate each end (see Figure 4-9):

(1) The termination of each signal shall consist of 220 ohms (+5%) to the TERMPWR line and 330 ohms (+5%) to ground. Using resistors with +1% tolerance improves noise margins.

(2) The termination of each signal shall meet these requirements:

(a) The terminators shall each supply a characteristic impedance between 100 and 132 ohms.

(b) The terminators shall be powered by the TERMPWR line and may receive additional power from other sources but shall not require such additional power for proper operation (see 4.4.3).

(c) The current available to any signal line driver shall not exceed 48 milliamps when the driver asserts the line and pulls it to 0.5 volts dc. Only 44.8 mA of this current shall be available from the two terminators.

(d) The voltage on all released signal lines shall be at least 2.5 volts dc when the TERMPWR line is within specified values (see 4.4.3).

(e) These conditions shall be met with any legal configuration of targets and initiators as long as at least one device is supplying TERMPWR.

The first termination method above is the same as in SCSI-1. The second termination method is recommended for better signal quality.

4.4.1.1. Output Characteristics

All signals shall use open-collector or three-state drivers. Each signal driven by an SCSI device shall have the following output characteristics when measured at the SCSI device's connector:

V_{OL} (Low-level output voltage) = 0.0 to 0.5 volts dc at 48 mA sinking

V_{OH} (High-level output voltage) = 2.5 to 5.25 volts dc (signal negation)

4.4.1.2. Input Characteristics

SCSI devices with power on shall meet the following electrical characteristics on each signal (including both receivers and passive drivers):

V_{IL} (Low-level input voltage) = 0.0 to 0.8 volts dc (signal true)

V_{IH} (High-level input voltage) = 2.0 to 5.25 volts dc (signal false)

I_{IL} (Low-level input current) = -0.4 to 0.0 mA at $V_I = 0.5$ volts dc

I_{IH} (High-level input current) = 0.0 to 0.1 mA at $V_I = 2.7$ volts dc

Minimum input hysteresis = 0.2 volts dc

Maximum input capacitance = 25 pF (measured at the device connector closest to the stub, if any, within the device)

It is recommended that SCSI devices with power off also meet the above I_{IL} and I_{IH} electrical characteristics on each signal.

To achieve maximum noise immunity and to assure proper operation with complex cable configurations, it is recommended that the nominal switching threshold be approximately 1.4 volts.

4.4.2. Differential Alternative

All signals consist of two lines denoted +SIGNAL and -SIGNAL. A signal is true when +SIGNAL is more positive than -SIGNAL, and a signal is false when -SIGNAL is more positive than +SIGNAL. All assigned signals of the A and B cables described in 4.6 shall be terminated at each end of the cable with a terminator network as shown in Figure 4-10. Resistor tolerances in the terminator network shall be +5% or less.

The DIFFSENS signal of the connector is used as an active high enable for the differential drivers. If a single-ended device or terminator is inadvertently connected, this signal is grounded, disabling the differential drivers (see Figure 4-11).

4.4.2.1. Output Characteristics

Each signal driven by an SCSI device shall have the following output characteristics when measured at the SCSI device's connector:

V_{OL} (Low-level output voltage) = 1.7 V maximum at I_{OL} (Low-level output current) = 55 mA.

V_{OH} (High-level output voltage) = 2.7 V minimum at I_{OH} (High-level output current) = -55 mA.

V_{OD} (Differential output voltage) = 1.0 V minimum with common-mode voltage ranges from -7 to +12 volts dc.

V_{OL} and V_{OH} shall be as measured between the output terminal and the SCSI device's logic ground reference.

The output characteristics shall additionally conform to ISO 8482.

4.4.2.2. Input Characteristics

SCSI devices shall meet the following electrical characteristics on each signal (including both receivers and passive drivers):

I_I (Input current on either input) = ± 2.0 mA maximum.

Maximum input capacitance = 25 pF.

The I_I requirement shall be met with the input voltage varying between -7 and +12 volts dc, with power on or off, and with the hysteresis equaling 35 millivolts, minimum.

The input characteristics shall additionally conform to ISO 8482.

4.4.3. Terminator Power

SCSI initiators shall supply terminator power to the TERMPWR contact(s) and, if it implements the wide SCSI option, to the TERMPWRB contacts. This power shall be supplied through a diode or similar semiconductor that prevents backflow of power to the SCSI device. Targets and SCSI devices that become temporary initiators (e.g., targets which implement the COPY command or asynchronous event notification) are not required to supply terminator power. Any SCSI device may supply terminator power. Interface error rates are lower if the termination voltage is maintained at the extreme ends of the cable.

All terminators independent of location shall be powered from the TERMPWR and TERMPWRB contact(s). The use of keyed connectors is recommended in SCSI devices that provide terminator power to prevent accidental grounding or misconnection of terminator power.

IMPLEMENTORS NOTE: Regulatory agencies may require limiting maximum (short circuit) current to the terminator power lines. Recommended current limiting is 1.5 amperes for TERMPWR and 2 amperes for TERMPWRB. For systems utilizing multiple initiators, the initiators may be configured with option straps or current limiting devices. Maximum available current should not exceed 5 amperes.

SCSI devices shall sink no more than 1.0 mA from TERMPWR and no more than 1.0 mA from TERMPWRB except to power an optional internal terminator.

Single-ended SCSI devices providing terminator power on cable A shall have the following characteristics:

V_{Term} = 4.25 to 5.25 volts dc

Term 900 mA minimum source drive capability

Differential SCSI devices providing terminator power on cable A shall have the following characteristics:

V_{Term} = 4.0 to 5.25 volts dc

Term 600 mA minimum source drive capability

Single-ended SCSI devices providing terminator power on cable B shall have the following characteristics:

V_{Term} = 4.25 to 5.25 volts dc

Term 1500 mA minimum source drive capability

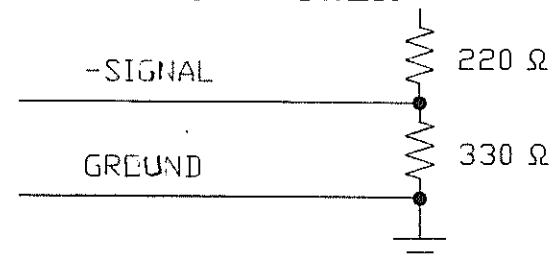
Differential SCSI devices providing terminator power on cable B shall have the following characteristics:

V_{Term} = 4.0 to 5.25 volts dc

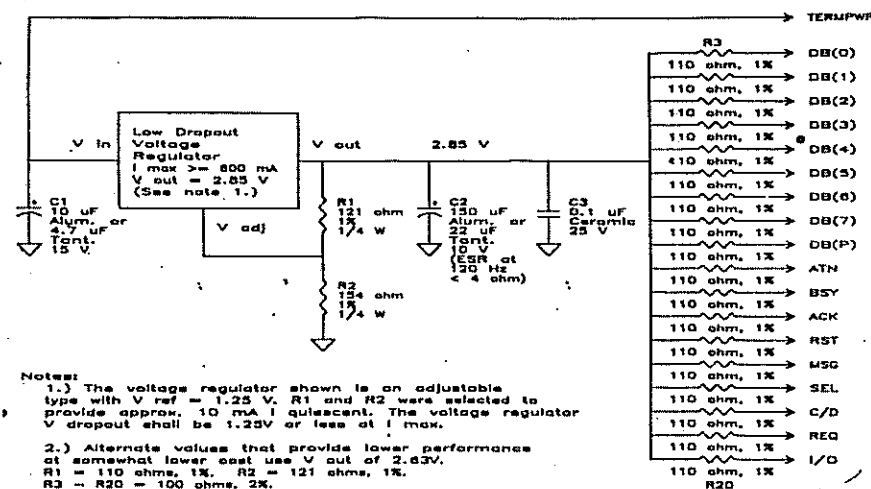
Term 1000 mA minimum source drive capability

IMPLEMENTORS NOTE: It is recommended that the terminator power lines be decoupled at each terminator with at least a 2.2 microfarad high-frequency capacitor to improve signal quality.

TERMINATOR POWER



Alternative 1 Termination



Alternative 2 Termination

Figure 4-9: Termination for Single-Ended Devices

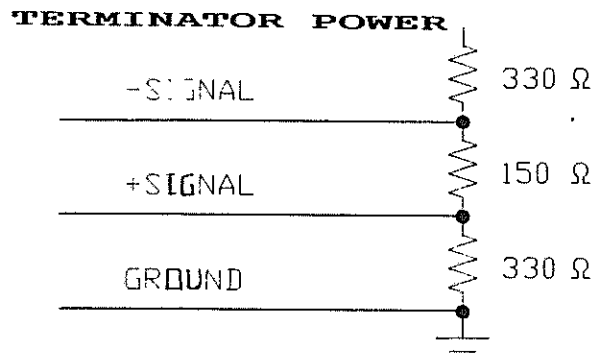


Figure 4-10: Termination for Differential Devices

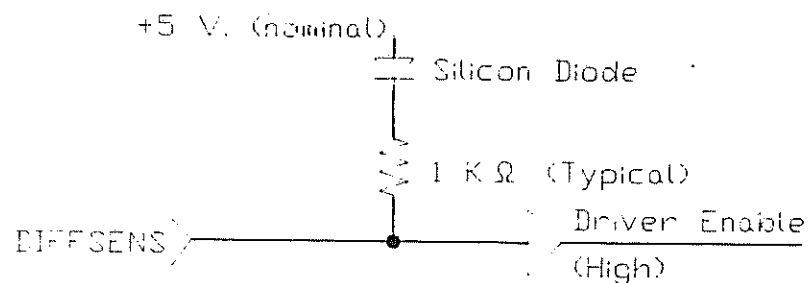


Figure 4-11: Differential Driver Protection Circuit

4.4.4. RESERVED Lines

The lines labeled RESERVED in the A cable contact assignment tables (Table 4-2 and Table 4-4) shall be connected to ground in the bus terminator assemblies or in the end devices on the SCSI cable. The RESERVED lines should be open in the other SCSI devices, but may be connected to ground.

4.5. SCSI Bus

Communication on the SCSI bus is allowed between only two SCSI devices at any given time. There is a maximum of eight SCSI devices. Each SCSI device has an SCSI ID bit assigned as shown in Figure 4-12. Three sample system configurations are shown in Figure 4-13.

When two SCSI devices communicate on the SCSI bus, one acts as an initiator and the other acts as a target. The initiator originates an operation and the target performs the operation. An SCSI device usually has a fixed role as an initiator or target, but some devices may be able to assume either role.

An initiator may address up to eight peripheral devices that are connected to a target. The target may be physically housed within the peripheral device in which case the peripheral device is referred to as an embedded SCSI device.

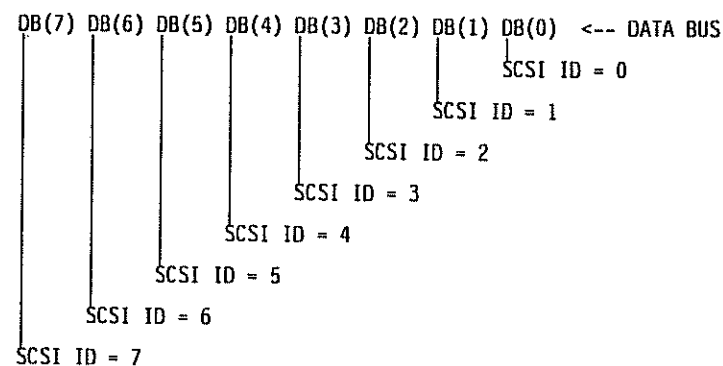


Figure 4-12: SCSI ID Bits.

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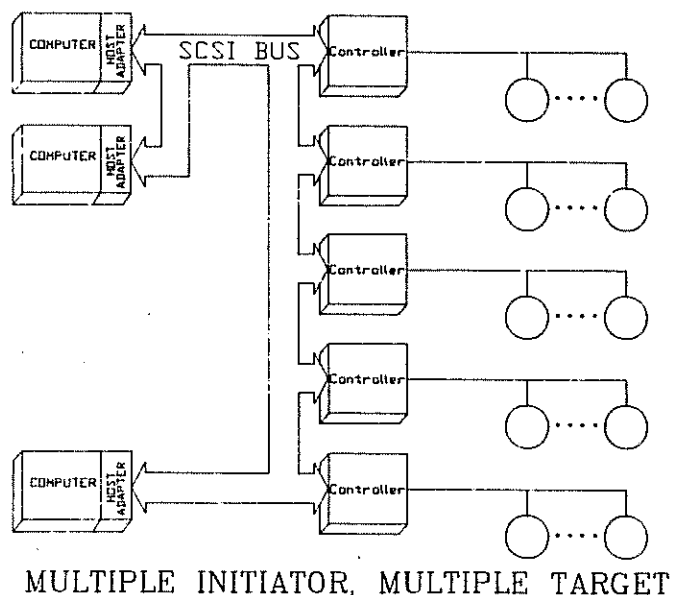
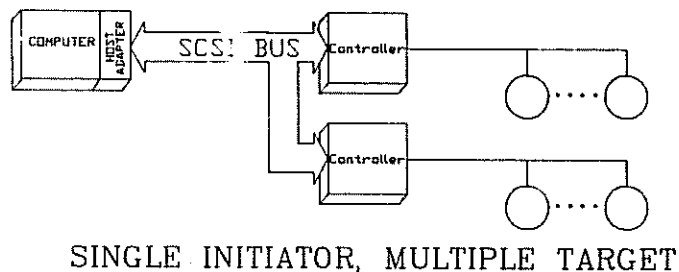
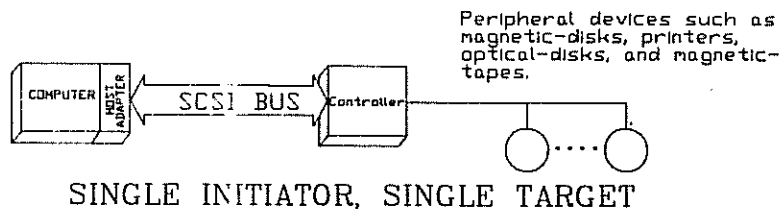


Figure 4-13: Sample SCSI Configurations

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Up to eight SCSI devices can be supported on the SCSI bus. They can be any combination of initiators and targets provided there is at least one of each.

Certain SCSI bus functions are assigned to the initiator and certain SCSI bus functions are assigned to the target. The initiator may arbitrate for the SCSI bus and select a particular target. The target may request the transfer of COMMAND, DATA, STATUS, or other information on the DATA BUS, and in some cases it may arbitrate for the SCSI bus and reselect an initiator for the purpose of continuing an operation.

Information transfers on the DATA BUS are asynchronous and follow a defined REQ/ACK handshake protocol. One byte of information may be transferred with each handshake on the A cable and, if the wide data transfer option is implemented, one or three bytes of information may be transferred with each handshake on the B cable. An option is defined for synchronous data transfer.

4.6. SCSI Bus Signals

There are a total of 18 signals on the A cable and 29 signals on the B cable. A total of 11 signals are used for control and 36 are used for data (messages, commands, status, and data), including parity. These signals are described as follows:

BSY (BUSY). An "OR-tied" signal that indicates that the bus is being used.

SEL (SELECT). An "OR-tied" signal used by an initiator to select a target or by a target to reselect an initiator.

IMPLEMENTORS NOTE: The SEL signal was not defined as "OR-tied" in SCSI-1. It has been defined as "OR-tied" in SCSI-2 in anticipation of needing another "OR-tied" signal for future standardization. This does not cause an operational problem in mixing SCSI-1 and SCSI-2 devices.

C/D (CONTROL/DATA). A signal driven by a target that indicates whether CONTROL or DATA information is on the DATA BUS. True indicates CONTROL.

I/O (INPUT/OUTPUT). A signal driven by a target that controls the direction of data movement on the DATA BUS with respect to an initiator. True indicates input to the initiator. This signal is also used to distinguish between SELECTION and RESELECTION phases.

MSG (MESSAGE). A signal driven by a target during the MESSAGE phase.

REQ (REQUEST). A signal driven by a target on the A cable to indicate a request for a REQ/ACK data transfer handshake.

REQB (REQUEST). A signal driven by a target on the B cable to indicate a request for a REQ/ACKB data transfer handshake.

ACK (ACKNOWLEDGE). A signal driven by an initiator on the A cable to indicate an acknowledgment for a REQ/ACK data transfer handshake.

ACKB (ACKNOWLEDGE). A signal driven by an initiator on the B cable to indicate an acknowledgment for a REQ/ACKB data transfer handshake.

ATN (ATTENTION). A signal driven by an initiator to indicate the ATTENTION condition.

RST (RESET). An "OR-tied" signal that indicates the RESET condition.

DB(7-0,P) (DATA BUS). Eight data-bit signals, plus a parity-bit signal that form a DATA BUS. DB(7) is the most significant bit and has the highest priority during the ARBITRATION phase. Bit number, significance, and priority decrease downward to DB(0). A data bit is defined as one when the signal value is true and is defined as zero when the signal value is false. Data parity DB(P) shall be odd. Parity is undefined during the ARBITRATION phase.

DB(31-8,P1,P2,P3) (DATA BUS). Twenty-four data-bit signals, plus three parity-bit signals that form an extension to the DATA BUS. DB(P1,P2,P3) are parity bits for DB(15-8), DB(23-16), and DB(31-24) respectively. A data bit is defined as one when the signal value is true and is defined as zero when the signal value is false. Data parity DB(Px) shall be odd.

4.6.1. Signal Values

Signals may assume true or false values. There are two methods of driving these signals. In both cases, the signal shall be actively driven true, or asserted. In the case of OR-tied drivers, the driver does not drive the signal to the false state, rather the bias circuitry of the bus terminators pulls the signal false whenever it is released by the drivers at every SCSI device. If any driver is asserted, then the signal is true. In the case of non-OR-tied drivers, the signal may be actively driven false. In this standard, wherever the term negated is used, it means that the signal may be actively driven false, or may be simply released (in which case the bias circuitry pulls it false), at the option of the implementor. The advantage to actively driving signals false during information transfer is that the transition from true to false occurs more quickly and the noise margin is much higher than if the signal is simply released. This facilitates reliable data transfer at high rates, especially at the longer cable lengths used with differential drivers.

4.6.2. OR-Tied Signals

The BSY, SEL, and RST signals shall be OR-tied only. In the ordinary operation of the bus, the BSY and RST signals may be simultaneously driven true by several drivers. No signals other than BSY, RST, and DB(P) are simultaneously driven by two or more drivers, and any signal other than BSY, SEL, and RST may employ OR-tied or non-OR-tied drivers. DB(P) shall not be driven false during the ARBITRATION phase but may be driven false in other phases. There is no operational problem in mixing OR-tied and non-OR-tied drivers on signals other than BSY and RST.

4.6.3. Signal Sources

Table 4-6 indicates which type of SCSI device is allowed to source each signal. No attempt is made to show if the source is driving asserted, driving negated, or is passive. All SCSI device drivers that are not active sources shall be in the passive state. The RST signal may be asserted by any SCSI device at any time.

Table 4-6: Signal Sources

Bus Phase	A Cable Signals					B Cable Signals		
	BSY	SEL	C/D, I/O, MSG, REQ	ACK, ATN	DB(7-0) DB(P)	REQB	ACKB	DB(31-8), DB(P1,P2,P3)
BUS FREE	None	None	None	None	None	None	None	None
ARBITRATION	All	Win	None	None	S ID	None	None	None
SELECTION	I&T	Init	None	Init	Init	None	None	None
RESELECTION	I&T	Targ	Targ	Init	Targ	None	None	None
COMMAND	Targ	None	Targ	Init	Init	None	None	None
DATA IN	Targ	None	Targ	Init	Targ	Targ	Init	Targ
DATA OUT	Targ	None	Targ	Init	Init	Targ	Init	Init
STATUS	Targ	None	Targ	Init	Targ	None	None	None
MESSAGE IN	Targ	None	Targ	Init	Targ	None	None	None
MESSAGE OUT	Targ	None	Targ	Init	Init	None	None	None

- All:** The signal shall be driven by all SCSI devices that are actively arbitrating.
- S ID:** A unique data bit (the SCSI ID) shall be driven by each SCSI device that is actively arbitrating; the other seven data bits shall be released (i.e., not driven) by this SCSI device. The parity bit (DB(P)) may be released or driven to the true state, but shall never be driven to the false state during this phase.
- I&T:** The signal shall be driven by the initiator, target, or both, as specified in the SELECTION phase and RESELECTION phase.
- Init:** If driven, this signal shall be driven only by the active initiator.
- None:** The signal shall be released; that is, not be driven by any SCSI device. The bias circuitry of the bus terminators pulls the signal to the false state.
- Win:** The signal shall be driven by the one SCSI device that wins arbitration.
- Targ:** If the signal is driven, it shall be driven only by the active target.

4.7. SCSI Bus Timing

Unless otherwise indicated, the delay-time measurements for each SCSI device, shown in Table 4-7, shall be calculated from signal conditions existing at that SCSI device's own SCSI bus connection. Thus, these measurements (except cable skew delay) can be made without considering delays in the cable. The timing characteristics of each signal are described in the following paragraphs.

Table 4-7: SCSI Bus Timing Values

Arbitration Delay	2.4 microseconds
Assertion Period	90 nanoseconds
Bus Clear Delay	800 nanoseconds
Bus Free Delay	800 nanoseconds
Bus Set Delay	1.8 microseconds
Bus Settle Delay	400 nanoseconds
Cable Skew Delay	10 nanoseconds
Data Release Delay	400 nanoseconds
Deskew Delay	45 nanoseconds
Disconnection Delay	200 microseconds
Hold Time	45 nanoseconds
Negation Period	90 nanoseconds
Power-On to Selection Time	10 seconds recommended
Reset to Selection Time	250 milliseconds recommended
Reset Hold Time	25 microseconds
Selection Abort Time	200 microseconds
Selection Time-out Delay	250 milliseconds recommended
Transfer Period	set during an SDTR message
Fast Assertion Period	30 nanoseconds
Fast Cable Skew Delay	5 nanoseconds
Fast Deskew Delay	20 nanoseconds
Fast Hold Time	10 nanoseconds
Fast Negation Period	30 nanoseconds

4.7.1. Arbitration Delay

The minimum time an SCSI device shall wait from asserting BSY for arbitration until the DATA BUS can be examined to see if arbitration has been won. There is no maximum time.

4.7.2. Assertion Period

The minimum time that a target shall assert REQ (or REQB) while using synchronous data transfers. Also, the minimum time that an initiator shall assert ACK (or ACKB) while using synchronous data transfers. REQB and ACKB timings only apply to optional wide data transfers.

4.7.3. Bus Clear Delay

The maximum time for an SCSI device to stop driving all bus signals after:

- (1) The BUS FREE phase is detected (BSY and SEL both false for a bus settle delay)
- (2) SEL is received from another SCSI device during the ARBITRATION phase
- (3) The transition of RST to true.

For the first condition above, the maximum time for an SCSI device to clear the bus is 1200 nanoseconds from BSY and SEL first becoming both false. If an SCSI device requires more than a bus settle delay to detect BUS FREE phase, it shall clear the bus within a bus clear delay minus the excess time.

4.7.4. Bus Free Delay

The minimum time that an SCSI device shall wait from its detection of the BUS FREE phase (BSY and SEL both false for a bus settle delay) until its assertion of BSY when going to the ARBITRATION phase.

4.7.5. Bus Set Delay

The maximum time for an SCSI device to assert BSY and its SCSI ID bit on the DATA BUS after it detects BUS FREE phase (BSY and SEL both false for a bus settle delay) for the purpose of entering the ARBITRATION phase.

4.7.6. Bus Settle Delay

The minimum time to wait for the bus to settle after changing certain control signals as called out in the protocol definitions.

4.7.7. Cable Skew Delay

The maximum difference in propagation time allowed between any two SCSI bus signals measured between any two SCSI devices.

4.7.8. Data Release Delay

The maximum time for an initiator to release the DATA BUS signals following the transition of the I/O signal from false to true.

4.7.9. Deskew Delay

The minimum time required for deskew of certain signals.

4.7.10. Disconnection Delay

The minimum time that a target shall wait after releasing BSY before participating in an ARBITRATION phase when honoring a DISCONNECT message from the initiator.

4.7.11. Hold Time

The minimum time added between the assertion of REQ (or REQB) or ACK (or ACKB) and the changing of the data lines to provide hold time in the initiator or target while using synchronous data transfers. REQB and ACKB timings only apply to optional wide data transfers.

4.7.12. Negation Period

The minimum time that a target shall negate REQ (or REQB) while using synchronous data transfers. Also, the minimum time that an initiator shall negate ACK (or ACKB) while using synchronous data transfers. REQB and ACKB timings only apply to optional wide data transfers.

4.7.13. Power-On to Selection Time

The recommended maximum time from power application until an SCSI target is able to respond with appropriate status and sense data to the TEST UNIT READY, INQUIRY, and REQUEST SENSE commands.

4.7.14. Reset to Selection Time

The recommended maximum time after a hard RESET condition until an SCSI target is able to respond with appropriate status and sense data to the TEST UNIT READY, INQUIRY, and REQUEST SENSE commands.

4.7.15. Reset Hold Time

The minimum time for which RST is asserted. There is no maximum time.

4.7.16. Selection Abort Time

The maximum time that a target (or initiator) shall take from its most recent detection of being selected (or reselected) until asserting a BSY response. This time-out is required to ensure that a target (or initiator) does not assert BSY after a SELECTION (or RESELECTION) phase has been aborted. This is not the selection time-out period; see 5.1.3.1 and 5.1.4.2 for a complete description.

4.7.17. Selection Time-out Delay

The minimum time that an initiator (or target) should wait for a BSY response during the SELECTION (or RESELECTION) phase before starting the time-out procedure. Note that this is only a recommended time period.

4.7.18. Transfer Period

The Transfer Period specifies the minimum time allowed between the leading edges of successive REQ pulses and of successive ACK pulses while using synchronous data transfers. (See 5.1.5.2 and 5.6.21.)

4.8. Fast Synchronous Transfer Option

When devices negotiate a synchronous data transfer period between 100 ns and 200 ns they are said to be using "fast synchronous data transfers". Devices which negotiate a synchronous data transfer period greater than 200 ns use timing parameters specified in 4.7. When a fast synchronous data transfer period is negotiated, those specific times redefined in this section are used; those not redefined remain the same. The minimum synchronous data transfer period is 100 ns.

4.8.1. Fast Assertion Period

This value is the minimum time that a target shall assert REQ (or REQB) while using fast synchronous data transfers. Also, the minimum time that an initiator shall assert ACK (or ACKB) while using fast synchronous data transfers. REQB and ACKB timings only apply to optional wide data transfers.

4.8.2. Fast Cable Skew Delay

This value is the maximum difference in propagation time allowed between any two SCSI bus signals measured between any two SCSI devices while using fast synchronous data transfers.

4.8.3. Fast Deskew Delay

This value is the minimum time required for deskew of certain signals while using fast synchronous data transfers.

4.8.4. Fast Hold Time

This value is the minimum time added between the assertion of REQ (or REQB) or ACK (or ACKB) and the changing of the data lines to provide hold time in the initiator or target, respectively, while using fast synchronous data transfers. REQB and ACKB timings only apply to optional wide data transfers.

4.8.5. Fast Negation Period

This value is the minimum time that a target shall negate REQ (or REQB) while using fast synchronous data transfers. Also, the minimum time that an initiator shall negate ACK (or ACKB) while using fast synchronous data transfers. REQB and ACKB timings only apply to optional wide data transfers.

5. Logical Characteristics

5.1. SCSI Bus Phases

The SCSI architecture includes eight distinct phases:

BUS FREE phase	
ARBITRATION phase	
SELECTION phase	
RESELECTION phase	
COMMAND phase	
DATA phase	} These phases are collectively termed the information transfer phases.
STATUS phase	
MESSAGE phase	

The SCSI bus can never be in more than one phase at any given time. In the following descriptions, signals that are not mentioned shall not be asserted.

5.1.1. BUS FREE Phase

The BUS FREE phase is used to indicate that no SCSI device is actively using the SCSI bus and that it is available. In some cases a target reverts to BUS FREE phase to indicate an error condition that it has no other way to handle. This is called as an unexpected disconnect.

SCSI devices shall detect the BUS FREE phase after the SEL and BSY signals are both false for at least a bus settle delay.

SCSI devices shall release all SCSI bus signals within a bus clear delay after the BSY and SEL signals become continuously false for a bus settle delay. If an SCSI device requires more than a bus settle delay to detect the BUS FREE phase then it shall release all SCSI bus signals within a bus clear delay minus the excess time to detect the BUS FREE phase. The total time to clear the SCSI bus shall not exceed a bus settle delay plus a bus clear delay.

Initiators normally do not expect BUS FREE phase to begin because of the target's release of the BSY signal except after one of the following occurrences:

- (1) after a reset condition is detected.
- (2) after an ABORT message is successfully received by a target.
- (3) after a BUS DEVICE RESET message is successfully received by a target.
- (4) after a DISCONNECT message is successfully transmitted from a target (see 5.6.6)
- (5) after a COMMAND COMPLETE message is successfully transmitted from a target (see 5.6.5).
- (6) after a RELEASE RECOVERY message is successfully received by a target.
- (7) after an ABORT TAG message is successfully received by a target.
- (8) after a CLEAR QUEUE message is successfully received by a target.

The BUS FREE phase may also be entered after an unsuccessful selection or reselection, although in this case it is the release of the SEL signal rather than the release of the BSY signal that first establishes the BUS FREE phase.

If an initiator detects the release of the BSY signal by the target at any other time, the target is indicating an error condition to the initiator. The

target may perform this transition to the BUS FREE phase independent of the state of the ATN signal. The initiator shall manage this condition as an unsuccessful I/O process termination. The target terminates the I/O process by clearing all pending data and status information for the affected logical unit or target routine. The target may optionally prepare sense data that may be retrieved by a REQUEST SENSE command. When an initiator detects an unexpected disconnect, it is recommended that a REQUEST SENSE command be attempted to obtain any valid sense data that may be available.

5.1.2. ARBITRATION Phase

The ARBITRATION phase allows one SCSI device to gain control of the SCSI bus so that it can initiate or resume an I/O process.

The procedure for an SCSI device to obtain control of the SCSI bus is as follows:

- (1) The SCSI device shall first wait for the BUS FREE phase to occur. The BUS FREE phase is detected whenever both the BSY and SEL signals are simultaneously and continuously false for a minimum of a bus settle delay.

IMPLEMENTORS NOTE: This bus settle delay is necessary because a transmission line phenomenon known as a "wired-OR glitch" may cause the BSY signal to briefly appear false, even though it is being driven true.

- (2) The SCSI device shall wait a minimum of a bus free delay after detection of the BUS FREE phase (i.e. after the BSY and SEL signals are both false for a bus settle delay) before driving any signal.

- (3) Following the bus free delay in Step (2), the SCSI device may arbitrate for the SCSI bus by asserting both the BSY signal and its own SCSI ID, however the SCSI device shall not arbitrate (i.e. assert the BSY signal and its SCSI ID) if more than a bus set delay has passed since the BUS FREE phase was last observed.

IMPLEMENTORS NOTE: There is no maximum delay before asserting the BSY signal and the SCSI ID following the bus free delay in Step (2) as long as the bus remains in the BUS FREE phase. However, SCSI devices that delay longer than a bus settle delay plus a bus set delay from the time when the BSY and SEL signals first become false may fail to participate in arbitration when competing with faster SCSI devices.

- (4) After waiting at least an arbitration delay (measured from its assertion of the BSY signal) the SCSI device shall examine the DATA BUS. If a higher priority SCSI ID bit is true on the DATA BUS (DB(7) is the highest), then the SCSI device has lost the arbitration and the SCSI device may release its signals and return to Step (1). If no higher priority SCSI ID bit is true on the DATA BUS, then the SCSI device has won the arbitration and it shall assert the SEL signal. Any SCSI device other than the winner has lost the arbitration and shall release the BSY signal and its SCSI ID bit within a bus clear delay after the SEL signal becomes true. An SCSI device that loses arbitration may return to Step (1).

IMPLEMENTORS NOTE: It is recommended that new implementations wait for the SEL signal to become true before releasing the BSY signal and SCSI ID bit when arbitration is lost.

(5) The SCSI device that wins arbitration shall wait at least a bus clear delay plus a bus settle delay after asserting the SEL signal before changing any signals.

NOTE: The SCSI ID bit is a single bit on the DATA BUS that corresponds to the SCSI device's unique SCSI address. All other seven DATA BUS bits shall be released by the SCSI device. Parity is not valid during the ARBITRATION phase. During the ARBITRATION phase, DB(P) may be released or asserted, but shall not be actively driven false.

5.1.3. SELECTION Phase

The SELECTION phase allows an initiator to select a target for the purpose of initiating some target function (e.g., READ or WRITE command). During the SELECTION phase the I/O signal is negated so that this phase can be distinguished from the RESELECTION phase.

The SCSI device that won the arbitration has both the BSY and SEL signals asserted and has delayed at least a bus clear delay plus a bus settle delay before ending the ARBITRATION phase. The SCSI device that won the arbitration becomes an initiator by not asserting the I/O signal.

The initiator shall set the DATA BUS to a value which is the OR of its SCSI ID bit and the target's SCSI ID bit and it shall assert the ATN signal (indicating that a MESSAGE OUT phase is to follow the SELECTION phase). The initiator shall then wait at least two deskew delays and release the BSY signal. The initiator shall then wait at least a bus settle delay before looking for a response from the target.

The target shall determine that it is selected when the SEL signal and its SCSI ID bit are true and the BSY and I/O signals are false for at least a bus settle delay. The selected target may examine the DATA BUS in order to determine the SCSI ID of the selecting initiator. The selected target shall then assert the BSY signal within a selection abort time of its most recent detection of being selected; this is required for correct operation of the selection time-out procedure.

The target shall not respond to a selection if bad parity is detected. Also, if more than two SCSI ID bits are on the DATA BUS, the target shall not respond to selection.

IMPLEMENTORS NOTE: If a target chooses to support the single initiator option of SCSI-1, it should not be an error to have only the target ID bit present during selection. In this case the target may respond as an SCSI-1 device.

No less than two deskew delays after the initiator detects the BSY signal is true, it shall release the SEL signal and may change the DATA BUS. The target shall wait until the SEL signal is false before asserting the REQ signal to enter an information transfer phase.

5.1.3.1. SELECTION Time-out Procedure

Two optional selection time-out procedures are specified for clearing the SCSI bus if the initiator waits a minimum of a selection time-out delay and there has been no BSY signal response from the target:

- (1) Optionally, the initiator shall assert the RST signal (see 5.2.2).
- (2) Optionally, the initiator shall continue asserting the SEL and ATN signals and shall release the DATA BUS. If the initiator has not detected the BSY signal to be true after at least a selection abort time plus two deskew delays, the initiator shall release the SEL and ATN signals allowing the SCSI bus to go to the BUS FREE phase. SCSI devices shall ensure that when responding to selection that the selection was still valid within a selection abort time of their assertion of the BSY signal. Failure to comply with this requirement could result in an improper selection (two targets connected to the same initiator, wrong target connected to an initiator, or a target connected to no initiator).

5.1.4. RESELECTION Phase

RESELECTION is an optional phase that allows a target to reconnect to an initiator for the purpose of continuing some operation that was previously started by the initiator but was suspended by the target, (i.e., the target disconnected by allowing a BUS FREE phase to occur before the operation was complete).

5.1.4.1. RESELECTION

Upon completing the ARBITRATION phase, the winning SCSI device has both the BSY and SEL signals asserted and has delayed at least a bus clear delay plus a bus settle delay. The winning SCSI device becomes a target by asserting the I/O signal. The winning SCSI device shall also set the DATA BUS to a value that is the logical OR of its SCSI ID bit and the initiator's SCSI ID bit. The target shall wait at least two deskew delays and release the BSY signal. The target shall then wait at least a bus settle delay before looking for a response from the initiator.

The initiator shall determine that it is reselected when the SEL and I/O signals and its SCSI ID bit are true and the BSY signal is false for at least a bus settle delay. The reselected initiator may examine the DATA BUS in order to determine the SCSI ID of the reselecting target. The reselected initiator shall then assert the BSY signal within a selection abort time of its most recent detection of being reselected; this is required for correct operation of the time-out procedure. The initiator shall not respond to a RESELECTION phase if bad parity is detected. Also, the initiator shall not respond to a RESELECTION phase if other than two SCSI ID bits are on the DATA BUS.

After the target detects the BSY signal is true, it shall also assert the BSY signal and wait at least two deskew delays and then release the SEL signal. The target may then change the I/O signal and the DATA BUS. After the reselected initiator detects the SEL signal is false, it shall release the BSY signal. The target shall continue asserting the BSY signal until it relinquishes the SCSI bus.

NOTE: When the target is asserting the BSY signal, a transmission line phenomenon known as a "wired-OR glitch" may cause the BSY signal to appear false for up to a round-trip propagation delay following the release of the

BSY signal by the initiator. This is the reason why the BUS FREE phase is recognized only after both the BSY and SEL signals are continuously false for a minimum of a bus settle delay. Cables longer than 25 meters should not be used even if the chosen driver, receiver, and cable provide adequate noise margins, because they increase the duration of the glitch and could cause SCSI devices to inadvertently detect the BUS FREE phase.

5.1.4.2. RESELECTION Time-out Procedure

Two optional RESELECTION time-out procedures are specified for clearing the SCSI bus during a RESELECTION phase if the target waits a minimum of a selection time-out delay and there has been no BSY signal response from the initiator:

- (1) Optionally, the target shall assert the RST signal (see 5.2.2).
- (2) Optionally, the target shall continue asserting the SEL and I/O signals and shall release all DATA BUS signals. If the target has not detected the BSY signal to be true after at least a selection abort time plus two deskew delays, the target shall release the SEL and I/O signals allowing the SCSI bus to go to the BUS FREE phase. SCSI devices that respond to the RESELECTION phase shall ensure that the reselection was still valid within a selection abort time of their assertion of the BSY signal. Failure to comply with this requirement could result in an improper reselection (two initiators connected to the same target or the wrong initiator connected to a target).

5.1.5. Information Transfer Phases

NOTE: The COMMAND, DATA, STATUS, and MESSAGE phases are all grouped together as the information transfer phases because they are all used to transfer data or control information via the DATA BUS. The actual content of the information is beyond the scope of this section.

The C/D, I/O, and MSG signals are used to distinguish between the different information transfer phases (see Table 5-1). The target drives these three signals and therefore controls all changes from one phase to another. The initiator can request a MESSAGE OUT phase by asserting the ATN signal, while the target can cause the BUS FREE phase by releasing the MSG, C/D, I/O, and BSY signals.

The information transfer phases use one or more REQ/ACK handshakes to control the information transfer. Each REQ/ACK handshake allows the transfer of one byte of information. During the information transfer phases the BSY signal shall remain true and the SEL signal shall remain false. Additionally, during the information transfer phases, the target shall continuously envelope the REQ/ACK handshake(s) with the C/D, I/O, and MSG signals in such a manner that these control signals are valid for a bus settle delay before the assertion of the REQ signal of the first handshake and remain valid until after the negation of the ACK signal at the end of the handshake of the last transfer of the phase.

IMPLEMENTORS NOTES:

- (1) After the negation of the ACK signal of the last transfer of the phase, the target may prepare for a new phase by asserting or negating the C/D, I/O, and MSG signals. These signals may be changed together or individually. They may be changed in any order and may be changed more than once. It is desirable that each line change only once. A new phase does not begin until the REQ signal is asserted for the first byte of the new phase.
- (2) A phase is defined as ending when the C/D, I/O, or MSG signals change after the negation of the ACK signal. The time between the end of a phase and the assertion of the REQ signal beginning a new phase is undefined. An initiator is allowed to anticipate a new phase based on the previous phase, the expected new phase, and early information provided by changes in the C/D, I/O, and MSG signals. However, the anticipated phase is not valid until the REQ signal is asserted at the beginning of a the next phase.

Table 5-1: Information Transfer Phases

Signal			Phase Name	Direction Of Transfer	Comment
MSG	C/D	I/O			
0	0	0	DATA OUT	Initiator to target	Data Phase
0	0	1	DATA IN	Initiator from target	
0	1	0	COMMAND	Initiator to target	
0	1	1	STATUS	Initiator from target	Message Phase
1	0	0	*		
1	0	1	*		
1	1	0	MESSAGE OUT	Initiator to target	
1	1	1	MESSAGE IN	Initiator from target	

Key: 0 = False, 1 = True, * = Reserved for future standardization.

5.1.5.1. Asynchronous Information Transfer

The target shall control the direction of information transfer by means of the I/O signal. When the I/O signal is true, information shall be transferred from the target to the initiator. When the I/O signal is false, information shall be transferred from the initiator to the target.

If the I/O signal is true (transfer to the initiator), the target shall first drive the DB(7-0,P) signals to their desired values, delay at least one deskew delay plus a cable skew delay, then assert the REQ signal. The DB(7-0,P) signals shall remain valid until the ACK signal is true at the target. The initiator shall read the DB(7-0,P) signals after the REQ signal is true, then indicate its acceptance of the data by asserting the ACK signal. When the ACK signal becomes true at the target, the target may change or release the DB(7-0,P) signals and shall negate the REQ signal. After the REQ signal is false the initiator shall then negate the ACK signal. After the ACK signal is false the target may continue the transfer by driving the DB(7-0,P) signals and asserting the REQ signal, as described above.

If the I/O signal is false (transfer to the target) the target shall request information by asserting the REQ signal. The initiator shall drive the DB(7-

0,P) signals to their desired values, delay at least one deskew delay plus a cable skew delay and assert the ACK signal. The initiator shall continue to drive the DB(7-0,P) signals until the REQ signal is false. When the ACK signal becomes true at the target, the target shall read the DB(7-0,P) signals then negate the REQ signal. When the REQ signal becomes false at the initiator, the initiator may change or release the DB(7-0,P) signals and shall negate the ACK signal. The target may continue the transfer by asserting the REQ signal, as described above.

5.1.5.2. Synchronous Data Transfer

Synchronous data transfer is optional and is only used in data phases. It shall be used in a data phase if a synchronous data transfer agreement has been established (see 5.6.21). The agreement specifies the REQ/ACK offset and the minimum transfer period.

The REQ/ACK offset specifies the maximum number of REQ pulses that can be sent by the target in advance of the number of ACK pulses received from the initiator, establishing a pacing mechanism. If the number of REQ pulses exceeds the number of ACK pulses by the REQ/ACK offset, the target shall not assert the REQ signal until after the leading edge of the next ACK pulse is received. A requirement for successful completion of the data phase is that the number of ACK and REQ pulses be equal.

The target shall assert the REQ signal for a minimum of an assertion period. The target shall then wait at least the greater of a transfer period from the last transition of the REQ signal to true or a minimum of a negation period from the last transition of the REQ signal to false before again asserting the REQ signal.

The initiator shall send one pulse on the ACK signal for each REQ pulse received. The ACK signal may be asserted as soon as the leading edge of the corresponding REQ pulse has been received. The initiator shall assert the ACK signal for a minimum of an assertion period. The initiator shall wait at least the greater of a transfer period from the last transition of the ACK signal to true or for a minimum of a negation period from the last transition of the ACK signal to false before asserting the ACK signal.

If the I/O signal is true (transfer to the initiator), the target shall first drive the DB(7-0,P) signals to their desired values, wait at least one deskew delay plus one cable skew delay, then assert the REQ signals. The DB(7-0,P) signals shall be held valid for a minimum of one deskew delay plus one cable skew delay plus one hold time after the assertion of the REQ signal. The target shall assert the REQ signal for a minimum of an assertion period. The target may then negate the REQ signal and change or release the DB(7-0,P) signals. The initiator shall read the value on the DB(7-0,P) signals within one hold time of the transition of the REQ signal to true. The initiator shall then respond with an ACK pulse.

If the I/O signal is false (transfer to the target), the initiator shall transfer one byte for each REQ pulse received. After receiving the leading edge of a REQ pulse, the initiator shall first drive the DB(7-0,P) signals to their desired values, delay at least one deskew delay plus one cable skew delay, then assert the ACK signal. The initiator shall hold the DB(7-0,P) signals valid for at least one deskew delay plus one cable skew delay plus one hold time after the assertion of the ACK signal. The initiator shall assert

the ACK signal for a minimum of an assertion period. The initiator may then negate the ACK signal and may change or release the DB(7-0,P) signals. The target shall read the value of the DB(7-0,P) signals within one hold time of the transition of the ACK signal to true.

IMPLEMENTORS NOTE: The description in SCSI-1 allowed some implementors to presume that the leading edge of the first REQ pulse beyond the REQ/ACK offset agreement would not occur until after the trailing edge of the last ACK pulse within the agreement. Devices implemented with this understanding may be subject to data destruction when in synchronous data transfer mode with devices that issue the leading edge of the next REQ pulse, at the boundary of the agreement, as soon as the leading edge of the last ACK pulse within the agreement is received. Implementors using devices of the former type in initiator designs may insure data integrity by restricting the synchronous offset agreement to values smaller than the maximum nominally offered by their device.

5.1.5.3. Wide Data Transfer

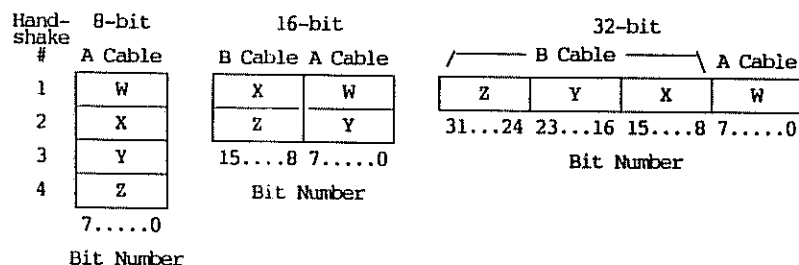
Wide data transfer is optional and may be used in the DATA phase only if a nonzero wide data transfer agreement is in effect (see WIDE DATA TRANSFER REQUEST message, 5.6.23). The messages determine the use of wide mode by both SCSI devices and establish a data path width to be used during the DATA phase.

Wide data transfers of 16- or 32-bits may be established. Although not mandatory, it is recommended that targets and initiators that support 32-bit wide transfers also support 16-bit wide transfers as well. All SCSI devices shall support 8-bit data transfers.

During 16-bit wide data transfers, the first logical data byte for each data phase shall be transferred across the DB(7-0,P) signals on the A cable and the second logical data byte shall be transferred across the DB(15-8,P1) signals on the B cable. Subsequent pairs of data bytes are likewise transferred in parallel across the A and B cables (see Figure 5-1).

During 32-bit wide data transfers, the first logical data byte for each data phase shall be transferred across the DB(7-0,P) signals on the A cable and the second, third, and fourth logical data bytes shall be transferred across the DB(15-8,P1), DB(23-16,P2), and DB(31-24,P3) signals, respectively, on the B cable. Subsequent groups of four data bytes are likewise transferred in parallel across the A and B cables (see Figure 5-1).

When transferring bytes W, X, Y and Z across the three bus widths, they are transferred as shown below:



NOTE: This figure does not represent how these bytes are stored in the initiator's memory, which may be different.

Figure 5-1: Wide SCSI Byte Ordering

If the last data byte transferred for a command does not fall on the DB(15-8,P1) signals for a 16-bit wide transfer or the DB(31-24,P3) signals for a 32-bit wide transfer, then the values of the remaining higher-numbered bits are undefined. However, parity bits for these undefined bytes shall be valid for whatever data is placed on the bus.

To ensure proper data integrity, certain sequence requirements shall be met between the REQ/ACK handshakes on the A cable and the REQB/ACKB handshakes on the B cable:

(1) The REQB and ACKB signals shall only be asserted during data phases while a nonzero wide data transfer agreement is in effect. These signals shall not be asserted during other phases.

(2) The same information transfer mode (asynchronous or synchronous) shall be used for both the A cable and the B cable. If synchronous data transfer mode is in effect, the same REQ/ACK offset and transfer period shall be used for both cables.

(3) The information transfer procedures defined in 5.1.5.1 and 5.1.5.2 for the A cable (the REQ, ACK, and DB(7-0,P) signals) shall also apply to the B cable (the REQB, ACKB, and DB(31-8,P1,P2,P3) signals). The only means available for a target to manage the timing relationship between the signals on the two cables is its management of the REQ and REQB signals. Similarly, the only means for the initiator to manage the timing between the two cables is its management of the ACK and ACKB signals.

(4) The target shall ensure that the number of REQ/ACK handshakes and the number of REQB/ACKB handshakes in a data phase are equal before it changes to another phase. The target shall not change the phase until the ACK and ACKB signals have both become false for the last REQ/ACK handshake and the last REQB/ACKB handshake.

IMPLEMENTORS NOTE: If any violations of these rules are detected by the target, the target may attempt to end the data phase and return CHECK CONDITION status. If it is impossible to correctly terminate the data phase, the target may abnormally terminate the I/O process by an unexpected disconnect. If any violations of these rules are detected by the initiator, the initiator may attempt to send an INITIATOR DETECTED ERROR message to the target. If the initiator is unable to terminate the I/O process normally, it may generate the reset condition.

5.1.6. COMMAND Phase

The COMMAND phase allows the target to request command information from the initiator.

The target shall assert the C/D signal and negate the I/O and MSG signals during the REQ/ACK handshake(s) of this phase.

5.1.7. Data Phase

The data phase is a term that encompasses both the DATA IN phase and the DATA OUT phase.

5.1.7.1. DATA IN Phase

The DATA IN phase allows the target to request that data be sent to the initiator from the target.

The target shall assert the I/O signal and negate the C/D and MSG signals during the REQ/ACK handshake(s) of this phase.

5.1.7.2. DATA OUT Phase

The DATA OUT phase allows the target to request that data be sent from the initiator to the target.

The target shall negate the C/D, I/O, and MSG signals during the REQ/ACK handshake(s) of this phase.

5.1.8. STATUS Phase

The STATUS phase allows the target to request that status information be sent from the target to the initiator.

The target shall assert the C/D and I/O signals and negate the MSG signal during the REQ/ACK handshake of this phase.

5.1.9. Message Phase

The message phase is a term that references either a MESSAGE IN, or a MESSAGE OUT phase. Multiple messages may be sent during either phase. The first byte transferred in either of these phases shall be either a single-byte message or the first byte of a multiple-byte message. Multiple-byte messages shall be wholly contained within a single message phase.

5.1.9.1. MESSAGE IN Phase

The MESSAGE IN phase allows the target to request that message(s) be sent to the initiator from the target.

The target shall assert the C/D, I/O, and MSG signals during the REQ/ACK handshake(s) of this phase.

5.1.9.2. MESSAGE OUT Phase

The MESSAGE OUT phase allows the target to request that message(s) be sent from the initiator to the target. The target invokes this phase in response to the attention condition created by the initiator (see 5.2.1).

The target shall assert the C/D and MSG signals and negate the I/O signal during the REQ/ACK handshake(s) of this phase. The target shall handshake byte(s) in this phase until the ATN signal is negated, except when rejecting a message.

If the target detects one or more parity error(s) on the message byte(s) received, it may indicate its desire to retry the message(s) by asserting the REQ signal after detecting the ATN signal has gone false and prior to changing to any other phase. The initiator, upon detecting this condition, shall re-send all of the previous message byte(s) in the same order as previously sent during this phase. When re-sending more than one message byte, the initiator shall assert the ATN signal at least two deskew delays prior to asserting the ACK signal on the first byte and shall maintain the ATN signal asserted until the last byte is sent as described in 5.2.1.

The target may act on messages as received as long as no parity error is detected and may ignore all remaining messages sent under one ATN condition after a parity error is detected. When a sequence of messages is re-sent by an initiator because of a target detected parity error, the target shall not act on any message which it acted on the first time received.

If the target receives all of the message byte(s) successfully (i.e., no parity errors), it shall indicate that it does not wish to retry by changing to any information transfer phase other than the MESSAGE OUT phase and transfer at least one byte. The target may also indicate that it has successfully received the message byte(s) by changing to the BUS FREE phase (e.g., ABORT or BUS DEVICE RESET messages).

5.1.10. Signal Restrictions Between Phases

When the SCSI bus is between two information transfer phases, the following restrictions shall apply to the SCSI bus signals:

- (1) The BSY, SEL, REQ, REQB, ACK and ACKB signals shall not change.
- (2) The C/D, I/O, MSG, and DATA BUS signals may change. When switching the DATA BUS direction from out (initiator driving) to in (target driving), the target shall delay driving the DATA BUS by at least a data release delay plus a bus settle delay after asserting the I/O signal and the initiator shall release the DATA BUS no later than a data release delay after the transition of the I/O signal to true. When switching the DATA BUS direction from in (target driving) to out (initiator driving), the target shall release the DATA BUS no later than a deskew delay after negating the I/O signal.
- (3) The ATN and RST signals may change as defined under the descriptions for the attention condition (see 5.2.1) and reset condition (see 5.2.2).

5.2. SCSI Bus Conditions

The SCSI bus has two asynchronous conditions; the attention condition and the reset condition. These conditions cause the SCSI device to perform certain actions and can alter the phase sequence.

Furthermore, an SCSI device may not all be powered on at the same time. This standard does not address power sequencing issues. However, each SCSI device, as it is powered on, should perform appropriate internal reset operations and internal test operations. It is recommended that following a power-on to selection time after power is applied, SCSI targets be able to respond with appropriate status and sense data to the TEST UNIT READY, INQUIRY, and REQUEST SENSE commands.

5.2.1. Attention Condition

The attention condition allows an initiator to inform a target that the initiator has a message ready. The target may get this message by performing a MESSAGE OUT phase.

The initiator creates the attention condition by asserting ATN at any time except during the ARBITRATION or BUS FREE phases.

The initiator shall assert the ATN signal at least two deskew delays before negating the ACK signal for the last byte transferred in a bus phase for the attention condition to be honored before transition to a new bus phase. Asserting the ATN signal later might not be honored until a later bus phase and then may not result in the expected action. The initiator shall negate the ATN signal at least two deskew delays before asserting the ACK signal while transferring the last byte of the messages indicated with a "Yes" in Table 5-2. If the target detects that the initiator failed to meet this requirement, then the target shall go to BUS FREE phase (see unexpected disconnect, 5.1.1).

A target shall respond with MESSAGE OUT phase as follows:

(1) If the ATN signal becomes true during a COMMAND phase, the target shall enter MESSAGE OUT phase after transferring part or all of the command descriptor block bytes.

(2) If the ATN signal becomes true during a DATA phase, the target shall enter MESSAGE OUT phase at the target's earliest convenience (often, but not necessarily on a logical block boundary). The initiator shall continue REQ/ACK handshakes until it detects the phase change.

(3) If the ATN signal becomes true during a STATUS phase, the target shall enter MESSAGE OUT phase after the status byte has been acknowledged by the initiator.

(4) If the ATN signal becomes true during a MESSAGE IN phase, the target shall enter MESSAGE OUT phase before it sends another message. This permits a MESSAGE PARITY ERROR message from the initiator to be associated with the appropriate message.

(5) If the ATN signal becomes true during a SELECTION phase and before the initiator releases the BSY signal, the target shall enter MESSAGE OUT phase immediately after that SELECTION phase.

(6) If the ATN signal becomes true during a RESELECTION phase, the target shall enter MESSAGE OUT phase after the target has sent its IDENTIFY message for that RESELECTION phase.

The initiator shall keep the ATN signal asserted if more than one byte is to be transferred. The initiator may negate the ATN signal at any time except it shall not negate the ATN signal while the ACK signal is asserted during a MESSAGE OUT phase. Normally, the initiator negates the ATN signal while the REQ signal is true and the ACK signal is false during the last REQ/ACK handshake of the MESSAGE OUT phase.

5.2.2. Reset Condition

The reset condition is used to immediately clear all SCSI devices from the bus. This condition shall take precedence over all other phases and conditions. Any SCSI device may create the reset condition by asserting the RST signal for a minimum of a reset hold time. During the reset condition, the state of all SCSI bus signals other than the RST signal is not defined.

All SCSI devices shall release all SCSI bus signals (except the RST signal) within a bus clear delay of the transition of the RST signal to true. The BUS FREE phase always follows the reset condition.

The effect of the reset condition on I/O processes which have not completed, SCSI device reservations, and SCSI device operating modes is determined by whether the SCSI device has implemented the hard reset alternative or the soft reset alternative (one of which shall be implemented) as defined in 5.2.2.1 and 5.2.2.2. The hard and soft reset alternatives are mutually exclusive within a system. A facility for targets to report which reset alternative is implemented is provided in the SftRe bit of the INQUIRY data (7.2.5).

IMPLEMENTORS NOTE: Environmental conditions (e.g., static discharge) may generate brief glitches on the RST signal. It is recommended that SCSI devices not react to these glitches. The manner of rejecting glitches is vendor specific. The bus clear delay following a RST signal transition to true is measured from the original transition of the RST signal, not from the time that the signal has been confirmed. This limits the time to confirm the RST signal to a maximum of a bus clear delay.

5.2.2.1. Hard Reset Alternative

SCSI devices that implement the hard reset alternative, upon detection of the reset condition, shall:

- (1) Clear all I/O processes including queued I/O processes.
- (2) Release all SCSI device reservations.
- (3) Return any SCSI device operating modes to their appropriate initial conditions, similar to those conditions that would be found after a normal power-on reset. MODE SELECT conditions shall be restored to their last saved values if saved values have been established. MODE SELECT conditions for which no values have been saved shall be returned to their default values.
- (4) Unit attention condition shall be set (See 6.9).

It is recommended that following a reset to selection time after a hard reset condition ends, SCSI targets be able to respond with appropriate status and sense data to the TEST UNIT READY, INQUIRY, and REQUEST SENSE commands.

5.2.2.2. Soft Reset Alternative

SCSI devices that implement the soft reset alternative, upon detection of the reset condition, shall:

- (1) Attempt to complete any I/O processes which have not completed and that were fully identified
- (2) Preserve all SCSI device reservations
- (3) Preserve any SCSI device operating modes (MODE SELECT, PREVENT/ALLOW MEDIUM REMOVAL commands, etc.)
- (4) Preserve all the information required to continue normal dispatching of I/O processes queued prior to the reset condition.

The soft reset alternative allows an initiator to reset the SCSI bus with minimum disruption to the operation of other initiators in a multiple initiator system. To ensure proper operation the following conditions shall be met:

- (1) An initiator shall not consider an I/O process to be fully identified until the IDENTIFY message (and queue tag message, if any) is sent to the target and the target responds by changing to any other information transfer phase and requests that at least one byte be transferred.
- (2) A target shall consider an I/O process to be fully identified when it successfully receives the IDENTIFY message and any queue tag message and the initiator negates the ATN signal.
- (3) If an initiator selects a logical unit for which there already is an active I/O process with the same queue tag (if any) for the same initiator, the target shall clear the original I/O process and perform the new I/O process.
- (4) If a target reselects an initiator to continue an I/O process for which the initiator has no record, the initiator shall abort that I/O process by sending the ABORT or ABORT TAG message, depending on whether the reselecting I/O process is a tagged I/O process.
- (5) An initiator shall consider an I/O process to be completed when it negates ACK for a successfully received COMMAND COMPLETE message.
- (6) A target shall consider an I/O process to be completed when it detects the transition of ACK to false for the COMMAND COMPLETE message with the ATN signal false.
- (7) An initiator shall not negate the ACK signal for the SAVE DATA POINTER message until it has actually saved the data pointer for the I/O process.

(8) A target shall consider the data pointer to be saved when it detects the transition of the ACK signal to false for the SAVE DATA POINTER message with the ATN signal false.

(9) If the reset condition occurs between the time that the target asserts the REQ signal for the SAVE DATA POINTER message and it detects the transition of the ACK signal to false, the target shall terminate the I/O process with CHECK CONDITION status. The target shall set the sense key to ABORTED COMMAND. This is necessary because the target cannot determine whether the data pointer has actually been saved.

NOTE: If the ATN signal is true in conditions (6) or (8), the target would normally switch to MESSAGE OUT phase and attempt to transfer a message byte. If the reset condition occurs before the target successfully receives the message byte, it may assume that the initiator has not successfully received the COMMAND COMPLETE message or the SAVE DATA POINTER message. In the case of COMMAND COMPLETE message, the target may reselect the initiator and attempt to send the COMMAND COMPLETE message again. In the case of the SAVE DATA POINTER message, the target may reselect the initiator and terminate the I/O process as described in condition (9).

5.3. SCSI Bus Phase Sequences

The order in which phases are used on the SCSI bus follows a prescribed sequence.

The reset condition can abort any phase and is always followed by the BUS FREE phase. Also any other phase can be followed by the BUS FREE phase but many such instances are error conditions (see unexpected disconnect, 5.1.1).

The additional allowable sequences shall be as shown in Figure 5-2. The normal progression is from the BUS FREE phase to ARBITRATION, from ARBITRATION to SELECTION or RESELECTION, and from SELECTION or RESELECTION to one or more of the information transfer phases (COMMAND, DATA, STATUS, or MESSAGE). The final information transfer phase is normally the MESSAGE IN phase where a DISCONNECT, or COMMAND COMPLETE message is transferred, followed by the BUS FREE phase.

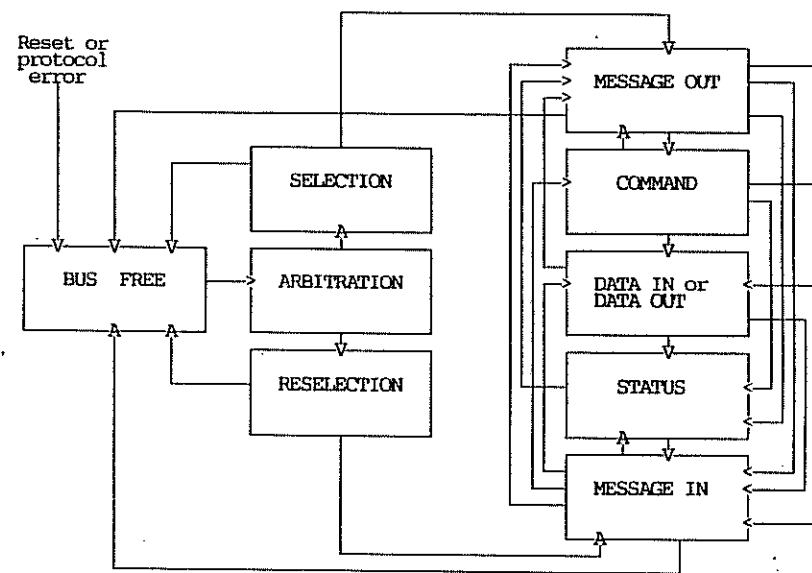


Figure 5-2: Phase Sequences

5.4. SCSI Pointers

Consider the system shown in Figure 5-3 in which an initiator and target communicate on the SCSI bus in order to execute an I/O process.

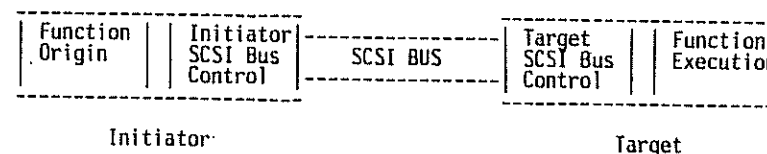


Figure 5-3: Simplified SCSI System

The SCSI architecture provides for a set of three pointers for each I/O process, called the saved pointers. The set of three pointers consist of one for the command, one for the data, and one for the status. When an I/O process becomes active, its three saved pointers are copied into the initiator's set of three active pointers. There is only one set of active pointers in each initiator. The active pointers point to the next command, data, or status byte to be transferred between the initiator's memory and the target. The saved and active pointers reside in the initiator.

The saved command pointer always points to the start of the command descriptor block (see 6.2) for the I/O process. The saved status pointer always points to the start of the status area for the I/O process. The saved data pointer points to the start of the data area until the target sends a SAVE DATA POINTER message (see 5.6.20) for the I/O process.

In response to the SAVE DATA POINTER message, the initiator stores the value of the active data pointer into the saved data pointer for that I/O process. The target may restore the active pointers to the saved pointer values for the active I/O process by sending a RESTORE POINTERS message (see 5.6.19) to the initiator. The initiator then copies the set of saved pointers into the set of active pointers. Whenever a target disconnects from the bus, only the set of saved pointers are retained. The set of active pointers is restored from the set of saved pointers upon reconnection of the I/O process.

IMPLEMENTORS NOTE: Since the data pointer value may be modified by the target before the I/O process ends, it should not be used to test for actual transfer length.

5.5. Message System Description

The message system allows communication between an initiator and target for the purpose of interface management. A message may be one, two, or multiple bytes in length. One or more messages may be sent during a single MESSAGE phase, but a message may not be split over MESSAGE phases. The initiator is required to end the MESSAGE OUT phase (by negating ATN) when it sends certain messages identified in Table 5-2.

One-byte, Two-byte, and extended message formats are defined. The first byte of the message determines the format as follows:

Value	Message Format
00h	One-Byte Message (COMMAND COMPLETE)
01h	Extended Messages
02h - 1Fh	One-Byte Messages
20h - 2Fh	Two-Byte Messages
30h - 7Fh	Reserved
80h - FFh	One-Byte Message (IDENTIFY)

One-byte messages consist of a single byte transferred during a MESSAGE phase. The value of the byte determines which message is to be performed as defined in Table 5-2.

Table 5-2: Message Codes

Code	Support	Message Name	Direction	Negate ATN
Init	Targ			Before last ACK
06h	O	M ABORT	Out	Yes
0Dh	O	O ABORT TAG (Note 1)	Out	Yes
0Ch	O	M BUS DEVICE RESET	Out	Yes
0Eh	O	O CLEAR QUEUE (Note 1)	Out	Yes
00h	M	M COMMAND COMPLETE	In	---
04h	O	O DISCONNECT	In	---
04h	O	O DISCONNECT	Out	Yes
80h+	M	O IDENTIFY	In	---
80h+	M	M IDENTIFY	Out	No
23h	O	O IGNORE WIDE RESIDUE (Two Bytes)	In	---
0Fh	O	O INITIATE RECOVERY	In	---
0Fh	O	O INITIATE RECOVERY (Note 2)	Out	Yes
05h	M	M INITIATOR DETECTED ERROR	Out	Yes
0Ah	O	O LINKED COMMAND COMPLETE	In	---
0Bh	O	O LINKED COMMAND COMPLETE (WITH FLAG)	In	---
09h	M	M MESSAGE PARITY ERROR	Out	Yes
07h	M	M MESSAGE REJECT	In	Out
***	O	O MODIFY DATA POINTER	In	---
08h	M	M NO OPERATION	Out	Yes
Queue Tag Messages (Two Bytes)				
21h	O	O HEAD OF QUEUE TAG	Out	No
22h	O	O ORDERED QUEUE TAG	Out	No
20h	O	O SIMPLE QUEUE TAG	In	Out
10h	O	O RELEASE RECOVERY	Out	Yes
03h	O	O RESTORE POINTERS	In	---
02h	O	O SAVE DATA POINTER	In	---
***	O	O SYNCHRONOUS DATA TRANSFER REQUEST	In	Out
***	O	O WIDE DATA TRANSFER REQUEST	In	Out
11h	O	O TERMINATE I/O PROCESS	Out	Yes
12h - 1Fh		Reserved		
24h - 2Fh		Reserved for two-byte messages		
30h - 7Fh		Reserved		

Key: M = Mandatory support, O = Optional support.
 In = Target to initiator, Out = Initiator to target.
 Yes = Initiator shall negate ATN before last ACK of message.
 No = Initiator may or may not negate ACK before last ACK of message. (see attention condition, 5.2.1.)
 --- = Not Applicable
 *** = Extended message (see Tables 5-3 and 5-4)
 80h+ = Codes 80h through FFh are used for IDENTIFY messages (see Table 5-7).

NOTES:

- (1) The ABORT TAG and CLEAR QUEUE messages are required if tagged queuing is implemented.
- (2) Outbound INITIATE RECOVERY messages are only valid during the asynchronous event notification protocol.

Two-byte messages consist of two consecutive bytes transferred during a MESSAGE phase. The value of the first byte determines which message is to be performed as defined in Table 5-2. The second byte is a parameter byte which is used as defined in the message description (see 5.6).

A value of one in the first byte of a message indicates the beginning of a multiple-byte extended message. The minimum number of bytes sent for an extended message is three. The extended message format and the extended message codes are shown in Tables 5-3 and 5-4, respectively.

Table 5-3: Extended Message Format

Byte	Value	Description
0	01h	Extended message
1	n	Extended message length
2	y	Extended message code
3 - n+1	x	Extended message arguments

The extended message length specifies the length in bytes of the extended message code plus the extended message arguments to follow. Therefore, the total length of the message is equal to the extended message length plus two. A value of zero for the extended message length indicates 256 bytes follow.

The extended message codes are listed in Table 5-4. The extended message arguments are specified within the extended message descriptions.

Table 5-4: Extended Message Codes

Code (y)	Description
02h	Reserved (See Note)
00h	MODIFY DATA POINTER
01h	SYNCHRONOUS DATA TRANSFER REQUEST
03h	WIDE DATA TRANSFER REQUEST
04h - 7Fh	Reserved
80h - FFh	Vendor Unique

NOTE: Extended message code 02h was used for the EXTENDED IDENTIFY message in SCSI-1.

The first message sent by the initiator after the SELECTION phase shall be an IDENTIFY, ABORT, or BUS DEVICE RESET message. If a target receives any other message it shall go to BUS FREE phase (see unexpected disconnect, 5.1.1).

If the first message is an IDENTIFY message, then it may be immediately followed by other messages, such as the first of a pair of SYNCHRONOUS DATA TRANSFER REQUEST messages. If tagged queuing is used the queue tag message immediately follows the IDENTIFY message (see 5.6.7). The IDENTIFY message establishes a logical connection between the initiator and the specified logical unit or target routine within the target known as an I T L nexus or I T R nexus. After the RESELECTION phase, the target's first message shall be IDENTIFY. This allows the I T L nexus or I T R nexus to be re-established. Only one logical unit or target routine shall be identified for any connection; if a target receives a second IDENTIFY message with a different logical unit number or target routine number during a connection, it shall go to BUS FREE phase (see unexpected disconnect, 5.1.1). The treatment of other logical unit addressing errors is described in 6.5.

All initiators shall implement the mandatory messages tabulated in the "Init" column of Table 5-2. All targets shall implement the mandatory messages tabulated in the "Targ" column of Table 5-2.

Whenever an I T L nexus or I T R nexus is established by an initiator that is allowing disconnection, the initiator shall ensure that the active pointers are equal to the saved pointers for that particular logical unit or target routine. An implied restore pointers operation shall occur as a result of a reconnection.

5.6. Messages

The SCSI messages are defined in this section.

5.6.1. ABORT

The ABORT message is sent from the initiator to the target to clear the present I/O process plus any queued I/O process for the I T x nexus. The target shall go to the BUS FREE phase following successful receipt of this message. Pending data, status, and queued I/O processes for any other I T x nexus shall not be cleared.

If only an I T nexus has been established, the target shall go to the BUS FREE phase. No status or message shall be sent for the current I/O process and no other I/O process shall be affected.

It is not an error to issue this message to an I T x nexus that does not have an active or queued I/O process.

Previously established conditions, including MODE SELECT parameters, reservations, and extended contingent allegiance shall not be changed by the ABORT message.

IMPLEMENTORS NOTES: The BUS DEVICE RESET, CLEAR QUEUE, ABORT, and ABORT TAG messages provide a means to clear one or more I/O processes prior to normal termination. The BUS DEVICE RESET message clears all I/O processes for all initiators on all logical units and all target routines of the target. The CLEAR QUEUE message clears all I/O processes for all initiators on the specified logical unit or target routine of the target. The ABORT message clears all I/O processes for the selecting initiator on the specified logical unit or target routine of the target. The ABORT TAG message clears

the ~~active~~ ^{current} I/O process only.

5.6.2. ABORT TAG

The ABORT TAG message shall be implemented if tagged queuing is implemented. The target shall go to the BUS FREE phase following successful receipt of this message. The target shall clear the current I/O process. If the target has already started execution of the I/O process, the execution shall be halted. The medium contents may have been modified before the execution was halted. In either case, any pending status or data for the I/O process shall be cleared and no status or ending message shall be sent to the initiator. Pending status, data, and commands for other active or queued I/O processes shall not be affected. Execution of other I/O processes queued for the I_T_x nexus shall not be aborted.

Previously established conditions, including MODE SELECT parameters, reservations, and extended contingent allegiance shall not be changed by the ABORT TAG message.

5.6.3. BUS DEVICE RESET

The BUS DEVICE RESET message is sent from an initiator to direct a target to clear all I/O processes on that SCSI device. This message forces a hard reset condition to the selected SCSI device. The target shall go to the BUS FREE phase following successful receipt of this message. The target shall create a unit attention condition for all initiators (see 6.9).

5.6.4. CLEAR QUEUE

The CLEAR QUEUE message shall be implemented if tagged queuing is implemented and may be implemented if untagged queuing is implemented. The target shall go to the BUS FREE phase following successful receipt of this message. The target shall perform an action equivalent to receiving a series of ABORT messages from each initiator. All I/O processes, from all initiators, in the queue for the specified logical unit or target routine shall be cleared from the queue. All active I/O processes shall be terminated. The medium may have been altered by partially executed commands. All pending status and data for that logical unit or target routine for all initiators shall be cleared. No status or message shall be sent for any of the I/O processes. A unit attention condition shall be generated for all other initiators with I/O processes that either were active or were queued for that logical unit or target routine. When reporting the unit attention condition the additional sense code shall be set to COMMANDS CLEARED BY ANOTHER INITIATOR.

Previously established conditions, including MODE SELECT parameters, reservations, and extended contingent allegiance shall not be changed by the CLEAR QUEUE message.

5.6.5. COMMAND COMPLETE

The COMMAND COMPLETE message is sent from a target to an initiator to indicate that the execution of an I/O process has completed and that valid status has been sent to the initiator. After successfully sending this message, the target shall go to the BUS FREE phase by releasing the BSY

signal. The target shall consider the message transmission to be successful when it detects the negation of ACK for the COMMAND COMPLETE message with the ATN signal false.

IMPLEMENTORS NOTE: The I/O process may have completed successfully or unsuccessfully as indicated in the status.

5.6.6. DISCONNECT

The DISCONNECT message is sent from a target to inform an initiator that the present connection is going to be broken (the target plans to disconnect by releasing the BSY signal), but that a later reconnect will be required in order to complete the current I/O process. This message shall not cause the initiator to save the data pointer. After successfully sending this message, the target shall go to the BUS FREE phase by releasing the BSY signal. The target shall consider the message transmission to be successful when it detects the negation of the ACK signal for the DISCONNECT message with the ATN signal false.

Targets which break data transfers into multiple connections shall end each successful connection (except possibly the last) with a SAVE DATA POINTER - DISCONNECT message sequence.

This message may also be sent from an initiator to a target to instruct the target to disconnect from the SCSI bus. If this option is supported, and after the DISCONNECT message is received, the target shall switch to MESSAGE IN phase, send the DISCONNECT message to the initiator (possibly preceded by SAVE DATA POINTER message), and then disconnect by releasing BSY. After releasing the BSY signal, the target shall not participate in another ARBITRATION phase for at least a disconnection delay. If this option is not supported or the target cannot disconnect at the time when it receives the DISCONNECT message from the initiator, the target shall respond by sending a MESSAGE REJECT message to the initiator.

5.6.7. IDENTIFY

The IDENTIFY message (Table 5-5) is sent by either the initiator or the target to establish an I_T_L nexus or an I_T_R nexus.

IMPLEMENTORS NOTE: Use of the IDENTIFY message to establish an I_T_R nexus allows connection to one of up to eight target routines or functions in the target itself. These target routines are expected to be used for maintenance and diagnostic purposes.

Table 5-5: IDENTIFY Message Format

Bit	7	6	5	4	3	2	1	0
Byte								
0	Identify	DiscPriv	LUNAR	Reserved	Reserved		LUNTRN	

The identify bit shall be set to one to specify that this is an IDENTIFY message.

A disconnect privilege (DiscPriv) bit of one specifies that the initiator has granted the target the privilege of disconnecting. A DiscPriv bit of zero specifies that the target shall not disconnect. This bit is not defined and shall be set to zero when an IDENTIFY message is sent by a target.

A logical unit target (LUNTAR) bit of zero specifies that the I/O process is directed to or from a logical unit. A LUNTAR bit of one specifies that the I/O process is directed to or from a target routine.

The logical unit number target routine number (LUNTRN) field specifies a logical unit number if the LUNTAR bit is zero. The LUNTRN field specifies a target routine number if the LUNTAR bit is one. The response to an invalid value in the LUNTRN field is described in 6.5.3. Only the INQUIRY and REQUEST SENSE commands are valid for target routines. If a target receives any other command for a target routine, it shall return CHECK CONDITION status and shall set the sense key to ILLEGAL REQUEST.

An IDENTIFY message is invalid if a reserved bit is set to one or if the LUNTAR bit is set to one and the target does not implement target routines. A device may respond to an invalid IDENTIFY message by immediately sending a MESSAGE REJECT message or by returning CHECK CONDITION status. If a CHECK CONDITION status is returned, the sense key shall be set to ILLEGAL REQUEST and the additional sense code shall be set to INVALID BITS IN IDENTIFY MESSAGE FIELD.

Only one logical unit number or target routine number shall be identified per I/O process. The initiator may send one or more IDENTIFY messages during a connection. A second IDENTIFY message with a different value in either the LUNTAR bit or LUNTRN field shall not be issued before a BUS FREE phase has occurred; if a target receives a second IDENTIFY message with a different value in either of these fields, it shall go to BUS FREE phase (see unexpected disconnect, 5.1.1). Thus an initiator may change the DiscPriv bit, but may not attempt to switch to another I/O process. (See the DTDC field of the disconnect-reconnect page (7.3.3.2) for additional controls over disconnection.)

An implied RESTORE POINTERS message shall be performed by the initiator prior to the assertion of the ACK signal on the next phase for an inbound IDENTIFY message sent during reconnection.

5.6.8. IGNORE WIDE RESIDUE

Table 5-6: IGNORE WIDE RESIDUE Message Format

Bit	7	6	5	4	3	2	1	0
Byte	Message Code (23h)							
0	Message Code (23h)							
1	Ignore							

The IGNORE WIDE RESIDUE message (Table 5-6) shall be sent from a target to indicate that the number of valid bytes sent during the last REQ/ACK handshake and REQB/ACKB handshake of a DATA IN phase is less than the negotiated transfer width. The ignore field indicates the number of invalid data bytes transferred. This message shall be sent immediately following the DATA IN phase and prior to any other messages. The ignore field is defined as follows:

Ignore	Invalid Data Bits	
	32-bit Transfers	16-bit Transfers
00h	Reserved	Reserved
01h	DB(31-24)	DB(15-8)
02h	DB(31-16)	Reserved
03h	DB(31-8)	Reserved
04h to FFh	Reserved	Reserved

Even though a byte is invalid its corresponding parity bit shall be valid for the value transferred. For 16-bit transfers, DB(31-16) are always invalid and the corresponding parity bits are also invalid.

5.6.9. INITIATE RECOVERY

A target that supports extended contingent allegiance shall inform the initiator it is entering this condition by sending an INITIATE RECOVERY message immediately following a CHECK CONDITION or COMMAND TERMINATED status. The extended contingent allegiance condition remains in effect until terminated as described in 6.7.

If an asynchronous event occurs, the target may enter an extended contingent allegiance condition by becoming a temporary initiator and sending the INITIATE RECOVERY message following the IDENTIFY message and any queue tag message and before the COMMAND phase of the SEND command that is used to perform the asynchronous event notification (see 6.5.5). The successful transmission of this message establishes the extended contingent allegiance condition which remains in effect until terminated as described in 6.7.

IMPLEMENTORS NOTE: If the target notifies multiple initiators of the asynchronous event, it should include the INITIATE RECOVERY message in only one of the notifications.

A MESSAGE REJECT response to an INITIATE RECOVERY message indicates that an extended contingent allegiance condition shall not be established. The enabled or disabled state of an extended contingent allegiance (see the DQue bit of the control mode page, 7.3.3.1) is not changed by the rejection of an INITIATE RECOVERY message.

5.6.10. INITIATOR DETECTED ERROR

The INITIATOR DETECTED ERROR message is sent from an initiator to inform a target that an error has occurred that does not preclude the target from retrying the operation. The source of the error may either be related to previous activities on the SCSI bus or may be internal to the initiator and unrelated to any previous SCSI bus activity. Although present pointer

integrity is not assured, a RESTORE POINTERS message or a disconnect followed by a reconnect, shall cause the pointers to be restored to their defined prior state.

5.6.11. LINKED COMMAND COMPLETE

The LINKED COMMAND COMPLETE message is sent from a target to an initiator to indicate that the execution of a linked command has completed and that status has been sent. The initiator shall then set the pointers to the initial state for the next linked command.

5.6.12. LINKED COMMAND COMPLETE (WITH FLAG)

The LINKED COMMAND COMPLETE (WITH FLAG) message is sent from a target to an initiator to indicate that the execution of a linked command (with the flag bit set to one) has completed and that status has been sent. The initiator shall then set the pointers to the initial state of the next linked command. Typically this message would be used to cause an interrupt in the initiator between two linked commands.

5.6.13. MESSAGE PARITY ERROR

The MESSAGE PARITY ERROR message is sent from the initiator to the target to indicate that the last message byte it received had a parity error.

In order to indicate its intentions of sending this message, the initiator shall assert the ATN signal prior to its release of the ACK signal for the REQ/ACK handshake of the message byte that has the parity error. This provides an interlock so that the target can determine which message byte has the parity error. If the target receives this message under any other circumstance, it shall signal a catastrophic error condition by releasing the BSY signal without any further information transfer attempt (see 5.1.1).

If the target returns to the MESSAGE IN phase before switching to some other phase, after receiving the MESSAGE PARITY ERROR message, the target shall re-send the entire message that had the parity error.

5.6.14. MESSAGE REJECT

The MESSAGE REJECT message is sent from either the initiator or target to indicate that the last message or message byte it received was inappropriate or has not been implemented.

In order to indicate its intentions of sending this message, the initiator shall assert the ATN signal prior to its release of the ACK signal for the REQ/ACK handshake of the message byte that is to be rejected. If the target receives this message under any other circumstance, it shall reject this message.

When a target sends this message, it shall change to MESSAGE IN phase and send this message prior to requesting additional message bytes from the initiator. This provides an interlock so that the initiator can determine which message byte is rejected.

After a target sends a MESSAGE REJECT message and if the ATN signal is still

asserted, then it shall return to the MESSAGE OUT phase. The subsequent MESSAGE OUT phase shall begin with the first byte of a message.

5.6.15. MODIFY DATA POINTER Message

Table 5-7: MODIFY DATA POINTER

Byte	Value	Description
0	01h	Extended message
1	05h	Extended message length
2	00h	MODIFY DATA POINTER code
3	x	Argument (Most Significant Byte)
4	x	Argument
5	x	Argument
6	x	Argument (Least Significant Byte)

The MODIFY DATA POINTER message (Table 5-7) is sent from the target to the initiator and requests that the signed argument be added (two's complement) to the value of the current data pointer.

5.6.16. NO OPERATION

The NO OPERATION message is sent from an initiator in response to a target's request for a message when the initiator does not currently have any other valid message to send.

For example, if the target does not respond to the attention condition until a later phase and at that time the original message is no longer valid the initiator may send the NO OPERATION message when the target enters the MESSAGE OUT phase.

5.6.17. Queue Tag Messages

Table 5-8: Queue Tag Message Format

Bit	7	6	5	4	3	2	1	0
Byte								
0	Message Code (20h or 21h or 22h)							
1	Queue Tag							

Table 5-8 defines the format for the queue tag messages. If the target implements tagged queuing, all of the queue tag messages are mandatory: HEAD OF QUEUE TAG, ORDERED QUEUE TAG, and SIMPLE QUEUE TAG. Tagged queuing is only defined for logical units, not target routines.

If a target does not implement tagged queuing and a queue tag message is received or if a queue tag message is received for a target routine, it shall respond with a MESSAGE REJECT message and accept the I/O process as if it were untagged.

The queue tag messages are used to specify an identifier, called a queue tag, for an I/O process which establishes the I T L Q nexus. The queue tag field is an 8-bit unsigned integer assigned by the Initiator during an initial connection. The queue tag for every I/O process for each I T L nexus should be unique. If the target receives a queue tag that is currently in use for the I T L nexus, then it shall respond as defined in 6.5.2. A queue tag becomes available for re-assignment when the I/O process ends. The numeric value of a queue tag has no effect on the order of execution.

IMPLEMENTORS NOTE: For each logical unit on each target, each initiator has up to 256 queue tags to assign to I/O processes. Thus a target with eight logical units could have up to 14336 I/O processes concurrently in existence if there were seven initiators on the bus.

Whenever an initiator connects to a target, the appropriate queue tag message shall be sent immediately following the IDENTIFY message and within the same MESSAGE OUT phase to establish the I T L Q nexus for the I/O process. Only one I T L Q nexus may be established during a connection. If a queue tag message is not sent, then only an I_T_x nexus is established for the I/O process (untagged command).

Whenever a target reconnects to an initiator to continue a tagged I/O process, the SIMPLE QUEUE TAG message shall be sent immediately following the IDENTIFY message and within the same MESSAGE IN phase to revive the I T L Q nexus for the I/O process. Only one I T L Q nexus may be revived during a reconnection. If the SIMPLE QUEUE TAG message is not sent, then only an I_T_x nexus is revived for the I/O process (untagged command).

If a target attempts to reconnect using an invalid queue tag, then the initiator should respond with an ABORT TAG message.

5.6.17.1. HEAD OF QUEUE TAG

The HEAD OF QUEUE TAG message specifies that the I/O process be placed first in that logical unit's command queue. An I/O process already being executed by the target shall not be pre-empted. A subsequent I/O process received with a HEAD OF QUEUE TAG message shall be placed at the head of the command queue for execution in last-in, first-out order.

5.6.17.2. ORDERED QUEUE TAG

The ORDERED QUEUE TAG message specifies that the I/O process be placed in that logical unit's command queue for execution in the order received. All queued I/O processes for the logical unit received prior to this I/O process shall be executed before this I/O process is executed. All queued I/O processes received after this I/O process shall be executed after this I/O process, except for I/O processes received with a HEAD OF QUEUE TAG message.

5.6.17.3. SIMPLE QUEUE TAG

The SIMPLE QUEUE TAG message specifies that the I/O process be placed in that logical unit's command queue. The order of execution is described in 6.8.

5.6.18. RELEASE RECOVERY

The RELEASE RECOVERY message is sent from an initiator to a target to terminate an extended contingent allegiance condition previously established by an INITIATE RECOVERY message. This message shall be sent immediately following the IDENTIFY message in the same MESSAGE OUT phase. The extended contingent allegiance condition ends upon successful receipt of the RELEASE RECOVERY message. The target shall go to the BUS FREE phase following successful receipt of this message.

If a RELEASE RECOVERY message is received by a target that implements extended contingent allegiance when no extended contingent allegiance condition is active, the message shall not be rejected and the target shall go to the BUS FREE phase.

5.6.19. RESTORE POINTERS

The RESTORE POINTERS message is sent from a target to direct the initiator to copy the most recently saved command, data, and status pointers for the I/O process to the corresponding active pointers. The command and status pointers shall be restored to the beginning of the present command and status areas. The data pointer shall be restored to the value at the beginning of the data area in the absence of a SAVE DATA POINTER message or to the value at the point at which the last SAVE DATA POINTER message occurred for that nexus.

5.6.20. SAVE DATA POINTER

The SAVE DATA POINTER message is sent from a target to direct the initiator to copy the active data pointer to the saved data pointer for the current I/O process. (See 5.4 for a definition of pointers.)

5.6.21. SYNCHRONOUS DATA TRANSFER REQUEST Message

Table 5-9: SYNCHRONOUS DATA TRANSFER REQUEST

Byte	Value	Description
0	01h	Extended message
1	03h	Extended message length
2	01h	SYNCHRONOUS DATA TRANSFER REQUEST code
3	m	Transfer period (m times 4 nanoseconds)
4	x	REQ/ACK offset

A SYNCHRONOUS DATA TRANSFER REQUEST (SDTR) message (Table 5-9) exchange shall be initiated by an SCSI device whenever a previously-arranged data transfer agreement may have become invalid. The agreement becomes invalid after any condition which may leave the data transfer agreement in an indeterminate state such as:

- (1) after a hard reset condition
- (2) after a BUS DEVICE RESET message and
- (3) after a power cycle.

In addition, an SCSI device may initiate an SDTR message exchange whenever it is appropriate to negotiate a new data transfer agreement (either synchronous or asynchronous). SCSI devices that are capable of synchronous data transfers shall not respond to an SDTR message with a MESSAGE REJECT message.

IMPLEMENTORS NOTES:

- (1) Re-negotiation at every selection is not recommended, since a significant performance impact is likely.
- (2) Due to historical problems with early host adapters that could not accept an SDTR message, some targets may not initiate synchronous negotiation after a power cycle as required by this standard. Host adapters that support synchronous mode may avoid the ensuing failure modes when the target is independently power cycled by initiating a synchronous negotiation on each REQUEST SENSE and INQUIRY command.

The SDTR message exchange establishes the permissible transfer periods and the REQ/ACK offsets for all logical units and target routines on the two devices. This agreement only applies to data phases.

The transfer period is the minimum time allowed between leading edges of successive REQ pulses and of successive ACK pulses to meet the device requirements for successful reception of data.

The REQ/ACK offset is the maximum number of REQ pulses allowed to be outstanding before the leading edge of its corresponding ACK pulse is received at the target. This value is chosen to prevent overflow conditions in the device's reception buffer and offset counter. A REQ/ACK offset value of zero

shall indicate asynchronous data transfer mode; a value of FFh shall indicate unlimited REQ/ACK offset.

The originating device (the device that sends the first of the pair of SDTR messages) sets its values according to the rules above to permit it to receive data successfully. If the responding device can also receive data successfully with these values (or smaller transfer periods or larger REQ/ACK offsets or both), it returns the same values in its SDTR message. If it requires a larger transfer period, a smaller REQ/ACK offset, or both in order to receive data successfully, it substitutes values in its SDTR message as required, returning unchanged any value not required to be changed. Each device when transmitting data shall respect the limits set by the other's SDTR message, but it is permitted to transfer data with larger transfer periods, smaller REQ/ACK offsets, or both than specified in the other's SDTR message. The successful completion of an exchange of SDTR messages implies an agreement as follows:

Responding Device SDTR response	Implied Agreement
(1) Non-zero REQ/ACK offset	Each device transmits data with a transfer period equal to or greater than and a REQ/ACK offset equal to or less than the values received in the other device's SDTR message.
(2) REQ/ACK offset equal to zero	Asynchronous transfer
(3) MESSAGE REJECT message	Asynchronous transfer

If the initiator recognizes that negotiation is required, it asserts the ATN signal and sends a SDTR message to begin the negotiating process. After successfully completing the MESSAGE OUT phase, the target shall respond with the proper SDTR message. If an abnormal condition prevents the target from returning an appropriate response, both devices shall go to asynchronous data transfer mode for data transfers between the two devices.

Following target response (1) above, the implied agreement for synchronous operation shall be considered to be negated by both the initiator and the target if the initiator asserts the ATN signal and the first message out is either MESSAGE PARITY ERROR or MESSAGE REJECT. In this case, both devices shall go to asynchronous data transfer mode for data transfers between the two devices. For the MESSAGE PARITY ERROR case, the implied agreement shall be reinstated if a re-transmission of the second of the pair of messages is successfully accomplished. After a vendor-specific number of retry attempts (greater than zero), if the target receives a MESSAGE PARITY ERROR message, it shall terminate the retry activity. This may be done either by changing to any other information transfer phase and transferring at least one byte of information or by going to the BUS FREE phase (see 5.1.1). The initiator shall accept such action as aborting the negotiation, and both devices shall go to asynchronous data transfer mode for data transfers between the two devices.

If the target recognizes that negotiation is required, it sends an SDTR message to the initiator. Prior to releasing the ACK signal on the last byte of the SDTR message from the target, the initiator shall assert the ATN signal

and respond with its SDTR message or with a MESSAGE REJECT message. If an abnormal condition prevents the initiator from returning an appropriate response, both devices shall go to asynchronous data transfer mode for data transfers between the two devices.

Following an initiator's responding SDTR message, an implied agreement for synchronous operation shall not be considered to exist until the target leaves the MESSAGE OUT phase, indicating that the target has accepted the negotiation. After a vendor-specific number of retry attempts (greater than zero), if the target has not received the initiator's responding SDTR message, it shall go to the BUS FREE phase without any further information transfer attempt (see 5.1.1). This indicates that a catastrophic error condition has occurred. Both devices shall go to asynchronous data transfer mode for data transfers between the two devices.

If, following an initiator's responding SDTR message, the target shifts to MESSAGE IN phase and the first message in is MESSAGE REJECT, the implied agreement shall be considered to be negated and both devices shall go to asynchronous data transfer mode for data transfers between the two devices.

The implied synchronous agreement shall remain in effect until a BUS DEVICE RESET message is received, until a hard reset condition occurs, or until one of the two SCSI devices elects to modify the agreement. The default data transfer mode is asynchronous data transfer mode. The default data transfer mode is entered at power on, after a BUS DEVICE RESET message, or after a hard reset condition.

5.6.22. TERMINATE I/O PROCESS

not defined
↓

The TERMINATE I/O PROCESS message is sent from the initiator to the target to advise the target to terminate the current I/O process without corrupting the medium. Upon successful receipt of this message the target shall terminate the current I/O process as soon as possible and return COMMAND TERMINATED status. The sense key shall be set to NO SENSE and the additional sense code and qualifier shall be set to ~~NO SENSE~~ I/O PROCESS TERMINATED. The TERMINATE I/O PROCESS message shall not affect pending status, data, and commands for other queued or executing I/O processes. However, continued execution and status of other I/O processes queued for the I-T nexus may be affected by the queue error recovery option specified in the control mode page (see 7.3.3.1).

If the I/O process that is being terminated has a data transfer associated with it (i.e., DATA IN or DATA OUT phase), the valid bit in the sense data shall be set to one and the information field shall be set as follows:

(1) If the command descriptor block specifies an allocation length or parameter list length in bytes, the information field shall be set to the difference (residue) between the transfer length and the number of bytes successfully transferred.

(2) If the command descriptor block specifies a transfer length field, the information field shall be set as defined in the REQUEST SENSE command (see 7.2.14).

If the I/O process being terminated has no data transfer associated with it the target shall set the valid bit in the sense data to zero and terminate the I/O process with COMMAND TERMINATED status. The sense key shall be set to NO SENSE and the additional sense code and qualifier shall be set to I/O PROCESS TERMINATED.

When any error condition is detected for an I/O process the target shall ignore the TERMINATE I/O PROCESS message and terminate the I/O process with the appropriate error status and sense data for the error condition.

If the target completes all processing for a command (i.e., all data has been read, written, or processed) and a TERMINATE I/O PROCESS message is received before the I/O process is terminated, the target shall ignore the TERMINATE I/O PROCESS message and terminate the I/O process in the normal manner.

If the target receives a TERMINATE I/O PROCESS message before the command descriptor block is transferred (i.e., before the COMMAND phase) or the message is issued to an I-T nexus that does not have an active or queued I/O process, the target shall set the valid bit in the sense data to zero and terminate the I/O process with COMMAND TERMINATED status. The sense key shall be set to NO SENSE and the additional sense code and qualifier shall be set to I/O PROCESS TERMINATED.

If the current I/O process is in the command queue (I-T nexus for untagged queuing or I-T-L-Q nexus for tagged queuing) and has not started execution, the target shall either terminate the I/O process immediately or disconnect and wait until the command is at the head of the queue (started executing) then terminate the I/O process. In either case, the target shall terminate the I/O process with COMMAND TERMINATED status. The sense key shall be set to

NO SENSE and the additional sense code and qualifier shall be set to I/O PROCESS TERMINATED. The valid bit shall be set to zero.

If the target does not support this message or is unable to stop the current I/O process for the I T x nexus, it shall respond by sending a MESSAGE REJECT message to the initiator and continuing the I/O process in a normal manner.

IMPLEMENTORS NOTE: The TERMINATE I/O PROCESS message provides a means for the initiator to request the target to reduce the transfer length of the current command to the amount that has already been transferred. The initiator can use the sense data to determine the actual number of bytes or blocks that have been transferred. This message is normally used by the initiator to stop a lengthy read, write, or verify operation when a higher-priority command is available to be executed. It is up to the initiator to complete the terminated command at a later time, if required.

5.6.23. WIDE DATA TRANSFER REQUEST Message

Table 5-10: WIDE DATA TRANSFER MESSAGE

Byte	Value	Description
0	01h	Extended message
1	02h	Extended message length
2	03h	WIDE DATA TRANSFER REQUEST code
3	m	Transfer Width (2*m bytes)

A WIDE DATA TRANSFER REQUEST (WDTR) message (Table 5-10) exchange shall be initiated by an SCSI device whenever a previously-arranged transfer width agreement may have become invalid. The agreement becomes invalid after any condition which may leave the data transfer agreement in an indeterminate state such as:

- (1) after a hard reset condition
- (2) after a BUS DEVICE RESET message and
- (3) after a power cycle.

In addition, an SCSI device may initiate an WDTR message exchange whenever it is appropriate to negotiate a new transfer width agreement. SCSI devices that are capable of wide data transfers (greater than eight bits) shall not respond to an WDTR message with a MESSAGE REJECT message.

IMPLEMENTORS NOTE: Re-negotiation at every selection is not recommended, since a significant performance impact is likely.

The WDTR message exchange establishes an agreement between two SCSI devices on the width of the data path to be used for DATA phase transfers between the two devices. This agreement applies to DATA IN and DATA OUT phases only. All other information transfer phases shall use an eight-bit data path.

If an SCSI device implements both wide data transfer option and synchronous data transfer option, then it shall negotiate the wide data transfer agreement prior to negotiating the synchronous data transfer agreement. If a synchronous data transfer agreement is in effect, then an SCSI device that accepts a WDTR message shall reset the synchronous agreement to asynchronous mode.

The transfer width that is established applies to all logical units on both SCSI devices. Valid transfer widths are 8 bits (m = 00h), 16 bits (m = 01h), and 32 bits (m = 02h). Values of m greater than 02h are reserved.

The originating SCSI device (the SCSI device that sends the first of the pair of WDTR messages) sets its transfer width value to the maximum data path width it elects to accommodate. If the responding SCSI device can also accommodate this transfer width, it returns the same value in its WDTR message. If it requires a smaller transfer width, it substitutes the smaller value in its WDTR message. The successful completion of an exchange of WDTR messages implies an agreement as follows:

Responding Device WDTR Response	Implied Agreement
(1) Non-zero transfer width	Each device transmits and receives data with a transfer width equal to the responding SCSI device's transfer width.
(2) Transfer width equal to zero	Eight-bit Data Transfer
(3) MESSAGE REJECT message	Eight-bit Data Transfer

If the initiator recognizes that negotiation is required, it asserts the ATN signal and sends a WDTR message to begin the negotiating process. After successfully completing the MESSAGE OUT phase, the target shall respond with the proper WDTR message. If an abnormal condition prevents the target from returning an appropriate response, both devices shall go to eight-bit data transfer mode for data transfers between the two devices.

Following target response (1) above, the implied agreement for wide data transfers shall be considered to be negated by both the initiator and the target if the initiator asserts ATN and the first message out is either MESSAGE PARITY ERROR or MESSAGE REJECT. In this case, both devices shall go to eight-bit data transfer mode for data transfers between the two devices. For the MESSAGE PARITY ERROR case, the implied agreement shall be reinstated if a re-transmission of the second of the pair of messages is successfully accomplished. After a vendor-specific number of retry attempts (greater than zero), if the target receives a MESSAGE PARITY ERROR message, it shall terminate the retry activity. This may be done either by changing to any other information transfer phase and transferring at least one byte of information or by going to the BUS FREE phase (see 5.1.1). The initiator shall accept such action as aborting the negotiation, and both devices shall go to eight-bit data transfer mode for data transfers between the two devices.

If the target recognizes that negotiation is required, it sends a WDTR message to the initiator. Prior to releasing the ACK signal on the last byte of the WDTR message from the target, the initiator shall assert the ATN signal and respond with its WDTR message or with a MESSAGE REJECT message. If an abnormal condition prevents the initiator from returning an appropriate response, both devices shall go to eight-bit data transfer mode for data transfers between the two devices.

Following an initiator's responding WDTR message, an implied agreement for wide data transfer operation shall not be considered to exist until the target leaves the MESSAGE OUT phase, indicating that the target has accepted the negotiation. After a vendor-specific number of retry attempts (greater than zero), if the target has not received the initiator's responding WDTR message, it shall go to the BUS FREE phase without any further information transfer attempt (see 5.1.1). This indicates that a catastrophic error condition has occurred. Both devices shall go to eight-bit data transfer mode for data transfers between the two devices.

If, following an initiator's responding WDTR message, the target shifts to MESSAGE IN phase and the first message in is MESSAGE REJECT, the implied agreement shall be considered to be negated and both devices shall go to eight-bit data transfer mode for data transfers between the two devices.

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The implied transfer width agreement shall remain in effect until a BUS DEVICE RESET message is received, until a hard reset condition occurs, or until one of the two SCSI devices elects to modify the agreement. The default data transfer width is eight-bit data transfer mode. The default data transfer mode is entered at power on, after a BUS DEVICE RESET message, or after a hard reset condition.

6. SCSI Commands and Status

This section defines the SCSI command and status structures and gives several examples.

By keeping to a minimum the functions essential to communicate via this protocol, a wide range of peripheral devices of varying capability can operate in the same environment. Because subsets of the full architecture may be implemented, optional functions are noted.

6.1. Command Implementation Requirements

The first byte of all SCSI commands shall contain an operation code as defined in this standard. Targets shall implement all commands with a mandatory operation code (see 6.1.2) both in section 7 and in the appropriate section for their device type.

6.1.1. Reserved

Reserved bits, fields, bytes, and code values are set aside for future standardization. Their use and interpretation may be specified by future extensions to this standard. A reserved bit, field, or byte shall be set to zero, or in accordance with a future extension to this standard. A target that receives a reserved bit, field, or byte that is not zero or receives a reserved code value shall terminate the command with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST. It shall also be acceptable for a target to interpret a bit, field, byte, or code value in accordance with a future extension to this standard.

6.1.2. Operation Code Types

Operation Code Type	Description
M	Mandatory - Commands so designated shall be implemented in order to meet the minimum requirement of this standard.
O	Optional - Commands so designated, if implemented, shall be implemented as defined in this standard.
V	Vendor specific - Operation codes so designated are available for vendor defined commands. See the vendor specifications where compatibility is desired.
R	Reserved - Operation codes so designated shall not be used. They are reserved for future extensions to this standard.

6.2. Command Descriptor Block

A command is communicated by sending a command descriptor block to the target. For several commands, the command descriptor block is accompanied by a list of parameters sent during the DATA OUT phase. See the specific commands for detailed information.

The command descriptor block always has an operation code as its first byte and a control byte as its last byte.

For all commands, if there is an invalid parameter in the command descriptor block, then the target shall terminate the command without altering the medium.

Table 6-1: Typical Command Descriptor Block for Six-byte Commands

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code							
1	Logical Unit Number			(MSB)				
2	Logical Block Address (if required)							
3	(LSB)							
4	Transfer Length (if required) Parameter List Length (if required) Allocation Length (if required)							
5	Control							

Table 6-2: Typical Command Descriptor Block for Ten-byte Commands

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code							
1	Logical Unit Number				Reserved			
2	(MSB)							
3								
4	Logical Block Address (if required)							
5								
6	(LSB)							
7	Reserved							
8	(MSB)							
9	Transfer Length (if required) Parameter List Length (if required) Allocation Length (if required)							
9	(LSB)							
9	Control							

Table 6-3: Typical Command Descriptor Block for Twelve-byte Commands

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code							
1	Logical Unit Number			Reserved				
2	(MSB)							
3								
4	Logical Block Address (if required)							
5								
6	(MSB)							
7								
8	Transfer Length (if required) Parameter List Length (if required) Allocation Length (if required)							
9								
10	(LSB)							
11	Reserved							
	Control							

6.2.1. Operation Code

The operation code (Table 6-4) of the command descriptor block has a group code field and a command code field. The three-bit group code field provides for eight groups of command codes. The five-bit command code field provides for thirty-two command codes in each group. Thus, a total of 256 possible operation codes exist. Operation codes are defined in the subsequent Sections of this document.

The group code specifies one of the following groups:

- Group 0 - six-byte commands (see Table 6-1)
- Group 1 - ten-byte commands (see Table 6-2)
- Group 2 - ten-byte commands (see Table 6-2)
- Group 3 - reserved
- Group 4 - reserved
- Group 5 - twelve-byte commands (see Table 6-3)
- Group 6 - vendor specific
- Group 7 - vendor specific

Table 6-4: Operation Code

Bit	7	6	5	4	3	2	1	0
Group Code				Command Code				

6.2.2. Logical Unit Number

The logical unit number is defined in the IDENTIFY message (5.6.7). The target shall ignore the logical unit number specified within the command descriptor block if an IDENTIFY message was received. It is recommended that the logical unit number in the command descriptor block be set to zero.

NOTICE: The logical unit number field is included in the command descriptor block for compatibility with some SCSI-1 devices. This field may be reclaimed in SCSI-3. New implementations should use the outbound IDENTIFY message, which is mandatory in SCSI-2, to establish the I_T_L nexus.

6.2.3. Logical Block Address

The logical block address on logical units or within a partition on device volumes shall begin with block zero and be contiguous up to the last logical block on that logical unit or within that partition.

A six-byte command descriptor block contains a 21-bit logical block address. The ten-byte and the twelve-byte command descriptor blocks contain 32-bit logical block addresses. Logical block addresses in additional parameter data have their length specified for each occurrence. See the specific command descriptions.

6.2.4. Transfer Length

The transfer length field specifies the amount of data to be transferred, usually the number of blocks. For several commands the transfer length indicates the requested number of bytes to be sent as defined in the command description. For these commands the transfer length field may be identified by a different name. See the following descriptions and the individual command descriptions for further information.

Commands that use one byte for the transfer length allow up to 256 blocks of data to be transferred by one command. A transfer length value of 1 to 255 indicates the number of blocks that shall be transferred. A value of zero indicates 256 blocks.

In commands that use multiple bytes for the transfer length, a transfer length of zero indicates that no data transfer shall take place. A value of one or greater indicates the number of blocks that shall be transferred.

Refer to the specific command description for further information.

6.2.5. Parameter List Length

The parameter list length is used to specify the number of bytes sent during the DATA OUT phase. This field is typically used in command descriptor blocks for parameters that are sent to a target (e.g., mode parameters, diagnostic parameters, log parameters, etc.). A parameter length of zero indicates that no data shall be transferred. This condition shall not be considered as an error.

6.2.6. Allocation Length

The allocation length field specifies the maximum number of bytes that an initiator has allocated for returned data. An allocation length of zero indicates that no data shall be transferred. This condition shall not be considered as an error. The target shall terminate the DATA IN phase when allocation length bytes have been transferred or when all available data have been transferred to the initiator, whichever is less.

The allocation length is used to limit the maximum amount of data (e.g., sense data, mode data, log data, diagnostic data, etc.) returned to an initiator.

6.2.7. Control Field

The control field is the last byte of every command descriptor block. The control field is defined in Table 6-5.

Table 6-5: Control Field

Bit	7	6	5	4	3	2	1	0
	Vendor specific			Reserved			Flag	Link

The flag bit specifies which message the target shall return to the initiator if the link bit is one and the command completes without error. Implementation of the flag bit is optional.

The flag bit should be set to zero if the link bit is zero. If link bit is zero and the flag bit is one, the target shall return CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST.

If the flag bit is zero and the link bit is one, and if the command completes successfully, the target shall send the LINKED COMMAND COMPLETE message. If the flag bit is one and the link bit is one, and if the command completes successfully, the target shall send the LINKED COMMAND COMPLETE (WITH FLAG) message.

IMPLEMENTORS NOTE: The flag bit is typically used to cause an interrupt in the initiator between linked commands.

The link bit is used to continue the I/O process across multiple commands. Implementation of the link bit is optional.

A link bit of one indicates that the initiator requests a continuation of the I/O process and that the target should enter the command phase upon successful completion of the current command.

If the link bit is one, and if the command completes successfully, the target shall return INTERMEDIATE or INTERMEDIATE-CONDITION MET status and shall then send one of the two messages defined by the flag bit.

If either of the link and flag bits are set to one, and the target does not implement linked commands, it shall return CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST.

6.3. Status

The status byte and status byte code are specified in Tables 6-6 and 6-7. A status byte shall be sent from the target to the initiator during the STATUS phase at the completion of each command unless the command is terminated by one of the following events:

- (1) an ABORT message
- (2) an ABORT TAG message
- (3) a BUS DEVICE RESET message
- (4) a CLEAR QUEUE message
- (5) a hard reset condition
- (6) an unexpected disconnect (see 5.1.1)

The STATUS phase normally occurs at the end of a command but in some case may occur prior to transferring the command descriptor block.

Table 6-6: Status Byte

Bit	7	6	5	4	3	2	1	0
	Reserved			Status Byte Code				Reserved

Table 6-7: Status Byte Code

Bits of Status Byte								Status
7	6	5	4	3	2	1	0	
R	R	0	0	0	0	0	R	GOOD
R	R	0	0	0	0	1	R	CHECK CONDITION
R	R	0	0	0	1	0	R	CONDITION MET
R	R	0	0	1	0	0	R	BUSY
R	R	0	1	0	0	0	R	INTERMEDIATE
R	R	0	1	0	1	0	R	INTERMEDIATE-CONDITION MET
R	R	0	1	1	0	0	R	RESERVATION CONFLICT
R	R	1	0	0	0	1	R	COMMAND TERMINATED
R	R	1	0	1	0	0	R	QUEUE FULL
All Other Codes								Reserved

Key: R - Reserved bit

A definition of the status byte codes is given below.

GOOD. This status indicates that the target has successfully completed the command.

CHECK CONDITION. This status indicates that a contingent allegiance condition has occurred (see 6.6).

CONDITION MET. This status or INTERMEDIATE-CONDITION MET is returned whenever the requested operation is satisfied (see the SEARCH DATA and PRE-FETCH commands).

BUSY. This status indicates that the target is busy. This status shall be returned whenever a target is unable to accept the command from an otherwise acceptable initiator (i.e., no reservation conflicts). The recommended initiator recovery action is to issue the command again at a later time.

INTERMEDIATE. This status or INTERMEDIATE-CONDITION MET shall be returned for every successfully completed command in a series of linked commands (except the last command), unless the command is terminated with CHECK CONDITION, RESERVATION CONFLICT, or COMMAND TERMINATED status. If INTERMEDIATE or INTERMEDIATE-CONDITION MET status is not returned, the series of linked commands is terminated and the I/O process is ended.

INTERMEDIATE-CONDITION MET. This status is the combination of the CONDITION MET and INTERMEDIATE statuses.

RESERVATION CONFLICT. This status shall be returned whenever an initiator attempts to access a logical unit or an extent within a logical unit that is reserved with a conflicting reservation type for another SCSI device (see the RESERVE and RESERVE UNIT commands). The recommended initiator recovery action is to issue the command again at a later time.

COMMAND TERMINATED. This status shall be returned whenever the target terminates the current I/O process after receiving a TERMINATE I/O PROCESS message (see 5.6.22). This status also indicates that a contingent allegiance condition has occurred (see 6.6).

QUEUE FULL. This status shall be implemented if tagged queuing is implemented. This status is returned when a SIMPLE QUEUE TAG, ORDERED QUEUE TAG, or HEAD OF QUEUE TAG message is received and the command queue is full. The I/O process is not placed in the command queue.

6.4. Command Examples

The following sections give examples of typical command processing in the SCSI environment.

6.4.1. Single Command Example

An I/O process containing one untagged READ command is used in this section to illustrate a simple I/O process on the SCSI bus. This example does not include error or exception conditions.

The initiator has one set of active pointers that includes a command pointer, a data pointer, and a status pointer. In addition, the initiator has one set of saved pointers for each I/O process that it is able to concurrently manage. The initiator sets up the saved pointers to point to the appropriate bytes for the I/O process and copies the saved pointers to the active pointers. It then arbitrates for the SCSI bus, and upon winning arbitration, selects the target. Once the target is selected, the target assumes control of the I/O process.

During the SELECTION phase, the initiator asserts the ATN signal to inform the target that the initiator wishes to send a message. The target enters the MESSAGE OUT phase and transfers the IDENTIFY message from the initiator. This message informs the target of which logical unit is to be used. At this point, an I T L nexus has been established for the I/O process. This nexus associates the initiator's pointers with the I/O process.

The target switches to the COMMAND phase and transfers the command descriptor block from the initiator. In this case, the command descriptor block contains a READ command. The target interprets the command and switches to the DATA IN phase, transfers the data, switches to STATUS phase, sends GOOD status, switches to MESSAGE IN phase, and transfers a COMMAND COMPLETE message. After successfully sending the COMMAND COMPLETE message, the target goes to the BUS FREE phase by releasing the BSY signal and the I/O process ends.

6.4.2. Disconnect Example

In the above single command example, the length of time necessary to obtain the data may require a time-consuming physical positioning operation. In order to improve system throughput, the target may disconnect from the initiator, thereby freeing the SCSI bus to allow other I/O process to occur.

After the target has received the READ command (and has determined that there will be a delay), it disconnects from the SCSI bus by sending a DISCONNECT message and by going to the BUS FREE phase.

After the target retrieves the requested data from the peripheral device it arbitrates for the SCSI bus. Upon winning arbitration, it reselects the initiator and sends an IDENTIFY message to the initiator via the MESSAGE IN phase. This revives the I T L nexus so that the initiator can retrieve the correct set of pointers for the I/O process. The initiator restores the active pointers to their most recent saved values (which, in this case, are the initial values) and the target continues (as in the single command example) to finish the I/O process.

If target wishes to disconnect after transferring part of the data (e.g., while crossing a cylinder boundary), it may do so by sending a SAVE DATA POINTER message and a DISCONNECT message to the initiator and then disconnecting. When reconnection is completed, the current data pointer is restored to its value immediately prior to the SAVE DATA POINTER message.

On those occasions when an error or exception condition occurs and the target elects to repeat the information transfer, the target may repeat the transfer by either issuing a RESTORE POINTERS message or by disconnecting without issuing a SAVE DATA POINTER message. When reconnection is completed, the most recent saved pointer values are restored.

6.4.3. Linked Command Example

An I/O process may contain multiple commands "linked" together. Upon completing a linked command successfully, the target automatically proceeds to the next linked command for the I/O process. All commands in a series of linked commands are addressed to the same nexus and are part of a single I/O process.

The commands are not entirely independent. When using the relative address bit (see 8.1.10), the address of the last logical block accessed by one of the commands is available to the next command. Thus one can search for a particular data pattern using a SEARCH DATA command and then read the logical block containing the data pattern with a READ command linked to the SEARCH DATA command. One can also read a logical block at a specified displacement from the block containing the data pattern.

A LINKED COMMAND COMPLETE or LINKED COMMAND COMPLETE (WITH FLAG) message is sent from the target to the initiator to indicate that a linked command completed. The initiator then updates the saved pointers for the nexus so that subsequent transfers from the target reference the next command of the series. Command processing of linked and single commands is similar except that relative addressing is permitted in linked commands.

For example, a successful completion of a SEARCH DATA EQUAL command causes the target to continue with the linked READ command from the initiator. If the relative address bit in the READ command has been set to one, and the address field of the READ command is set to zero, the target transfers the successfully searched block to the initiator.

6.5. Command Processing Considerations and Exception Conditions

The following sections describe some exception conditions and errors associated with command processing and the sequencing of commands.

6.5.1. Programmable Operating Definition

Some applications require that the operating definition of a logical unit be modified to meet the special requirements of a particular initiator. The program-controlled modification of the operating definition is typically provided to allow operating systems to change the operating definition of a more recently developed targets to one which is more compatible with the operating system. This ability requires that the system comply with the low-level hardware definitions of SCSI-2.

The parameters that can be changed by modifying the operating definition of a logical unit include the vendor identification, the device type, the device model, the SCSI compliance level, the SCSI specification level, the command set, and other parameters. The low-level hardware parameters including signal timing and parity definitions cannot be changed by modifying the operating definition. The present operating definition of a logical unit with respect to an initiator can be determined at any time by execution of an INQUIRY command. In some vendor-specific cases, it may also be necessary to perform other commands including MODE SENSE and READ CAPACITY.

Each logical unit begins at a particular operating definition. If the logical unit supports the CHANGE DEFINITION command, the present operating definition can be changed to any other operating definition supported by the logical unit. The actual details of the operating definition of a logical unit are vendor-specific. If the operating definition is changed to one that does not include the CHANGE DEFINITION command, the target should continue to accept the CHANGE DEFINITION command.

If an error occurs during execution of a CHANGE DEFINITION command, the original operating definition remains in effect after the command is executed. The new operating definition becomes active only after successful execution of the CHANGE DEFINITION command.

Since new operating definitions may preclude the execution of I/O processes that are already in progress, the target may disconnect to allow completion of any I/O processes that are in progress. Operating definition changes that may cause conflicts with the normal operation from other initiators shall be indicated to those initiators by generating a unit attention condition for each other initiator. The additional sense code shall be set to CHANGED OPERATING DEFINITION.

An initiator may request a list of the operating definitions that the target supports and descriptive text for each operating definition using the INQUIRY command.

6.5.2. Incorrect Initiator Connection

An incorrect initiator connection occurs on a reconnection if:

- (1) an initiator attempts to reconnect to an I/O process, and
- (2) a soft reset condition has not occurred, and
- (3) the initiator does not send an ABORT, ABORT TAG, BUS DEVICE RESET, CLEAR QUEUE, or TERMINATE I/O PROCESS message during the same MESSAGE OUT phase as the IDENTIFY message.

An incorrect initiator connection also occurs on an initial connection when an initiator:

- (1) attempts to establish an I_T_L_Q nexus when an I_T_L nexus already exists from a previous connection, or
- (2) attempts to establish an I_T_L nexus when an I_T_L_Q nexus already exists unless there is a contingent allegiance or extended contingent allegiance condition present for the logical unit or target routine.

A target that detects an incorrect initiator connection shall abort all I/O processes for the initiator on the logical unit or target routine and shall return CHECK CONDITION status. The sense key shall be set to ABORTED COMMAND and the additional sense code shall be set to OVERLAPPED COMMANDS ATTEMPTED.

If an initiator reconnects to an I/O process and a soft reset condition has occurred, the target shall meet the requirements of 5.2.2.2.

IMPLEMENTORS NOTES:

(1) An incorrect initiator connection may be indicative of a serious error and, if not detected, could result in an I/O process operating with a wrong set of pointers. This is considered a catastrophic failure on the part of the initiator. Therefore, vendor-specific error recovery procedures may be

required to guarantee the data integrity on the medium. The target may return additional sense data to aid in this error recovery procedure (e.g., sequential-access devices may return the residue of blocks remaining to be written or read at the time the second command was received).

(2) Some targets may not detect an incorrect initiator connection until after the command descriptor block has been received.

6.5.3. Selection of an Invalid Logical Unit

The target's response to selection of a logical unit which is not valid is described in the following paragraphs.

The logical unit may not be valid because:

(1) the target does not support the logical unit (e.g., some targets support only one peripheral device). In response to an INQUIRY command the target shall return the INQUIRY data with the peripheral qualifier set to the value required in Table 7-16. In response to any other command except REQUEST SENSE the target shall terminate the command with CHECK CONDITION status. In response to a REQUEST SENSE command the target shall return sense data. The sense key shall be set to ILLEGAL REQUEST and the additional sense code shall be set to LOGICAL UNIT NOT SUPPORTED.

(2) the target supports the logical unit, but the peripheral device is not currently attached to the target. In response to an INQUIRY command the target shall return the INQUIRY data with the peripheral qualifier set to the value required in Table 7-16. In response to any other command except REQUEST SENSE the target shall terminate the command with CHECK CONDITION status. In response to a REQUEST SENSE command the target shall return sense data. The sense key shall be set to ILLEGAL REQUEST and the additional sense code shall be set to LOGICAL UNIT NOT SUPPORTED.

(3) the target supports the logical unit and the peripheral device is attached, but not operational. In response to an INQUIRY command the target shall return the INQUIRY data with the peripheral qualifier set to the value required in Table 7-16. The target's response to any command other than INQUIRY and REQUEST SENSE is vendor specific.

(4) the target supports the logical unit but is incapable of determining if the peripheral device is attached or is not operational when it is not ready. In response to an INQUIRY command the target shall return the INQUIRY data with the peripheral qualifier set to the value required in Table 7-16. In response to a REQUEST SENSE command the target shall return the REQUEST SENSE data with a sense key of NO SENSE unless a prior error condition exists. The target's response to any other command is vendor specific.

6.5.4. Parameter Rounding

Certain parameters sent to a target with various commands contain a range of values. Targets may choose to implement only selected values from this range. When the target receives a value that it does not support, it either rejects the command (CHECK CONDITION status with ILLEGAL REQUEST sense key) or it rounds the value received to a supported value. The target shall reject unsupported values unless rounding is permitted in the description of the parameter.

Rounding of parameter values, when permitted, shall be performed as follows: A target that receives a parameter value that is not an exact supported value

shall adjust the value to one that it supports and shall return CHECK CONDITION status with a sense key of RECOVERED ERROR. The additional sense code shall be set to ROUNDED PARAMETER. The initiator is responsible to issue an appropriate command to learn what value the target has selected.

IMPLEMENTORS NOTE: Generally, the target should adjust maximum-value fields down to the next lower supported value than the one specified by the initiator. Minimum-value fields should be rounded up to the next higher supported value than the one specified by the initiator. In some cases, the type of rounding (up or down) is explicitly specified in the description of the parameter.

6.5.5. Asynchronous Event Notification

Implementation of asynchronous event notification is optional. This protocol can be used to inform processor devices that an asynchronous event has occurred. A SEND command with an AEN bit of one is issued to a processor device with a subsequent data phase that includes the REQUEST SENSE information. SCSI devices that respond to an INQUIRY command as a processor device type with asynchronous event notification capability may be notified of asynchronous events using this protocol. An SCSI device has to be capable of acting as an initiator in order to perform asynchronous event notification.

IMPLEMENTORS NOTE: Asynchronous event notification cannot be used with a device that acts as a temporary initiator (e.g., devices executing COPY commands) since they are not identified as a processor device.

Parameters affecting the use of asynchronous event notification are contained in the control mode page (see 7.3.3.1).

The four uses of asynchronous event notification are:

- (1) informing a processor of an error condition encountered after command completion
- (2) informing all processor devices that a newly initialized device is available
- (3) informing all processor devices of other unit attention conditions
- (4) informing all processor devices of other asynchronous events.

Other uses of asynchronous event notification are not prohibited, however this protocol is not intended to be used while an I_T_L or I_T_L_Q nexus exists between the processor device (i.e., the initiator of the nexus) and the other SCSI device (i.e., the target of the nexus). Asynchronous event notification is not intended for use with target routines (i.e., an I_T_R or I_T_R_Q nexus).

An example of the first case above is a device that implements a write cache. If the target is unable to write cached data to the medium, it may use an asynchronous event notification to inform the initiator of the failure. An extended contingent allegiance condition may also be established during the same I_T_L nexus used for the asynchronous event notification (see 5.6.9).

An example of the third case above is a device that supports removable media. Asynchronous event notification may be used to inform an initiator of a not-ready-to-ready transition (medium changed) or of an operator initiated event (e.g., activating a write protect switch or activating a start or stop

switch).

An example of the fourth case above is a sequential-access device performing a REWIND command with the immediate bit set to one. Asynchronous event notification may be used to inform an initiator that the beginning of medium has been reached. Completion of a CD-ROM AUDIO PLAY command started in the immediate mode is another example of this case.

Notification of an asynchronous event is performed using the SEND command with the AEN bit set to one. The information identifying the condition being reported shall be returned during the DATA OUT phase of the SEND command (see 11.2.2.). The data sent is shown in Table 11-4.

An error condition or unit attention condition shall be reported once per occurrence of the event causing it. The target may choose to use an asynchronous event notification or to return CHECK CONDITION status on a subsequent command, but not both. Notification of command-related error conditions shall be sent only to the processor that initiated the I/O process.

The asynchronous event notification protocol can be used to notify processor devices that a system resource has become available. If a target chooses to use this method, the sense key in the sense data sent to the processor device shall be set to UNIT ATTENTION.

The asynchronous event notification protocol shall be used only to SCSI devices that return processor device type with an AENC bit of one in response to an INQUIRY command. The INQUIRY command should be issued to logical unit zero of each SCSI device responding to selection. This procedure shall be conducted prior to the first asynchronous event notification and shall be repeated whenever the device deems it appropriate or when an event occurs that may invalidate the current information. (See 5.6.21, SYNCHRONOUS DATA TRANSFER REQUEST message, for examples of these events.)

Each SCSI device that returns processor device type with an AENC bit of one shall be issued a TEST UNIT READY command to determine that the SCSI device is ready to receive an asynchronous event notification. An SCSI device returning CHECK CONDITION status is issued a REQUEST SENSE command. This clears any pending unit attention condition. An SCSI device that returns processor device type with an AENC bit of one and returns GOOD status when issued a TEST UNIT READY command shall accept a SEND command with an AEN bit of one.

IMPLEMENTORS NOTE: An SCSI device which can use asynchronous event notification at initialization time should provide means to defeat these notifications. This can be done with a switch or jumper wire. Devices which implement saved parameters may alternatively save the asynchronous event notification permissions either on a per SCSI device basis or as a system wide option. In any case, a device conducts a survey with INQUIRY commands to be sure that the devices on the SCSI bus are appropriate destinations for SEND commands with an AEN bit of one. (The devices on the bus or the SCSI ID assignments may have changed.)

6.6. Contingent Allegiance Condition

The contingent allegiance condition shall exist following the return of CHECK CONDITION or COMMAND TERMINATED status and may optionally exist following an unexpected disconnect. The contingent allegiance condition shall be preserved for the I T x nexus until it is cleared. The contingent allegiance condition shall be cleared upon the generation of a hard reset condition, or by an ABORT message, a BUS DEVICE RESET message, or any subsequent command for the I T x nexus. While the contingent allegiance condition exists the target shall preserve the sense data for the initiator.

While the contingent allegiance condition exists, if the target is unable to maintain separate sense data, the target shall respond to any other requests for access to the logical unit or target routine from another initiator with a BUSY status. Execution of queued commands for the logical unit or target routine for which the contingent allegiance condition exists shall be suspended until the contingent allegiance condition is cleared.

6.7. Extended Contingent Allegiance Condition

Implementation of extended contingent allegiance is optional. The extended contingent allegiance condition extends the contingent allegiance condition for an I T x nexus. This condition is generated by the target sending an INITIATE RECOVERY message to the initiator following CHECK CONDITION or COMMAND TERMINATED status and prior to the COMMAND COMPLETE message. This condition shall be preserved until it is cleared by a BUS DEVICE RESET message, a RELEASE RECOVERY message, or a hard reset condition.

While the extended contingent allegiance condition exists the target shall respond to any other request for access to the logical unit from another initiator with BUSY status. Execution of queued commands for the logical unit for which the extended contingent allegiance condition exists shall be suspended until the RELEASE RECOVERY message is received by the target.

IMPLEMENTORS NOTES:

- (1) It is not required to generate an extended contingent allegiance condition for every CHECK CONDITION or COMMAND TERMINATED status that occurs. Simple errors not requiring an extended recovery may be dealt with by using the contingent allegiance protocol.
- (2) During the existence of the extended contingent allegiance condition, appropriate error recovery sequences may be executed. Such commands can correct data, modify or delete queued commands, perform LOG SENSE commands and obtain diagnostic information. Extended contingent allegiance is recommended for error conditions that may require execution of multiple-step error-recovery protocols without interference from other initiators.

An extended contingent allegiance condition may also be generated using an asynchronous event notification protocol. When the event is detected, the bus device which detected the event assumes the initiator role and transmits a SEND command with an AEN bit of one to the appropriate device(s) (see 11.2.2.).

If the device wishes to generate an extended contingent allegiance condition during an asynchronous event notification, it shall send an INITIATE RECOVERY message following the IDENTIFY message (and following any queue tag message) and prior to the COMMAND phase of the SEND command. An extended contingent

allegiance condition can be generated for only one I T L nexus at a time. The extended contingent allegiance condition is cleared when a RELEASE RECOVERY message is received from the device to which the INITIATE RECOVERY message was sent. The generation of a hard reset condition, or receipt of a BUS DEVICE RESET message shall also clear the extended contingent allegiance condition.

During an extended contingent allegiance only untagged I/O processes from the SCSI device to which the INITIATE RECOVERY message was sent shall be executed by the target. If the initiator sends a queue tag message the target shall respond with QUEUE FULL status. After the extended contingent allegiance condition is cleared any commands remaining in the command queue shall be executed.

6.8. Queued I/O Processes

The implementation of queuing for I/O processes is optional. Queuing of I/O processes allows a target to accept multiple I/O processes for execution at a later time.

There are two methods for implementation of queuing, tagged and untagged. Tagged queuing allows a target to accept multiple I/O processes from each initiator for each logical unit or target routine. Untagged queuing allows a target to accept one I/O process from each initiator for each logical unit. A target may be capable of both methods queuing, but only one method may be used at a time.

Untagged queuing may be supported by SCSI-1 or SCSI-2 devices. Tagged queuing is new in SCSI-2.

6.8.1. Untagged Queuing

Untagged queuing allows a target to accept a command from an initiator for a logical unit or target routine while a command from another initiator is being executed. Only one command for each I_T_x nexus shall be accepted at a time.

An I/O process may be initiated any time the BUS FREE phase exists even if an I/O process from a different initiator is active. If the disconnect privilege is not granted in the IDENTIFY message of the new I/O process, the target may either suspend the active I/O process or it may return BUSY status to the new I/O process.

The I_T_x nexus sufficiently specifies the relationship so that the target can reconnect to the initiator to restore the pointers for I/O process as long as only one I/O process per I_T_x nexus is issued. It is the responsibility of the initiator to assure that only one such I/O process is issued at any time (see 6.5.2).

6.8.2. Tagged Queuing

Tagged queuing allows a target to accept multiple I/O processes from the same or different initiators until the logical unit's command queue is full. If an I/O process is received and the command queue is full, the target shall terminate the I/O process with QUEUE FULL status.

The command queue is setup by the target for each supported logical unit and target routine. Initiators may add or delete I/O processes to the queue. When adding an I/O process, the initiator may specify fixed order of execution, allow the target to define the order of execution, or specify that the I/O process is to be executed next.

If the disconnect privilege is not granted in the IDENTIFY message for a tagged I/O process the target shall return BUSY status.

The queue tag messages (see 5.6.17) allow the initiator to establish a unique I T L Q nexus to identify each I/O process. The I T L Q nexus allows the target to reconnect to a specific I/O process and allows the initiator to restore the set of pointers for that I/O process. An initiator may have several I/O processes ongoing to the same or different logical units or target routines as long as each has a unique nexus.

If only SIMPLE QUEUE TAG messages are used, the target may execute the commands in any order that is deemed desirable within the constraints of the queue management algorithm specified in the control mode page (see 7.3.3.1).

If ORDERED QUEUE TAG messages are used, the target shall execute the commands in the order received with respect to other commands received with ORDERED QUEUE TAG messages. All commands received with an SIMPLE QUEUE TAG message prior to a command received with an ORDERED QUEUE TAG message, regardless of initiator, shall be executed before that command with the ORDERED QUEUE TAG message. All commands received with an SIMPLE QUEUE TAG message after a command received with an ORDERED QUEUE TAG message, regardless of initiator, shall be executed after that command with the ORDERED QUEUE TAG message.

A command received with a HEAD OF QUEUE TAG message is placed first in the queue, to be executed next. A command received with a HEAD OF QUEUE TAG message shall not suspend an I/O process for which the target has begun execution. Consecutive commands received with HEAD OF QUEUE TAG messages are executed in a last-in-first-out order.

An I/O process received without a queue tag message while there are any tagged I/O commands in the command queue is an incorrect initiator connection (see 6.5.2) unless there is a contingent allegiance or extended contingent allegiance condition.

A series of linked commands are a single I/O process, and are assigned the queue tag established in the initial connection. A command received with a HEAD OF QUEUE TAG message shall not suspend a series of linked commands for which the target has begun execution.

The RESERVE and RELEASE commands should be sent with an ORDERED QUEUE TAG message. Use of the HEAD OF QUEUE TAG message with these commands could result in reservation conflicts with previously issued commands.

The TEST UNIT READY and INQUIRY commands are often sent with a HEAD OF QUEUE TAG message, since the information returned has no effect on the condition of the logical unit.

Two error recovery options are allowed. The error recovery option to be used is specified in the control mode page (see 7.3.3.1).

The first post-recovery option is to continue execution of commands in the queue after the contingent allegiance or extended contingent allegiance condition has cleared. The target returns BUSY status to other initiators while the contingent allegiance or extended contingent allegiance condition exists. During this time all commands in the queue are suspended. All commands used for recovery operations shall be untagged commands. The queue may be modified by removing all or selected commands in the queue as part of the recovery procedure.

If commands are combined by the queuing algorithm such that the error affects more than one command, then a contingent allegiance condition shall be generated for all affected commands.

The second recovery option clears the queue after the contingent allegiance or extended contingent allegiance condition has been cleared. When the queue is cleared because of this recovery option, a unit attention condition shall be generated for all other initiators and the additional sense code shall be set to COMMANDS CLEARED BY ANOTHER INITIATOR.

Deferred errors are normally related to a command that has already completed. As such, there is no attempt to return the queue tag value assigned to the original command.

A device that does not support tagged queuing for any reason (e.g., not implemented, disabled by the DQue bit in the control mode page, or it has been switched to an operating definition that does not include tagged queuing) shall respond to any queue tag message with a MESSAGE REJECT message. The target shall continue the I/O process as if it was an untagged I/O process.

Tagged queuing may also be temporarily disabled internal to the SCSI device during certain initialization periods or to control internal resource utilization. During these periods the target may elect to return QUEUE FULL status or it may respond to any queue tag message with a MESSAGE REJECT message.

Several messages may be used to clear part or all of the command queue. Please refer to the ABORT, ABORT TAG, BUS DEVICE RESET, and CLEAR QUEUE messages in Section 5 for details.

6.8.3. Example of Queued I/O Process

An example of I/O process queuing benefits from the consideration of the execution of a number of commands. After each command, the state of the queue kept in the target is shown to indicate the function actually performed by the queuing.

6.8.3.1. Typical Sequences for Tagged Queuing

An I/O process using tagged queuing uses the following sequences for normal execution. The initiator first arbitrates for the SCSI bus, and after successfully obtaining the SCSI bus, selects the appropriate SCSI device. The ATN signal is asserted during the SELECTION phase to indicate that a MESSAGE OUT phase is requested by the initiator. The first message byte transferred is an IDENTIFY message. The ATN signal continues to be asserted during the MESSAGE OUT phase to indicate that the initiator has another message. The second message byte transferred is the first byte of the appropriate queue tag message, in this case a SIMPLE QUEUE TAG message. The third and last message byte is transmitted containing the second byte of the queue tag message, the queue tag. As it is transferred, the ATN signal is negated to indicate that no more message bytes are available. The target then transfers the command descriptor block. Assuming the command requires disconnection, the target transmits a DISCONNECT message to the initiator and then enters the BUS-FREE phase. The target places the command, identified by the I_T_L_Q nexus, at the appropriate place in the command queue.

When the target removes I/O processes from the queue for execution, a physical latency period may occur. At the end of this period, when the target is prepared to transfer the appropriate data, the target begins an ARBITRATION phase and, upon winning, enters a RESELECTION phase. After a successful reselection, the target sends the IDENTIFY message followed by a SIMPLE QUEUE TAG message with the queue tag value originally sent by the initiator. The initiator uses the I_T_L_Q nexus to identify the correct set of pointers and control blocks associated with the I/O process and to establish the necessary conditions for data transfer. The target begins data transfer. When the data transfer is successfully completed, the target returns GOOD status and terminates the I/O process with a COMMAND COMPLETE message.

6.8.3.2. Example of Tagged Queuing

An example of the execution of five queued I/O processes is described to demonstrate how tagged queuing operates. All tagged I/O processes are from one initiator to a single logical unit of a single target. The five I/O processes are defined in Table 6-8. The target is a direct-access device. At the time the I/O processes are first being executed, it is assumed that the actuator is in position to access logical block 10000.

Table 6-8: Commands in Order Received by Target

Command	Queue Tag Message	Queue Tag Value	Logical Block Address	Transfer Length	Status
READ	SIMPLE	01h	10000	1000	Queued
READ	SIMPLE	02h	100	1	Queued
READ	ORDERED	03h	1000	1000	Queued
READ	SIMPLE	04h	10000	1	Queued
READ	SIMPLE	05h	2000	1000	Queued

The optimum order would require that those blocks close to the actuator position be the first blocks accessed, followed by those increasingly far from the actuator position. However, the command with queue tag 03h is an ordered I/O process, so that all simple I/O processes transferred previously must be executed before, while all simple I/O processes transferred after the ordered I/O process must be executed after the ordered I/O process.

If a target supports an optimizing algorithm the actual order in which the I/O processes are executed could be as shown in Table 6-9.

Table 6-9: Commands in Order of Execution

Command	Queue Tag Message	Queue Tag Value	Logical Block Address	Transfer Length	Status
READ	SIMPLE	01h	10000	1000	Queued
READ	SIMPLE	02h	100	1	Queued
READ	ORDERED	03h	1000	1000	Queued
READ	SIMPLE	05h	2000	1000	Queued
READ	SIMPLE	04h	10000	1	Queued

I/O processes with queue tag values 01h and 02h are executed in the order received since the actuator is already in position to execute I/O process 01h. I/O process 02h must be executed before I/O process 04h or 05h because the ordered I/O process 03h was transmitted after I/O processes 01h and 02h but before I/O processes 04h and 05h. I/O process 03h is then executed after I/O process 02h. The I/O processes 04h and 05h are executed after the ordered I/O process 03h. I/O process 05h is executed before I/O process 04h because the actuator is in position to access block 2000 after executing I/O process 03h. I/O process 04h is executed last.

As an example of the operation of the HEAD OF QUEUE TAG I/O process, consider that a new I/O process, identified by a HEAD OF QUEUE TAG message with a queue tag of 08h, is transmitted to the target while the ordered I/O process 03h is being executed. The I/O process 03h continues execution, but the new HEAD OF QUEUE TAG I/O process is placed in the queue for execution before all subsequent I/O processes. In this case, the queue for execution after the ordered I/O process 03h was executed would appear as shown in Table 6-10.

Table 6-10: Modified by HEAD OF QUEUE TAG Message

Command	Queue Tag Message	Queue Tag Value	Logical Block Address	Transfer Length	Status
READ	ORDERED	03h	1000	1000	Executing
READ	HEAD OF QUEUE	08h	0	8	Queued
READ	SIMPLE	05h	2000	1000	Queued
READ	SIMPLE	04h	10000	1	Queued

To obtain maximum performance gains using tagged queuing requires careful implementation of the queuing algorithms in the target. In addition, initiators should allow a maximum number of simple I/O processes to be executed with a minimum number of ordered I/O processes. RESERVE and RELEASE commands, SET LIMITS commands, and appropriate software locking conventions should be used to guarantee the proper relationship between the commands executed and the data stored on the peripheral devices. These conventions are not defined by this standard.

6.9. Unit Attention Condition

The target shall generate a unit attention condition for each initiator on each valid logical unit whenever the target has been reset by a BUS DEVICE RESET message, a hard reset condition, or by a power-on reset. The target shall also generate a unit attention condition on the affected logical unit(s) for each initiator whenever one of the following events occurs:

- (1) A removable medium may have been changed.
- (2) The mode parameters in effect for this initiator have been changed by another initiator.
- (3) The version or level of microcode has been changed.
- (4) Tagged commands queued for this initiator were cleared by another initiator.
- (5) INQUIRY data has been changed.
- (6) The mode parameters in effect for this initiator have been restored from non-volatile memory.
- (7) A change in the condition of a synchronized spindle.
- (8) Any other event occurs that requires the attention of the initiator.

IMPLEMENTORS NOTES:

- (1) Targets may queue unit attention conditions on logical units. After the first unit attention condition is cleared, another unit attention condition may exist (e.g., a power on condition followed by a microcode change condition). The initiator can clear all pending unit attention conditions by repeatedly sending the REQUEST SENSE command until a sense key other than UNIT ATTENTION is returned by the target.
- (2) See 6.5.3 for requirements concerning selection of an invalid logical unit.

The unit attention condition shall persist on the logical unit for each initiator until that initiator clears the condition as described in the following paragraphs.

If an INQUIRY command is received from an initiator to a logical unit with a pending unit attention condition (before the target generates the contingent allegiance condition), the target shall perform the INQUIRY command and shall not clear the unit attention condition. If the INQUIRY command is received after the target has generated the contingent allegiance condition for a pending unit attention condition, then the unit attention condition on the logical unit shall be cleared, and the target shall perform the INQUIRY command.

If any other command is received after the target has generated the contingent allegiance condition for a pending unit attention condition, the unit attention condition on the logical unit shall be cleared, and if no other unit attention condition is pending the target shall perform the command. If another unit attention condition is pending the target shall not perform the command and shall generate another contingent allegiance condition.

If a REQUEST SENSE command is received from an initiator with a pending unit attention condition (before the target generates the contingent allegiance condition), then the target may either:

- (1) report any pending sense data and preserve the unit attention condition on the logical unit or
- (2) report the unit attention condition, may discard any pending sense data, and clear the unit attention condition on the logical unit for that initiator.

(This page is intentionally blank.)

If the target has already generated the contingent allegiance condition for the unit attention condition, the target shall perform the second action listed above.

If an initiator issues a command other than INQUIRY or REQUEST SENSE while a unit attention condition exists for that initiator (prior to generating the contingent allegiance condition for the unit attention condition), the target shall not perform the command and shall report CHECK CONDITION status unless a higher priority status as defined by the target is also pending (e.g., BUSY or RESERVATION CONFLICT).

If after generating the contingent allegiance condition for a pending unit attention condition, the next command received from that initiator on the logical unit is not REQUEST SENSE, then that command shall be performed and the unit attention condition shall be cleared for that initiator on the logical unit and the sense data is lost (see 6.6).

If a target becomes a temporary initiator to issue a SEND command with an AEN bit of one, which informs the initiator (temporary target) of the unit attention condition, and the SEND command completes with GOOD status, then the target shall clear the unit attention condition for that initiator on the logical unit (see 6.5.5).

7. All Device Types

7.1. Model for All Device Types

This model describes some of the general characteristics expected of most SCSI devices. It is not intended to define any requirements nor is it intended to alter any requirements defined elsewhere in this standard. Section 6 of this standard also contains model information pertaining to all device types.

7.1.1. SCSI Addresses

There are two levels of addresses within the SCSI architecture: the SCSI device address and the logical unit number or target routine number.

7.1.1.1. SCSI Device Address

SCSI devices occupy (i.e., respond to) one address on the SCSI bus. Generally the SCSI device provides a means (e.g., switches, jumpers) to select one of the eight available addresses (0 through 7). This address is used during bus arbitration and selection or reselection of SCSI devices. Each device on the SCSI bus is assigned an unique address.

Normally, the SCSI device address is set when the system is configured and it remains static thereafter. Some systems and devices provide vendor-specific means to alter this address at other times.

7.1.1.2. Logical Units

Each target has one or more logical units, beginning with logical unit zero. There is a maximum of eight logical units. These logical units are usually mapped directly to peripheral devices, but they may be a portion of a peripheral device or may comprise multiple peripheral devices.

An initiator can determine whether a target implements a logical unit by issuing an INQUIRY command and examining the returned peripheral qualifier and peripheral device type.

The concept of a logical unit is not defined for an initiator. (An SCSI device may implement both the initiator role and the target role. In this case logical unit(s) are defined only for the target role.)

7.1.1.3. Target Routines

An optional feature of the SCSI architecture permits each target to have one or more target routines, beginning with target routine number zero. There is a maximum of eight target routines. These target routines are processes that execute directly on the target and are not associated with a particular logical unit or peripheral device. Target routines are addressed using the LUNAR bit of the IDENTIFY message (see 5.6.7).

Target routines are principally intended to return information about the target and the only valid commands are INQUIRY and REQUEST SENSE.

7.1.2. Commands Implemented by all SCSI Devices

This standard defines four commands that all SCSI-2 targets implement: INQUIRY, REQUEST SENSE, SEND DIAGNOSTIC, and TEST UNIT READY. These commands are used to configure the system, to test targets, and to return important information concerning errors and exception conditions.

7.1.2.1. Using the INQUIRY Command

The INQUIRY command may be used by a system to determine the configuration of the SCSI bus. Target devices respond with information that includes their type and standard level and may include the vendor's identification, model number and other useful information. It is recommended that SCSI targets be capable of returning this information (or whatever part of it that is available) upon completing power-on initialization. An SCSI device may take longer to get certain portions of this information, especially if it retrieves the information from the medium.

7.1.2.2. Using the REQUEST SENSE Command

Whenever a contingent allegiance condition (6.6) is established, the initiator that received the error should issue a REQUEST SENSE command to receive the sense data describing what caused the contingent allegiance condition. If the initiator issues some other command, the sense data is lost.

7.1.2.3. Using the SEND DIAGNOSTIC Command

The SEND DIAGNOSTIC command provides a means to request the target to perform a self test. While the test is target specific, the means of requesting the test is standardized and the response is simply GOOD status if all is well or CHECK CONDITION status if the test fails.

The SEND DIAGNOSTIC command also provides other powerful features when used in conjunction with the RECEIVE DIAGNOSTIC RESULTS command, but this capability is optional.

7.1.2.4. Using the TEST UNIT READY Command

The TEST UNIT READY command is useful in that it allows an initiator to poll a logical unit until it is ready without the need to allocate space for returned data. It is especially useful to check cartridge status of logical units with removable media. Targets are expected to respond promptly to indicate the current status of the device (i.e., a target should avoid lengthy disconnections in an attempt to respond with GOOD status).

7.2. Commands for All Device Types

The operation codes for commands that apply to all device types are listed in Table 7-1.

Table 7-1: Commands for All Device Types

Command Name	Operation Code	Type	Section
CHANGE DEFINITION	40h	O	7.2.1
COMPARE	39h	O	7.2.2
COPY	18h	O	7.2.3
COPY AND VERIFY	3Ah	O	7.2.4
INQUIRY	12h	M	7.2.5
LOG SELECT	4Ch	O	7.2.6
LOG SENSE	4Dh	O	7.2.7
MODE SELECT(6)	15h	Z	7.2.8
MODE SELECT(10)	55h	Z	7.2.9
MODE SENSE(6)	1Ah	Z	7.2.10
MODE SENSE(10)	5Ah	Z	7.2.11
READ BUFFER	3Ch	O	7.2.12
RECEIVE DIAGNOSTIC RESULTS	1Ch	O	7.2.13
REQUEST SENSE	03h	M	7.2.14
SEND DIAGNOSTIC	1Dh	O	7.2.15
TEST UNIT READY	00h	M	7.2.16
WRITE BUFFER	3Bh	O	7.2.17

Key: M = Command implementation is mandatory.
 O = Command implementation is optional.
 Z = Command implementation is device type specific.

7.2.1. CHANGE DEFINITION Command

Table 7-2: CHANGE DEFINITION Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (40h)							
1	Logical Unit Number				Reserved			
2	Reserved							Save
3	Reserved		Definition Parameter					
4	Reserved							
5	Reserved							
6	Reserved							
7	Reserved							
8	Parameter Data Length							
9	Control							

The CHANGE DEFINITION command (Table 7-2) modifies the operating definition of the selected logical unit or target with respect to commands from the selecting initiator or for all initiators.

A save control bit (Save) of zero indicates that the target shall not save the operating definition. A Save bit of one indicates that the target shall save the operating definition to non-volatile memory.

The definition parameter field is defined in Table 7-3.

Table 7-3: Definition Parameter Field

Value	Meaning of Definition Parameter
00h	Use Current Operating Definition
01h	SCSI X3.131-1986 Operating Definition
02h	CCS Operating Definition
03h	SCSI-2 X3.131-198X Operating Definition
04-3Fh	Reserved
40-7Fh	Vendor Specific

IMPLEMENTORS NOTE: The current operating definition parameter values establish operating definitions compatible with the appropriate SCSI specification. Vendor-specific values are available for those applications where more complex operation definition changes are required. Definitions supported by a device are returned in the implemented operating definition page (see 7.3.4.3).

The parameter data length field specifies the length in bytes of the parameter data that shall be transferred from the initiator to the target. A parameter data length of zero indicates that no data shall be transferred. This condition shall not be considered as an error. Parameter data lengths greater than zero indicate the number of bytes of parameter data that shall be transferred.

The parameter data is vendor specific.

IMPLEMENTORS NOTE: The parameter data may be used to specify a password to validate an operating definition change.

The CHANGE DEFINITION command causes one of the operating definition modifications listed below:

- (1) Change the operating definition of a logical unit relative to the initiator that issued the command. In this case, the target is capable of maintaining an unique operating definition for each logical unit relative to each initiator in the system.
- (2) Change the operating definition of the target relative to the initiator that issued the command. In this case, the target is capable of maintaining an unique operating definition, for each initiator in the system, that applies to all logical units of the target.
- (3) The operating definition of a logical unit relative to all initiators in the system. In this case, the target is capable of maintaining an unique operating definition for each logical unit relative to all initiators in the system.
- (4) The operating definition of the target relative to all initiators in the system. In this case, the target is capable of maintaining only one operating definition.

IMPLEMENTORS NOTE:

- (1) This standard does not provide a direct means to determine which of the above four methods has been implemented by the target. An indirect means of determining which method is implemented exists in that the target is required to inform affected initiators of operating definition changes via the unit attention condition.
- (2) Cases (3) and (4), above, may result in incompatibilities if there are other initiators in the system operated below the SCSI-2 level.

The operating definition is modified after successful completion of the command. A target shall consider the command successfully completed when it detects the assertion of the ACK signal for the COMMAND COMPLETE message. The initiator should verify the new operating definition by issuing an INQUIRY command requesting the implemented operating definition page (see Table 7-75).

It is permissible for an SCSI-2 device that has its definition changed to an SCSI-1 device to accept a CHANGE DEFINITION command.

IMPLEMENTORS NOTE: The method of changing the operating definition is implementation dependent. Some implementations may require the target's operating mode be re-initialized as if a power-up or hard-reset had occurred. Other implementations may modify only those operating definitions that are affected by the CHANGE DEFINITION command.

If the CHANGE DEFINITION command is not executed successfully for any reason, the operating definition shall remain the same as it was before the CHANGE DEFINITION command was attempted. If it is impossible to return to the previous operating definition, an unit attention condition shall be generated by the target.

IMPLEMENTORS NOTE: The present operating definition of the target may always be interrogated through the INQUIRY command. A change in the operating definition may change the vendor identifier, the device type, the device model, the SCSI implementation level, the command set, and any other operating characteristics. The logical unit responds to any command according to the requirements of its present operating definition.

After a power-on condition or a hard RESET condition, the target shall set its initial operating definition to the last saved value, if saving is implemented, or its default value, if saving is not implemented.

7.2.2. COMPARE Command

Table 7-4: COMPARE Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (39h)							
1	Logical Unit Number			Reserved				Pad
2	Reserved							
3	(MSB)							
5	Parameter List Length							(LSB)
6								
8	Reserved							
9	Control							

The COMPARE command (Table 7-4) provides the means to compare data from one logical unit with another or the same logical unit in a manner similar to the COPY command.

This command functions in the same manner as the COPY command, except that the data from the source is compared on a byte-by-byte basis with the data from the destination. The parameter list transferred to the target is the same as for the COPY command. This parameter list contains the information to identify the logical units involved in the comparison and the length of the comparison. (See 7.2.3 for additional information about the COPY command.)

If the comparison is unsuccessful, the command shall be terminated with CHECK CONDITION status and the sense key shall be set to MISCOMPARE. The remaining fields in the sense data shall be set as documented in the COPY command.

7.2.3. COPY Command

Table 7-5: COPY Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (18h)							
1	Logical Unit Number			Reserved				Pad
2	(MSB)							
4	Parameter List Length							(LSB)
5	Control							

The COPY command (Table 7-5) provides a means to copy data from one logical unit to another or the same logical unit. The logical unit that receives and performs the COPY command is called the copy manager. The copy manager is responsible for copying data from a logical unit (source device) to a logical unit (destination device). These logical units may reside on different SCSI devices or the same SCSI device (in fact all three may be the same logical unit). Some SCSI devices that implement this command may not support copies to or from another SCSI device, or may not support third party copies (i.e., both the source and the destination logical units reside on other SCSI devices).

The pad bit (7.2.3.7) is used in conjunction with the Cat bit (7.2.3.7) in the segment descriptors to define what action should be taken when a segment of the copy does not fit exactly into an integer number of destination blocks.

The parameter list length field specifies the length in bytes of the parameters that shall be sent during the DATA OUT phase of the command. A parameter list length of zero indicates that no data shall be transferred. This condition shall not be considered as an error.

The COPY parameter list (Table 7-6) begins with a four-byte header that contains the COPY function code and priority. Following the header is one or more segment descriptors.

Table 7-6: COPY Parameter List

Bit Byte	7	6	5	4	3	2	1	0
0	COPY Function Code					Priority		
1	Vendor Specific							
2	Reserved							
3	Reserved							
	Segment Descriptor(s)							
0 - xx	Segment Descriptor 0 (See specific table for length.)							
	:							
	:							
0 - xx	Segment Descriptor n (See specific table for length.)							

The COPY function code field defines a specific format for the segment descriptors. The COPY function codes are defined in Table 7-7. A target need not support all function codes for its device type.

The priority field of the COPY parameter list establishes the relative priority of this COPY command to other commands being executed by the same target. All other commands are assumed to have a priority of 1. Priority 0 is the highest priority with increasing values indicating lower priorities.

The segment descriptor formats are determined by the COPY function code. The segment descriptor format used for write-once devices, CD-ROM devices, and optical-memory devices shall be the same as for direct-access devices. The segment descriptor format used for printer devices, communications, and processor devices, shall be the same as for sequential-access devices. Thus a COPY from a write-once device to a printer device uses the same segment descriptor format as for a COPY from a direct-access device to a sequential-access device (see Table 7-7). The segment descriptor formats are described in Tables 7-8 through 7-11. A maximum of 256 segment descriptors are permitted. The segment descriptors are identified by ascending numbers beginning with zero.

Table 7-7: COPY Function Codes

Peripheral Device Type		COPY Function Code	Segment Descriptor Table	Comments
Source	Destination			
Direct-Access (0,4,5,7)	Sequential-Access (1,2,3,9)	0	7-8	
Sequential-Access (1,3,9)	Direct-Access (0,4,7)	1	7-8	
Direct-Access (0,4,5,7)	Direct-Access (0,4,7)	2	7-9	
Sequential-Access (1,3,9)	Sequential-Access (1,2,3,9)	3	7-10	
Sequential-Access (1)	Sequential-Access (1)	4	7-11	Image Copy

The numbers in parenthesis are the peripheral device type codes (see Table 7-17).

COPY function code: 00h Direct-access to sequential-access
 01h Sequential-access to writable direct-access
 02h Direct-access to writable direct-access
 03h Sequential-access to sequential-access
 04h Sequential-access image copy
 05h - 0Fh Reserved
 10h - 1Fh Vendor specific

For the COMPARE command the destination direct-access device does not have to be a writable device.

7.2.3.1. Errors Detected by the Managing SCSI Device

Two classes of exception conditions may occur during execution of a COPY command. The first class consists of those exception conditions detected by the SCSI device that received the COPY command and is managing the execution of the command. These conditions include parity errors while transferring the COPY command and status byte, invalid parameters in the COPY command, invalid segment descriptors, and inability of the SCSI device controlling the COPY functions to continue operating. In the event of such an exception condition, the SCSI device managing the COPY shall:

- (1) Terminate the COPY command with CHECK CONDITION status.
- (2) The valid bit in the sense data shall be set to one. The segment number shall contain the number of the segment descriptor being processed at the time the exception condition is detected. The sense key shall contain the sense key code describing the exception condition (i.e., not COPY ABORTED). The information field shall contain the difference between the number of blocks field in the segment descriptor being processed at the time of the failure and the number of blocks successfully copied. This number is the residue of unprocessed blocks remaining for the segment descriptor.

7.2.3.2. Errors Detected by a Target

The second class of errors consists of exception conditions detected by the SCSI device transferring data at the request of the SCSI device managing the transfer. The SCSI device managing the COPY command detects exception conditions by receiving CHECK CONDITION status from one of the SCSI devices it is managing. It then shall recover the sense data associated with the exception condition.

The SCSI device managing the COPY command may also be the source or destination SCSI device (or both). It shall distinguish between a failure of the management of the COPY and a failure of the data transfer being requested. It shall then create the appropriate sense data internally.

After recovering the sense data associated with the detected error, the SCSI device managing the COPY command shall:

- (1) Terminate the COPY command with CHECK CONDITION status.
- (2) The valid bit in the sense data shall be set to one. The segment number shall contain the number of the segment descriptor being processed at the time the exception condition is detected. The sense key shall be set to COPY ABORTED. The information field shall contain the difference between the number of blocks field in the segment descriptor being processed at the time of the failure and the number of blocks successfully copied. This number is the residue of unprocessed blocks remaining for the segment descriptor.

The first byte of the command-specific information field shall specify the starting byte number, relative to the first byte of sense data, of an area that contains (unchanged) the source logical unit's status byte and sense data. A zero value indicates that no status byte or sense data is being returned for the source logical unit.

The second byte of the command-specific information field shall specify the starting byte number, relative to the first byte of sense data, of an area that contains (unchanged) the destination logical unit's status byte and sense data. A zero value indicates that no status byte or sense data is being returned for the destination logical unit.

7.2.3.3. COPY Function Code 00h and 01h

The format for the segment descriptors for COPY transfers between direct-access and sequential-access devices is specified in Table 7-8. This format is required for COPY function codes 00h or 01h. The segment descriptor may be repeated up to 256 times within the parameter list length specified in the command descriptor block.

Table 7-8: Segment Descriptor for COPY Function Codes 00h and 01h

Bit Byte	7	6	5	4	3	2	1	0
0	Source Address			Reserved	Cat	Source LUN		
1	Destination Address			Reserved		Destination LUN		
2	(MSB)							
3	Sequential-Access Device Block Length							(LSB)
4	(MSB)							
7	Direct-Access Device Number of Blocks							(LSB)
8	(MSB)							
11	Direct-Access Device Logical Block Address							(LSB)

The source address and source LUN fields specify the SCSI bus ID and logical unit of the device to copy the data from for this segment of the COPY command. The destination address and destination LUN fields specify the SCSI bus ID and logical unit to copy the data to for this segment of the COPY command. Some SCSI devices may not support third-party COPY in which the copying SCSI device is not the source or destination device. Some SCSI devices only support COPY within the SCSI device and not to other SCSI devices. If an unsupported COPY operation is requested, the command shall be terminated with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST with an additional sense code of INVALID FIELD IN PARAMETER LIST (see 7.2.3.1).

A catenate (Cat) bit (optional) of one indicates that the COPY manager shall catenate the last source block of a segment with the first source block of the next segment if the last source block does not end exactly at the end of the destination block. The definition of a cat bit of zero depends on the setting of the pad bit in the command descriptor block (see 7.2.3.7).

The sequential-access device block-length field specifies the block length to be used on the sequential-access logical unit during this segment of the COPY command. If the SCSI device managing the COPY knows this block length is not supported, the command shall be terminated with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST with an additional sense code of INVALID FIELD IN PARAMETER LIST. If the block length is found to be invalid while executing a read or write operation to the sequential-access device, the command shall be terminated with CHECK CONDITION status and the sense key shall be set to COPY ABORTED (see 7.2.3.2).

The direct-access device number of blocks field specifies the number of blocks in the current segment to be copied. A value of zero indicates that no blocks shall be transferred in this segment.

The direct-access device logical block address field specifies the starting logical block address on the logical unit for this segment.

7.2.3.4. COPY Function Code 02h

The format for the segment descriptors for COPY transfers among direct-access devices is specified in Table 7-9. This format is required for COPY function code 02h. The segment descriptor may be repeated up to 256 times within the parameter list length specified in the command descriptor block.

Table 7-9: Segment Descriptor for COPY Function Code 02h

Bit Byte	7	6	5	4	3	2	1	0
0	Source Address			DC	Cat	Source LUN		
1	Destination Address			Reserved		Destination LUN		
2	Reserved							
3	Reserved							
4	(MSB)							
7	Number of Blocks (LSB)							
8	(MSB)							
11	Source Logical Block Address (LSB)							
12	(MSB)							
15	Destination Logical Block Address (LSB)							

See 7.2.3.3 for definitions of the source address, the source LUN, the destination address, the destination LUN, and CAT fields.

A destination count (DC) bit of zero indicates that the number of blocks field refers to the source logical unit. A DC bit of one indicates that the number of blocks field refers to the destination logical unit.

The number of blocks field specifies the number of blocks to be transferred to or from (depending on the DC bit) the device during this segment. A value of zero indicates that no blocks shall be transferred.

The source logical block address field specifies the starting logical block address on the source device.

The destination logical block address field specifies the starting logical block address on the destination device.

7.2.3.5. COPY Function Code 03h

The format for the segment descriptors for COPY transfers among sequential-access devices is specified by Table 7-10. This format is required for COPY function code 03h. The segment descriptor may be repeated up to 256 times within the parameter list length specified in the command descriptor block.

Table 7-10: Segment Descriptor for COPY Function Code 03h

Bit Byte	7	6	5	4	3	2	1	0
0	Source Address			DC	Cat	Source LUN		
1	Destination Address			Reserved		Destination LUN		
2	Reserved							
3	Reserved							
4	(MSB)							
5	Source Block Length (LSB)							
6	(MSB)							
7	Destination Block Length (LSB)							
8	(MSB)							
11	Number of Blocks (LSB)							

See 7.2.3.3 for definitions of the source address, the source LUN, the destination address, the destination LUN, and CAT fields.

A destination count (DC) bit of zero indicates that the number of blocks field refers to the source logical unit. A DC bit of one indicates that the number of blocks field refers to the destination logical unit.

The source block length field specifies the block-length of the source device for this segment of the COPY. A zero in this field indicates variable block-length. For non-zero values, this field shall match the logical unit's actual block-length.

If block-length mismatches are detected prior to the beginning of the read operation by the SCSI device managing the COPY, the command shall be terminated with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST and the additional sense code shall be set to INVALID FIELD IN PARAMETER LIST (see 7.2.3.1).

If the mismatches are detected during the read operation by the COPY manager, the command shall be terminated with CHECK CONDITION status. The sense key shall be set to COPY ABORTED (see 7.2.3.2), and the additional sense code shall be set to INVALID FIELD IN PARAMETER LIST.

The destination block-length field specifies the block length to be used on the destination logical unit during the COPY. Destination block length mismatches are handled in an analogous manner as source block length mismatches.

The number of blocks field specifies the number of blocks to be transferred to or from (depending on the DC bit) the device during this segment. A value of zero indicates that no blocks shall be transferred.

7.2.3.6. COPY Function Code 04h

The format for the segment descriptors for image COPY transfers between sequential-access devices is specified in Table 7-11. This format is required for COPY function code 04h. The segment descriptor may be repeated up to 256 times within the parameter list length specified in the command descriptor block.

Table 7-11: Segment Descriptor for COPY Function Code 04h

Bit Byte	7	6	5	4	3	2	1	0
0	Source Address			Reserved		Source LUN		
1	Destination Address			Reserved		Destination LUN		
2	Count							
3								
7	-	-	Reserved					-
8								
11	-	-	Vendor Specific					-

See 7.2.3.3 for definitions of the source address, the source LUN, the destination address, the destination LUN, and CAT fields.

The image mode COPY command copies an exact image of the source device medium to the destination device medium, beginning at their current positions. The copy function terminates when the source device:

- (1) encounters an end-of-partition as defined by the source device
- (2) encounters an end-of-data as defined by the source device (i.e., BLANK CHECK sense key)
- (3) has copied the number of consecutive filemarks specified in the count field from the source device to the destination device
- (4) has copied the number of consecutive setmarks specified in the count field from the source device to the destination device, if the RSmk bit in the device configuration page (see 9.3.3.1) is one.

A count field of zero indicates that the COPY command shall not terminate due to any number of consecutive filemarks or setmarks. Other error or exception conditions (e.g., early-warning end-of-partition on the destination device) may cause the COPY command to terminate prior to completion. In such cases, it is not possible to calculate a residue, so the information field in the sense data shall be set to zero.

7.2.3.7. Copies With Unequal Block Lengths

When copying data between two devices with unequal block lengths, it is possible for the last source block to not completely fill the last destination block for one or more segments in the COPY command. Two optional bits are defined to assist in controlling the copy manager's actions in this circumstance. The Pad bit (in the command descriptor block) and the Cat bit (in each applicable segment descriptor) are defined in Table 7-12.

Table 7-12: Pad and Cat Bit Definition

Pad	Cat	COPY Manager's Action
0	0	On inexact segments, it is device specific whether the COPY manager rejects the COPY command with CHECK CONDITION status and ILLEGAL REQUEST sense key, the COPY manager writes or accepts short blocks (variable-block mode on sequential-access devices), or the COPY manager adds pad characters (00h) to the destination block or strips pad characters from the source block.
1	0	On inexact segments, the COPY manager shall add pad characters (00h) to the destination block to completely fill the block or it shall strip pad characters from the source block, always stopping at the end of a complete block.
X	1	The COPY manager shall always write or read complete blocks. On inexact segments, the remainder of the block contains data from the next segment. This code is not valid in the last segment of the COPY command.

IMPLEMENTORS NOTE: Use of pad characters is intended to assist in managing COPY commands between devices of different block lengths where partial-block residues may occur. The initiator which issued the COPY command is responsible for management of these pad areas (i.e., remembering where they

are). One possible method is to write the COPY parameter list information to the destination medium prior to issuing the COPY command for backup and to read this information prior to issuing the COPY command for restore.

7.2.4. COPY AND VERIFY Command

Table 7-13: COPY AND VERIFY Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (3Ah)							
1	Logical Unit Number			Reserved			BytChk	Pad
2	Reserved							
3	(MSB)							
5	Parameter List Length							(LSB)
6	Reserved							
8	Reserved							
9	Control							

The COPY AND VERIFY command (Table 7-13) performs the same function as the COPY command, except that a verification of the data written to the destination logical unit is performed after the data is written. The parameter list transferred to the target is the same as for the COPY command. This parameter list contains the information to identify the logical units involved in the copy and the length of the copy. See 7.2.3 for additional information about the COPY command.

A byte check (BytChk) bit of zero causes a medium verification to be performed with no data comparison. A BytChk bit of one causes a byte-by-byte compare of data written on the destination medium and the data transferred from the source medium. If the compare is unsuccessful for any reason, the copy manager shall return CHECK CONDITION status with the sense key set to MISCOMPARE. The remaining fields in the sense data shall be set as documented in the COPY command.

7.2.5. INQUIRY Command

Table 7-14: INQUIRY Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (12h)							
1	Logical Unit Number			Reserved				EVPD
2	Page Code							
3	Reserved							
4	Allocation Length							
5	Control							

The INQUIRY command (Table 7-14) requests that information regarding parameters of the target and its attached peripheral device(s) be sent to the initiator. An option allows the initiator to request additional information about the target or logical unit (see 7.2.5.2).

An enable vital product data (EVPD) bit of one specifies that the target shall return the optional vital product data specified by the page code field. If the target does not support vital product data and this bit is set to one, the target shall return CHECK CONDITION status with the sense key set to ILLEGAL REQUEST and an additional sense code of INVALID FIELD IN CDB.

An EVPD bit of zero specifies that the target shall return the standard INQUIRY data. If the page code field is not zero, the target shall return CHECK CONDITION status with the sense key set to ILLEGAL REQUEST and an additional sense code of INVALID FIELD IN CDB.

The page code field specifies which page of vital product data information the target shall return (see 7.3.4).

The INQUIRY command shall return CHECK CONDITION status only when the target cannot return the requested INQUIRY data.

IMPLEMENTORS NOTE: It is recommended that the INQUIRY data be returned even though the peripheral device may not be ready for other commands.

If an INQUIRY command is received from an initiator with a pending unit attention condition (i.e., before the target reports CHECK CONDITION status), the target shall perform the INQUIRY command and shall not clear the unit attention condition (see 6.9).

IMPLEMENTORS NOTES:

(1) The INQUIRY command is typically used by the initiator after a reset or power-up condition to determine the device types for system configuration. To minimize delays after a reset or power-up condition the standard INQUIRY data should be available without incurring any media access delays. If the target does store some of the INQUIRY data on the device it may return zeros or ASCII spaces (20h) in those fields until the data is available from the device.

(2) The INQUIRY data may change as the target executes its initialization sequence or in response to a CHANGE DEFINITION command. For example, the target may contain a minimum command set in its non-volatile memory and may load its final firmware from the device when it becomes ready. After it has loaded the firmware it may support more options and therefore return different supported options information in the INQUIRY data.

7.2.5.1. Standard INQUIRY Data

The standard INQUIRY data (Table 7-15) contains 36 required bytes, followed by a variable number of vendor-specific parameters. Bytes 56 through 95, if returned, are reserved for future standardization.

The standard INQUIRY data format is shown in Table 7-15.

Table 7-15: Standard INQUIRY Data Format

Bit Byte	7	6	5	4	3	2	1	0
0	Peripheral Qualifier			Peripheral Device Type				
1	RMB	Device-Type Modifier						
2	ISO Version		ECMA Version			ANSI-Approved Version		
3	AENC	TrmIOP	Reserved		Response Data Format			
4	Additional Length (n-4)							
5	Reserved							
6	Reserved							
7	RelAdr	WBus32	WBus16	Sync	Linked	Reserved	CmdQue	SftRe
8	(MSB)							
15	Vendor Identification (LSB)							
16	(MSB)							
31	Product Identification (LSB)							
32	(MSB)							
35	Product Revision Level (LSB)							
36								
55	Vendor Specific							
56								
95	Reserved							
Vendor-Specific Parameters								
96 to n	Vendor-Specific Parameter Bytes							

The peripheral qualifier and peripheral device-type fields identify the device currently connected to the logical unit. If the target is not capable of supporting a device on this logical unit, this field shall be set to 7Fh (peripheral qualifier set to 011b and peripheral device type set to 1Fh). The peripheral qualifier is defined in Table 7-16 and the peripheral device type is defined in Table 7-17.

Table 7-16: Peripheral Qualifier

Qualifier	Description
000b	The specified peripheral device type is currently connected to this logical unit. If the target cannot determine whether or not a physical device is currently connected it shall also use this peripheral qualifier when returning the INQUIRY data. Note: This peripheral qualifier does not imply that the device is ready for access by the initiator.
001b	The target is capable of supporting the specified peripheral device type on this logical unit, however the physical device is not currently connected to this logical unit.
010b	Reserved
011b	The target is not capable of supporting a physical device on this logical unit. For this peripheral qualifier the peripheral device type shall be set to 1Fh to provide compatibility with previous versions of SCSI. All other peripheral device type values are reserved for this peripheral qualifier.
1XXb	Vendor specific

Table 7-17: Peripheral Device Type

Code	Description
00h	Direct-access device (e.g., magnetic disk)
01h	Sequential-access device (e.g., magnetic tape)
02h	Printer device
03h	Processor device
04h	Write-once read-multiple device (e.g., some optical disks)
05h	CD-ROM device
06h	Scanner device
07h	Optical memory device (e.g., some optical disks)
08h	Medium Changer device (e.g., jukeboxes)
09h	Communications device
0Ah - 0Bh	Defined by ASC IT8 (Graphics/Pre-Press Devices)
0Ch - 1Eh	Reserved
1Fh	Unknown or no device type

A removable medium (RMB) bit of zero indicates that the medium is not removable. A RMB bit of one indicates that the medium is removable.

The device-type modifier field was defined in SCSI-1 to permit vendor-specific qualification codes of the device type. This field is retained for compatibility with SCSI-1. Targets that do not support this field should return a value of zero.

The usage of non-zero code values in the ISO version and ECMA version fields are defined by the International Organization for Standardization and the European Computer Manufacturers Association, respectively. A zero code value in these fields shall indicate that the target does not claim compliance to the ISO version of SCSI (ISO IS 9316) or the ECMA version of SCSI (ECMA-111). It is possible to claim compliance to more than one of these SCSI standards.

The ANSI-approved version field indicates the implemented version of this standard and is defined in Table 7-18.

Table 7-18: ANSI-Approved Version

Code	Description
0h	The device might or might not comply to an ANSI-approved standard.
1h	The device complies to ANSI X3.131-1986 (SCSI-1).
2h	The device complies to this version of SCSI. This code is reserved to designate this standard upon approval by ANSI.
3h - 7h	Reserved

The asynchronous event notification capability (AENC) bit indicates that the device supports the asynchronous event notification capability as defined in 6.5.5.

- (1) Processor device-type definition: An AENC bit of one indicates that the processor device is capable of accepting asynchronous event notifications. An AENC bit of zero indicates that the processor device does not support asynchronous event notifications.
- (2) All other device-types: This bit is reserved.

A terminate I/O process (TrmIOP) bit of one indicates that the device supports the TERMINATE I/O PROCESS message as defined in 5.6.22. A value of zero indicates that the device does not support the TERMINATE I/O PROCESS message.

A response data format value of zero indicates the INQUIRY data format is as specified in SCSI-1. A response data format value of one indicates compatibility with some products that were designed prior to the development of this standard (i.e., CCS). A response data format value of two indicates that the data shall be in the format specified in this standard. Response data format values greater than two are reserved.

The additional length field shall specify the length in bytes of the parameters. If the allocation length of the command descriptor block is too small to transfer all of the parameters, the additional length shall not be adjusted to reflect the truncation.

A relative addressing (RelAdr) bit of one indicates that the device supports the relative addressing mode for this logical unit. If this bit is set to one the linked command (Linked) bit shall also be set to one since relative addressing can only be used with linked commands. A RelAdr bit of zero indicates the device does not support relative addressing for this logical unit.

A wide bus 32 (WBus32) bit of one indicates that the device supports 32-bit wide data transfers. A value of zero indicates that the device does not support 32-bit wide data transfers.

A wide bus 16 (WBus16) bit of one indicates that the device supports 16-bit wide data transfers. A value of zero indicates that the device does not support 16-bit wide data transfers.

NOTE: If the values of both the WBus16 and WBus32 bits are zero the device only supports 8-bit wide data transfers.

A synchronous transfer (Sync) bit of one indicates that the device supports synchronous data transfer. A value of zero indicates the device does not support synchronous data transfer.

A linked command (Linked) bit of one indicates that the device supports linked commands for this logical unit. A value of zero indicates the device does not support linked commands for this logical unit.

A command queuing (CmdQue) bit of one indicates that the device supports tagged command queuing for this logical unit. A value of zero indicates the device does not support tagged command queuing for this logical unit.

A soft reset (SftRe) bit of zero indicates that the device responds to the RESET condition with the hard RESET alternative (see 5.2.2.1). A SftRe bit of one indicates that the device responds to the RESET condition with the soft RESET alternative (see 5.2.2.2).

ASCII data fields shall contain only graphic codes (i.e., code values 20h through 7Eh). Left-aligned fields shall place any unused bytes at the end of the field (highest offset) and the unused bytes shall be filled with space characters (20h). Right-aligned fields shall place any unused bytes at the start of the field (lowest offset) and the unused bytes shall be filled with space characters (20h).

The vendor identification field contains eight bytes of ASCII data identifying the vendor of the product. The data shall be left aligned within this field.

NOTE: It is intended that this field provide a unique vendor identification of the manufacturer of the SCSI device. In the absence of a formal registration procedure, X3T9.2 maintains a list of vendor identification codes in use. Vendors are requested to voluntarily submit their identification codes to X3T9.2 to prevent duplication of codes (see Appendix J).

The product identification field contains sixteen bytes of ASCII data as defined by the vendor. The data shall be left-aligned within this field.

The product revision level field contains four bytes of ASCII data as defined by the vendor. The data shall be left-aligned within this field.

7.2.5.2. Vital Product Data

Implementation of vital product data is optional. The information returned consists of configuration data (e.g., vendor identification, product identification, model, serial number), manufacturing data (e.g., plant and date of manufacture), field replaceable unit data and other vendor- or device-specific data.

The initiator requests the vital product data information by setting the EVPD bit to one and specifying the page code of the desired vital product data (see Table 7-72). If the target does not implement the requested page it shall return CHECK CONDITION status. The a sense key shall be set to ILLEGAL REQUEST and the additional sense code shall be set to INVALID FIELD IN CDB.

IMPLEMENTORS NOTES:

(1) It is recommended that the target have the ability to execute the INQUIRY command even when a device error occurs which prohibits normal command execution. In such a case, CHECK CONDITION status would be returned for commands other than INQUIRY or REQUEST SENSE. The sense data returned may contain the field replaceable unit code. The vital product data would be obtained for the failing device using the INQUIRY command.

(2) This standard defines a format which allows device-independent initiator software to display the vital product data returned by the INQUIRY command. For example, the initiator may display the data associated for the field replaceable unit returned in the sense data. The contents of the data may be vendor-specific; therefore it may not be usable without detailed information about the device.

(3) This standard does not define the location or method of storing the vital product data. The retrieval of the data may require completion of initialization operations within the device which may induce delays before the data is available to the initiator. Time-critical requirements are an implementation consideration and are not addressed in this standard.

7.2.6. LOG SELECT Command

Table 7-19: LOG SELECT Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (4Ch)							
1	Logical Unit Number			Reserved			PCR	SP
2	PC		Reserved					
3	Reserved							
4	Reserved							
5	Reserved							
6	Reserved							
7	(MSB)							
8	Parameter List Length						(LSB)	
9	Control							

The LOG SELECT command (Table 7-19) provides a means for the initiator to manage statistical information maintained by the device about the device or its logical units. Targets that implement the LOG SELECT command shall also implement the LOG SENSE command. Structures in the form of log parameters within log pages are defined as a way to manage the log data. The LOG SELECT command provides for sending zero or more log pages during a DATA OUT phase. This standard defines the format of the log pages, but does not define the exact conditions and events which are logged.

A parameter code reset (PCR) bit of one and a parameter list length of zero shall cause all implemented parameters to be set to the target-defined default values (e.g., zero). If the PCR bit is one and a the parameter list length is greater than zero the command is terminated with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST and the additional sense code shall be set to INVALID FIELD IN CDB. A PCR bit of zero specifies that the log parameters shall not be reset.

A save parameters (SP) bit of one indicates that after performing the specified LOG SELECT operation the target shall save to non-volatile memory all parameters identified as savable by the DS bit in the log page (see 7.3.2). A SP bit of zero specifies that parameters shall not be saved.

Saving of log parameters is optional and indicated for each log parameter by the DS bit in the page. Log parameters may be saved at vendor-specific times

subject to the TSD bit (see 7.3.2) in the log parameter. If the target does not implement saved parameters for any log parameter and the SP bit is set to one, the command shall be terminated with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST, and the additional sense code set to INVALID FIELD IN CDB.

It is not an error to set the SP bit to one and to set the DS bit of a log parameter to one. In this case, the parameter value for that log parameter is not saved.

The page control (PC) field defines the type of parameter values to be selected. The page control field is defined in Table 7-20.

Table 7-20: Page Control Field

Value	Type of Parameter Values	
	LOG SENSE	LOG SELECT
00b	Current Threshold Values	Threshold Values
01b	Current Cumulative Values	Cumulative Values
10b	Default Threshold Values	Default Threshold Values
11b	Default Cumulative Values	Default Cumulative Values

The current cumulative values may be updated by the target or by the initiator using the LOG SELECT command to reflect the cumulative number of events experienced by the target. Fields in the parameter control byte (7.3.2) of each log parameter control the updating and saving of the current cumulative parameters.

The target shall set the current threshold parameters to the default threshold values in response to a LOG SELECT command with the PC field set to 10b and the parameter list length field set to zero.

The target shall set all cumulative parameters to their default values in response to a LOG SELECT command with the PC field set to 11b and the parameter list length field set to zero.

The current threshold value can only be modified by the initiator via the LOG SELECT command. If the initiator attempts to change current threshold values that are not available or not implemented for that log parameter, then the target shall terminate the LOG SELECT command with CHECK CONDITION status, the sense key set to ILLEGAL REQUEST, and the additional sense code set to INVALID FIELD IN PARAMETER LIST. The saving of current threshold parameters and the criteria for the current threshold being met are controlled by bits in the parameter control byte (7.3.2).

IMPLEMENTORS NOTE: Pages or log parameters that are not available may become available at some later time (e.g., after the device has become ready).

The parameter list length field specifies the length in bytes of the parameter list that shall be transferred from the initiator to the target during the DATA OUT phase. A parameter list length of zero indicates that no pages shall be transferred. This condition shall not be considered an error. If the initiator sends page codes or parameter codes within the parameter list that are reserved or not implemented by the target, the target shall terminate the LOG SELECT command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST and the additional sense code set to INVALID FIELD IN PARAMETER LIST.

If a parameter list length results in the truncation of any log parameter, the target shall terminate the command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST and the additional sense code set to INVALID FIELD IN CDB.

The initiator should send pages in ascending order by page code value if multiple pages are sent during a DATA OUT phase. If multiple log parameters within a page are sent during the DATA OUT phase then they should be sent in ascending order by parameter code value. The target shall return CHECK CONDITION status if the initiator sends pages out of order or parameter codes out of order. The sense key shall be set to ILLEGAL REQUEST and the additional sense code set to INVALID FIELD IN PARAMETER LIST.

IMPLEMENTORS NOTE: Initiators should issue LOG SENSE commands prior to issuing LOG SELECT commands to determine supported pages and page lengths.

The target may provide independent sets of log parameters for each logical unit or for each combination of logical units and initiators. If the target does not support independent sets of log parameters and any log parameters are changed that affect other initiators, then the target shall generate an unit attention condition for all initiators except the one that issued the LOG SELECT command (see 6.9). This unit attention condition is returned with an additional sense code of LOG PARAMETERS CHANGED.

If the initiator sends a log parameter that is not supported by the target, the target shall terminate the command with CHECK CONDITION status, set the sense key to ILLEGAL REQUEST, and set the additional sense code to INVALID FIELD IN PARAMETER LIST.

7.2.7. LOG SENSE Command

Table 7-21: LOG SENSE Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (4Dh)							
1	Logical Unit Number			Reserved			PPC	SP
2	PC		Page Code					
3	Reserved							
4	Reserved							
5	(MSB)							
6	Parameter Pointer						(LSB)	
7	(MSB)							
8	Allocation Length						(LSB)	
9	Control							

The LOG SENSE command (Table 7-21) provides a means for the initiator to retrieve statistical information maintained by the device about the device or its logical units. It is a complementary command to the LOG SELECT command.

The parameter pointer control (PPC) bit controls the type of parameters requested from the target:

(1) A PPC bit of one indicates that the target shall return a log page with parameter code values which have changed since the last LOG SELECT or LOG SENSE command. The target shall return only those parameter codes following the parameter pointer field. (See 7.3.2 for a description of the time ordering of list parameters.) The number returned shall be a 16-bit unsigned integer.

(2) A PPC bit of zero indicates that the log parameter requested from the target shall begin with the parameter code specified in the parameter pointer field and return the number of bytes specified by the allocation length field in ascending order of parameter codes from the specified log page. A PPC bit of zero and a parameter pointer field of zero shall cause all available log parameters for the specified log page to be returned to the initiator subject to the specified allocation length.

Saving parameters is an optional function of the LOG SENSE command. If the target does not implement saving log parameters and if the save parameters (SP) bit is one, then the target shall return CHECK CONDITION status, set the sense key to ILLEGAL REQUEST, and set the additional sense code to INVALID FIELD IN CDB.

A SP bit of zero indicates the target shall perform the specified LOG SENSE command and shall not save any log parameters. If saving log parameters is implemented, a SP bit of one indicates that the target shall perform the specified LOG SENSE command and shall save all log parameters identified as savable by the DS bit (7.3.2) to a non-volatile vendor-specific location.

The page control (PC) field defines the type of parameter values to be selected (see 7.2.6 for the definition of the page control field). The parameter values returned by a LOG SENSE command are determined as follows:

(1) The specified parameter values at the last update (in response to a LOG SELECT or LOG SENSE command or done automatically by the target for cumulative values).

(2) The saved values if an update has not occurred since the last power-on, hard RESET condition, or BUS DEVICE RESET message and saved parameters are implemented.

(3) The default values if an update has not occurred since the last power-on, hard RESET condition, or BUS DEVICE RESET message and saved values are not available or not implemented.

The page code field identifies which page of data is being requested (see Table 7-53). If the page code is reserved or not implemented, the target shall terminate the command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST with the additional sense code set to INVALID FIELD IN CDB.

The parameter pointer field allows the initiator to request parameter data beginning from a specific parameter code to the maximum allocation length or the maximum parameter code supported by the target, whichever is less. If the value of the parameter pointer field is larger than the largest available parameter code that can be returned by the target on the specified page, the target shall terminate the command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST and the additional sense code shall be set to INVALID FIELD IN CDB.

Log parameters within the specified log page shall be transferred in ascending order according to parameter code.

7.2.8. MODE SELECT(6) Command

Table 7-22: MODE SELECT(6) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (15h)							
1	Logical Unit Number			PF	Reserved			SP
2	Reserved							
3	Reserved							
4	Parameter List Length							
5	Control							

The MODE SELECT(6) command (Table 7-22) provides a means for the initiator to specify medium, logical unit, or peripheral device parameters to the target. Targets that implement the MODE SELECT command shall also implement the MODE SENSE command. Initiators should issue MODE SENSE prior to MODE SELECT to determine supported pages, page lengths, and other parameters.

If a target supports saved pages, it may save only one copy of the page for each logical unit and have it apply to all initiators or it may save separate copies for each initiator for each logical unit. If separate copies are saved, the target shall maintain separate current values for each I T L nexus. Pages which are common to all initiators are not required to have multiple copies.

If an initiator sends a MODE SELECT command that changes any parameters that apply to other initiators, the target shall generate a unit attention condition for all initiators except the one that issued the MODE SELECT command (see 6.9). The target shall set the additional sense code to MODE PARAMETERS CHANGED.

The target may provide for independent sets of parameters for each attached logical unit or for each combination of logical unit and initiator. If independent sets of parameters are implemented, and a third party reservation is requested the target transfers the set of parameters in effect for the initiator of the RESERVE command to the parameters used for commands from the third party device (see 8.2.12.3 and 9.2.10.1).

A page format (PF) bit of zero indicates that the MODE SELECT parameters are as specified in SCSI-1 (i.e., all parameters after the block descriptors are vendor-specific. A PF bit of one indicates that the MODE SELECT parameters following the header and block descriptor(s) are structured as pages of related parameters and are as specified in this standard.

A save pages (SP) bit of zero indicates the target shall perform the specified MODE SELECT operation, and shall not save any pages. A SP bit of one indicates that the target shall perform the specified MODE SELECT operation, and shall save to a non-volatile vendor-specific location all the savable pages including any sent during the DATA OUT phase. The SP bit is optional, even when mode pages are supported by the target. Pages which are saved are identified by the PS bit that is returned in the page header by the MODE SENSE command (see 7.2.10.4). If the target does not implement saved pages and the SP bit is set to one, the command shall be terminated with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST, and the additional sense code shall be set to INVALID FIELD IN CDB.

The parameter list length field specifies the length in bytes of the MODE SELECT parameter list that shall be transferred from the initiator to the target during the DATA OUT phase. A parameter list length of zero indicates that no data shall be transferred. This condition shall not be considered as an error. A parameter list length that results in the truncation of any descriptor, header or page of parameters shall cause the target to terminate the command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST, and the additional sense code shall be set to PARAMETER LIST LENGTH ERROR.

The parameter list for the MODE SELECT and MODE SENSE commands is defined in 7.3.3. Parts of each parameter list are uniquely defined for each device-type.

IMPLEMENTORS NOTE: In some situations where there is a conflict between the information in the command descriptor block and the information in the parameter list, one of several additional sense codes may apply. As a guide, INVALID FIELD IN CDB or PARAMETER LIST LENGTH ERROR should be used if the error is detected prior to any operations that alter the logical unit's mode parameters. INVALID FIELD IN PARAMETER LIST should be used if the logical unit's mode parameters have been altered.

The target shall terminate the MODE SELECT command with CHECK CONDITION status, set the sense key to ILLEGAL REQUEST and set the additional sense code to INVALID FIELD IN PARAMETER LIST for the following conditions:

- (1) If the initiator attempts to change any field that is not changeable as reported by the target. In this case, no parameters shall be changed by this command.
- (2) If the initiator attempts to send an unsupported value or a non-zero value to a reserved field in the MODE SELECT header, block descriptor, or any page header.
- (3) If an initiator attempts to send a page with a length not equal to the parameter length reported for that page by the MODE SENSE command.
- (4) If the initiator attempts to send a value for a parameter that is outside the range supported by the target and rounding is not implemented for that parameter.

If the initiator sends a value for a parameter that is outside the range supported by the target and rounding is implemented for that parameter, the target may either:

- (1) round the parameter to an acceptable value and terminate the command as described in 6.5.4.
- (2) terminate the command with CHECK CONDITION status, the sense key set to ILLEGAL REQUEST, and the additional sense code set to INVALID FIELD IN PARAMETER LIST.

7.2.9. MODE SELECT(10) Command

Table 7-23: MODE SELECT(10) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (55h)							
1	Logical Unit Number			PF	Reserved			SP
2	Reserved							
3	Reserved							
4	Reserved							
5	Reserved							
6	Reserved							
7	(MSB)							
8	Parameter List Length							(LSB)
9	Control							

The MODE SELECT(10) command (Table 7-23) provides a means for the initiator to specify medium, logical unit, or peripheral device parameters to the target. See the MODE SELECT(6) command (7.2.8) for a description of the fields in this command. Initiators should issue MODE SENSE prior to MODE SELECT to determine supported pages, page lengths, and other parameters.

Targets that implement the MODE SELECT(10) command shall also implement the MODE SENSE(10) command.

7.2.10. MODE SENSE(6) Command

Table 7-24: MODE SENSE(6) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (1Ah)							
1	Logical Unit Number			Reserved	DBD	Reserved		
2	PC		Page Code					
3	Reserved							
4	Allocation Length							
5	Control							

The MODE SENSE(6) command (Table 7-24) provides a means for a target to report parameters to the initiator. It is a complementary command to the MODE SELECT(6) command.

A disable block descriptors (DBD) bit of zero indicates that the target may return zero or more block descriptors in the returned MODE SENSE data (see 7.3.3), at the target's discretion. A DBD bit of one specifies that the target shall not return any block descriptors in the returned MODE SENSE data.

The page control (PC) field defines the type of parameter values to be returned. The page control field is defined in Table 7-25.

Table 7-25: Page Control Field

Bit 7	Bit 6	Type of Parameter Values	Section
0	0	Current Values	7.2.10.1
0	1	Changeable Values	7.2.10.2
1	0	Default Values	7.2.10.3
1	1	Saved Values	7.2.10.4

The page code specifies which page or pages to return. Page code usage is defined in Table 7-26.

Table 7-26: Page Code Usage for All Devices

Page Code	Description
00h	Vendor Specific (does not require page format)
01h - 1Fh	See specific device-types
20h - 3Eh	Vendor Specific (page format required)
3Fh	Return all pages

An initiator may request any one or all of the supported pages of a target. If an initiator attempts a MODE SENSE command with a page code value not implemented by the target, the target shall return CHECK CONDITION status and shall set the sense key to ILLEGAL REQUEST and the additional sense code to INVALID FIELD IN CDB.

A page code of 3Fh indicates that all pages implemented by the target shall be returned to the initiator. Page 00h, if implemented, shall be returned after all other pages. Targets that implement more than 256 bytes of mode page parameter data and block descriptors shall return CHECK CONDITION status to a MODE SENSE(6) request of 3Fh in the page code field. The sense key shall be set to ILLEGAL REQUEST and the additional sense code shall be set to INVALID FIELD IN CDB.

IMPLEMENTORS NOTE: Targets that support more than 256 bytes of block descriptors and parameter pages should consider implementing the MODE SELECT(10) and MODE SENSE(10) commands. Targets might not support an entire mode page as defined in this standard. Typically, the mode pages are returned in ascending order (except for page 00h).

The parameter list for MODE SELECT and MODE SENSE is defined in 7.3.3. Parts of the parameter lists are specifically defined for each device type.

7.2.10.1. Current Values

A PC field value of 0h requests the target to return the current parameter values for the specified page code for the logical unit. The current values returned are:

- (1) The parameters set in the last successful MODE SELECT command.
- (2) The saved values if a MODE SELECT command has not been executed since the last power-on, hard RESET condition, or BUS DEVICE RESET message.
- (3) The default values if saved values are not available or not supported.

7.2.10.2. Changeable Values

A PC field value of 1h requests the target to return the changeable parameter mask for the page code specified. The page requested shall be returned containing information that indicates which parameters are changeable. All bits of parameters that are changeable shall be set to one. All bits of parameters that are target defined (not changeable by the initiator) shall be set to zero.

NOTE: An attempt to change a target-defined parameter (via MODE SELECT) results in an error condition (see 7.2.8).

IMPLEMENTORS NOTE: The initiator should issue a MODE SENSE command with the PC field set to 1h and the page code field set to 3Fh to determine which pages are supported, which parameters within the pages are changeable, and the supported length of each page prior to issuing any MODE SELECT commands.

7.2.10.3. Default Values

A PC field value of 2h requests the target return the default values for the page code specified. The page requested shall be returned with each supported parameter set to its default value. Parameters not supported by the target shall be set to zero.

7.2.10.4. Saved Values

A PC field value of 3h requests the target return the saved values for the page code specified. Implementation of saved page parameters is optional. The page requested shall be returned with the parameters set to their saved values. Parameters not supported by the target shall be set to zero. If saved values are not implemented, the command shall be terminated with CHECK CONDITION status with the sense key set to ILLEGAL REQUEST and the additional sense code set to SAVING PARAMETERS NOT SUPPORTED.

IMPLEMENTORS NOTE: The method of saving parameters is vendor-specific. The parameters are preserved in such a manner that they are retained when the target is powered down. All savable pages can be considered saved when a MODE SELECT command issued with the SP bit set to one has returned a GOOD status or after the successful completion of a FORMAT UNIT command.

7.2.10.5. Initial Responses

After a power-up condition or hard reset condition, the target shall respond in the following manner:

- (1) If default values are requested, report the default values.
- (2) If saved values are requested, report valid restored parameters, or restore the parameters and report them. If the saved parameters are not able to be accessed from the non-volatile vendor-specific location, terminate the command with CHECK CONDITION status and the sense key set to NOT READY. If saved parameters are not implemented respond as defined in 7.2.10.4.
- (3) If current values are requested and current values have not been sent by the initiator (via a MODE SELECT command), the target may return either the default or saved parameters (if implemented), as defined above. If current values have been sent, the current values shall be reported.

7.2.11. MODE SENSE(10) Command

Table 7-27: MODE SENSE(10) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (5Ah)							
1	Logical Unit Number			Reserved	DBD	Reserved		
2	PC		Page Code					
3	Reserved							
4	Reserved							
5	Reserved							
6	Reserved							
7	(MSB)							
8	Allocation Length						(LSB)	
9	Control							

The MODE SENSE(10) command (Table 7-27) provides a means for a target to report parameters to the initiator. It is a complementary command to the MODE SELECT(10) command. If the MODE SELECT(10) command is implemented the MODE SENSE(10) command shall be implemented. See the MODE SENSE(6) command for a description of the fields in this command.

7.2.12. READ BUFFER

Table 7-28: READ BUFFER Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (3Ch)							
1	Logical Unit Number			Reserved		Mode		
2	Buffer ID							
3	(MSB)							
5	Buffer Offset							
6	(MSB)							
8	Allocation Length							
9	(LSB)							
	Control							

The READ BUFFER command (Table 7-28) is used in conjunction with the WRITE BUFFER command as a diagnostic function for testing target memory and the SCSI bus integrity. This command shall not alter the medium.

The function of this command and the meaning of fields within the command descriptor block depend on the contents of the mode field. The mode field is defined in Table 7-29.

Table 7-29: READ BUFFER Mode Field

Mode	Description	Type
000b	Combined header and data	Optional
001b	Vendor-specific	Vendor specific
010b	Data	Optional
011b	Descriptor	Optional
100b	Reserved	Reserved
101b	Reserved	Reserved
110b	Reserved	Reserved
111b	Reserved	Reserved

IMPLEMENTORS NOTE: Modes 000b and 001b are included for compatibility with products that were designed prior to the generation of this standard. Some products that were designed prior to the generation of this standard restrict the available length to 65535 bytes.

7.2.12.1. Combined Header and Data Mode (000b)

In this mode, a four-byte header followed by data bytes are returned to the initiator during the DATA IN phase. The buffer ID and the buffer offset fields are reserved.

The four-byte READ BUFFER header (Table 7-30) is followed by data bytes from the target's data buffer.

Table 7-30: READ BUFFER Header

Bit Byte	7	6	5	4	3	2	1	0
0	Reserved							
1	(MSB)							
3	Buffer Capacity							
	(LSB)							

The buffer capacity field specifies the total number of data bytes that are available in the target's data buffer. This number is not reduced to reflect the allocation length nor is it reduced to reflect the actual number of bytes written using the WRITE BUFFER command. Following the READ BUFFER header, the target shall transfer data from its data buffer. The target terminates the DATA IN phase when allocation length bytes of header plus data have been transferred or when all available header and buffer data have been transferred to the initiator, whichever is less.

7.2.12.2. Vendor-Specific Mode (001b)

In this mode, the meaning of the buffer ID, buffer offset, and allocation length fields are not specified by this standard.

7.2.12.3. Data Mode (010b)

In this mode, the DATA IN phase contains buffer data. The buffer ID field identifies a specific buffer within the target from which data shall be transferred. The vendor assigns buffer ID codes to buffers within the target. Buffer ID zero shall be supported. If more than one buffer is supported, additional buffer ID codes shall be assigned contiguously, beginning with one. Buffer ID code assignments for the READ BUFFER command shall be the same as for the WRITE BUFFER command. If an unsupported buffer ID code is selected, the target shall return CHECK CONDITION status, shall set the sense key to ILLEGAL REQUEST and set the additional sense code to ILLEGAL FIELD IN CDB. The target terminates the DATA IN phase when allocation length bytes have been transferred or when all the available data from the buffer has been transferred to the initiator, whichever is less.

The buffer offset field contains the byte offset within the specified buffer where data shall be transferred from. The initiator should conform to the offset boundary requirements returned in the READ BUFFER descriptor (see 7.2.12.4). If the target is unable to accept the specified buffer offset, it shall return CHECK CONDITION status, shall set the sense key to ILLEGAL REQUEST and set the additional sense code to ILLEGAL FIELD IN CDB.

7.2.12.4. Descriptor Mode (011b)

In this mode, a maximum of four bytes of READ BUFFER descriptor information are returned. The target shall return the descriptor information for the buffer specified by the buffer ID (see the description of the buffer ID in 7.2.12.3). If there is no buffer associated with the specified buffer ID, the target shall return all zeros in the READ BUFFER descriptor. The buffer offset field is reserved in this mode. The allocation length should be set to four or greater. The target shall transfer the lesser of the allocation length or four bytes of READ BUFFER descriptor. The READ BUFFER descriptor is defined as shown in Table 7-31.

Table 7-31: READ BUFFER Descriptor

Bit Byte	7	6	5	4	3	2	1	0
0	Offset Boundary							
1	(MSB)							
3	Buffer Capacity (LSB)							

The offset boundary field returns the boundary alignment within the selected buffer for subsequent WRITE BUFFER and READ BUFFER commands. The value contained in the offset boundary field shall be interpreted as a power of two.

The value contained in the buffer offset field of subsequent WRITE BUFFER and READ BUFFER commands should be a multiple of 2^{offset boundary} as shown in Table 7-32.

Table 7-32: Buffer Offset Boundary

Offset Boundary	Offset Boundary	Buffer Offsets
0	2 ⁰ = 1	Byte boundaries
1	2 ¹ = 2	Even-byte boundaries
2	2 ² = 4	Four-byte boundaries
3	2 ³ = 8	Eight-byte boundaries
4	2 ⁴ = 16	16-byte boundaries
.	.	.
FFh	Not Applicable	0 is the only supported buffer offset.

The buffer capacity field shall return the size of the selected buffer in bytes.

IMPLEMENTORS NOTE: In a multi-tasking system, a buffer may be altered between the WRITE BUFFER and READ BUFFER commands by another task. Buffer testing applications may wish to insure that only a single task is active. Use of reservations (to all logical units on the device) or linked commands may also be helpful in avoiding buffer alteration between these two commands.

7.2.13. RECEIVE DIAGNOSTIC RESULTS Command**Table 7-33: RECEIVE DIAGNOSTIC RESULTS Command**

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (1Ch)							
1	Logical Unit Number			Reserved				
2	Reserved							
3	(MSB)							
4	Allocation Length (LSB)							
5	Control							

The RECEIVE DIAGNOSTIC RESULTS command (Table 7-33) requests analysis data be sent to the initiator after completion of a SEND DIAGNOSTIC command (see 7.2.15). If the target supports the optional page format the page code field sent in the previous SEND DIAGNOSTIC command specifies the format of the returned data.

IMPLEMENTORS NOTES:

(1) To insure that the diagnostic command information is not destroyed by a command sent from another initiator, the SEND DIAGNOSTIC command should either be linked to the RECEIVE DIAGNOSTIC RESULTS command or the logical unit should be reserved.

(2) Although diagnostic software is generally device-specific, this command and the SEND DIAGNOSTIC command provide a means to isolate the operating system software from the device-specific diagnostic software. Hence the operating system can remain device-independent. This also allows diagnostic software to be more easily transferred to other operating systems.

See 7.3.1 for RECEIVE DIAGNOSTIC RESULTS page format definitions.

7.2.14. REQUEST SENSE Command

Table 7-34: REQUEST SENSE Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (03h)							
1	Logical Unit Number				Reserved			
2	Reserved							
3	Reserved							
4	Allocation Length							
5	Control							

The REQUEST SENSE command (Table 7-34) requests that the target transfer sense data to the initiator.

The sense data:

- (1) shall be available if the previous command to the specified I_T_x nexus terminated with CHECK CONDITION or COMMAND TERMINATED status
- (2) shall be available if other information (e.g., medium position) is available in any field
- (3) may be available if the previous command to the specified I_T_x nexus ended with an unexpected BUS FREE error (5.1.1).

If the target has no sense data available to return, it shall return a sense key of NO SENSE and an additional sense code of NO ADDITIONAL SENSE INFORMATION.

The sense data shall be preserved by the target for the initiator until retrieved by the REQUEST SENSE command or until the receipt of any other command for the same I_T_x nexus (see 6.6). Sense data shall be cleared upon receipt of any subsequent command (including REQUEST SENSE) to the same I_T_x nexus.

IMPLEMENTORS NOTE: Some target implementations do not update sense data except on commands that return CHECK CONDITION or COMMAND TERMINATED status. Thus when polling for a logical unit to become ready, the initiator should issue TEST UNIT READY commands until GOOD status is returned. If desired, the initiator may issue REQUEST SENSE commands after the TEST UNIT READY commands that return CHECK CONDITION or COMMAND TERMINATED status to obtain the sense data.

The target shall return CHECK CONDITION status for a REQUEST SENSE command only to report errors specific to the command itself. For example:

- (1) A non-zero reserved bit is detected in the command descriptor block.
- (2) An unrecovered parity error is detected on the data bus.
- (3) A target malfunction prevents return of the sense data.

If a recovered error occurs during the execution of the REQUEST SENSE command, the target shall return the sense data with GOOD status. If a target returns CHECK CONDITION status for a REQUEST SENSE command the sense data may be invalid.

IMPLEMENTORS NOTE: The sense data appropriate to the selection of an invalid logical unit is defined in 6.5.3.

Targets shall be capable of returning eighteen bytes of data in response to a REQUEST SENSE command. If the allocation length is eighteen or greater and a target returns less than eighteen bytes of data the initiator should assume that the bytes not transferred would have been zeros had the target returned those bytes. Initiators can determine how much sense data has been returned by examining the allocation length parameter in the command descriptor block and the additional sense length in the sense data. Targets shall not adjust the additional sense length to reflect truncation if the allocation length is less than the sense data available.

Table 7-35: Error Codes 70h and 71h Sense Data Format

Bit Byte	7	6	5	4	3	2	1	0						
0	Valid	Error Code (70h or 71h)												
1	Segment Number													
2	Filemark	EOM	ILI	Reserved	Sense Key									
3	(MSB)	Information						(LSB)						
6														
7	Additional Sense Length (n-7)													
8	(MSB)	Command-Specific Information						(LSB)						
11														
12	Additional Sense Code													
13	Additional Sense Code Qualifier													
14	Field Replaceable Unit Code													
15 to	SKSV	Sense-Key Specific												
17														
18 to	Additional Sense Bytes													
n														

A valid bit of zero indicates that the information field is not as defined in this standard. A valid bit of one indicates the information field contains valid information as defined in this standard. Targets shall implement the valid bit.

The sense data format for error codes 70h (current errors) and 71h (deferred errors) are defined in Table 7-35. Error code values of 72h to 7Eh are reserved. Error code 7Fh is for a vendor-specific sense data format. Targets shall implement error code 70h; implementation of error code 71h is optional. Error code values of 00h to 6Fh are not defined by this standard and their use is not recommended.

The segment number field contains the number of the current segment descriptor if the REQUEST SENSE command is in response to a COPY, COMPARE, or COPY AND VERIFY command. Up to 256 segments are supported beginning with segment zero.

The filemark bit is mandatory for sequential-access devices and this bit is reserved for all other device types. A filemark bit of one indicates that the current command has read a filemark or setmark. The additional sense code field may be used to indicate whether a filemark or setmark was read. Reporting of setmarks is optional and indicated by the RSmk bit for sequential-access devices in the configuration parameters page (see 9.3.3).

The end-of-medium (EOM) bit is mandatory for sequential-access and printer devices and this bit is reserved for all other device types. An EOM bit of one indicates that an end-of-medium condition (end-of-partition, beginning-of-partition, out-of-paper, etc.) exists. For sequential-access devices, this bit indicates that the unit is at or past the early-warning if the direction was forward or that the command could not be completed because beginning-of-partition was encountered if the direction was reverse.

An incorrect length indicator (ILI) bit of one usually indicates that the requested logical block length did not match the logical block length of the data on the medium.

The sense key, additional sense code and additional sense code qualifier provide a hierarchy of information. The intention of the hierarchy is to provide a top-down approach for an initiator to determine information relating to the error and exception conditions. The sense key provides generic categories in which error and exception conditions can be reported. Initiators would typically use sense keys for high level error recovery procedures. Additional sense codes provide further detail describing the sense key. Additional sense code qualifiers add further detail to the additional sense code. The additional sense code and additional sense code qualifier can be used by initiators where sophisticated error recovery procedures require detailed information describing the error and exception conditions.

The sense key field is mandatory and indicates generic information describing an error or exception condition. The sense keys are defined in Tables 7-39 and 7-40.

The contents of the information field is device-type or command specific and is defined within the appropriate section for the device type or command of interest. Targets shall implement the information field. Unless specified otherwise, this field contains:

(1) The unsigned logical block address associated with the sense key, for direct-access devices (Device Type 0), write-once devices (Device Type 4), CD-ROM devices (Device Type 5), and optical memory devices (Device Type 7).

(2) The difference (residue) of the requested length minus the actual length in either bytes or blocks, as determined by the command, for sequential-access devices (Device Type 1), printer devices (Device Type 2), processor devices (Device Type 3) and some direct access device commands, except as defined for (4) below. (Negative values are indicated by two's complement notation.)

(3) The difference (residue) of the requested number of blocks minus the actual number of blocks copied or compared for the current segment descriptor of a COPY, COMPARE, or COPY AND VERIFY command.

(4) For sequential-access devices operating in buffered modes 1h or 2h that detect an unrecoverable write error when unwritten data blocks, filemarks, or setmarks remain in the buffer, the value of the information field for all commands shall be:

- (a) the total number of data blocks, filemarks, and setmarks in the buffer if the device is in fixed block mode (block length field of the MODE SENSE block descriptor is non-zero and the fixed bit of the WRITE command is one).
- (b) the number of bytes in the buffer, including filemarks and setmarks, if the device is in variable mode (the fixed bit of the WRITE command is zero).

The additional sense length field indicates the number of additional sense bytes to follow. If the allocation length of the command descriptor block is too small to transfer all of the additional sense bytes, the additional sense length is not adjusted to reflect the truncation.

The command-specific information field contains information that depends on the command which was executed. Further meaning for this field is defined within the command description. The command-specific information field is mandatory if the target supports any of the following commands: COPY, COMPARE, COPY AND VERIFY, SEARCH DATA, and REASSIGN BLOCKS.

The additional sense code field indicates further information related to the error or exception condition reported in the sense key field. Targets shall support the additional sense code field. Support of the additional sense codes not explicitly required by this standard is optional. A list of additional sense codes is in Table 7-41. If the target does not have further information related to the error or exception condition, the additional sense code is set to NO ADDITIONAL SENSE INFORMATION.

The additional sense code qualifier indicates detailed information related to the additional sense code. The additional sense code qualifier is optional. If the error or exception condition is reportable by the device the value returned shall be as specified in Table 7-41. If the target does not have detailed information related to the error or exception condition, the additional sense code qualifier is set to zero.

Non-zero values in the field replaceable unit code field are used to define a device-specific mechanism or unit that has failed. A value of zero in this field shall indicate that no specific mechanism or unit has been identified to have failed or that the data is not available. The field replaceable unit code field is optional. The format of this information is not specified by this standard. Additional information about the field replaceable unit may be available in the ASCII information page (see 7.3.4.2), if supported by the target.

The sense-key specific bytes are described in 7.2.14.1, below.

The additional sense bytes field may contain command specific data, peripheral device specific data, or vendor-specific data that further defines the nature of the CHECK CONDITION status.

7.2.14.1. Sense-Key Specific

The sense-key specific field is defined by this standard when the value of the sense-key specific valid (SKSV) bit is one. The sense-key specific valid bit and sense-key specific field are optional. The definition of this field is determined by the value of the sense key field. This field is reserved for sense keys not described below. An SKSV value of zero indicates that this field is not defined by this standard.

If the sense key field is set to ILLEGAL REQUEST and the SKSV bit is set to one, the sense-key specific field shall be defined as shown in Table 7-36. These Field Pointer field indicates which illegal parameters in command descriptor blocks or data parameters are in error.

Table 7-36: Field Pointer Bytes

Bit Byte	7	6	5	4	3	2	1	0
15	SKSV	C/D	Reserved	Reserved	BPV	Bit Pointer		
16	(MSB)							
17	Field Pointer							(LSB)

A command data (C/D) bit of one indicates that the illegal parameter is in the command descriptor block. A C/D bit of zero indicates that the illegal parameter is in the data parameters sent by the initiator during the DATA OUT phase.

A bit pointer valid (BPV) bit of zero indicates that the value in the bit pointer field is not valid. A BPV bit of one indicates that the bit pointer field specifies which bit of the byte designated by the field pointer field is in error. When a multiple-bit field is in error, the bit pointer field shall point to the most-significant (left-most) bit of the field.

The field pointer field indicates which byte of the command descriptor block or of the parameter data was in error. Bytes are numbered starting from zero, as shown in the tables describing the commands and parameters. When a multiple-byte field is in error, the pointer shall point to the most-significant (left-most) byte of the field.

IMPLEMENTORS NOTE: Bytes identified as being in error are not necessarily the place that has to be changed to correct the problem.

If the sense key is RECOVERED ERROR, HARDWARE ERROR or MEDIUM ERROR and if the SKSV bit is one, the sense-key specific field shall be defined as shown in Table 7-37.

Table 7-37: Actual Retry Count Bytes

Bit	7	6	5	4	3	2	1	0
Byte	15	SKSV	Reserved					
16	(MSB)							
17	Actual Retry Count							
	(LSB)							

The actual retry count field returns implementation-specific information on the actual number of retries of the recovery algorithm used in attempting to recover an error or exception condition.

IMPLEMENTORS NOTE: It is recommended that this field relate to the retry count fields within the error recovery page of the MODE SELECT command.

If the sense key is NOT READY and the SKSV bit is one, the sense-key specific field shall be defined as shown in Table 7-38. These fields are only defined for the FORMAT UNIT command with the Immed bit set to one.

Table 7-38: Format Progress Indication Bytes

Bit	7	6	5	4	3	2	1	0
Byte								
15	SKSV	Reserved						
16	(MSB)	Progress Indication						
17		(LSB)						

The progress indication field is a percent complete indication in which the returned value is the numerator that has 65536 (10000h) as its denominator. The progress indication shall be based upon the total format operation including any certification or initialization operations.

IMPLEMENTORS NOTE: It is intended that the progress indication be time related. However, since format time varies with the number of defects encountered, etc., it is reasonable for the target to assign values to various steps within the process. The granularity of these steps should be small enough to provide reasonable assurances to the initiator that progress is being made.

7.2.14.2. Deferred Errors

Error code 70h indicates that the CHECK CONDITION or COMMAND TERMINATED status returned is the result of an error or exception condition on the command that returned the CHECK CONDITION or COMMAND TERMINATED status or an unexpected bus free condition. This includes errors generated during execution of the command by the actual execution process. It also includes errors not related to any command that are first observed during execution of a command. Examples of this latter type of error include disk servo-mechanism off-track errors and power-up test errors.

Error Code 71h (deferred error) indicates that the CHECK CONDITION status returned is the result of an error or exception condition that occurred during execution of a previous command for which GOOD status has already been returned. Such commands are associated with use of the immediate bit, with some forms of caching, and with multiple command buffering. Targets that implement these features are required to implement deferred error reporting.

The deferred error indication may be sent at a time selected by the target through the asynchronous event notification process (see 6.5.5) if AEN is supported by both the initiator and target.

If AEN is not supported, the deferred error may be indicated by returning CHECK CONDITION status to the appropriate initiator as described below. The subsequent execution of a REQUEST SENSE command shall return the deferred error sense information.

If CHECK CONDITION status for a deferred error is returned, the current command has not performed any storage operations or output operations to the media. After the target detects a deferred error condition on a logical unit, it shall return a deferred error according to the rules described below:

- (1) If a deferred error can be recovered with no external system intervention, a deferred error indication shall not be posted unless required by the error handling parameters of the MODE SELECT command. The occurrence of the error may be logged if statistical or error logging is supported.
- (2) If a deferred error can be associated with a causing initiator and with a particular function or a particular subset of data, and the error is either unrecovered or required to be reported by the mode parameters, a deferred error indication shall be returned to the causing initiator. If an initiator other than the causing initiator attempts access to the particular function or subset of data associated with the deferred error, a BUSY status shall be returned to that initiator in response to the command attempting the access.

NOTE: Not all devices may be sufficiently sophisticated to identify the function or data that has failed. Those that cannot should treat the error in the following manner.

- (3) If a deferred error cannot be associated with a causing initiator or with a particular subset of data, a deferred error indication shall be returned on behalf of the failing logical unit to each initiator. If multiple deferred errors have accumulated for some initiators, only the last error shall be returned.
- (4) If a deferred error cannot be associated with a particular logical unit, it shall be returned to the appropriate initiator for all logical units supported by the target.

(5) If a current command has not yet started executing, and a deferred error occurs, the command shall be terminated with CHECK CONDITION status and deferred error information posted in the sense data. By convention, the current command is considered to have started execution if the target has changed phase from the COMMAND phase to the next normal phase of the command sequence. If a deferred error occurs while a current command is executing and the current command has been affected by the error, the command shall be terminated by CHECK CONDITION status and current error information shall be returned in the sense data. In this case, if the current error information does not adequately define the deferred error condition, a deferred error may be returned after the current error information has been recovered. If a deferred error occurs while a current command is executing and the current command completes successfully, the target may choose to return the deferred error information after the completion of the current command.

IMPLEMENTORS NOTE: Deferred errors may indicate that an operation was unsuccessful long after the command performing the data transfer returned GOOD status. If data that cannot be replicated or recovered from other sources is being stored using such buffered write operations, synchronization commands should be performed before the critical data is destroyed in the host initiator. This is necessary to be sure that recovery actions can be taken if deferred errors do occur in the storing of the data. If AEN is not implemented, the synchronizing process should provide the necessary commands to allow returning CHECK CONDITION status and subsequent returning of deferred error sense information after all buffered operations are guaranteed to be complete.

Table 7-39: Sense Key (0h-7h) Descriptions

Sense Key	Description
0h	NO SENSE. Indicates that there is no specific sense key information to be reported for the designated logical unit. This would be the case for a successful command or a command that received CHECK CONDITION or COMMAND TERMINATED status because one of the filemark, EOM, or ILI bits is set to one.
1h	RECOVERED ERROR. Indicates that the last command completed successfully with some recovery action performed by the target. Details may be determinable by examining the additional sense bytes and the information field. When multiple recovered errors occur during one command, the choice of which error to report (first, last, most severe, etc.) is device specific.
2h	NOT READY. Indicates that the logical unit addressed cannot be accessed. Operator intervention may be required to correct this condition.
3h	MEDIUM ERROR. Indicates that the command terminated with a non-recovered error condition that was probably caused by a flaw in the medium or an error in the recorded data. This sense key may also be returned if the target is unable to distinguish between a flaw in the medium and a specific hardware failure (sense key 4h).
4h	HARDWARE ERROR. Indicates that the target detected a non-recoverable hardware failure (for example, controller failure, device failure, parity error, etc.) while performing the command or during a self test.
5h	ILLEGAL REQUEST. Indicates that there was an illegal parameter in the command descriptor block or in the additional parameters supplied as data for some commands (FORMAT UNIT, SEARCH DATA, etc.). If the target detects an invalid parameter in the command descriptor block, then it shall terminate the command without altering the medium. If the target detects an invalid parameter in the additional parameters supplied as data, then the target may have already altered the medium. This sense key may also indicate that an invalid IDENTIFY message was received (5.6.7).
6h	UNIT ATTENTION. Indicates that the removable medium may have been changed or the target has been reset. See 6.9 for more detailed information about the unit attention condition.
7h	DATA PROTECT. Indicates that a command that reads or writes the medium was attempted on a block that is protected from this operation. The read or write operation is not performed.

Table 7-40: Sense Key (8h-Fh) Descriptions

Sense Key	Description
8h	BLANK CHECK. Indicates that a write-once device or a sequential-access device encountered blank medium or format-defined end-of-data indication while reading or a write-once device encountered a non-blank medium while writing.
9h	Vendor Specific. This sense key is available for reporting vendor specific conditions.
Ah	COPY ABORTED. Indicates a COPY, COMPARE, or COPY AND VERIFY command was aborted due to an error condition on the source device, the destination device, or both. (See 7.2.3.2 for additional information about this sense key.)
Bh	ABORTED COMMAND. Indicates that the target aborted the command. The initiator may be able to recover by trying the command again.
Ch	EQUAL. Indicates a SEARCH DATA command has satisfied an equal comparison.
Dh	VOLUME OVERFLOW. Indicates that a buffered peripheral device has reached the end-of-partition and data may remain in the buffer that has not been written to the medium. A RECOVER BUFFERED DATA command(s) may be issued to read the unwritten data from the buffer.
Eh	MISCOMPARE. Indicates that the source data did not match the data read from the medium.
Fh	RESERVED.

Table 7-41: ASC and ASCQ Assignments

BYTE	12	13	DTLPWRSOMC	DESCRIPTION	COMMENTS
	13	00	D W 0	ADDRESS MARK NOT FOUND FOR DATA FIELD	
	12	00	D W 0	ADDRESS MARK NOT FOUND FOR ID FIELD	
	00	11	R	AUDIO PLAY OPERATION IN PROGRESS	
	00	12	R	AUDIO PLAY OPERATION PAUSED	
	00	14	R	AUDIO PLAY OPERATION STOPPED DUE TO ERROR	
	00	13	R	AUDIO PLAY OPERATION SUCCESSFULLY COMPLETED	
	00	04	T S	BEGINNING-OF-PARTITION/MEDIUM DETECTED	
	14	04	T	BLOCK SEQUENCE ERROR	
	30	02	DT WR 0	CANNOT READ MEDIUM - INCOMPATIBLE FORMAT	
	30	01	DT WR 0	CANNOT READ MEDIUM - UNKNOWN FORMAT	
	52	00	T	CARTRIDGE FAULT	
	3F	02	DTLPWRSOMC	CHANGED OPERATING DEFINITION	
	11	06	WR 0	CIRC UNRECOVERED ERROR	(L-EC UNAVAILABLE OR OFF)
	30	03	DT	CLEANING CARTRIDGE INSTALLED	
	4A	00	DTLPWRSOMC	COMMAND PHASE ERROR	
	2C	00	DTLPWRSOMC	COMMAND SEQUENCE ERROR	
	2B	00	DTLPWRSO C	COPY CANNOT EXECUTE SINCE HOST CANNOT DISCONNECT	
	41	00	D	DATA PATH FAILURE	SHOULD USE 40 NN
	4B	00	DTLPWRSOMC	DATA PHASE ERROR	
	11	07	W 0	DATA RESYNCHRONIZATION ERROR	
	16	00	D W 0	DATA SYNCHRONIZATION MARK ERROR	
	19	00	D 0	DEFECT LIST ERROR	
	19	03	D 0	DEFECT LIST ERROR IN GROWN LIST	
	19	02	D 0	DEFECT LIST ERROR IN PRIMARY LIST	
	19	01	D 0	DEFECT LIST NOT AVAILABLE	
	1C	00	D 0	DEFECT LIST NOT FOUND	
	32	01	D W 0	DEFECT LIST UPDATE FAILURE	
	40	NN	DTLPWRSOMC	DIAGNOSTIC FAILURE ON COMPONENT NN (80H-FFH)	USE ASCQ CODES 80H-FFH
	63	00	R	END OF USER AREA ENCOUNTERED ON THIS TRACK	
	00	05	T S	END-OF-DATA DETECTED	
	14	03	T	END-OF-DATA NOT FOUND	
	00	02	T S	END-OF-PARTITION/MEDIUM DETECTED	
	51	00	T 0	ERASE FAILURE	

Table 7-41: ASC and ASCQ Assignments (continued)

=====			
0A 00	DTLPWRSOMC	ERROR LOG OVERFLOW	
11 02	DT W SO	ERROR TOO LONG TO CORRECT	
03 02	T	EXCESSIVE WRITE ERRORS	
3B 07	L	FAILED TO SENSE BOTTOM-OF-FORM	
3B 06	L	FAILED TO SENSE TOP-OF-FORM	
00 01	T	FILEMARK DETECTED	
14 02	T	FILEMARK OR SETMARK NOT FOUND	
09 02	WR O	FOCUS SERVO FAILURE	
31 01	D L O	FORMAT COMMAND FAILED	
58 00	O	GENERATION DOES NOT EXIST	
1C 02	D O	GROWN DEFECT LIST NOT FOUND	
10 00	D W O	ID CRC OR ECC ERROR	
22 00	D	ILLEGAL FUNCTION FOR DEVICE TYPE	
SHOULD USE 20 00, 24 00, OR 26 00			
64 00	R	ILLEGAL MODE FOR THIS TRACK	
28 01	M	IMPORT / EXPORT ELEMENT ACCESSED	
30 00	DT WR OM	INCOMPATIBLE MEDIUM INSTALLED	
11 08	T	INCOMPLETE BLOCK READ	(POSTAMBLE NOT FOUND)
48 00	DTLPWRSOMC	INITIATOR DETECTED ERROR MESSAGE RECEIVED	
3F 03	DTLPWRSOMC	INQUIRY DATA HAS CHANGED	
44 00	DTLPWRSOMC	INTERNAL TARGET FAILURE	
3D 00	DTLPWRSOMC	INVALID BITS IN IDENTIFY MESSAGE	
2C 02	S	INVALID COMBINATION OF WINDOWS SPECIFIED	
20 00	DTLPWRSOMC	INVALID COMMAND OPERATION CODE	
21 01	M	INVALID ELEMENT ADDRESS	
24 00	DTLPWRSOMC	INVALID FIELD IN CDB	CHECK FIELD POINTER IN SENSE DATA
26 00	DTLPWRSOMC	INVALID FIELD IN PARAMETER LIST	CHECK FIELD POINTER IN SENSE DATA
=====			
49 00	DTLPWRSOMC	INVALID MESSAGE ERROR	
11 05	WR O	L-EC UNCORRECTABLE ERROR	(L-EC CODES PRESENT AND L-EC ON)
60 00	S	LAMP FAILURE	
5B 02	DTLPWRSOMC	LOG COUNTER AT MAXIMUM	
5B 00	DTLPWRSOMC	LOG EXCEPTION	
5B 03	DTLPWRSOMC	LOG LIST CODES EXHAUSTED	
2A 02	DTL WRSOMC	LOG PARAMETERS CHANGED	
21 00	DT WR OM	LOGICAL BLOCK ADDRESS OUT OF RANGE	
08 00	DTL WRSOMC	LOGICAL UNIT COMMUNICATION FAILURE	
08 02	DTL WRSOMC	LOGICAL UNIT COMMUNICATION PARITY ERROR	
08 01	DTL WRSOMC	LOGICAL UNIT COMMUNICATION TIME-OUT	
05 00	DTL WRSOMC	LOGICAL UNIT DOES NOT RESPOND TO SELECTION	
4C 00	DTLPWRSOMC	LOGICAL UNIT FAILED SELF-CONFIGURATION	
3E 00	DTLPWRSOMC	LOGICAL UNIT HAS NOT SELF-CONFIGURED YET	
04 01	DTLPWRSOMC	LOGICAL UNIT IS IN PROCESS OF BECOMING READY	
04 00	DTLPWRSOMC	LOGICAL UNIT NOT READY, CAUSE NOT REPORTABLE	
04 04	DTL O	LOGICAL UNIT NOT READY, FORMAT IN PROGRESS	
04 02	DTLPWRSOMC	LOGICAL UNIT NOT READY, INITIALIZING COMMAND REQUIRED	
04 03	DTLPWRSOMC	LOGICAL UNIT NOT READY, MANUAL INTERVENTION REQUIRED	
25 00	DTLPWRSOMC	LOGICAL UNIT NOT SUPPORTED	
15 01	DTL WRSOMC	MECHANICAL POSITIONING ERROR	
53 00	DTL WRSOMC	MEDIA LOAD/EJECT FAILED	
=====			

Table 7-41: ASC and ASCQ Assignments (continued)

=====			
3B 0D	M	MEDIUM DESTINATION ELEMENT FULL	
31 00	DT W O	MEDIUM FORMAT CORRUPTED	
3A 00	DTL WRSOMC	MEDIUM NOT PRESENT	
53 02	DT WR OM	MEDIUM REMOVAL PREVENTED	
3B 0E	M	MEDIUM SOURCE ELEMENT EMPTY	
43 00	DTLPWRSOMC	MESSAGE ERROR	
3F 01	DTLPWRSOMC	MICROCODE HAS BEEN CHANGED	
1D 00	D W O	MISCOMPARE DURING VERIFY OPERATION	
11 0A	DT O	MISCORRECTED ERROR	
2A 01	DTL WRSOMC	MODE PARAMETERS CHANGED	
07 00	DTL WRSOMC	MULTIPLE PERIPHERAL DEVICES SELECTED	
11 03	DT W SO	MULTIPLE READ ERRORS	
00 00	DTLPWRSOMC	NO ADDITIONAL SENSE INFORMATION	
00 15	R	NO CURRENT AUDIO STATUS TO RETURN	
32 00	D W O	NO DEFECT SPARE LOCATION AVAILABLE	
11 09	T	NO GAP FOUND	
01 00	D W O	NO INDEX/SECTOR SIGNAL	
06 00	D WR OM	NO REFERENCE POSITION FOUND	(TRACK ZERO OR EQUIVALENT)
02 00	D WR OM	NO SEEK COMPLETE	
03 01	T	NO WRITE CURRENT	
28 00	DTLPWRSOMC	NOT READY TO READY TRANSITION	MEDIUM MAY HAVE CHANGED
5A 01	DT WR OM	OPERATOR MEDIUM REMOVAL REQUEST	
5A 00	DTLPWRSOMC	OPERATOR REQUEST OR STATE CHANGE INPUT (UNSPECIFIED)	
5A 03	DT W O	OPERATOR SELECTED WRITE PERMIT	
5A 02	DT W O	OPERATOR SELECTED WRITE PROTECT	
61 02	S	OUT OF FOCUS	
4E 00	DTLPWRSOMC	OVERLAPPED COMMANDS ATTEMPTED	
2D 00	T	OVERWRITE ERROR ON UPDATE IN PLACE	
3B 05	L	PAPER JAM	
1A 00	DTLPWRSOMC	PARAMETER LIST LENGTH ERROR	
26 01	DTLPWRSOMC	PARAMETER NOT SUPPORTED	CHECK FIELD POINTER IN SENSE DATA
26 02	DTLPWRSOMC	PARAMETER VALUE INVALID	CHECK FIELD POINTER IN SENSE DATA
2A 00	DTL WRSOMC	PARAMETERS CHANGED	
03 00	DTL W SO	PERIPHERAL DEVICE WRITE FAULT	
3B 0C	S	POSITION PAST BEGINNING OF MEDIUM	
3B 0B	S	POSITION PAST END OF MEDIUM	
15 02	DT WR O	POSITIONING ERROR DETECTED BY READ OF MEDIUM	
29 00	DTLPWRSOMC	POWER ON, RESET, OR BUS DEVICE RESET OCCURRED	
42 00	D	POWER-ON OR SELF-TEST FAILURE	SHOULD USE 40 NN
1C 01	D O	PRIMARY DEFECT LIST NOT FOUND	
40 00	D	RAM FAILURE	SHOULD USE 40 NN
15 00	DTL WRSOMC	RANDOM POSITIONING ERROR	
3B 0A	S	READ PAST BEGINNING OF MEDIUM	
3B 09	S	READ PAST END OF MEDIUM	
11 01	DT W SO	READ RETRIES EXHAUSTED	
14 01	DT WR O	RECORD NOT FOUND	
14 00	DTL WRSO	RECORDED ENTITY NOT FOUND	
=====			

Table 7-41: ASC and ASCQ Assignments (continued)

17 05 D	WR 0	RECOVERED DATA USING PREVIOUS SECTOR ID
18 03	R	RECOVERED DATA WITH CIRC
18 02 D	WR 0	RECOVERED DATA WITH ECC AND/OR RETRIES. DATA AUTO-REALLOCATED
18 01 D	WR 0	RECOVERED DATA WITH ERROR CORRECTION AND RETRIES APPLIED
18 00 DT	WR 0	RECOVERED DATA WITH ERROR CORRECTION APPLIED
18 04	R	RECOVERED DATA WITH LEC
17 03 DT	WR 0	RECOVERED DATA WITH NEGATIVE HEAD OFFSET
17 00 DT	WRSO	RECOVERED DATA WITH NO ERROR CORRECTION APPLIED
17 02 DT	WR 0	RECOVERED DATA WITH POSITIVE HEAD OFFSET
17 01 DT	WRSO	RECOVERED DATA WITH RETRIES
17 04	WR 0	RECOVERED DATA WITH RETRIES AND/OR CIRC APPLIED
17 06 D	W 0	RECOVERED DATA WITHOUT ECC. DATA AUTO-REALLOCATED.
1E 00 D	W 0	RECOVERED ID WITH ECC CORRECTION
3B 08 T		REPOSITION ERROR
36 00 L		RIBBON/INK/TONER FAILURE
37 00 DTL	WRSOMC	ROUNDED PARAMETER
5C 00 D	O	RPL STATUS CHANGE
39 00 DTL	WRSOMC	SAVING PARAMETERS NOT SUPPORTED
62 00	S	SCAN HEAD POSITIONING ERROR
47 00 DTL	PWRSOMC	SCSI PARITY ERROR
54 00	P	SCSI TO HOST SYSTEM INTERFACE FAILURE
45 00 DTL	PWRSOMC	SELECT/RESELECT FAILURE
3B 00 TL		SEQUENTIAL POSITIONING ERROR
00 03 T		SETMARK DETECTED
3B 04 L		SLEW FAILURE
09 03	WR 0	SPINDLE SERVO FAILURE
5C 02 D	O	SPINDLES NOT SYNCHRONIZED
5C 01 D	O	SPINDLES SYNCHRONIZED
1B 00 DTL	PWRSOMC	SYNCHRONOUS DATA TRANSFER ERROR
55 00	P	SYSTEM RESOURCE FAILURE
2F 00 DTL	PWRSOMC	COMMANDS CLEARED BY ANOTHER INITIATOR
33 00 T		TAPE LENGTH ERROR
3B 03 L		TAPE OR ELECTRONIC VERTICAL FORMS UNIT NOT READY
3B 01 T		TAPE POSITION ERROR AT BEGINNING-OF-MEDIUM
3B 02 T		TAPE POSITION ERROR AT END-OF-MEDIUM
3F 00 DTL	PWRSOMC	TARGET OPERATING CONDITIONS HAVE CHANGED
5B 01 DTL	PWRSOMC	THRESHOLD CONDITION MET
26 03 DTL	PWRSOMC	THRESHOLD PARAMETERS NOT SUPPORTED
50 02 T		POSITION ERROR RELATED TO TIMING
2C 01	S	TOO MANY WINDOWS SPECIFIED
09 00 DT	WR 0	TRACK FOLLOWING ERROR
09 01	WR 0	TRACKING SERVO FAILURE
61 01	S	UNABLE TO ACQUIRE VIDEO
57 00	R	UNABLE TO RECOVER TABLE-OF-CONTENTS
53 01 T		UNLOAD TAPE FAILURE
11 00 DT	WRSO	UNRECOVERED READ ERROR
11 04 D	W 0	UNRECOVERED READ ERROR - AUTO REALLOCATE FAILED
46 00 DTL	PWRSOMC	UNSUCCESSFUL SOFT RESET
59 00	O	UPDATED BLOCK READ
61 00	S	VIDEO ACQUISITION ERROR

Table 7-41: ASC and ASCQ Assignments (continued)

50 00 T		WRITE APPEND ERROR
50 01 T		WRITE APPEND POSITION ERROR
0C 00 T	S	WRITE ERROR
0C 02 D	W 0	WRITE ERROR - AUTO REALLOCATION FAILED
0C 01 D	W 0	WRITE ERROR RECOVERED WITH AUTO REALLOCATION
27 00 DT	W 0	WRITE PROTECTED

80 XX	\	VENDOR UNIQUE.
THROUGH		
FF XX	/	
XX 80	\	VENDOR UNIQUE QUALIFICATION OF STANDARD ASC.
THROUGH		
XX FF	/	
ALL CODES NOT SHOWN ARE RESERVED.		

7.2.15. SEND DIAGNOSTIC Command

Table 7-42: SEND DIAGNOSTIC Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (1Dh)							
1	Logical Unit Number			PF	Reserved	SelfTest	DevOfL	UnitOfL
2	Reserved							
3	(MSB)							
4	Parameter List Length							(LSB)
5	Control							

The SEND DIAGNOSTIC command (Table 7-42) requests the target to perform diagnostic operations on itself, on the logical unit, or on both. The only mandatory implementation of this command is the self-test feature with the parameter list length of zero. Except when the self-test bit is one, this command is usually followed by a RECEIVE DIAGNOSTIC RESULTS command.

A page format (PF) bit of one specifies that the SEND DIAGNOSTIC parameters conform to the page structure as specified in this standard. The implementation of the PF bit is optional. See 7.3.1 for the definition of diagnostic pages. A PF bit of zero indicates that the SEND DIAGNOSTIC parameters are as specified in SCSI-1 (i.e., all parameters are vendor specific).

A self-test (SelfTest) bit of one directs the target to complete its default self-test. If the self-test successfully passes, the command shall be terminated with GOOD status; otherwise, the command shall be terminated with CHECK CONDITION status and the sense key shall be set to HARDWARE ERROR.

A self-test bit of zero requests that the target perform the diagnostic operation specified in the parameter list. The diagnostic operation might or might not require a target to return data which contains diagnostic results. If the return of data is not required, the return of GOOD status indicates successful completion of the diagnostic operation. If the return of data is required the target shall either:

(1) perform the requested diagnostic operation, prepare the data to be returned and indicate completion by returning GOOD status. The initiator issues a RECEIVE DIAGNOSTIC RESULTS command to recover the data.

(2) accept the parameter list and if no errors are detected in the parameter list return GOOD status. The requested diagnostic operation and the preparation of the data to be returned is performed upon receipt of a RECEIVE DIAGNOSTIC RESULTS command.

IMPLEMENTORS NOTE: To insure that the diagnostic command information is not destroyed by a command sent from another initiator, the SEND DIAGNOSTIC command should either be linked to the RECEIVE DIAGNOSTIC RESULTS command or the logical unit should be reserved.

The device off-line (DevOfL) and unit off-line (UnitOfL) bits are generally set by operating system software, while the parameter list is prepared by diagnostic application software. These bits grant permission to perform vendor-specific diagnostic operations on the target which may be visible to attached initiators. Thus, by preventing operations that are not enabled by these bits, the target assists the operating system in protecting its resources.

A UnitOfL bit of one grants permission to the target to perform diagnostic operations that may affect the user accessible medium on the logical unit, e.g., write operations to the user accessible medium, or repositioning of the medium on sequential access devices. The implementation of the UnitOfL bit is optional. A UnitOfL bit of zero prohibits any diagnostic operations that may be detected by subsequent I/O processes.

A DevOfL bit of one grants permission to the target to perform diagnostic operations that may affect all the logical units on a target, e.g., alteration of reservations, log parameters, or sense data. The implementation of the DevOfL bit is optional. A DevOfL bit of zero prohibits diagnostic operations that may be detected by subsequent I/O processes.

The parameter list length field specifies the length in bytes of the parameter list that shall be transferred from the initiator to the target. A parameter list length of zero indicates that no data shall be transferred. This condition shall not be considered as an error. If the specified parameter list length results in the truncation of one or more pages (PF bit set to one) the target shall return CHECK CONDITION status with a sense key of ILLEGAL REQUEST and an additional sense code of INVALID FIELD IN CDB.

See the implementors note under the RECEIVE DIAGNOSTIC RESULTS command in 7.2.13.

7.2.16. TEST UNIT READY Command

Table 7-43: TEST UNIT READY Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (00h)							
1	Logical Unit Number			Reserved				
2	Reserved							
3	Reserved							
4	Reserved							
5	Control							

The TEST UNIT READY command (Table 7-43) provides a means to check if the logical unit is ready. This is not a request for a self-test. If the logical unit would accept an appropriate medium-access command without returning CHECK CONDITION status, this command shall return a GOOD status. If the logical unit cannot become operational or is in a state such that an initiator action (e.g., START UNIT command) is required to make the unit ready the target shall return CHECK CONDITION status with a sense key of NOT READY.

Table 7-44 defines the preferred responses to the TEST UNIT READY command. Higher-priority responses (e.g., BUSY or RESERVATION CONFLICT) are also permitted.

Table 7-44: Preferred TEST UNIT READY Responses

Status	Sense Key	Additional Sense Code and Additional Sense Code Qualifier
GOOD	NO SENSE	NO ADDITIONAL SENSE INFORMATION or other valid additional sense code.
CHECK CONDITION	ILLEGAL REQUEST	LOGICAL UNIT NOT SUPPORTED
CHECK CONDITION	NOT READY	LOGICAL UNIT DOES NOT RESPOND TO SELECTION
CHECK CONDITION	NOT READY	MEDIUM NOT PRESENT
CHECK CONDITION	NOT READY	LOGICAL UNIT NOT READY, CAUSE NOT REPORTABLE
CHECK CONDITION	NOT READY	LOGICAL UNIT IS IN PROCESS OF BECOMING READY
CHECK CONDITION	NOT READY	LOGICAL UNIT NOT READY, INITIALIZING COMMAND REQUIRED
CHECK CONDITION	NOT READY	LOGICAL UNIT NOT READY, MANUAL INTERVENTION REQUIRED
CHECK CONDITION	NOT READY	LOGICAL UNIT NOT READY, FORMAT IN PROGRESS

7.2.17. WRITE BUFFER Command

Table 7-45: WRITE BUFFER Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (3Bh)							
1	Logical Unit Number			Reserved		Mode		
2	Buffer ID							
3	(MSB)							
5	Buffer Offset							(LSB)
6	(MSB)							
8	Parameter List Length							(LSB)
9	Control							

The WRITE BUFFER command (Table 7-45) is used in conjunction with the READ BUFFER command as a diagnostic for testing target memory and the SCSI bus integrity. Additional modes are provided for downloading microcode and for downloading and saving microcode.

This command shall not alter any medium of the target when the data mode or the combined header and data mode is specified.

The function of this command and the meaning of fields within the command descriptor block depend on the contents of the mode field. The mode field is defined in Table 7-46.

Table 7-46: WRITE BUFFER Mode Field

Mode	Description	Implementation Requirements
000b	Write combined header and data	Optional
001b	Vendor specific	Vendor specific
010b	Write data	Optional
011b	Reserved	Reserved
100b	Download Microcode	Optional
101b	Download Microcode and Save	Optional
110b	Reserved	Reserved
111b	Reserved	Reserved

IMPLEMENTORS NOTE: Modes 000b and 001b are included for compatibility with CCS products that were designed prior to the generation of this standard. These products restrict the maximum transfer length to 65535 bytes.

7.2.17.1. Combined Header and Data Mode (000b)

In this mode, data to be transferred is preceded by a four-byte header. The four-byte header consists of all reserved bytes. The buffer ID and the buffer offset fields shall be zero. The parameter list length field specifies the maximum number of bytes that shall be transferred during the DATA OUT phase. This number includes four bytes of header, so the data length to be stored in the target's buffer is parameter list length minus four. The initiator should attempt to ensure that the parameter list length is not greater than four plus the available length that is returned in the header of the READ BUFFER command (mode 00b). If the parameter list length exceeds the available length plus four, the target shall return CHECK CONDITION status and shall set the sense key to ILLEGAL REQUEST.

7.2.17.2. Vendor-Specific Mode (001b)

In this mode, the meaning of the buffer ID, buffer offset, and parameter list length fields are not specified by this standard.

7.2.17.3. Data Mode (010b)

In this mode, the DATA OUT phase contains buffer data. The buffer ID field identifies a specific buffer within the target. The vendor assigns buffer ID codes to buffers within the target. Buffer ID zero shall be supported. If more than one buffer is supported, additional buffer ID codes shall be assigned contiguously, beginning with one. If an unsupported buffer ID code is selected, the target shall return CHECK CONDITION status and shall set the sense key to ILLEGAL REQUEST with an additional sense code of INVALID FIELD IN CDB.

Data is written to the target buffer starting at the location specified by the buffer offset. The initiator should conform to the offset boundary requirements returned in the READ BUFFER descriptor. If the target is unable to accept the specified buffer offset, it shall return CHECK CONDITION status and it shall set the sense key to ILLEGAL REQUEST with an additional sense code of INVALID FIELD IN CDB.

The parameter list length specifies the maximum number of bytes that shall be transferred during the DATA OUT phase to be stored in the specified buffer beginning at the buffer offset. The initiator should attempt to ensure that the parameter list length plus the buffer offset does not exceed the capacity of the specified buffer. (The capacity of the buffer can be determined by the buffer capacity field in the READ BUFFER descriptor.) If the buffer offset and parameter list length fields specify a transfer that would exceed the buffer capacity, the target shall return CHECK CONDITION status and shall set the sense key to ILLEGAL REQUEST with an additional sense code of INVALID FIELD IN CDB.

7.2.17.4. Download Microcode Mode (100b)

In this mode, vendor-specific microcode or control information shall be transferred to the control memory space of the target. After a power-cycle or reset, the device operation shall revert to a vendor-specific condition. The meanings of the buffer ID, buffer offset, and parameter list length fields are not specified by this standard and are not required to be zero-filled. When the microcode download has completed successfully the target shall generate a unit attention condition for all initiators except the one that issued the WRITE BUFFER command (see 6.9). The additional sense code shall be set to MICROCODE HAS BEEN CHANGED.

7.2.17.5. Download Microcode and Save Mode (101b)

In this mode, vendor-specific microcode or control information shall be transferred to the target and, if the WRITE BUFFER command is completed successfully, also shall be saved in a non-volatile memory space (semiconductor, disk, or other). The downloaded code shall then be effective after each power-cycle and reset until it is supplanted in another download microcode and save operation. The meanings of the buffer ID, buffer offset, and parameter list length fields are not specified by this standard and are not required to be zero-filled. When the download microcode and save command has completed successfully the target shall generate a unit attention condition for all initiators except the one that issued the WRITE BUFFER command. When reporting the unit attention condition, the target shall set the additional sense code to MICROCODE HAS BEEN CHANGED.

7.3. Parameters for All Device Types

7.3.1. Diagnostic Parameters

This section describes the diagnostic page structure and the diagnostic pages that are applicable to all SCSI devices. Pages specific to each device type are described in the third subsection of each device-type section (i.e., 8.3, 9.3, etc.).

A SEND DIAGNOSTIC command with a PF bit of one specifies that the SEND DIAGNOSTIC parameter list consists of zero or more diagnostic pages and that the data returned by the subsequent RECEIVE DIAGNOSTIC RESULTS command shall use the diagnostic page format (Table 7-47) described in this standard.

Each diagnostic page defines a function or operation that the target shall perform. The page contains a page header followed by the analysis data which is formatted according to the page code specified in the previous SEND DIAGNOSTIC command.

Targets that implement diagnostic pages are only required to accept a single diagnostic page per command.

Table 7-47: Diagnostic Page Format

Bit Byte	7	6	5	4	3	2	1	0
0	Page Code							
1	Reserved							
2	(MSB)							
3	Page Length (n-3)							
4	(LSB)							
n	Diagnostic Parameters							

The page code field identifies which diagnostic page is being sent or returned. The page codes are defined in Table 7-48.

The page length field specifies the length in bytes of the diagnostic parameters which follow this field. If the initiator sends a page length that results in the truncation of any parameter, the target shall terminate the command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST with the additional sense code set to INVALID FIELD IN PARAMETER LIST.

The diagnostic parameters are defined for each page code. The diagnostic parameters within a page may be defined differently in a SEND DIAGNOSTIC command than in a RECEIVE DIAGNOSTIC RESULTS command.

Table 7-48: Diagnostic Page Codes

Page Code	Description	Section
00h	Supported Diagnostics Pages	7.3.1.1
01h - 3Fh	Reserved (for all device type pages)	
40h - 7Fh	See specific device type for definition	
80h - FFh	Vendor specific pages	

7.3.1.1. Supported Diagnostic Pages

The supported diagnostics page (Table 7-49) returns the list of diagnostic pages implemented by the target. This page shall be implemented if the target implements the page format option of the SEND DIAGNOSTIC and RECEIVE DIAGNOSTIC RESULTS commands.

Table 7-49: Supported Diagnostic Pages

Bit Byte	7	6	5	4	3	2	1	0
0	Page Code (00h)							
1	Reserved							
2	(MSB)							
3	Page Length (n-3)							
4	(LSB)							
n	Supported Page List							

The definition of this page for the SEND DIAGNOSTIC command includes only the first four bytes. If the page length field is not zero, the target shall terminate the SEND DIAGNOSTIC command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST with an additional sense code of INVALID FIELD IN PARAMETER LIST. This page instructs the target to make available the list of all supported diagnostic pages to be returned by a subsequent RECEIVE DIAGNOSTIC RESULTS command.

The definition of this page for the RECEIVE DIAGNOSTIC RESULTS command includes the list of diagnostic pages supported by the target.

The page length field specifies the length in bytes of the following supported page list.

The supported page list field shall contain a list of all diagnostic page codes implemented by the target in ascending order beginning with page code 00h.

7.3.2. Log Parameters

This section describes the log page structure and the log pages that are applicable to all SCSI devices. Pages specific to each device type are described in the third subsection of each device-type section (i.e., 8.3.2, 9.3.2, etc.). The LOG SELECT command supports the ability to send zero or more log pages. The LOG SENSE command shall return a single log page specified in the page code field of the command descriptor block (see Table 7-21).

Each log page begins with a four-byte page header followed by zero or more variable-length log parameters defined for that page. The log page format is defined in Table 7-50.

Table 7-50: Log Page Format

Bit Byte	7	6	5	4	3	2	1	0
0	Reserved			Page Code				
1	Reserved							
2	(MSB)							
3	Page Length (n-3)						(LSB)	
	Log Parameters(s)							
4 - x+3	Log Parameter (First) (Length x)							
	⋮							
n-y n	Log Parameter (Last) (Length y)							

The page code field identifies which log page is being transferred.

The page length field specifies the length in bytes of the following log parameters. If the initiator sends a page length that results in the truncation of any parameter, the target shall terminate the command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST with the additional sense code set to INVALID FIELD IN PARAMETER LIST.

Most log pages contain one or more special data structures called log parameters (see Table 7-51). Log parameters may be data counters which record a count of a particular event (or events) or log parameters may be list parameters (strings) which contain a description of a particular event.

Table 7-51: Log Parameter

Bit Byte	7	6	5	4	3	2	1	0
0	(MSB)							
1	Parameter Code						(LSB)	
2	DU	DS	TSD	ETC	TMC		Reserved	LP
3	Parameter Length (n-3)							
4								
n	Parameter Value							

Each log parameter (Table 7-51) begins with a four-byte parameter header followed by one or more bytes of parameter value data.

The parameter code field identifies which log parameter is being transferred for that log page.

The DU, DS, TSD, ETC, TMC, and LP fields are collectively referred to as the parameter control byte. These fields are described below.

For cumulative log parameter values (indicated by the PC field of the LOG SELECT and LOG SENSE command descriptor block), the disable update (DU) bit is defined as follows:

- (1) A zero value indicates that the target shall update the log parameter value to reflect all events that should be noted by that parameter.
- (2) A one value indicates that the target shall not update the log parameter value except in response to a LOG SELECT command that specifies a new value for the parameter.

IMPLEMENTORS NOTE: When updating cumulative log parameter values, a target may use volatile memory to hold these values until a LOG SELECT or LOG SENSE command is received with an SP bit of one (or a target-defined event occurs). Thus the updated cumulative log parameter values may be lost if a power cycle occurs.

The DU bit is not defined for threshold values (indicated by the PC field of the LOG SENSE command descriptor block) nor for list parameters (indicated by the LP bit). The target shall ignore the value of any DU bits in a LOG SELECT command.

A disable save (DS) bit of zero indicates that the target supports saving for that log parameter. The target shall save the current cumulative or the current threshold parameter value (depending on the value in the PC field of the command descriptor block) in response to a LOG SELECT or LOG SENSE command with a SP bit of one. A DS bit of one indicates that the target does not support saving that log parameter in response to a LOG SELECT or LOG SENSE command with a SP bit of one.

A target save disable (TSD) bit of zero indicates that the target provides a target-defined method for saving log parameters. This implicit saving operation shall be done frequently enough to insure that the cumulative parameter values retain statistical significance (i.e., across power cycles). A TSD bit of one indicates that either the target does not provide a target-defined method for saving log parameters or the target-defined method has been disabled by the initiator. If the initiator sets both the DS and the TSD bits set to one, the target shall terminate the command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST with the additional sense code set to INVALID FIELD IN PARAMETER LIST.

An enable threshold comparison (ETC) bit of one indicates that a comparison to the threshold value is performed whenever the cumulative value is updated. An ETC bit of zero indicates that a comparison is not performed. The value of the ETC bit is the same for cumulative and threshold parameters.

The threshold met criteria (TMC) field (Table 7-52) defines the basis for comparison of the cumulative and threshold values. The TMC field is valid only if the ETC bit is one. The value of the TMC field is the same for cumulative and threshold parameters.

If the ETC bit is one and the result of the comparison is true, a unit attention condition shall be generated for all initiators. When reporting the unit attention condition, the target shall set the sense key shall be set to UNIT ATTENTION; the additional sense code shall be set to LOG EXCEPTION; and the additional sense code qualifier set to THRESHOLD CONDITION MET.

Table 7-52: Threshold Met Criteria

Code	Basis For Comparison
00b	Every update of the cumulative value
01b	Cumulative value equal threshold value
10b	Cumulative value not equal threshold value
11b	Cumulative value greater than threshold value

The list parameter (LP) bit indicates the format of the log parameter. If an initiator attempts to set the value of the LP bit to a value other than the one returned for the same parameter in the LOG SENSE command, the target shall terminate the command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST with the additional sense code set to INVALID FIELD IN PARAMETER LIST.

An LP bit of zero indicates that the parameter is a data counter. Data counters are associated with one of more events. The data counter is updated whenever one of these events occurs by incrementing of the counter value provided the DU bit is zero. Each data counter has associated with it a target-defined maximum value. Upon reaching this maximum value, the data counter shall not be incremented (i.e., it does not wrap). When a data counter reaches its maximum value, the target shall set the associated DU bit to one. If the data counter is at or reaches its maximum value during the execution of a command, the target shall complete the command. If the command completes correctly (except for the data counter being at its maximum value) and if the RLEC bit of the control mode page (7.3.3.1) is set to one, then the target shall terminate the command with CHECK CONDITION status and set the sense key to RECOVERED ERROR with the additional sense code set to LOG COUNTER AT MAXIMUM.

An LP bit of one indicates that the parameter is a list parameter. List parameters are not counters and thus the ET and TMC fields shall be set to zero. A list parameter is a string of ASCII graphic codes (i.e., code values 20h through 7eh).

If more than one list parameter is defined in a single log page, the following rules apply to assigning parameter codes:

- (1) The parameter updated last shall have a higher parameter code than the previous parameter, except as defined in rule (2).
- (2) When the maximum parameter code value supported by the target is reached, the target shall assign the lowest parameter code value to the next log parameter (i.e., wrap-around parameter codes). If the associated command completes correctly (except for the parameter code being at its maximum value) and if the RLEC bit of the control mode page (7.3.3.1) is set to one, then the target shall terminate the command with CHECK CONDITION status and set the sense key to RECOVERED ERROR with the additional sense code set to LOG LIST CODES EXHAUSTED.

IMPLEMENTORS NOTE: List parameters can be used to store the locations of defective blocks in the following manner. When a defective block is identified, a list parameter is updated to reflect the location and cause of the defect. When the next defect is encountered, the list parameter with the next higher parameter code is updated to record this defect. The size of the page can be made target specific to accommodate memory limitations. It is recommended that one or more data counter parameters be defined for the page to keep track of the number of valid list parameters and the parameter code of the parameter with the oldest recorded defect. This technique can be adapted to record other types of information.

The parameter length field specifies the length in bytes of the following parameter value. If the initiator sends a parameter length value that results in the truncation of the parameter value, the target shall terminate the command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST with the additional sense code set to INVALID FIELD IN PARAMETER LIST.

If the initiator sends a log parameter value that is outside the range supported by the target and rounding is implemented for that parameter, the target may either:

- (1) round to an acceptable value and terminate the command as described in 6.5.4.
- (2) terminate the command with CHECK CONDITION status, set the sense key to ILLEGAL REQUEST with the additional sense code set to INVALID FIELD IN PARAMETER LIST.

When any counter in a log page reaches its maximum value, incrementing of all counters in that log page shall cease until re-initialized by the initiator via a LOG SELECT command. If the RLEC bit of the control mode page is one, then the target shall report the exception condition as described in 7.3.3.1.

The page code assignments for the log pages are listed in Table 7-53.

Table 7-53: Log Page Codes

Page Code	Description	Section
01h	Buffer Over-Run/Under-Run Page	7.3.2.1
03h	Error Counter Page (Read) Page	7.3.2.2
04h	Error Counter Page (Read Reverse) Page	7.3.2.2
05h	Error Counter Page (Verify) Page	7.3.2.2
02h	Error Counter Page (Write) Page	7.3.2.2
07h	Last n Error Events Page	7.3.2.3
06h	Non-Medium Error Page	7.3.2.4
00h	Supported Log Pages	7.3.2.5
08h - 2Fh	Reserved	
3Fh	Reserved	
30h - 3Eh	Vendor Specific	

7.3.2.1. Buffer Over-Run/Under-Run Page

The buffer over-run/under-run page (page code 01h) defines 24 data counters that may be used to record the number of buffer over-runs or under-runs for the logical unit. A target that implements this page may implement one or more of the defined data counters.

A buffer over-run or under-run can occur when an initiator does not transmit data to or from the target's buffer fast enough to keep up with reading or writing the media. This can be caused by a slow transfer rate across the SCSI bus or by a high SCSI bus utilization that prevents reconnection by the target. A buffer over-run condition can occur during a read operation when a buffer full condition prevents continued transfer of data from the media to the buffer. A buffer under-run condition can occur during a write operation when a buffer empty condition prevents continued transfer of data to the media from the buffer. Most devices incur a delay at this point while the media is repositioned.

Table 7-54 defines the parameter code field for the buffer over-run/under-run counters.

Table 7-54: Parameter Code Field for Buffer Over-Run/Under-Run Counters

Bit	7	6	5	4	3	2	1	0
Byte	Reserved							
0	Reserved							
1	Count Basis				Cause		Type	

The parameter code field for buffer over-run/under-run counters is a 16-bit value comprised of eight reserved bits, a three-bit count basis field (Table 7-55), a four-bit cause field (Table 7-56), and a one-bit type field. These are concatenated to determine the value of the parameter code for that log parameter. For example, a parameter code value of 0023h specifies a count basis of 001b, a cause of 0001b, and a type of 1b which is a counter that is incremented once per command that experiences an over-run due to the SCSI bus being busy.

The count basis field defines the criteria for incrementing the counter. The following criteria are defined:

Table 7-55: Count Basis Definition

Count Basis	Description
000b	Undefined
001b	Per Command
010b	Per Failed Reconnect
011b	Per Unit of Time
100b-111b	Reserved

IMPLEMENTORS NOTE: The per unit of time count basis is device type specific. Direct-access devices typically use a latency period (i.e., one revolution of the medium) as the unit of time.

The cause field indicates the reason that the over-run or under-run occurred. The following causes are defined:

Table 7-56: Cause Field Definition

Cause	Description
0h	Undefined
1h	SCSI Bus Busy
2h	Transfer Rate Too Slow
3h-Fh	Reserved

The type field indicates whether the counter records under-runs or over-runs. A value of zero specifies a buffer under-run condition and a value of one specifies a buffer over-run condition.

The counters contain the total number of times buffer over-run or under-run conditions have occurred since the last time the counter was cleared. The counter shall be incremented for each occurrence of an under-run or over-run condition and can be incremented more than once for multiple occurrences during the execution of a single command.

7.3.2.2. Error Counter Pages

This section defines the optional error counter pages for write errors (page code 02h), read errors (page code 03h), read reverse errors (page code 04h) and verify errors (page code 05h). Table 7-50 defines the page format for these pages. A page can return one or more log parameters which record events defined by the parameter codes (Table 7-57).

Table 7-57 defines the parameter codes for the error counter pages. Support of each log parameter is optional.

Table 7-57: Parameter Codes for Error Counter Pages

Parameter Code	Description
0000h	Errors corrected without substantial delay
0001h	Errors corrected with possible delays
0002h	Total (e.g., re-writes or re-reads)
0003h	Total errors corrected
0004h	Total times correction algorithm processed
0005h	Total bytes processed
0006h	Total uncorrected errors
0007-7FFFh	Reserved
8000-FFFFh	Vendor specific

IMPLEMENTORS NOTE: The exact definition of the error counters is not part of this standard. These counters should not be used to compare products because the products may define errors differently.

7.3.2.3. Last n Error Events Page

Log page (07h) provides for a number of error-event records using the list parameter format of the log page. The number of these error-event records supported, n, is device-specific. Each error-event record contains device-specific diagnostic information for a single error encountered by the device. The parameter code associated with error-event record indicates the relative time at which the error occurred. A higher parameter code indicates that the error event occurred later in time.

The content of the parameter value field of each log parameter is an ASCII character string which may describe the error event. The exact contents of the character string is not defined by this standard.

When the last supported parameter code is used by an error-event record, the recording on this page of all subsequent error information shall cease until one or more of the list parameters with the highest parameter codes have been re-initialized. If the RLEC bit of the common mode (7.3.3.1) is set to one, the target shall return CHECK CONDITION status with the sense key set to RECOVERED ERROR and the additional sense code set to LOG LIST CODES EXHAUSTED. Alternatively, the target may report this condition via asynchronous event notification (see 6.5.5).

7.3.2.4. Non-Medium Error Page

This page (page code 06h) provides for summing the occurrences of recoverable error events other than write, read, or verify failures. No discrimination among the various types of events is provided by parameter code (Table 7-58). Vendor-specific discrimination may be provided through the vendor-specific parameter codes.

Table 7-58: Non-Medium Error Event Parameter Codes

Parameter Code	Description
0000h	Non-Medium Error Count
0001h - 7FFFh	Reserved
8000h - FFFFh	Vendor-Specific Error Counts

7.3.2.5. Supported Log Pages

The supported log page (Table 7-59) returns the list of log pages implemented by the target. Targets that implement the LOG SENSE command shall implement this log page.

Table 7-59: Supported Log Pages

Bit Byte	7	6	5	4	3	2	1	0
0	Reserved		Page Code (00h)					
1	Reserved							
2	(MSB)		Page Length (n-3)					
3							(LSB)	
4								
n	-		Supported Page List				-	

This page is not defined for the LOG SELECT command. This log page returns the list of supported log pages for the specified logical unit.

The page length field specifies the length in bytes of the following supported page list.

The supported page list field shall contain a list of all log page codes implemented by the target in ascending order beginning with page code 00h.

7.3.3. Mode Parameters

This section describes the block descriptors and the pages used with MODE SELECT and MODE SENSE commands that are applicable to all SCSI devices. Pages specific to each device type are described in the third subsection of each device-type section (i.e., 8.3, 9.3, etc.).

The mode parameter list shown in Table 7-60 contains a header, followed by zero or more block descriptors, followed by zero or more variable-length pages. Parameter lists are defined for each device type.

Table 7-60: Mode Parameter List

Bit Byte	7	6	5	4	3	2	1	0
0 - n	Mode Parameter Header							
0 - n	Block Descriptor(s)							
0 - n	Page(s)							

The six-byte command descriptor block parameter header is defined in Table 7-61. The ten-byte command descriptor block parameter header is defined in Table 7-62.

Table 7-61: Mode Parameter Header(6)

Bit Byte	7	6	5	4	3	2	1	0
0	Mode Data Length							
1	Medium Type							
2	Device-Specific Parameter							
3	Block Descriptor Length							

The block length specifies the length in bytes of each logical block described by the block descriptor. For sequential-access devices, a block length of zero indicates that the logical block size written to the medium is specified by the transfer length field in the command descriptor block (see 9.2.5 and 9.2.15)

The mode page format is defined in Table 7-64.

Table 7-64: Mode Page Format

Bit	7	6	5	4	3	2	1	0
Byte								
0	PS	Reserved	Page Code					
1	Page Length							
2 - n	Mode Parameters							

Each mode page contains a page code, a page length, and a set of mode parameters. The page codes are defined in Table 7-65 and in the mode parameter sub-sections of the specific device type.

When using the MODE SENSE command, a parameters savable (PS) bit of one indicates that the mode page can be saved by the target in a non-volatile, vendor-specific location. A PS bit of zero indicates that the supported parameters cannot be saved. When using the MODE SELECT command, the PS bit is reserved.

The page code field identifies the format and parameters defined for that mode page. Some page codes are defined as applying to all device types (Table 7-65) and other page codes are defined for the specific device type.

When using the MODE SENSE command, if page code 00h (vendor-specific page) is implemented, the target shall return that page last in response to a request to return all pages (page code 3Fh). When using the MODE SELECT command, this page should be sent last.

The page length field specifies the length in bytes of the mode parameters that follow. If the initiator does not set this value to the value that is returned for the page by the MODE SENSE command, the target shall terminate the command with CHECK CONDITION status. The sense key shall be set to to ILLEGAL REQUEST with the additional sense code set to INVALID FIELD IN PARAMETER LIST. The target is permitted to implement a mode page that is less than the full page length defined by this standard, provided no field is truncated and the page length field correctly specifies the actual length implemented.

The mode parameters for each page are defined in the following sub-sections or in the mode parameters sub-section for the specific device type. Mode parameters not implemented by the target shall be set to zero.

Table 7-65 defines the mode pages that are applicable to all device types that include the MODE SELECT and MODE SENSE commands.

Table 7-65: Mode Page Codes

Page Code	Description	Section
0Ah	Control Mode Page	7.3.3.1
02h	Disconnect-Reconnect Page	7.3.3.2
09h	Peripheral Device Page	7.3.3.3
01h	(See specific device type)	
03h - 08h	(See specific device type)	
0Bh - 1Fh	(See specific device type)	
00h	Vendor specific (does not require page format)	
20h - 3Eh	Vendor specific (page format required)	
3Fh	Return all pages (valid only for the MODE SENSE command)	

7.3.3.1. Control Mode Page

Table 7-66: Control Mode Page

Bit Byte	7	6	5	4	3	2	1	0	
0	PS	Reserved		Page Code (0Ah)					
1	Page Length (06h)								
2	Reserved							RLEC	
3	Queue Algorithm Modifier				Reserved		QErr	DQue	
4	EECA	Reserved			RAENP		UAAENP	EAENP	
5	Reserved								
6									
7	---				Ready AEN Holdoff Period				---

The control mode page (Table 7-66) provides controls over several SCSI-2 features which are applicable to all device types such as tagged queuing, extended contingent allegiance, asynchronous event notification, and error logging.

A report log exception condition (RLEC) bit of one specifies that the target shall report log exception conditions as described in 7.3.2. A RLEC bit of zero specifies that the target shall not report log exception conditions.

The queue algorithm modifier field (Table 7-67) specifies restrictions on the algorithm used for re-ordering commands that are tagged with the SIMPLE QUEUE TAG message.

Table 7-67: Queue Algorithm Modifier

Value	Definition
0h	Guaranteed data integrity
1h	Unrestricted re-ordering allowed
2h - 7h	Reserved
8h - Fh	Vendor Specific

A value of zero in this field specifies that the target shall order the actual execution sequence of the queued commands such that data integrity is guaranteed at any time. This means that, if the transmission of new commands was halted at any time, the final value of all data observable on the medium shall have exactly the same value as it would have if the commands had been executed in the same received sequence without tagged queuing. The guaranteed data integrity value shall be the default value.

A value of one in this field specifies that the target may re-order the actual execution sequence of the queued commands in any manner it selects. Any data integrity exposures related to command sequence order are explicitly handled by the initiator through the selection of appropriate commands and queue tag messages.

A queue error management (QErr) bit of zero specifies that those commands still queued after the target has entered the contingent allegiance or extended contingent allegiance conditions shall continue execution in a normal manner when that condition has terminated (see 6.8). A QErr bit of one specifies that those commands still queued after the target has entered the contingent allegiance or extended contingent allegiance conditions shall be aborted when that condition has terminated. A unit attention condition shall be generated for each initiator which had commands in the queue except the initiator that received the original INITIATE RECOVERY message. When reporting the unit attention condition, the target shall set the additional sense code to TAGGED COMMANDS CLEARED BY ANOTHER INITIATOR.

A disable queuing (DQue) bit of zero specifies that tagged queuing shall be enabled if the target supports tagged queuing. A DQue bit of one specifies that tagged queuing shall be disabled. Any queued commands for that I T x nexus shall be aborted. Any subsequent queue tag message received shall be rejected with a MESSAGE REJECT message and the I/O process shall be executed as an untagged command (see 6.8.1).

An enable extended contingent allegiance (EECA) bit of one specifies that extended contingent allegiance is enabled (see 6.7). An EECA bit of zero specifies that extended contingent allegiance is disabled.

The RAENP, UAAENP, and EAENP bits enable specific events to be reported via the asynchronous event notification protocol. When all three bits are zero, the target shall not create asynchronous event notifications.

A ready AEN permission (RAENP) bit of one specifies that the target may issue an asynchronous event notification upon completing its initialization sequence instead of generating a unit attention condition. A RAENP bit of zero specifies that the target shall not issue an asynchronous event notification upon completing its initialization sequence.

IMPLEMENTORS NOTE: If the target's default value for the RAENP bit is one and it does not implement saved parameters or include a hardware switch, then it may not be possible to disable the initialization sequence asynchronous event notification.

A unit attention AEN permission (UAAENP) bit of one specifies that the target may issue an asynchronous event notification instead of creating a unit attention condition upon detecting an event which would cause a unit attention condition (other than upon completing an initialization sequence). A UAAENP bit of zero specifies that the target shall not issue an asynchronous event notification instead of creating a unit attention condition.

An error AEN permission (EAENP) bit of one specifies that the target may issue an asynchronous event notification upon detecting a deferred error condition instead of waiting to report the deferred error on the next command. An EAENP bit of zero specifies that the target shall not report deferred error conditions via an asynchronous event notification.

The ready AEN holdoff period field specifies the minimum time in milliseconds after the target starts its initialization sequence that it shall delay before attempting to issue an asynchronous event notification. This value may be rounded up as defined in 6.5.4.

Table 7-62: Mode Parameter Header(10)

Bit Byte	7	6	5	4	3	2	1	0
0	(MSB)	Mode Data Length						(LSB)
1		Medium Type						
2		Device-Specific Parameter						
3		Reserved						
4		Reserved						
5		Block Descriptor Length						
6	(MSB)							(LSB)
7								

When using the MODE SENSE command, the mode data length field specifies the length in bytes of the following data that is available to be transferred. The mode data length does not include itself. When using the MODE SELECT command, this field is reserved.

IMPLEMENTORS NOTE: Targets that support more than 256 bytes of block descriptors and pages may need to implement ten-byte mode commands. The mode data length field in the six-byte command descriptor block header limits the returned data to 256 bytes.

Medium types are unique for each device type. Refer to the mode parameters section of the specific device type for definition of these values. Some device types reserve this field.

The device specific parameter is unique for each device type. Refer to the mode parameters section of the specific device type for definition of this field. Some device types reserve all or part of this field.

The block descriptor length specifies the length in bytes of all the block descriptors. It is equal to the number of block descriptors times eight and does not include pages or vendor-specific parameters, if any, that may follow the last block descriptor. A block descriptor length of zero indicates that no block descriptors are included in the mode parameter list. This condition shall not be considered an error.

The mode parameter block descriptor is shown in Table 7-63.

Table 7-63: Mode Parameter Block Descriptor

Bit Byte	7	6	5	4	3	2	1	0
0	Density Code							
1	(MSB)	Number of Blocks						(LSB)
2		Reserved						
3		Block Length						
4								
5	(MSB)							(LSB)
6								
7								

Block descriptors specify some of the medium characteristics for all or part of a logical unit. Support for block descriptors is optional. Each block descriptor contains a density code field, a number of blocks field, and a block length field. Block descriptor values are always current (i.e., saving is not supported). A unit attention condition (see 6.9) shall be generated when any block descriptor values are changed.

The density code field is unique for each device type. Refer to the mode parameters section of the specific device type for definition of this field. Some device types reserve all or part of this field.

The number of blocks field specifies the number of logical blocks on the medium to which the density code and block length fields apply. A value of zero indicates that all of the remaining logical blocks of the logical unit shall have the medium characteristics specified.

IMPLEMENTORS NOTES:

(1) There may be implicit association between parameters defined in the pages and block descriptors. For direct-access devices, the block length affects the optimum values (the values that achieve best performance) for the sectors per track, bytes per physical sector, track skew factor, and cylinder skew factor fields in the format parameters page. In this case, the target may change parameters not explicitly sent with the MODE SELECT command. A subsequent MODE SENSE command would reflect these changes.

(2) The number of remaining logical blocks may be unknown for some device types.

7.3.3.2. Disconnect-Reconnect Page

Table 7-68: Disconnect-Reconnect Page

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved	Page Code (02h)					
1	Page Length (0Eh)							
2	Buffer Full Ratio							
3	Buffer Empty Ratio							
4	(MSB)	Bus Inactivity Limit						(LSB)
5								
6	(MSB)	Disconnect Time Limit						(LSB)
7								
8	(MSB)	Connect Time Limit						(LSB)
9								
10	(MSB)	Maximum Burst Size						(LSB)
11								
12	Reserved				DTDC			
13	Reserved							
14	Reserved							
15	Reserved							

The disconnect-reconnect page (Table 7-68) provides the initiator the means to tune the performance of the SCSI bus.

The buffer full ratio field indicates to the target, on read operations, how full the buffer should be prior to attempting a reselection. Targets that do not implement the requested ratio should round down to the nearest implemented ratio as defined in 6.5.4.

The buffer empty ratio field indicates to the target, on write operations, how empty the buffer should be prior to attempting a reselection. Targets that do not implement the requested ratio should round down to the nearest implemented ratio as defined in 6.5.4.

The buffer full and buffer empty ratios are numerators of a fractional multiplier that has 256 as its denominator. A value of zero indicates that the target determines when to initiate reselection consistent with the disconnect time limit parameter. These parameters are advisory to the target.

IMPLEMENTORS NOTE: As an example, consider a target with ten 512-byte buffers and a specified buffer full ratio of 3Fh. The formula is: $\text{INTEGER}((\text{ratio}/256) * \text{number of buffers})$. Thus $\text{INTEGER}(3Fh/256 * 10) = 2$. The target should attempt to reselect the initiator on read operations whenever two or more buffers are full.

The bus inactivity limit field indicates the maximum time in 100 microsecond increments that the target is permitted to assert the BSY signal without a REQ/ACK handshake. If the bus inactivity limit is exceeded the target shall attempt to disconnect if the initiator has granted the disconnect privilege (see 5.6.7) and it is not restricted by DTDC. This value may be rounded as defined in 6.5.4. A value of zero indicates that there is no bus inactivity limit.

The disconnect time limit field indicates the minimum time in 100 microsecond increments that the target shall wait after releasing the SCSI bus before attempting reselection. This value may be rounded as defined in 6.5.4. A value of zero indicates that there is no disconnect time limit.

The connect time limit field indicates the maximum time in 100 microsecond increments that the target is allowed to use the SCSI bus before disconnecting if the initiator has granted the disconnect privilege (see 5.6.7) and it is not restricted by DTDC. This value may be rounded as defined in 6.5.4. A value of zero indicates that there is no connect time limit.

The maximum burst size field indicates the maximum amount of data that the target shall transfer during a data phase before disconnecting if the initiator has granted the disconnect privilege. This value is expressed in increments of 512 bytes (e.g., a value of one means 512 bytes, two means 1024 bytes, etc.). A value of zero indicates there is no limit on the amount of data transferred per connection.

The data transfer disconnect control (DTDC) field (Table 7-69) defines further restrictions on when a disconnect is permitted.

Table 7-69: Data Transfer Disconnect Control

DTDC	Description
00b	Data transfer disconnect control is not used. Disconnect is controlled by the other fields in this page.
01b	A target shall not attempt to disconnect once the data transfer of a command has started until all data the command is to transfer has been transferred. The connect time limit and bus inactivity limit are ignored during the data transfer.
10b	Reserved
11b	A target shall not attempt to disconnect once the data transfer of a command has started until the command is complete. The connect time limit and bus inactivity limit are ignored once data transfer has started.

If DTDC is nonzero and the maximum burst size is nonzero the target shall return CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST and the additional sense code set to ILLEGAL FIELD IN PARAMETER LIST.

7.3.3.3. Peripheral Device Page

Table 7-70: Peripheral Device Page

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved	Page Code (09h)					
1	Page Length (n-1)							
2	(MSB)							
3	Interface Identifier (LSB)							
4	Reserved							
5	Reserved							
6	Reserved							
7	Reserved							
8	Reserved							
n	Vendor Specific							

The peripheral device page (Table 7-70) is used to pass vendor-specific information between an initiator and a peripheral interface below the target (i.e., between the target and the peripheral device). This standard does not define the format of this data, except to provide a standard header.

Interface identifier codes are defined in the Table 7-71.

Table 7-71: Interface Identifier Codes

Code Value	Interface	Reference Standard
0000h	Small Computer System Interface	X3.131
0001h	Storage Module Interface	X3.91M-1987
0002h	Enhanced Small Device Interface	X3.170
0003h	IPI-2	X3.130-1986; X3T9.3/87-002
0004h	IPI-3	X3.132-1987; X3.147-1988
0005h-7FFFh	Reserved	
8000h-FFFFh	Vendor Specific	

7.3.4. Vital Product Data Parameters

This section describes the optional vital product data page structure and the vital product data pages (Table 7-72) that are applicable to all SCSI devices. These pages are optionally returned by the INQUIRY command (7.2.5) and contain vendor-specific product information about a target or logical unit. The vital product data may include vendor identification, product identification, unit serial numbers, device operating definitions, manufacturing data, field replaceable unit information, and other vendor-specific information. This standard defines the structure of the vital product data, but not the contents.

Table 7-72: Vital Product Data Page Codes

Page Code	Description	Section
82h	ASCII Implemented Operating Definition Page	7.3.4.1
01h - 7Fh	ASCII Information Page	7.3.4.2
81h	Implemented Operating Definitions Page	7.3.4.3
00h	Supported Vital Product Data Pages	7.3.4.4
80h	Unit Serial Number Page	7.3.4.5
83h - BFh	Reserved	
COh - FFh	Vendor Specific	

7.3.4.1. ASCII Implemented Operating Definition Page

Table 7-73: ASCII Implemented Operating Definition

Bit Byte	7	6	5	4	3	2	1	0
0	Peripheral Qualifier			Peripheral Device Type				
1	Page Code (82h)							
2	Reserved							
3	Page Length (n-3)							
4	ASCII Operating Definition Description Length (m-4)							
5	ASCII Operating Definition Description Data							
m								
m+1	Vendor-Specific Description Data							
n								

This page (Table 7-73) contains operating definition description data for all operating definitions implemented by the target. The contents of this data is not defined by this standard.

The peripheral qualifier field is defined in Table 7-16 and the peripheral device type field is defined in Table 7-17.

The page length field specifies the length of the following page data. If the allocation length is less than the length of the data to be returned the page length shall not be adjusted to reflect the truncation.

The ASCII operating definition description length field specifies the length in bytes of the ASCII operating definition description data that follows. If the allocation length is less than the length of data to be returned, the ASCII operating definition description length shall not be adjusted to reflect the truncation. A value of zero in this field indicates that no ASCII operating definition description data is available.

The ASCII operating definition description data field contains the ASCII operating definition description data for the target or logical unit. The data in this field shall be formatted in lines (or character strings). Each line shall contain only graphic codes (i.e., code values 20h through 7Eh) and shall be terminated with a NULL (00h) character.

7.3.4.2. ASCII Information Page

Table 7-74: ASCII Information Page

Bit	7	6	5	4	3	2	1	0
Byte								
0	Peripheral Qualifier			Peripheral Device Type				
1				Page Code (01h - 7Fh)				
2				Reserved				
3				Page Length (n-3)				
4				ASCII Length (m-4)				
5								
m	-	-		ASCII Information				-
m+1								
n	-	-		Vendor-Specific Information				-

The ASCII information page (Table 7-74) returns information for the field replaceable unit code returned in the REQUEST SENSE sense data (see 7.2.14).

The peripheral qualifier field is defined in Table 7-16 and the peripheral device type field is defined in Table 7-17.

The page code field contains the same value as in the page code field of the INQUIRY command descriptor block (see 7.2.5) and is associated with the field replaceable unit code returned by the REQUEST SENSE command.

IMPLEMENTORS NOTE: The field replaceable unit field in the sense data provides for 255 possible codes, while the page code field provides for only 127 possible codes. Thus it is not possible to return ASCII information pages for the upper code values.

The page length field specifies the length of the following page data. If the allocation length of the command descriptor block is too small to transfer all of the page, the page length shall not be adjusted to reflect the truncation.

The ASCII length field specifies the length in bytes of the ASCII information that follows. If the allocation length is less than the length of the data to be returned the ASCII length shall not be adjusted to reflect the truncation. A value of zero in this field indicates that no ASCII information is available for the specified page code.

The ASCII information field contains ASCII information concerning the field replaceable unit identified by the page code. The data in this field shall be formatted in one or more lines (or character strings). Each line shall contain only graphic codes (i.e., code values 20h through 7Eh) and shall be terminated with a NULL (00h) character.

The contents of the vendor-specific information field is not defined in this standard.

7.3.4.3. Implemented Operating Definition Page

Table 7-75: Implemented Operating Definition Page

Bit	7	6	5	4	3	2	1	0
Byte								
0	Peripheral Qualifier			Peripheral Device Type				
1	Page Code (81h)							
2	Reserved							
3	Page Length (n-3)							
4	Reserved	Current Operating Definition						
5	SavImp	Default Operating Definition						
6	SavImp	Supported Operating Definition List						
n	SavImp							

The implemented operating definition page (Table 7-75) defines the current operating definition, the default operating definition, and which operating definitions are implemented by the target. These operating definition values are specified in the CHANGE DEFINITION command (see 7.2.1).

The peripheral qualifier field is defined in Table 7-16 and the peripheral device type field is defined in Table 7-17.

The page length field specifies the length of the following operating definitions. If the allocation length of the command descriptor block is too small to transfer all of the page, the page length shall not be adjusted to reflect the truncation.

For each operating definition there is an associated save implemented (SavImp) bit. A SavImp bit of zero indicates that the corresponding operating definition parameter cannot be saved. A SavImp bit of one indicates that the corresponding operating definition parameter can be saved.

All returned operating definitions use the codes defined in Table 7-3. The current operating definition field returns the value of the present operating definition. The default operating definition field returns the value of the operating definition the target uses when power is applied if no operating definition is saved. The supported operating definition list returns one or more operating definitions implemented by the target.

7.3.4.4. Supported Vital Product Data Pages

Table 7-76: Supported Vital Product Data Pages

Bit Byte	7	6	5	4	3	2	1	0
0	Peripheral Qualifier			Peripheral Device Type				
1	Page Code (00h)							
2	Reserved							
3	Page Length (n-3)							
4								
- -	Supported Page List							
n								

The supported vital product data pages are shown in Table 7-76.

The peripheral qualifier field is defined in Table 7-16 and the peripheral device type field is defined in Table 7-17.

The page code field shall be set to the value of the page code field in the INQUIRY command descriptor block (see 7.2.5).

The page length field specifies the length of the supported page list. If the allocation length is too small to transfer all of the page, the page length shall not be adjusted to reflect the truncation.

The supported page list field shall contain a list of all vital product data page codes implemented for the target or logical unit in ascending order beginning with page code 00h.

7.3.4.5. Unit Serial Number Page

Table 7-77: Unit Serial Number Page

Bit	7	6	5	4	3	2	1	0
Byte								
0	Peripheral Qualifier				Peripheral Device Type			
1	Page Code (80h)							
2	Reserved							
3	Page Length (n-3)							
8								
n	Product Serial Number							

This page (Table 7-77) provides a product serial number for the target or logical unit.

The peripheral qualifier field is defined in Table 7-16 and the peripheral device type field is defined in Table 7-17.

The page length field specifies the length of the product serial number. If the allocation length is too small to transfer all of the page, the page length shall not be adjusted to reflect the truncation.

The product serial number field contains ASCII data that is vendor specific. The least significant ASCII character of the serial number shall appear as the last byte of a successful data transfer. If the product serial number is not available, the target shall return ASCII spaces (20h) in this field.

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8. Direct-Access Devices

8.1. Direct-Access Device Model

Direct-access devices store blocks of data for later retrieval. Each block of data is stored at a unique logical block address. An initiator issues WRITE commands to store the blocks of data (write operations) and READ commands to retrieve the blocks of data (read operations). Other commands issued by the initiator to may also cause write and read operations to occur. A write operation causes a block of data to be written on the medium. A read operation causes a block of data to be read from the medium. A verify operation causes a block of data to be verified that it can be read without error from the medium.

Blocks of data are stored by a process that causes localized changes or transitions within the medium. The changes made to the medium to store the blocks of data may be volatile (i.e., not retained through power cycles) or non-volatile (retained through power cycles). The medium may be divided in parts that are used to data blocks, parts that are reserved as spares for defect handling, and parts that are reserved for use by the controller for the management of the device.

8.1.1. Removable Medium

The medium may be removable (i.e., a typically used in a floppy disk drive) or non-removable (i.e., typically used in a hard disk drive). Removable medium is contained within a cartridge (or jacket) to prevent damage to the recording surfaces. The combination of medium and cartridge is often called a volume.

A volume has an attribute of being mounted or demounted on a suitable transport mechanism. A volume is mounted when the direct-access device is capable of performing write or read operations to the medium. A mounted volume may be not be accessible by an initiator if it is reserved by another initiator. A volume is demounted at any other time (e.g., during loading, unloading, or storage).

An initiator may check whether a volume is mounted by issuing a TEST UNIT READY command. A volume that is loaded may need a START/STOP UNIT command issued to become accessible for write or read operations.

The PREVENT/ALLOW MEDIUM REMOVAL command allows an initiator to restrict the demounting of the volume. This is useful in maintaining system integrity. If the direct-access device implements cache memory, it must ensure that all logical blocks of the medium contain the most recent data prior to permitting demounting of the volume. If the initiator issues a START/STOP UNIT command to eject the cartridge, and the direct-access device is prevented from demounting by the PREVENT/ALLOW MEDIUM REMOVAL command, the START/STOP command is rejected by the direct-access device.

8.1.2. Logical Blocks

Blocks of data are stored on the medium along with additional information that the controller uses to manage the storage and retrieval. The format of the additional information is unique and is hidden from the initiator during normal read or write operations. This additional information is often used to identify the physical location of the blocks of data, the address of the logical block, and to provide protection against the loss of the user data.

The address of the first logical block is zero. The address of the last logical block is $[n-1]$, where $[n]$ is the number of logical blocks available on the medium. A READ CAPACITY command may be issued to determine the value of $[n-1]$. If a command is issued which requests access to a logical block not within the capacity of the medium the command is terminated with CHECK CONDITION.

The number of bytes of data contained in a logical block is known as the block length. Each logical block has a block length associated with it. The block length may be different for each logical block on the medium. However, in a typical device only one block length is used at a time. The block descriptor in the MODE SENSE data describes the block lengths that are used on the medium. An MODE SELECT command can be used to setup extents. An extent is a specified number of logical blocks that have the specified block length. The FORMAT UNIT command is typically required to change the block length of devices that support variable block lengths and make the extents that were setup active.

The location of a logical block on the medium does not have a relationship to the location of any other logical block. However, in a typical device the logical blocks are located in an ascending order. The time to access the logical block at address $[x]$ and then the logical block at address $[x+1]$ may not be less than time to access $[x]$ and then $[x+100]$. The READ CAPACITY with a PMI bit of one is useful in determining where longer access times occur.

8.1.3. Ready State

A direct-access device is ready when medium access commands can be executed. A device using removable media is usually not ready until a volume is mounted. Such a device normally returns CHECK CONDITION status and sets the sense key to NOT READY.

A direct-access device may also be switched between the ready state and the not ready state by using the START/STOP UNIT command. An initiator may need to issue a START UNIT to switch them to the ready state.

8.1.4. Initialization

Many direct-access devices must be initialized prior to write or read operations occurring. This initialization is usually performed by a FORMAT UNIT command. Parameters related to the geometry and performance characteristics can be set with the MODE SELECT command prior to the format operation. Some devices will be initialized by means not specified in this standard. The time at which this occurs is specific to the implementation of the direct-access device.

Devices using non-volatile medium typically save the parameters and only need to be initialized once. However, some mode parameters may need to be initialized after each power-on and reset. A catastrophic failure of the direct-access device may require the FORMAT UNIT command to be reissued.

Devices using volatile medium may need to be initialized at each power-on cycle prior read or write operations being performed. Mode parameters may also need initialization.

8.1.5. Medium Defects

Any medium has the potential for defects which can cause user data to be lost. Therefore, each logical block may contain information which allows the detection of changes to the user data caused by defects in the medium or other phenomena, and may also allow the data to be reconstructed following the detection of such a change. On some devices, the initiator has some control some through use of the mode parameters. Some devices may allow the initiator to examine and modify the additional information by using the READ LONG and WRITE LONG commands. Some media having a very low probability of defects in may not require these structures.

Defects may also be detected and managed during execution of the FORMAT UNIT command. The FORMAT UNIT command defines four sources of defect information. These defects may be reassigned or avoided during the initialization process so that they do not appear in a logical block.

Defects may also be avoided after initialization. The initiator issues a REASSIGN BLOCKS command to request that the specified logical block address be reassigned to a different part of the medium. This operation can be repeated if a new defect appears at a later time. The total number of defects that may be handled in this manner can be specified in the mode parameters.

8.1.6. Data Cache

Some direct-access devices implement some form of cache memory. A cache memory is usually an area of temporary storage on the direct-access with a fast access time that is used to enhance the performance of the device. It exists separately from the blocks of data stored and is normally not directly accessible by the initiator. Use of the cache memory for write or read operations typically reduces the access time to a logical block and can increase the overall data throughput.

During read operations the direct-access device uses the cache memory to keep blocks of data that the initiator may request at some future time. The algorithm used to manage the cache memory is not part of this standard. However, there are parameters provided by which the initiator may advise the direct-access device about its future requests, or restrict the use of cache memory for a particular request.

During write operations the direct-access device uses the cache memory to buffer data for later storage on the medium. This is typically called a write-back caching algorithm. Thus the command may complete prior to blocks of data being written to the medium. As a result of using write-back caching there is a period of time during which the blocks of data may be lost due to a power or hardware failure. There is also the possibility of an error

occurring during the write operation. If an error occurred during the write, it may be reported as a deferred error on a later command. However, there are parameters provided by which the the initiator may request the use of write-through caching thereby preventing these circumstances from arising.

When the cache memory fills up with blocks of data that are being kept for possible future access, new blocks of data that are to be kept must replace those currently in cache memory. The disable page out (DPO) bit is used to control replacement of logical blocks in the cache. For write operations, setting this bit to one advises the direct-access device to not replace existing blocks in the cache memory with the write data. For read operations, setting this bit to one causes blocks of data that are being read to not replace existing ones in the cache memory.

Sometimes the initiator may wish to have the blocks of data read from the medium instead of from the cache memory. The force unit access (FUA) bit is used to indicate that the direct-access device shall access the physical medium. For a write operation, setting FUA to one causes the direct-access device to complete the data write to the physical medium before completing the command. For a read operation, setting FUA to one causes the logical blocks to be retrieved from the physical medium.

Note that when the DPO and FUA bits are both set to one, write and read operations bypass the cache memory.

For the VERIFY command, an FUA is implied, since the blocks of data stored on the medium are being verified. Further, a SYNCHRONIZE CACHE operation (see below) is also implied to flush an unwritten blocks of data still in the cache memory. These blocks of data must be stored on the medium before the verify operation can begin. The DPO bit is still provided since the VERIFY command may still cause the replacement of blocks in the cache. The above also applies to the WRITE AND VERIFY command.

Certain special commands may be implemented by the direct-access device that allow the initiator to control other behavior of the cache memory:

- LOCK/UNLOCK CACHE controls whether certain logical blocks will be held in the data cache for future use. Locking a logical block prevents its replacement by a future access. Unlocking a logical block exposes it to possible replacement by a future access. (See 8.2.2)
- PRE-FETCH causes a set of logical blocks requested by the initiator to be read into the data cache for possible future access. The blocks fetched are subject to later replacement unless they are locked (see 8.2.3)
- SYNCHRONIZE CACHE forces any pending write data in the requested set of logical blocks to be stored in the physical medium. This command can be used to ensure that the data was written and any errors reported. (see 8.2.18)
- The MODE SELECT command defines a page for the control of cache behavior and handles certain basic elements of cache replacement algorithms (see 8.3.3.1).

8.1.7. Reservation

The access enabled or access disabled state determines when an initiator may store user data blocks in, and/or retrieve user data blocks from part or all of the physical medium. Access may be restricted for read operations, write operations, or both. This attribute may be controlled by an external mechanism or by the RESERVE and RELEASE commands (see 8.2.12 and 8.2.11).

The RESERVE and RELEASE commands define how different types of restricted access may be achieved, and to whom the access is restricted. This section will describe how commands behave relative to the initiator that has a reservation, and to the initiator(s) that are restricted.

In general, the initiator uses reservations to gain a mode of access to extents of logical blocks or the entire direct-access device. The initiator expects this mode to be retained until it is released by whatever mechanism that occurs. The direct-access device behaves in such a way as to ensure that the initiator with the reservation is able to access the reserved media with the operating parameters established by that initiator.

Unless otherwise noted below, if an initiator issues any command to a direct-access device that has a logical unit reservation by another initiator, a reservation conflict will occur.

CHANGE DEFINITION: If any initiator has an extent reservation on a direct-access device, no other initiator may affect the operating definition of that initiator by use of this command. If the direct-access device allows different operating definitions for each initiator, then there is no conflict; otherwise, a reservation conflict occurs.

COMPARE, COPY, COPY AND VERIFY: When a direct-access device is requested to copy to or from itself, those write and read accesses are evaluated for reservation conflict as if they were normal write and read operations. For example, if a COPY is issued to logical unit 0 that requests the direct-access device to copy from logical unit 0 to logical unit 1, access to logical unit 1 must also be evaluated for conflict.

FORMAT UNIT, PREVENT/ALLOW MEDIUM REMOVAL, REZERO UNIT, START/STOP UNIT: If any initiator has an extent reservation on a direct-access device, any of these commands from another initiator will result in a reservation conflict.

INQUIRY and REQUEST SENSE are not affected by any kind of reservation.

LOG SELECT, LOG SENSE, MODE SENSE, TEST UNIT READY, READ CAPACITY (PMI set to zero), READ BUFFER, WRITE BUFFER, and READ DEFECT DATA are not affected by extent reservations.

SEEK, LOCK/UNLOCK CACHE, PRE-FETCH, and SYNCHRONIZE CACHE are evaluated for reservation conflict as if they were normal write and read operations.

MODE SELECT: If an initiator has an extent reservation on a direct-access device, and another initiator attempts one of these command, a reservation conflict occurs if the command affects the manner in which access of the extent by the first initiator is performed. If the command does not affect access to the extent, or parameters are saved for each initiator, then a

conflict does not occur.

SEND DIAGNOSTIC, RECEIVE DIAGNOSTIC RESULTS: These commands conflict with an extent reservation only if they affect access to the extent (as with MODE SELECT).

REASSIGN BLOCKS: An initiator may not reassign a block that is extent reserved to another initiator.

SET LIMITS: The logical blocks specified in the command will conflict if it allows a conflicting access by the same initiator to an extent.

ALL OTHER COMMANDS are of the read or write access nature and are evaluated for conflict as described in the RESERVE command.

When a system is integrated with more than one initiator, there must be agreement between the initiators as to how media is reserved and released during operations, otherwise an initiator may be locked out of access to a target in the middle of an operation. For example, initiator 'A' is performing writes to a direct-access device that are being cached. Then, initiator 'B' issues a RESERVE command to the direct-access device. As a result, initiator 'A' is locked out of issuing a SYNCHRONIZE CACHE command to ensure the integrity of the writes. To prevent this from happening, initiator 'A' should issue a RESERVE prior to the cached write commands.

8.1.8. Seek and Rezero

The SEEK command provides a way for the initiator to position the device in preparation for access to a particular logical block at some later time. Since this positioning action is implicit in other commands, the SEEK command may not be useful with some direct-access devices.

The REZERO UNIT command is provided to bring the direct-access device to a known condition. This standard does not specify the condition. The REZERO UNIT command is used in some devices to position the actuator at cylinder zero. Some devices return GOOD status without attempting any action.

8.1.9. Notched Drives

A notched (also called partitioned or zoned) drive has areas of the medium in which the drive geometry changes. In the simplest case, the entire medium consists of a single notch. Multiple notches are often used to increase capacity of the drive. The notch page is used to indicate the notch for assignment of values to the parameters in the format device page. By sequencing the notch page through each notch, the format device parameters of each notch are set. This is usually done prior to initialization by the FORMAT UNIT command.

8.1.10. Relative Addressing

Relative addressing is a technique useful in accessing structured data in a uniform manner. Relative addressing is only allowed when commands are linked. Section 6.4.3 gives an example of relative addressing and linking for SEARCH DATA commands appropriate to direct-access devices.

The SET LIMITS command is provided to define the limits of a linked chain of relative addressing commands. This gives an additional protection against exceeding a particular set of blocks. The SET LIMITS command has no effect on any other initiator.

8.1.11. Error Reporting

If any of the following conditions occur during the execution of a command the target shall return CHECK CONDITION status. The appropriate sense key and additional sense code should be set. The following table illustrates some error conditions and the applicable sense keys. The table does not provide an exhaustive enumeration of all conditions that may cause the CHECK CONDITION status.

Condition	Sense Key
Invalid logical block address	ILLEGAL REQUEST
Unsupported option requested	ILLEGAL REQUEST
Target reset or medium change since last command from this initiator	UNIT ATTENTION
Self diagnostic failed	HARDWARE ERROR
Unrecovered read error	MEDIUM ERROR HARDWARE ERROR
Recovered read error	RECOVERED ERROR
Overrun or other error that might be resolved by repeating the command	ABORTED COMMAND
Attempt to write on write protected medium	DATA PROTECT

In the case of an invalid logical block address, the sense data information field shall be set to the logical block address of the first invalid address.

In the case of an attempt to read a blank or previously unwritten block, the information field shall be set to the logical block address of the first blank block encountered. The data read up to that block shall be transferred (optical memory and write-once devices only).

In the case of an attempt to write a previously written block and blank checking is enabled, the information field shall be set to the logical block address of the first non-blank block encountered (optical memory and write-once devices only).

8.1.12. Examples

The following examples show some typical variations of the direct-access device. Other variations are possible.

8.1.12.1. Rotating Media

The typical application of a direct-access device is a disk drive. The medium is a disk coated with a material in which flux changes may be induced. The disk drive allows direct and random access to the medium. This is done using an actuator which positions the read/write head, and a rotating disk. Data is stored and retrieved through the interaction of the read/write head and the disk.

The disk is typically divided into cylinders. Each cylinder is typically divided into tracks. Each track is typically divided into sectors. A cylinder is a set of tracks that can be accessed without movement of the actuator. A track is a recording path over which the read/write head travels during one rotation of the disk. A sector is a part of a track that contains the stored data blocks.

A logical block is stored in one or more sectors, or a sector may store more than one logical block. A sector is typically made up of a header, data and a trailer. The header contains a preamble used to synchronize read circuits to the data, an address field to identify the sector, flags to use for defect management, and a checksum that validates the header. The data contains the block of data. The trailer contains the checksum and the error correction information if it is used. The checksum or the error correction information allows the correction data for medium defects.

A disk drive is typically ready when the disks are rotating at the correct speed and the read-write circuitry is powered and ready to access the disks. Some disks, particularly removable disks, require the user to issue load or start commands to bring the disk drive to the ready state.

A disk drive will typically have to be formatted prior to the initial access. Exceptions to this are drives that are formatted at the factory and some optical drives with pre-formatted media (see 12.1). A disk drive format will typically create the headers for each sector and initialize the data field. The MODE SELECT command is often used at format time to establish the geometry (number of heads and tracks, sectors per track, etc.) and defect management scheme. Disk drives are usually non-volatile.

The defect management scheme of a disk drive is often shielded from the user, though some aspects can be evaluated and controlled by the initiator. The direct-access device will usually reserve some sectors and tracks for recording defect tables and for reassigning defective blocks. The READ LONG and WRITE LONG commands will typically access the user data and checksum portions of the data field so that defects may be induced by the initiator to test the defect detection logic of the direct-access device.

Notches find their most typical use in a rotating disk drive. On a disk, the inner tracks are physically shorter than the outer tracks. As a result, if each track is made to store the same number of data bits, the data is packed more densely on the inner tracks than the outer tracks. By using

notches, the outer tracks may be made to contain a different number of sectors than the inner tracks, balancing the data density. This results in increased capacity.

8.1.12.2. Sequential Media

Some tape devices are implemented as a direct access device so that they can be used in disk oriented operating system environments. These devices are sometimes referred to as 'random access tape' or 'floppy tape'. These devices might be thought of as a disk drive with one or a few very long tracks. Access time to a logical block is usually longer than for a disk drive, since the tape must be fast forwarded or rewound to the block. As a result, the SEEK command will often be more useful for a tape than for a disk. The only way an initiator may determine if a direct-access device is a tape is to use the medium type code returned by the MODE SENSE command.

8.1.12.3. Memory Media

Memory media includes devices which are traditionally used for primary storage within computer systems, such as solid state static or dynamic random access memories (SRAM or DRAM), or magnetic core or bubble memory. These devices are typically non-mechanical and therefore the entire physical medium may be accessed in virtually the same access time. The data is typically accessed as a bit or byte and this also speeds access time. Memory devices typically store less data than disks or tapes, and are usually volatile when not protected by battery backup.

8.2. Commands for Direct-Access Devices.

The commands for direct-access devices shall be as shown in Tables 8-1 and 8-2.

Table 8-1: Commands for Direct-Access Devices(1)

Command Name	Operation Code	Type	Section
CHANGE DEFINITION	40h	O	7.2.1
COMPARE	39h	O	7.2.2
COPY	18h	O	7.2.3
COPY AND VERIFY	3Ah	O	7.2.4
FORMAT UNIT	04h	M	8.2.1
INQUIRY	12h	M	7.2.5
LOCK-UNLOCK CACHE	36h	O	8.2.2
LOG SELECT	4Ch	O	7.2.6
LOG SENSE	4Dh	O	7.2.7
MODE SELECT(6)	15h	O	7.2.8
MODE SELECT(10)	55h	O	7.2.9
MODE SENSE(6)	1Ah	O	7.2.10
MODE SENSE(10)	5Ah	O	7.2.11
PRE-FETCH	34h	O	8.2.3
PREVENT-ALLOW MEDIUM REMOVAL	1Eh	O	8.2.4
READ(6)	08h	M	8.2.5
READ(10)	28h	M	8.2.6
READ BUFFER	3Ch	O	7.2.12
READ CAPACITY	25h	M	8.2.7
READ DEFECT DATA	37h	O	8.2.8
READ LONG	3Eh	O	8.2.9
REASSIGN BLOCKS	07h	O	8.2.10
RECEIVE DIAGNOSTIC RESULTS	1Ch	O	7.2.13
RELEASE	17h	M	8.2.11
REQUEST SENSE	03h	M	7.2.14
RESERVE	16h	M	8.2.12
REZERO UNIT	01h	O	8.2.13

Key: M = Command implementation is mandatory.
O = Command implementation is optional.

Table 8-2: Commands for Direct-Access Devices(2)

Command Name	Operation Code	Type	Section
SEARCH DATA EQUAL	31h	O	8.2.14.1
SEARCH DATA HIGH	30h	O	8.2.14.2
SEARCH DATA LOW	32h	O	8.2.14.3
SEEK(6)	08h	O	8.2.15
SEEK(10)	28h	O	8.2.15
SEND DIAGNOSTIC	1Dh	M	7.2.15
SET LIMITS	33h	O	8.2.16
START STOP UNIT	1Bh	O	8.2.17
SYNCHRONIZE CACHE	35h	O	8.2.18
TEST UNIT READY	00h	M	7.2.16
VERIFY	2Fh	O	8.2.19
WRITE(6)	0Ah	M	8.2.20
WRITE(10)	2Ah	M	8.2.21
WRITE AND VERIFY	2Eh	O	8.2.22
WRITE BUFFER	3Bh	O	7.2.17
WRITE LONG	3Fh	O	8.2.23
WRITE SAME	41h	O	8.2.24

Key: M = Command implementation is mandatory.
O = Command implementation is optional.

The following operation codes are vendor-specific: 02h, 05h, 06h, 09h, 0Ch, 0Dh, 0Eh, 0Fh, 10h, 11h, 13h, 14h, 19h, 20h, 21h, 22h, 23h, 24h, 26h, 27h, 29h, 2Ch, 2Dh and C0h through FFh. All remaining operation codes for direct-access devices are reserved for future standardization.

8.2.1. FORMAT UNIT Command

Table 8-3: FORMAT UNIT Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (04h)							
1	Logical Unit Number			FmtData	CmpLst	Defect List Format		
2	Vendor-Specific							
3	(MSB)							
4	Interleave							(LSB)
5	Control							

The FORMAT UNIT command (Table 8-3) formats the medium into initiator addressable logical blocks per the initiator defined options. In addition, the medium may be certified and control structures may be created for the management of the medium and defects. There is no guarantee that the medium has or has not been altered.

The simplest mandatory form of the FORMAT UNIT command (with no format data) accomplishes medium formatting with little initiator control over defect management. The target implementation determines the degree of defect management that is to be performed. Two additional mandatory forms of this command increase the initiator's control over defect management. Several optional forms of this command further increase the initiator's control over defect management, by allowing the initiator to specify which defect list(s) are to be used, to specify defect locations (in several formats), to enable target certification, and to specify what to do in the event that defect lists are not accessible.

The FORMAT UNIT command shall be rejected with RESERVATION CONFLICT status if the logical unit is reserved, or any extent reservation, from any initiator, is active in the specified logical unit.

IMPLEMENTORS NOTE: It is recommended that MODE SELECT parameters (if any) be set prior to issuing the FORMAT UNIT command.

During the execution of the FORMAT UNIT command, the target may perform a medium defect management algorithm (which can be controlled by the initiator, using optional forms of this command). Four sources of defect location information (hereafter called defects) are defined as follows:

Primary defect list (Plist). This is the list of defects, usually supplied by the original manufacturer of the device or medium, that are considered as permanent defects. The Plist is located outside of the initiator-accessible logical block space. The Plist is accessible by the target (to reference while formatting), but it is not normally accessible by the initiator except through the READ DEFECT DATA command. Once created, the original Plist shall not be subject to change.

Target certification list (Clist). This list includes defects detected by the target during an optional certification process executed during the FORMAT UNIT command. This list shall be added to the Glist.

Data defect list (Dlist). This list of defect descriptors may be supplied to the target by the initiator in the DATA OUT phase of the FORMAT UNIT command as shown in Table 8-4. This list shall be added to the Glist. The defect list length in the defect list header may be zero, in which case there is no Dlist.

Grown defect list (Glist). The Glist includes all defects sent by the initiator or detected by the target. The Glist does not include the Plist. If the CmpLst bit is zero the Glist shall include Dlists provided to the target during the previous and the current FORMAT UNIT commands. The Glist shall also include:

- (1) Defects detected by the format operation during medium certification.
- (2) Defects previously identified with a REASSIGN BLOCKS command.
- (3) Defects previously detected by the target and automatically reallocated.

A format data (FmtData) bit of one indicates that a DATA OUT phase shall take place during the command execution. The DATA OUT phase consists of a defect list header (see Table 8-4), followed by an initialization pattern descriptor (see Table 8-4), followed by zero or more defect descriptors (see Tables 8-7 to 8-9). Each defect descriptor identifies a location on the medium that the target shall map out of the user-accessible area.

A FmtData bit of zero indicates that a DATA OUT phase shall not occur. The source of defect information is not specified.

A complete list (CmpLst) bit of one indicates that the defect list sent by the initiator is a complete list of defects. Any existing defect list except the Plist shall be ignored by the target. As a result a new Glist is constructed which contains the Dlist (if it is sent by the initiator), and the Clist (if certification is enabled). The target may add any defects it detects during the format operation to this Dlist.

A CmpLst bit of zero indicates that the defect list sent by the initiator is an addition to existing list of defects. As a result a new Glist is constructed which contains the existing Glist, the Dlist (if it is sent by the initiator) and the Clist (if certification is enabled). The target may add any defects it detects during the format operation to this Dlist.

The defect list format field specifies which defect descriptor is used if the FmtData bit is one (see Table 8-5).

The interleave field specifies the interleave that is used when performing the format operation. This allows the logical blocks to be related in a way that facilitates matching the transfer rate between the initiator and the peripheral. An interleave of zero specifies that the target use its default interleave. An interleave of one specifies that consecutive logical blocks be placed in contiguous ascending order. All other values are vendor specific.

Table 8-4: FORMAT UNIT Parameter List

Bit Byte	7	6	5	4	3	2	1	0
Defect List Header								
Initialization Pattern Descriptor (if any)								
Defect Descriptor(s) (if any)								
Defect Descriptor 0 (See specific table for length.)								
Defect Descriptor n (See specific table for length.)								

Table 8-4A: Defect List Header

Bit Byte	7	6	5	4	3	2	1	0
0	Reserved							
1	FOV	DPRY	DCRT	STPF	IP	DSP	Immed	vs
2	(MSB)							
3	Defect List Length							(LSB)

The defect list header (TABLE 8-4A) provides several optional format control bits. Targets that implement these bits give the initiator additional control over the use of the four defect sources, and the formatting operation. If the initiator attempts to select any function not implemented by the target, the target shall terminate the command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST and the additional sense code shall be set to INVALID FIELD IN PARAMETER LIST.

A format options valid (FOV) bit of zero indicates that the target shall use its default settings for the DPRO, DCRT, STPF, IP and DSP bits (see below). The initiator shall set these bits to zero. If any of these bits are not zero, the target shall terminate the command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST and the additional sense code shall be set to INVALID FIELD IN PARAMETER LIST.

A FOV bit of one indicates that the target shall examine the setting of the DPRO, DCRT, STPF, IP and DSP bits. When FOV is one the DPRO, DCRT, STPF, IP and DSP bits are defined below.

A disable primary (DPRO) bit of zero indicates that the target shall not use portions of the medium identified as defective in the primary defect Plist for initiator addressable logical blocks. If the target cannot locate the Plist or it cannot determine whether a Plist exists, it shall perform the action specified by the STPF bit. A DPRO bit of one indicates that the target shall not use the Plist to identify defective areas of the medium. The Plist is not deleted.

A disable certification (DCRT) bit of zero indicates that the target shall perform a vendor-specific medium certification operation to generate a Clist. A DCRT bit of one indicates that the target shall not perform any vendor-specific medium certification process or format verification operation while executing the FORMAT UNIT command.

The stop format (STPF) bit controls the behavior of the target when one of the following events occurs:

(1) The target has been requested to use the primary defect list (DPRO is set to zero), or the grown defect list (Cmplst is set to zero) and the target cannot locate the list nor determine whether the list exists.

(2) The target has been requested to use the primary defect list (DPRO is set to zero) or the grown defect list (Cmplst is set to zero), and the target encounters an error while accessing the defect list.

A STPF bit of zero indicates that, if one or both of the above conditions occurs, the target shall continue to execute the FORMAT UNIT command. The target shall return CHECK CONDITION status at the completion of the FORMAT UNIT command. The sense key shall be set to RECOVERED ERROR and the additional sense code shall be set to either DEFECT LIST NOT FOUND if condition one occurred, or DEFECT LIST ERROR if condition two occurred.

A STPF bit of one indicates that, if one or both of the above conditions occurs, the target shall terminate the FORMAT UNIT command with CHECK CONDITION status. The sense key shall be set to MEDIUM ERROR and the additional sense code shall be set to either DEFECT LIST NOT FOUND if condition one occurred, or DEFECT LIST ERROR if condition two occurred.

IMPLEMENTORS NOTE: The use of the FmtData bit, the Cmplst bit, and the defect header allow the initiator to control the source of the defect lists used by the FORMAT UNIT command. Setting the defect list length to zero allows the initiator to control the use of Plist and Clist without having to specify a Dlist.

Table 8-5 defines the implementation requirements for the FORMAT UNIT command.

Table 8-5: FORMAT UNIT Defect Descriptor Format and Requirements

FmtData		Defect List Format	Defect List Length	Command Type	Comments
Cmplst					
0	0	000b	N/A	Mandatory	Target-defined use of defect sources.
BLOCK FORMAT:					
1	0	000b	Zero	Mandatory	(1),(3)
1	1	000b	Zero	Mandatory	(1),(4)
1	0	000b	>0	Optional	(2),(3)
1	1	000b	>0	Optional	(2),(4)
BYTES FROM INDEX FORMAT:					
1	0	100b	Zero	Optional	(1),(3)
1	1	100b	Zero	Optional	(1),(4)
1	0	100b	>0	Optional	(2),(3)
1	1	100b	>0	Optional	(2),(4)
PHYSICAL SECTOR FORMAT:					
1	0	101b	Zero	Optional	(1),(3)
1	1	101b	Zero	Optional	(1),(4)
1	0	101b	>0	Optional	(2),(3)
1	1	101b	>0	Optional	(2),(4)
1	0	110b	Vendor-Specific		
1	1	110b			
All other codes				Reserved	

NOTE: All options described above cause a new Glist to be created during the execution of the FORMAT UNIT command as described in the text above.

NOTES:

- (1) No Dlist is transferred to the target during the DATA OUT phase.
- (2) A Dlist is transferred to the target during the DATA OUT phase. Add the Dlist defects to the new Glist.
- (3) Use the existing Glist as a defect source. Add existing Glist defects to the new Glist.
- (4) Discard the existing Glist. Do not add existing Glist defects to the new Glist.

A disable saving parameters (DSP) bit of one, specifies that the target shall not save the MODE SELECT savable parameters to non-volatile memory during the format operation. A DSP bit of zero specifies that the target shall save all the MODE SELECT savable parameters for all initiators to non-volatile memory during the format operation.

An immediate (Immed) bit of zero indicates that status shall be returned after the format operation has completed. An Immed bit value of one indicates that the target shall return status as soon as the command descriptor block has been validated, and the entire defect list has been transferred. During the format operation, the target shall respond to commands as follows:

(1) In response to all commands except REQUEST SENSE and INQUIRY, the target shall return CHECK CONDITION status.

(2) In response to the INQUIRY command, the target shall respond as commanded.

(3) In response to the REQUEST SENSE command, unless an error has occurred, the target shall return a sense key of NOT READY and an additional sense code of LOGICAL UNIT NOT READY FORMAT IN PROGRESS, with the sense key specific bytes set for progress indication (as described in 7.2.14.1). Refer to 7.2.14.2 for a description of deferred error handling which may occur during the format operation.

The defect list length field in the defect list header specifies the total length in bytes of the defect descriptors that follow and does not include the initialization pattern descriptor or initialization pattern, if any. The length of the defect descriptors varies with the format of the defect list. The three formats for the defect descriptor(s) field in the defect lists are shown in Tables 8-8, 8-9, and 8-10. In Table 8-8, the defect list length is equal to four times the number of defect descriptors. In Tables 8-9 and 8-10, the defect list length is equal to eight times the number of defect descriptors.

8.2.1.1. Initialization Pattern Option

The initialization pattern option specifies that the logical blocks contain the specified initialization pattern. The initialization pattern descriptor (Table 8-6A) is sent to the target as part of the FORMAT UNIT parameter list.

Table 8-6A: Initialization Pattern Descriptor

0	IP Modifier	Reserved
1	Pattern Type	
2	(MSB)	Initialization Pattern Length
3		(LSB)
0 - n	Initialization Pattern	

IMPLEMENTORS NOTE: The initialization pattern option is not intended for media analysis or certification. This option may only initialize the initiator accessible area of the media to the specified pattern and may not write to any initiator inaccessible areas of the disk.

An initialization pattern (IP) bit of one indicates that an initialization pattern descriptor is included in the FORMAT UNIT parameter list immediately following the defect list header. An IP bit of zero indicates that an initialization pattern descriptor is not included and that the target shall use its default initialization pattern.

The IP modifier field specifies the type and location of a header that modifies the initialization pattern (Table 8-6).

Table 8-6: Initialization Pattern Modifier

IP Modifier	Description
0 0	No header. The target shall not modify the the initialization pattern.
0 1	The target shall overwrite the initialization pattern to write the logical block address in the first four bytes of the logical block. The logical block address shall be written with the most significant byte first.
1 0	The target shall overwrite the initialization pattern to write the logical block address in the first four bytes of each physical block contained within the logical block. The lowest numbered logical block or part thereof that occurs within the physical block is used. The logical block address shall be written with the most significant byte first.
1 1	Reserved.

The initialization pattern type field indicates the type of pattern the target shall use to initialize each logical block within the initiator accessible portion of the disk (Table 8-7). All bytes within a logical block shall be written with the initialization pattern. The initialization pattern is modified by the IP modifier field as described above in Table 8-6.

Table 8-7: Initialization Pattern Type

Pattern Type	Note	Description
00h	(1)	Use default pattern
01h	(2)	Repeat the initialization pattern as required to fill the logical block
02-7Fh	none	Reserved
80-FFh	none	Vendor-Specific

NOTES:

(1) If the initialization pattern length is not zero the target shall terminate the command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST and the additional sense code shall be set to INVALID FIELD IN PARAMETER LIST.

(2) If the initialization pattern length is zero the target shall terminate the command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST and the additional sense code shall be set to INVALID FIELD IN PARAMETER LIST.

The initialization pattern length field indicates the number of bytes contained in the initialization pattern. If the length exceeds the current logical block size the target shall terminate the command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST and the additional sense code shall be set to INVALID FIELD IN PARAMETER LIST. The pattern is modified by the IP modifier field.

8.2.1.2. Defect List Formats

This section describes the defect list formats used in the FORMAT UNIT, READ DEFECT DATA and translate page of the SEND DIAGNOSTIC and RECEIVE DIAGNOSTIC RESULTS commands.

IMPLEMENTORS NOTE: The selected reporting format must account for variables which can impact the information in the returned. For example, the specific location of a defect, while constant in angular and radial location on the device, may change in reported location a format operation with different geometry parameters is performed. It is the responsibility of the initiator to use a defect list format appropriate for the intended operation with the current or future geometry parameters. If the target is able to detect that the selected defect list format would provide inconsistent results, the target may return CHECK CONDITION status.

Table 8-8: Defect Descriptor - Block Format

Byte	Defect Descriptor
0	(MSB)
3	(LSB)

Each defect descriptor for the logical block format specifies a four-byte defective block address that contains the defect. The defect descriptors should be in ascending order. More than one physical or logical block may be affected by each defect descriptor.

A target may return CHECK CONDITION if the defect descriptors are not in ascending order.

Table 8-9: Defect Descriptor - Bytes From Index Format

Byte	Defect Descriptor
0	(MSB)
2	(LSB)
3	
4	(MSB)
7	(LSB)

Each bytes from index defect descriptor specifies the location of a defect which is no more than eight-bytes in length. Each descriptor is comprised of the cylinder number of defect, the head number of defect, and the defect bytes from index to the defect. The defect descriptors shall be in ascending order. The cylinder number of defect is the most significant part of the address and the defect bytes from index is the least significant part of the address. More than one physical or logical block may be affected by each defect.

A defect bytes from index of FFFFFFFFh indicates that the entire track shall be considered defective.

Table 8-10: Defect Descriptor - Physical Sector Format

Byte	Defect Descriptor
0	(MSB)
2	Cylinder Number of Defect (LSB)
3	Head Number of Defect
4	(MSB)
7	Defective Sector Number (LSB)

Each physical sector defect descriptor specifies the location of a defect that is the length of a sector. Each descriptor is comprised of a cylinder number of defect, the head number of defect, and the defective sector number. The defect descriptors shall be in ascending order. The cylinder number of defect is the most significant part of the address and the defective sector number is the least significant part of the address. More than one block may be affected by each defect descriptor.

A defective sector number of FFFFFFFFh indicates that the entire track shall be considered defective.

8.2.2. LOCK UNLOCK CACHE Command

Table 8-11: LOCK UNLOCK CACHE Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (36h)							
1	Logical Unit Number			Reserved			Lock	RelAdr
2	(MSB)							
5	Logical Block Address (LSB)							
6	Reserved							
7	(MSB)							
8	Number of Blocks (LSB)							
9	Control							

The LOCK UNLOCK CACHE command (Table 8-11) requests that the target disallow or allow logical blocks within the specified range to be removed from the cache memory by the target's cache replacement algorithm. Locked logical blocks may be written to the medium when modified, but a copy of the modified logical block shall remain in the cache memory.

A lock bit of one indicates that any logical block in the specified range that is currently present in the cache memory shall be locked into cache memory. Only logical blocks that are already present in the cache memory are actually locked. A lock bit of zero indicates that all logical blocks in the specified range that are currently locked into the cache memory shall be unlocked, but not necessarily removed.

A relative address (RelAdr) bit of one indicates that the logical block address field is a two's complement displacement. This negative or positive displacement is to be added to the logical block address last accessed on the logical unit to form the logical block address for this command. This feature is only available when linking commands. The feature requires that a previous command in the linked group have accessed a block of data on the logical unit.

A RelAdr bit of zero indicates that the logical block address field specifies the first logical block of the range of logical blocks to be operated on by this command.

The number of blocks specifies the total number of contiguous logical blocks within the range. A number of blocks field of zero indicates that all remaining logical blocks on the logical unit shall be within the range.

Multiple locks may be in effect from more than one initiator. Locks from different initiators may overlap. An unlock of an overlapped area does not release the lock of another initiator.

8.2.3. PRE-FETCH Command

Table 8-12: PRE-FETCH Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (34h)							
1	Logical Unit Number			Reserved			Immed	RelAdr
2	(MSB)							
5	Logical Block Address							(LSB)
6	Reserved							
7	(MSB)							
8	Transfer Length							(LSB)
9	Control							

The PRE-FETCH command (Table 8-12) requests the target to transfer the specified logical blocks to the cache memory. No data shall be transferred to the initiator.

An immediate (Immed) bit of one indicates that status shall be returned as soon as the command descriptor block has been validated. An Immed bit of zero indicates that status shall be returned after the operation is complete.

See 8.2.2 for a definition of the RelAdr bit and the logical block address field.

The transfer length field specifies the number of contiguous logical blocks of data that shall be transferred to the target's cache memory. A transfer length of zero indicates that the contiguous logical blocks up to and including the last logical block of the logical unit shall be transferred to the target's cache memory. Any other value indicates the number of logical blocks that shall be transferred. The target may elect to not transfer logical blocks that already are contained in the cache memory.

If the Immed bit is zero and the specified logical blocks were successfully transferred to the cache memory the target shall return CONDITION MET status. If the link bit (see 6.2.7) is one the target shall return INTERMEDIATE-CONDITION MET status.

If Immed is one and the unlocked cache memory has sufficient capacity to accept all of the specified logical blocks the target shall return CONDITION MET status. If the link bit (see 6.2.7) is one the target shall return INTERMEDIATE-CONDITION MET status.

If Immed is one and the unlocked cache memory does not have sufficient capacity to accept all of the specified logical blocks the target shall return GOOD status. The target shall transfer to cache memory as many logical blocks as will fit. If the link bit (see 6.2.7) is one the target shall return INTERMEDIATE status.

8.2.4. PREVENT ALLOW MEDIUM REMOVAL Command

Table 8-13: PREVENT ALLOW MEDIUM REMOVAL Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (1Eh)							
1	Logical Unit Number			Reserved				
2	Reserved							
3	Reserved							
4	Reserved							Prevent
5	Control							

The PREVENT ALLOW MEDIUM REMOVAL command (Table 8-13) requests that the target enable or disable the removal of the medium in the logical unit. This mechanism is independent of device reservations and the target shall not allow medium removal if any initiator currently has medium removal prevented.

The prevention of medium removal shall begin when any initiator issues a PREVENT ALLOW MEDIUM REMOVAL command with a prevent bit of one (medium removal prevented). The prevention of medium removal for the logical unit shall terminate:

- (1) after all initiators that have medium removal prevented issue PREVENT ALLOW MEDIUM REMOVAL commands with a prevent bit of zero,
- (2) upon the receipt of a BUS DEVICE RESET message from any initiator, or
- (3) upon a hard RESET condition.

While a prevention of medium removal condition is in effect the target shall inhibit mechanisms that normally allow removal of the medium by an operator.

Targets that contain cache memory shall perform a synchronize cache operation for the entire medium prior to allowing medium removal. If the synchronization operation fails, then the prevention of medium removal shall remain in effect.

8.2.5. READ(6) Command

Table 8-14: READ(6) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (08h)							
1	Logical Unit Number			(MSB)				
2	Logical Block Address							
3								(LSB)
4	Transfer Length							
5	Control							

The READ(6) command (Table 8-14) requests that the target transfer data to the initiator. The most recent data value written in the addressed logical block shall be returned.

The cache control bits (see 8.2.6) are not provided for this command. Targets with cache memory may have values for the cache control bits which may affect the READ(6) command, however no default value is defined by this standard. If explicit control is required, the READ(10) command should be used.

The logical block address field specifies the logical block at which the read operation shall begin.

The transfer length field specifies the number of contiguous logical blocks of data to be transferred. A transfer length of zero indicates that 256 logical blocks shall be transferred. Any other value indicates the number of logical blocks that shall be transferred.

8.2.6. READ(10) Command

Table 8-15: READ(10) Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (28h)							
1	Logical Unit Number			DPO	FUA	Reserved		RelAdr
2	(MSB)							
5	Logical Block Address							
	(LSB)							
6	Reserved							
7	(MSB)							
8	Transfer Length							
	(LSB)							
9	Control							

The READ(10) command (Table 8-15) requests that the target transfer data to the initiator. The most recent data value written in the addressed logical block shall be returned.

A disable page out (DPO) bit of one indicates that the target shall assign the logical blocks accessed by this command the lowest priority for being fetched into or retained by the cache. A DPO bit of one overrides any retention priority specified in the cache page (see 8.3.3.1). A DPO bit of zero indicates the priority shall be determined by the retention priority fields in the cache page. All other aspects of the algorithm implementing the cache memory replacement strategy are not defined by this standard.

IMPLEMENTORS NOTE: The DPO bit is used to control replacement of logical blocks in the cache memory when the host has information on the future usage of the logical blocks. If the the DPO bit is set to one, the host knows the logical blocks accessed by the command are not likely to be accessed again in the near future and should not be put in the cache memory nor retained by the cache memory. If the DPO bit is zero, the host expects that logical blocks accessed by this command are likely to be accessed again in the near future.

A force unit access (FUA) bit of one indicates that the target shall access the media in performing the command prior to returning GOOD status. Read commands shall access the specified logical blocks from the media (i.e., the data is not directly retrieved from the cache). In the case where the cache contains a more recent version of a logical block than the media, the logical block shall first be written to the media. Write commands shall not return GOOD status until the logical blocks have actually been written on the media (i.e., the data is not write cached).

An FUA bit of zero indicates that the target may satisfy the command by accessing the cache memory. For read operations, any logical blocks that are contained in the cache memory may be transferred to the initiator directly from the cache memory. For write operations, logical blocks may be transferred directly to the cache memory. GOOD status may be returned to the initiator prior to writing the logical blocks to the medium. Any errors which occur after the GOOD status is returned is a deferred error and information regarding the error is not reported until a subsequent command.

See 8.2.2 for a definition of the RelAdr bit and the logical block address field.

The transfer length field specifies the number of contiguous logical blocks of data that shall be transferred. A transfer length of zero indicates that no logical blocks shall be transferred. This condition shall not be considered an error. Any other value indicates the number of logical blocks that shall be transferred.

8.2.7. READ CAPACITY Command

Table 8-16: READ CAPACITY Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (25h)							
1	Logical Unit Number				Reserved			RelAdr
2	(MSB)							
5	Logical Block Address							(LSB)
6	Reserved							
7	Reserved							
8	Reserved							PMI
9	Control							

The READ CAPACITY command (Table 8-16) provides a means for the initiator to request information regarding the capacity of the logical unit.

See 8.2.2 for a definition of the RelAdr bit and the logical block address field.

The logical block address shall be zero if the PMI bit is zero. If the PMI bit is zero and the logical block address is not zero, the target shall return a CHECK CONDITION status, the sense key shall be set to ILLEGAL REQUEST and the additional sense code set to ILLEGAL FIELD IN CDB.

A partial medium indicator (PMI) bit of zero indicates that the returned logical block address and the block length in bytes are that of the last logical block of the logical unit.

A PMI bit of one indicates that the returned logical block address and block length in bytes are that of the logical block address after which a substantial delay in data transfer will be encountered. This returned logical block address shall be greater than or equal to the logical block address specified by the RelAdr and logical block address fields in the command descriptor block.

IMPLEMENTORS NOTE: This function is intended to assist storage management software in determining whether there is sufficient space on the current track, cylinder, etc. to contain a frequently accessed data structure such as a file directory or file index without incurring an access delay.

The READ CAPACITY data (Table 8-17) shall be sent during the DATA IN phase of the command.

Table 8-17: READ CAPACITY Data

Byte	Description
0 (MSB)	Returned Logical Block Address (LSB)
3	
4 (MSB)	
7	Block Length In Bytes (LSB)

8.2.8. READ DEFECT DATA Command

Table 8-18: READ DEFECT DATA Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (37h)							
1	Logical Unit Number				Reserved			
2	Reserved			PList	GList	Defect List Format		
3	Reserved							
6								
7	(MSB)							(LSB)
8	Allocation Length							
9	Control							

The READ DEFECT DATA command (Table 8-18) requests that the target transfer the medium defect data to the initiator. If the target is unable to access any medium defect data it shall terminate the command with CHECK CONDITION status. The sense key shall be set to either MEDIUM ERROR if a medium error occurred or NO SENSE if the list does not exist and the additional sense code shall be set to DEFECT LIST NOT FOUND.

IMPLEMENTORS NOTE: Some targets may not be able to return medium defect data until after a FORMAT UNIT command has been successfully completed.

A primary defect list (PList) bit of one requests that the target return the primary list of defects. A PList bit of zero requests that the target not return the primary list of defects.

A grown defect list (GList) bit of one requests that the target return the grown defect list. A GList bit of zero requests that the target not return the grown defect list.

A PList bit of one and a GList bit of one requests that the target return the primary and the grown defect lists. The order in which the lists are returned is vendor-specific. Whether the lists are merged or not is vendor-specific.

A PList bit of zero and a GList bit of zero requests that the target return only the defect list header.

The defect list format field is used by the initiator to indicate the preferred format for the defect list. This field is intended for those targets capable of returning more than one format, as defined in the FORMAT UNIT command (see 8.2.1.2, defect list format). A target unable to return the requested format shall return the defect list in its default format (see the defect list format field in the defect list header below).

If the requested defect list format and the returned defect list format are not the same the target shall transfer the defect data and then terminate the command with CHECK CONDITION status. The sense key shall be set to RECOVERED ERROR and the additional sense code shall be set to DEFECT LIST NOT FOUND.

The READ DEFECT DATA defect list (Table 8-19) contains a four-byte header, followed by zero or more defect descriptors.

Table 8-19: READ DEFECT DATA Defect List

Bit	7	6	5	4	3	2	1	0
Byte								
0	Reserved							
1	Reserved			PList	GList	Defect List Format		
2	(MSB)							
3	Defect List Length							
	(LSB)							
Defect Descriptors								
0-n								

A PList bit of one indicates that the data returned contains the primary defect list. A PList bit of zero indicates that the data returned does not contain the primary defect list.

A GList bit of one indicates that the data returned contains the grown defect list. A GList bit of zero indicates that the data returned does not contain the grown defect list.

The defect list format field indicates the format of the defect descriptors returned by the target. This field is defined in the FORMAT UNIT command (see 8.2.1.2).

IMPLEMENTORS NOTE: The use of the block format is not recommended. There is no universal model that sensibly defines the meaning of the logical block address of a defect. In the usual case a defect that has been reassigned no longer has a logical block address.

Defect descriptors returned in the block format are vendor-specific. Defect descriptors returned in the physical sector format may or may not include defects in areas not accessible to the initiator. Defect descriptors returned in bytes-from-index format shall comprise a complete list of defects. A complete list of defects may include defects in areas not within the capacity returned in the READ CAPACITY command.

The defect list length field specifies the length in bytes of the defect descriptors that follow. The defect list length is equal to four or eight times the number of defect descriptors depending on the format of the returned descriptors (see Tables 8-8, 8-9 and 8-10 in the FORMAT UNIT command).

If the allocation length is insufficient to transfer all of the defect descriptors, the defect list length shall not be adjusted to reflect the truncation. The target shall not create CHECK CONDITION status. The initiator is responsible for comparing the defect list length and the allocation length to ensure that a partial list was not received.

IMPLEMENTORS NOTE: The initiator can determine the length of the defect list by sending the READ DEFECT DATA command with an allocation length of four. The target will return the defect list header which contains the length of the defect list.

The defect descriptors may or may not be sent in ascending order. The initiator can determine the exact number of defects by dividing the defect list length by the length of a single defect descriptor for the returned format.

8.2.9. READ LONG Command

Table 8-20: READ LONG Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (3Eh)							
1	Logical Unit Number			Reserved			CORRCT	RelAdr
2	(MSB)							
5	Logical Block Address							(LSB)
6	Reserved							
7	(MSB)							
8	Byte Transfer Length							(LSB)
9	Control							

The READ LONG command (Table 8-20) requests that the target transfer data to the initiator. The data passed during the READ LONG command is implementation specific, but shall include the data bytes and the ECC bytes recorded on the medium. The most recent data written in the addressed logical block shall be returned.

IMPLEMENTORS NOTE: Any other bytes that can be corrected by ECC should be included (e.g., data synchronization mark within the area covered by ECC). It is not important for the ECC bytes to be at the end of the data bytes, however they should be in the same order as they are on the media.

A corrected (CORRCT) bit of zero causes a logical block to be read without any correction made by the target. A CORRCT bit of one causes the data to be corrected by ECC before being transferred to the initiator.

See 8.2.2 for a definition of the RelAdr bit and the logical block address field.

The byte transfer length field should exactly specify the number of bytes of data that are available for transfer. If a non-zero byte transfer length does not exactly match the available data length, the target shall terminate the command with CHECK CONDITION status, the sense key shall be set to ILLEGAL REQUEST and an additional sense code set to INVALID FIELD IN CDB. The valid and ILI bits shall be set to one and the information field shall be set to the difference (residue) of the requested length minus the actual length in bytes. Negative values shall be indicated by two's complement notation.

A byte transfer length of zero indicates that no bytes shall be transferred and shall not be considered an error.

8.2.10. REASSIGN BLOCKS Command

Table 8-21: REASSIGN BLOCKS Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (07h)							
1	Logical Unit Number				Reserved			
2	Reserved							
3	Reserved							
4	Reserved							
5	Control							

set aside

The REASSIGN BLOCKS command (Table 8-21) requests the target to reassign the defective logical blocks to an area ~~on the logical unit reserved~~ for this purpose and to record the defective logical blocks to the grown defect list if such a list is supported. More than one physical or logical block may be relocated by each defect descriptor sent by the initiator. This command does not alter the contents or location of the Plist (see 8.2.1, FORMAT UNIT command).

The initiator transfers a defect list that contains the logical block addresses to be reassigned. The target shall reassign the physical medium used for each logical block address in the list. The data contained in the logical blocks specified in the defect list may be altered, but the data in all other logical blocks on the medium shall be preserved.

IMPLEMENTORS NOTE: The effect of specifying a logical block to be reassigned that previously has been reassigned is to reassign the block again. Over the life of the medium, a logical block can be assigned to multiple physical addresses until no more spare locations remain on the medium.

The REASSIGN BLOCKS defect list (Table 8-22) contains a four-byte header followed by one or more defect descriptors. The length of each defect descriptor is four bytes.

Table 8-22: REASSIGN BLOCKS Defect List

Byte	Defect List Header
0	Reserved
1	Reserved
2	(MSB)
3	Defect List Length (LSB)
	Defect Descriptor(s)
0	(MSB)
3	Defect Logical Block Address (LSB)

The defect list length field specifies the total length in bytes of the defect descriptors that follow. The defect list length is equal to four times the number of defect descriptors and does not include the defect list header length.

The defect descriptor specifies a four-byte defect logical block address that contains the defect. The defect descriptors shall be in ascending order.

If the logical unit has insufficient capacity to reassign all of the logical blocks specified in the defect descriptors, the command shall terminate with CHECK CONDITION status, the sense key shall be set to HARDWARE ERROR and the additional sense code set to NO DEFECT SPARE LOCATION AVAILABLE.

If the logical unit is unable to successfully complete a REASSIGN BLOCKS command, the command shall terminate with CHECK CONDITION status with the appropriate sense information. The logical block address of the first defect descriptor not reassigned shall be returned in the command-specific information field of the sense data. If information about the first defect descriptor not reassigned is not available, or if all the defects have been reassigned, this field shall be set to FFFFFFFFh.

If the REASSIGN BLOCKS command failed due to an unexpected unrecoverable read error that would cause the loss of data in a block not specified in the defect list, the logical block address of the unrecoverable block shall be returned in the information field of the sense data and the valid bit shall be set to one.

IMPLEMENTORS NOTE: If the REASSIGN BLOCKS command returns CHECK CONDITION status and the sense data command-specific information field contains a valid logical block address, the initiator should remove all defect descriptors from the defect list prior to the one returned in the command-specific information field. If the sense key is MEDIUM ERROR and the valid bit is one (the information field contains the valid block address) the initiator should insert that new defective logical block address into the defect list and reissue the REASSIGN BLOCKS command with the new defect list. Otherwise, the initiator should perform any corrective action indicated by the sense data and then reissue the REASSIGN BLOCKS command with the new defect list.

8.2.11. RELEASE Command

Table 8-23: RELEASE Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (17h)							
1	Logical Unit Number			3rdPty	Third Party Device ID			Extent
2	Reservation Identification							
3	Reserved							
4	Reserved							
5	Control							

The RESERVE and RELEASE commands provide the basic mechanism for contention resolution in multiple-initiator systems. The RELEASE command (Table 8-23) is used to release a previously reserved logical unit, or, if the extent release option is implemented, to release previously reserved extents within a logical unit. It is not an error for an initiator to attempt to release a reservation that is not currently valid. In this case, the target returns GOOD status without altering any other reservation.

IMPLEMENTORS NOTE: The reservation queuing option in X3.131-1986 has been removed from SCSI-2.

8.2.11.1. Logical Unit Release (Mandatory)

If the extent bit is zero, this command shall cause the target to terminate all logical unit and extent reservations that are active from the initiator to the specified logical unit. The reservation ID field in the command descriptor block is ignored by the target.

8.2.11.2. Extent Release (Optional)

If the extent bit is one and the extent release option is not implemented, then the RELEASE command shall be terminated with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST. This option shall be implemented if the extent reservation option (see 8.2.12.2) is implemented.

If the extent bit is one and the extent release option is implemented, this command shall cause any reservation from the requesting initiator with a matching reservation identification to be terminated. Other reservations from the requesting initiator shall remain in effect.

8.2.11.3. Third-Party Release (Mandatory)

Third-party release allows an initiator to release a logical unit or extents within a logical unit that were previously reserved using third-party reservation (see 8.2.12.3). Third-party release shall be implemented and is intended for use in multiple-initiator systems that use the COPY command.

If the third-party (3rdPty) bit is zero, then a third-party release is not requested. If the 3rdPty bit is one then the target shall release the specified logical unit or extents, but only if the reservation was made using a third-party reservation by the initiator that is requesting the release for the same SCSI device as specified in the third-party device ID field.

If the 3rdPty bit is one the target shall not modify the mode parameters for commands received from the third-party device even if the target implements the transfer of mode parameters with a third-party RESERVE command.

IMPLEMENTORS NOTE: If a target implements independent storage of mode parameters for each initiator, a third-party RESERVE command copies the current mode parameters for the initiator that sent the RESERVE command to the current mode parameters for the initiator specified as the third-party device (usually a copy master device). A unit attention condition notifies the third-party of the changed mode parameters due to the reservation. A successful third-party RELEASE command does not return the third-party devices' current mode parameters back to their previous values. The third-party device can issue MODE SENSE and MODE SELECT commands to query and modify the mode parameters.

8.2.12. RESERVE Command

Table 8-24: RESERVE Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (16h)							
1	Logical Unit Number			3rdPty	Third Party Device ID			Extent
2	Reservation Identification							
3	(MSB)							
4	Extent List Length						(LSB)	
5	Control							

The RESERVE and RELEASE commands provide the basic mechanism for contention resolution in multiple-initiator systems. The RESERVE command (Table 8-24) is used to reserve a logical unit or, if the extent reservation option is implemented, extents within a logical unit. The third-party reservation allows logical units or extents to be reserved for another specified SCSI device.

IMPLEMENTORS NOTE: The reservation queuing option in X3.131-1986 has been removed from SCSI-2.

8.2.12.1. Logical Unit Reservation (Mandatory).

If the extent bit is zero, this command shall request that the entire logical unit be reserved for the exclusive use of the initiator until the reservation is superseded by another valid RESERVE command from the initiator that made the reservation or until released by a RELEASE command from the same initiator that made the reservation, by a BUS DEVICE RESET message from any initiator, by a hard RESET condition, or by a power on cycle. A logical unit reservation shall not be granted if the logical unit or any extent is reserved by another initiator. It shall be permissible for an initiator to reserve a logical unit that is currently reserved by that initiator. If the extent bit is zero, the reservation identification and the extent list length shall be ignored.

If the logical unit, or any extent within the logical unit is reserved for another initiator, the target shall return RESERVATION CONFLICT status.

If, after honoring the reservation, any other initiator attempts to perform any command on the reserved logical unit other than an INQUIRY command, a REQUEST SENSE command, an ALLOW MEDIUM REMOVAL, or a RELEASE command, which shall be ignored, then the command shall be rejected with RESERVATION CONFLICT status.

8.2.12.2. Extent Reservation (Optional)

The reservation identification field provides a means for an initiator to identify each extent reservation. This allows an initiator in a multiple tasking environment, to have multiple reservations outstanding. The reservation identification is used in the RELEASE command to specify which reservation is to be released. It is also used in superseding RESERVE commands to specify which reservation is to be superseded.

If the extent reservation option is implemented, then the extent release option (see 8.2.11.2) shall also be implemented. These options permit multiple extents within the logical unit to be reserved, each with a separate reservation type.

If the extent bit is one, and the extent reservation option is implemented, then the target shall process the reservation request as follows:

(1) The extent list shall be checked for the number of extents in the reservation request. If the extent list length is zero, no current reservations shall be changed, no new reservations shall be created, and this condition shall not be treated as an error. If the extent list contains more extents than are supported on the logical unit, the command shall be terminated with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST. If the extent list contains more extents than are currently available on the logical unit, then the target shall return a RESERVATION CONFLICT status.

(2) The extent list shall be checked for valid extent logical block addresses. If any logical block address is invalid for this logical unit, the command shall be terminated with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST. The extent list shall be checked for invalid extent overlaps (as defined by reservation type) with other extent descriptors in the extent list and if invalid overlaps are found, the command shall be terminated with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST.

(3) If the requested reservation does not conflict with an existing reservation, the extents specified shall be reserved until superseded by another valid RESERVE command from the initiator that made the reservation or until released by a RELEASE command from the same initiator, by a BUS DEVICE RESET message from any initiator, or by a hard RESET condition. If either of the last two conditions occur, the next command from each initiator shall be terminated with CHECK CONDITION status and the sense key shall be set to UNIT ATTENTION.

(4) If the reservation request conflicts with an existing reservation, then the target shall return a RESERVATION CONFLICT status.

If the extent bit is one, and the extent reservation option is not implemented, then the RESERVE command shall be rejected with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST.

Table 8-25: Data Format of Extent Descriptors

Bit	7	6	5	4	3	2	1	0
Byte								
0	Reserved					RelAdr	Reservation Type	
1	(MSB)					Number of Blocks		
3						(LSB)		
4	(MSB)					Logical Block Address		
7						(LSB)		

The size of the extent list shall be defined by the extent list length field. The extent list shall consist of zero or more descriptors as shown in Table 8-25. Each extent descriptor defines an extent beginning at the specified logical block address for the specified number of blocks. If the number of blocks is zero, the extent shall begin at the specified logical block address and continue through the last logical block address on the logical unit.

The reservation type field shall determine the type of reservation to be effected for each extent. Four types of reservations are possible as follows:

DB(1)	DB(0)	Reservation Type
1	0	Read Exclusive
0	1	Write Exclusive
1	1	Exclusive Access
0	0	Read Shared

Read Exclusive. While this reservation is active, no other initiator shall be permitted read operations to the indicated extent. This reservation shall not inhibit write operations from any initiator or conflict with a write exclusive reservation; however, read exclusive, exclusive access, and read shared reservations that overlap this extent shall conflict with this reservation.

Write Exclusive. While this reservation is active, no other initiator shall be permitted write operations to the indicated extent. This reservation shall not inhibit read operations from any initiator or conflict with a read exclusive reservation from any initiator. This reservation shall conflict with write exclusive, exclusive access, and read shared reservations that overlap this extent.

Exclusive Access. While this reservation is active, no other initiator shall be permitted any access to the indicated extent. All reservation types that overlap this extent shall conflict with this reservation.

Read Shared. While this reservation is active, no write operations shall be permitted by any initiator to the indicated extent. This reservation shall not inhibit read operations from any initiator or conflict with a read shared reservation. Read exclusive, write exclusive, and exclusive access reservations that overlap with this extent shall conflict with this reservation.

If the relative address bit is one, the logical block address in the extent descriptor shall be treated as a two's complement displacement. This displacement shall be added to the logical block address last accessed on the logical unit to form the logical block address for this extent. This feature is only available when linking commands and requires that a previous command in the linked group has accessed a logical block on the logical unit; if not, the RESERVE command shall be terminated with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST.

If an initiator attempts a command to a logical block that has been reserved and that access is prohibited by the reservation, the command shall not be performed and the command shall be terminated with a RESERVATION CONFLICT status. If a reservation conflict precludes any part of the command, none of the command shall be performed. COPY commands shall be terminated with CHECK CONDITION status and the sense key shall be set to DATA PROTECT if any part of the copy operation is prohibited by an active reservation. If any extent in a logical unit is reserved in any way, by any initiator, a FORMAT UNIT command shall be rejected with a RESERVATION CONFLICT status.

8.2.12.3. Third Party Reservation (Mandatory).

The third-party reservation for the RESERVE command allows an initiator to reserve a logical unit or extents within a logical unit for another SCSI device. This is intended for use in multiple-initiator systems that use the COPY command. Third-party reservation is required.

If the third-party (3rdPty) bit is zero, then a third-party reservation is not requested. If the 3rdPty bit is one then the target shall reserve the specified logical unit or extents for the SCSI device specified in the third-party device ID field. The target shall preserve the reservation until it is superseded by another valid RESERVE command from the initiator that made the reservation or until it is released by the same initiator, by a BUS DEVICE reset message from any initiator, or a hard reset condition. The target shall ignore any attempt to release the reservation made by any other initiator.

If independent sets of parameters are implemented, a third party reservation shall cause the target to transfer the set of parameters in effect for the initiator of the RESERVE command to the parameters used for commands from the third party device. Any subsequent command issued by the third-party device is executed according to the mode parameters in effect for the initiator that sent the RESERVE command.

IMPLEMENTORS NOTE: This transfer of the mode parameters is applicable to target devices which store mode information independently for different initiators. This mechanism allows an initiator to set the mode parameters of a target for the use of a copy master (i.e., the third-party device). The third-party copy master may subsequently issue a MODE SELECT command to modify the mode parameters.

8.2.12.4. Superseding Reservations (Mandatory)

An initiator that holds a current reservation (unit or extent) may modify that reservation by issuing another RESERVE command (unit or extent) to the same logical unit. The superseding RESERVE command shall release the previous reservation state (unit or extent) when the new reservation request is granted. If the superseding reservation is for an extent reservation and the current reservation is also an extent reservation, the current extent reservation identification value is used for the superseding reservation. The current reservation shall not be modified if the superseding reservation request cannot be granted. If the superseding reservation cannot be granted because of conflicts with a previous reservation (other than the reservation being superseded), then the target shall return RESERVATION CONFLICT status.

IMPLEMENTORS NOTE: Superseding reservations allow the SCSI device ID to be changed on a reservation using the third-party reservation option. This capability is necessary for certain situations when using COMPARE, COPY, and COPY AND VERIFY commands.

8.2.13. REZERO UNIT Command**Table 8-26: REZERO UNIT Command**

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (01h)							
1	Logical Unit Number				Reserved			
2	Reserved							
3	Reserved							
4	Reserved							
5	Control							

The REZERO UNIT command (Table 8-26) requests that the target set the logical unit to a specific state. See vendor specifications for details.

8.2.14. SEARCH DATA Commands**Table 8-27: SEARCH DATA Commands**

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (31h 30h 32h)							
1	Logical Unit Number			Invert	Reserved		SpnDat	RelAdr
2	(MSB)							
5	Logical Block Address							(LSB)
6	Reserved							
7	(MSB)							
8	Number of Blocks to Search							(LSB)
9	Control							

The SEARCH DATA commands (Table 8-27) search one or more logical blocks for equality or inequality to a data pattern. The concept of records within a logical block is used to allow multiple records within a logical block to be searched.

The invert bit determines whether the search condition is to be inverted. See 8.2.14.1 through 8.2.14.3 for a description of the search conditions for the individual SEARCH DATA commands.

A spanned data (SpnDat) bit of zero indicates that each record shall be wholly contained within a single block. Any space at the end of a block that is smaller than the record length is ignored by the SEARCH DATA commands. A SpnDat bit of one indicates that records span block boundaries (i.e., record may start in one block and end in the next or a subsequent block).

The number of blocks to search field specifies the maximum number of contiguous logical blocks to be searched. A value of zero indicates that no logical blocks shall be searched. This condition shall not be considered an error. Any other value indicates the maximum number of logical blocks that shall be searched.

A link bit (see 6.2.7) of zero indicates a non-linked command and if the search is satisfied, the command shall be terminated with a CONDITION MET status. A REQUEST SENSE command can then be issued to determine the logical block address and record offset of the matching record. If the search is not satisfied and no error occurs, the command shall be terminated with GOOD status.

A link bit (see 6.2.7) of one indicates a command is linked to the SEARCH DATA command and if the search is satisfied, INTERMEDIATE-CONDITION MET status is returned and the next command is executed. If the RelAdr bit in the next command is one, the logical block address of the next command is used as a displacement from the logical block address at which the search was satisfied. If a linked search is not satisfied, the command is terminated with CHECK CONDITION status. A REQUEST SENSE command may then be issued.

A REQUEST SENSE command following a satisfied SEARCH DATA command shall:

(1) Return a sense key of EQUAL if the search was satisfied by an exact match. If the search was satisfied by an inequality then a sense key of NO SENSE shall be returned.

(2) Return the valid bit set to one.

(3) Return the logical block address of the logical block containing the first matching record in the information field.

(4) Return the record offset of the matching record in the command-specific information field.

A REQUEST SENSE command following a SEARCH DATA command that is not satisfied shall:

(1) Return a sense key of NO SENSE, if no errors occurred during the command execution.

(2) Return the valid bit set to zero.

The SEARCH DATA parameter list (Table 8-28) contains a fourteen-byte header, followed by one or more search argument descriptors.

Table 8-28: SEARCH DATA Parameter List

Byte	Parameter List Header		
0	(MSB)	Logical Record Length	(MSB)
3	(LSB)		(LSB)
4	(MSB)	First Record Offset	(MSB)
7	(LSB)		(LSB)
8	(MSB)	Number of Records	(MSB)
11	(LSB)		(LSB)
12	(MSB)	Search Argument Length	(MSB)
13	(LSB)		(LSB)
Search Argument Descriptor(s)			
0	(MSB)	Displacement	(MSB)
3	(LSB)		(LSB)
4	(MSB)	Pattern Length	(MSB)
5	(LSB)		(LSB)
6 - n	Pattern		

The logical record length field specifies the record length in bytes.

The first record offset field specifies the number of bytes that shall be ignored in the first logical block before the search begins. If the value of the first record offset field shall is larger than the logical block length the target shall terminate the command with a CHECK CONDITION status, set the sense key to ILLEGAL REQUEST and set the additional sense code to INVALID FIELD IN PARAMETERS LIST. Subsequent logical blocks shall be searched beginning with the first byte in the logical block. This permits one or more records to be skipped initially.

The number of records field specifies the maximum number of records that shall be searched by this command. A search shall terminate when the search pattern is found or when the number of records is exhausted or when the number of blocks to search is exhausted.

The search argument length field specifies the length in bytes of all the search argument descriptors that follow.

IMPLEMENTORS NOTE: Since the pattern length can vary, there is no fixed multiple of the search argument descriptor to determine the search argument length.

The search argument descriptors specify one or more search conditions to execute within a single record in order to satisfy the search. Each search argument descriptor is made up of a displacement field, a pattern length field, and a pattern field.

The displacement field specifies the displacement in bytes of the first byte of the data to be compared from the start of the logical record.

The pattern length field specifies the length in bytes of the pattern that follows.

The pattern field specifies the data to compare to the logical record.

8.2.14.1. SEARCH DATA EQUAL Command

The SEARCH DATA EQUAL command (Table 8-27, operation code 31h) shall be satisfied by the first logical record searched that contains data that satisfies all of the search argument descriptor(s). If the invert bit in the command descriptor block is zero, the search argument descriptor(s) shall be satisfied by data in the logical record being equal to the data in the pattern. If the invert bit is one, the search argument descriptor(s) shall be satisfied by data in the logical record being not equal to the data in the pattern. (See 8.2.14.)

8.2.14.2. SEARCH DATA HIGH Command

The SEARCH DATA HIGH command (Table 8-27, operation code 30h) shall be satisfied by the first logical record searched that contains data that satisfies all of the search argument descriptor(s). If the invert bit in the command descriptor block is zero, the search argument descriptor(s) shall be satisfied by data in the logical record being greater than the data in the pattern. If the invert bit is one, the search argument descriptor(s) shall be satisfied by data in the logical record being less than or equal to the data in the pattern. (See 8.2.14.)

8.2.14.3. SEARCH DATA LOW Command

The SEARCH DATA LOW command (Table 8-27, operation code 32h) shall be satisfied by the first logical record searched that contains data that satisfies all of the search argument descriptor(s). If the invert bit in the command descriptor block is zero, the search argument descriptor(s) shall be satisfied by data in the logical record being less than the data in the pattern. If the invert bit is one, the search argument descriptor(s) shall be satisfied by data in the logical record being greater than or equal to the data in the pattern. (See 8.2.14.)

8.2.15. SEEK(6) and SEEK(10) Commands

Table 8-29: SEEK(6) Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (0Bh)							
1	Logical Unit Number			(MSB)				
2	Logical Block Address							
3	(LSB)							
4	Reserved							
5	Control							

Table 8-30: SEEK(10) Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (2Bh)							
1	Logical Unit Number			Reserved				
2	(MSB)							
5	Logical Block Address							(LSB)
6	Reserved							
7	Reserved							
8	Reserved							
9	Control							

The SEEK(6) (Table 8-29) and SEEK(10) (Table 8-30) commands request that the logical unit seek to the specified logical block address.

8.2.16. SET LIMITS Command

Table 8-31: SET LIMITS Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (33h)							
1	Logical Unit Number			Reserved			RdInh	WrInh
2	(MSB)							
5	Logical Block Address							(LSB)
6	Reserved							
7	(MSB)							
8	Number of Blocks							(LSB)
9	Control							

The SET LIMITS command (Table 8-31) defines the range within which subsequent linked commands may operate. A second SET LIMITS command may not be linked to a chain of commands in which a SET LIMITS command has already been issued.

The read inhibit (RdInh) bit of one indicates that read operations within the range shall be inhibited.

A write inhibit (WrInh) bit of one indicates that write operations within the range shall be inhibited.

The logical block address field specifies the starting address for the range.

The number of blocks field specifies the number of logical blocks within the range. A number of blocks of zero indicates that the range shall extend to the last logical block on the logical unit.

Any attempt to access outside of the restricted range or any attempt to perform an inhibited operation within the restricted range shall not be performed. The command shall be terminated with CHECK CONDITION status and the sense key shall be set to DATA PROTECT. A second SET LIMITS command within a linked list of commands shall be rejected with CHECK CONDITION status and the sense key shall be set to DATA PROTECT.

8.2.17. START STOP UNIT Command

Table 8-32: START STOP UNIT Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (1Bh)							
1	Logical Unit Number			Reserved				Immed
2	Reserved							
3	Reserved							
4	Reserved				LoEj		Start	
5	Control							

The START STOP UNIT command (Table 8-32) requests that the target enable or disable the logical unit for media access operations.

An immediate (Immed) bit of one indicates that status shall be returned as soon as the command descriptor block has been validated. An Immed bit of zero indicates that status shall be returned after the operation is completed.

A load eject (LoEj) bit of zero requests that no action be taken regarding loading or ejecting the medium. A LoEj bit of one requests that the medium shall be unloaded if the start bit is zero. A LoEj bit of one requests that the medium is to be loaded if the start bit is one.

A start bit of one requests the logical unit be made ready for use. A start bit of zero requests that the logical unit be stopped (media cannot be accessed by the initiator).

Targets that contain cache memory shall implicitly perform a SYNCHRONIZE CACHE command for the entire medium prior to executing the STOP UNIT command.

8.2.18. SYNCHRONIZE CACHE Command

Table 8-33: SYNCHRONIZE CACHE Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (35h)							
1	Logical Unit Number			Reserved			Immed	RelAdr
2	(MSB)							
5	Logical Block Address							(LSB)
6	Reserved							
7	(MSB)							
8	Number of Blocks							(LSB)
9	Control							

The SYNCHRONIZE CACHE command (Table 8-33) ensures that logical blocks in the cache memory, within the specified range, have their most recent data value recorded on the physical medium. If a more recent data value for a logical block within the specified range exists in the cache memory than on the physical medium, then the logical block from the cache memory shall be written to the physical medium. Logical blocks are not necessarily removed from the cache memory as a result of the synchronize cache operation.

An immediate (Immed) bit of one indicates that the target shall return status as soon as the command descriptor block has been validated. An Immed bit of zero indicates that the status shall not be returned until the operation has been completed. If the Immed bit is one and the target does not support it and the command shall terminate with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST and the additional sense code shall be set to INVALID FIELD IN CDB.

See 8.2.2 for a definition of the RelAdr bit and the logical block address field.

The number of blocks field specifies the total number of contiguous logical blocks within the range. A number of blocks of zero indicates that all remaining logical blocks on the logical unit shall be within the range.

A logical block within the specified range that is not in cache memory is not considered an error.

8.2.19. VERIFY Command

Table 8-34: VERIFY Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (2Fh)							
1	Logical Unit Number			DPO	Reserved	Reserved	BytChk	RelAdr
2	(MSB)							
5	Logical Block Address							(LSB)
6	Reserved							
7	(MSB)							
8	Verification Length							(LSB)
9	Control							

The VERIFY command (Table 8-34) requests that the target verify the data written on the medium.

If the MODE SELECT command is implemented, and the verify error recovery parameters page is also implemented, then the current settings in that page specifies the verification error criteria. If the verify error recovery parameters page is not implemented, then the verification criteria is vendor-specific.

A byte check (BytChk) bit of zero causes a medium verification to be performed with no data comparison. A BytChk bit of one causes a byte-by-byte compare of data written on the medium and the data transferred from the initiator. If the compare is unsuccessful for any reason, the target shall return CHECK CONDITION status with the sense key set to MISCOMPARE.

See 8.2.2 for a definition of the RelAdr bit and the logical block address field.

The verification length field specifies the number of contiguous logical blocks of data that shall be verified. A transfer length of zero indicates that no logical blocks shall be verified. This condition shall not be considered as an error. Any other value indicates the number of logical blocks that shall be verified.

See 8.2.6 for a description of the cache control bit (DPO).

Table 8-35: WRITE(6) Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (0Ah)							
1	Logical Unit Number			(MSB)				
2	Logical Block Address							
3								(LSB)
4	Transfer Length							
5	Control							

The WRITE(6) command (Table 8-35) requests that the target write the data transferred by the initiator to the medium.

The cache control bits (see 8.2.6) are not provided for this command. Targets with cache memory may have values for the cache control bits which may affect the WRITE(6) command, however no default value is defined by this standard. If explicit control is required, the WRITE(10) command should be used.

The logical block address field specifies the logical block at which the write operation shall begin.

The transfer length field specifies the number of contiguous logical blocks of data to be transferred. A transfer length of zero indicates that 256 logical blocks shall be transferred. Any other value indicates the number of logical blocks that shall be transferred.

Table 8-36: WRITE(10) Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (2Ah)							
1	Logical Unit Number				DPO	FUA	Reserved	RelAdr
2	(MSB)							
5	Logical Block Address							
6	(LSB)							
6	Reserved							
7	(MSB)							
8	Transfer Length							
9	(LSB)							
9	Control							

The WRITE(10) command (Table 8-36) requests that the target write the data transferred by the initiator to the medium.

See READ(10) command (8.2.6) for a definition of the cache control bits (DPO and FUA).

See 8.2.2 for a definition of the RelAdr bit and the logical block address field.

The transfer length field specifies the number of contiguous logical blocks of data that shall be transferred. A transfer length of zero indicates that no logical blocks shall be transferred. This condition shall not be considered an error and no data shall be written. Any other value indicates the number of logical blocks that shall be transferred.

8.2.22. WRITE AND VERIFY Command

Table 8-37: WRITE AND VERIFY Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (2Eh)							
1	Logical Unit Number			DPO	Reserved	Reserved	BytChk	RelAdr
2	(MSB)							
5	Logical Block Address							(LSB)
6	Reserved							
7	(MSB)							
8	Transfer Length							(LSB)
9	Control							

The WRITE AND VERIFY command (Table 8-37) requests that the target write the data transferred from the initiator to the medium and then verify that the data is correctly written. The data is only transferred once from the initiator to the target.

If the MODE SELECT command is implemented, and the verify error recovery page is also implemented (see 8.3.3.8), then the current settings in that page along with the AWRE bit from the read-write error recovery page specify the verification error criteria. If these pages are not implemented, then the verification criteria is vendor-specific.

A byte check (BytChk) bit of zero requests a medium verification to be performed with no data comparison. A BytChk bit of one requests a byte-by-byte compare of data written on the medium and the data transferred from the initiator. If the compare is unsuccessful for any reason, the target shall return CHECK CONDITION status with the sense key set to MISCOMPARE.

See 8.2.2 for a definition of the RelAdr bit and the logical block address field.

See WRITE(10) command (8.2.21) for a definition of the transfer length field.

See 8.2.6 for a description of the cache control bit (DPO).

IMPLEMENTORS NOTE: The WRITE AND VERIFY command specifically states that the data are not to be transferred twice (i.e., once for the write pass, and once for the verify pass) when performing a byte compare. If there is a need for two transfers to occur (e.g., to ensure the integrity of the path to the media), then the initiator should issue a WRITE command with a LINK bit of one followed by a VERIFY command with a BytCmp bit of one, transferring the same data on each command.

8.2.23. WRITE LONG Command

Table 8-38: WRITE LONG Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (3Fh)							
1	Logical Unit Number			Reserved				RelAdr
2	(MSB)							
5	Logical Block Address							(LSB)
6	Reserved							
7	(MSB)							
8	Byte Transfer Length							(LSB)
9	Control							

The WRITE LONG command (Table 8-38) requests that the target write the data transferred by the initiator to the medium. The data passed during the WRITE LONG command is implementation specific, but shall include the data bytes and the ECC bytes.

IMPLEMENTORS NOTE: Any other bytes that can be corrected by ECC should be included (e.g., a data synchronization mark within the area covered by ECC). The READ LONG command is usually issued before issuing a WRITE LONG command. The WRITE LONG data should be the same length and in the same order as the data returned by the READ LONG command.

See 8.2.2 for a definition of the RelAdr bit and the logical block address field.

The byte transfer length field should specify the number of bytes of data that the target would return for the READ LONG command. If a non-zero byte transfer length does not exactly match the data length the target would return for the READ LONG command, then the target shall terminate the command with CHECK CONDITION status and a sense key of ILLEGAL REQUEST and an additional sense code of INVALID FIELD IN CDB. The ILI and valid bits shall be set to one

and the information field shall be set to the difference (residue) of the requested length minus the actual length in bytes. Negative values shall be indicated by two's complement notation. A transfer length of zero indicates that no bytes shall be transferred and shall not be considered an error.

8.2.24. WRITE SAME Command

Table 8-39: WRITE SAME Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (41h)							
1	Logical Unit Number			Reserved		PBdata	LBdata	RelAdr
2	(MSB)							
5	Logical Block Address							(LSB)
6	Reserved							
7	(MSB)							
8	Number of Blocks							(LSB)
9	Control							

The WRITE SAME command (Table 8-39) requests that the target write the single block of data transferred by the initiator to the medium multiple times.

IMPLEMENTORS NOTE: This command is useful if large areas of the medium need to be written, prepared for certification, or otherwise initialized without the initiator having to transfer all the data.

A logical block data (LBdata) bit of one requests that the target replace the first four bytes of the data to be written to the current logical block with the logical block address of the block currently being written.

A physical block data (PBdata) bit of one requests that the target replace the first eight bytes of the data to be written to the current physical sector with the physical address of the sector currently being written using the physical sector format (see Table 8-10).

If PBdata and LBdata are one the command shall be terminated with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST.

See 8.2.2 for a definition of the RelAdr bit and the logical block address field.

The number of blocks field specifies the number of contiguous logical blocks to be written. A number of blocks field of zero requests that all the remaining logical blocks on the medium be written.

8.3. Parameters for Direct-Access Devices

8.3.1. Diagnostic Parameters

This section defines the descriptors and pages for diagnostic parameters used with direct-access devices.

The diagnostic page codes for direct-access devices are defined in Table 8-40.

Table 8-40: Diagnostic Page Codes

Page Code	Description	Section
00h	List of supported pages	7.3.1.1
01h - 3Fh	Reserved (for all device type pages)	
40h	Translate Address Page	8.3.1.1
41h - 7Fh	Reserved	
80h - FFh	Vendor-specific pages	

8.3.1.1. Translate Address Page - SEND DIAGNOSTIC

The translate address page allows the initiator to translate a logical block address, physical sector address or physical bytes from index address into any one of the other formats. The address to be translated is passed to the target with the SEND DIAGNOSTIC command and the results are returned to the initiator by the RECEIVE DIAGNOSTIC RESULTS command. The format of the translate address page - SEND DIAGNOSTIC is shown in Table 8-41. The translated address is returned in the translate address page - RECEIVE DIAGNOSTIC RESULTS (see Table 8-42).

Table 8-41: Translate Address Page - SEND DIAGNOSTIC

Bit	7	6	5	4	3	2	1	0
Byte								
0	Page Code (40h)							
1	Reserved							
2	(MSB)	Page Length (000Ah)						(LSB)
3								
4	Reserved				Supplied Format			
5	Reserved				Translate Format			
6								
13	Address to Translate							

The supplied format field specifies the format of address to translate field. Valid values for this field are defined in the FORMAT UNIT command (see Table 8-5). If the target does not support the requested format it shall terminate the SEND DIAGNOSTIC command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST and an additional sense code shall be set to INVALID FIELD IN PARAMETER LIST.

The translate format field specifies which format the initiator would like the address to be translated to. Valid values for this field are defined in the FORMAT UNIT command (see Table 8-5). If the target does not support the requested format it shall terminate the command with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST and an additional sense code shall be set to INVALID FIELD IN PARAMETER LIST.

The address to translate field contains a single address the initiator is requesting the target to translate. The format of this field depends on the value in the supplied format field. The formats are described in Tables 8-8, 8-9 and 8-10. If the logical block format is specified the block address shall be in the first four bytes of the field with the remaining bytes set to zero.

8.3.1.2. Translate Address Page - RECEIVE DIAGNOSTIC

The translate address page allows the initiator to translate a logical block address, physical sector address, or physical bytes from index address into any one of the other formats. The address to be translated is passed to the target with the SEND DIAGNOSTIC command and the results are returned to the initiator by the RECEIVE DIAGNOSTIC RESULTS command. The translated address is returned in the translate address page - RECEIVE DIAGNOSTIC (Table 8-42).

Table 8-42: Translate Address Page - RECEIVE DIAGNOSTIC

Bit Byte	7	6	5	4	3	2	1	0
0	Page Code (40h)							
1	Reserved							
2	(MSB)							
3	Page Length (LSB)							
4	Reserved				Supplied Format			
5	RAREA	ALTSEC	ALTTRK	Reserved		Translated Format		
6								
13	Translated Address 1							
14								
21	Translated Address 2 (if required)							
n								
nn	Translated Address n (if required)							

The translate address page contains a four byte page header which specifies the page code and length followed by two bytes which describe the translated address followed by zero or more translated address(s).

The page length field contains the number of parameter bytes which follow.

The supplied format field contains the value from the SEND DIAGNOSTIC command supplied format field (see 8.3.1.1).

A reserved area (RAREA) bit of one indicates that all or part of the translated address falls within a reserved area of the medium (e.g., speed tolerance gap, alternate sector, vendor reserved area, etc.). If the entire translated address falls within a reserved area the target may not return a translated address. An RAREA bit of zero indicates that no part of the translated address falls within a reserved area of the medium.

An alternate sector (ALTSEC) bit of one indicates that the translated address is physically located in an alternate sector of the medium. If the target cannot determine if all or part of the translated address is located in an alternate sector it shall set this bit to zero. An ALTSEC bit of zero indicates that no part of the translated address is located in an alternate sector of the medium or that the target is unable to determine this information.

An alternate track (ALTTRK) bit of one indicates that part or all of the translated address is located on an alternate track of the medium or the target cannot determine if all or part of the translated address is located on an alternate track. An ALTTRK bit of zero indicates that no part of the translated address is located on an alternate track of the medium.

The translated format field contains the value from the SEND DIAGNOSTIC command translate format field (see 8.3.1.1).

The translated address field contains the address(s) the target translated from the address supplied by the initiator in the SEND DIAGNOSTIC command. This field shall be in the format specified in the translate format field. The different formats are described in Tables 8-8, 8-9 and 8-10. If the logical block format is specified the block address shall be in the first four bytes of the field and the remaining bytes shall be set to zero.

If the returned data is in the logical block or physical sector format and the address to be translated covers more than one address after it has been translated (e.g., accounting for speed tolerance or multiple physical sectors within a single logical block or multiple logical blocks within a single physical sector) the target shall return all possible addresses which are contained in the area specified by the address to be translated.

If the returned data is in bytes from index format the target shall return a pair of translated values for each of the possible addresses which are contained in the area specified by the address to translate field. Of the pair of translated values returned, the first indicates the starting location and the second the ending location of the area.

8.3.2. Log Parameters

This section defines the descriptors and pages for log parameters used with direct-access devices.

The log page codes for direct-access devices are defined in Table 8-43.

Table 8-43: Log Page Codes

Page Code	Description	Section
01h	Buffer Over-Run/Under-Run Page	7.3.2.1
03h	Error Counter Page (Read) Page	7.3.2.2
05h	Error Counter Page (Verify) Page	7.3.2.2
02h	Error Counter Page (Write) Page	7.3.2.2
07h	Last n Error Events Page	7.3.2.3
06h	Non-Medium Error Page	7.3.2.4
00h	Supported Log Pages	7.3.2.5
04h	Reserved	
08h - 2Fh	Reserved	
3Fh	Reserved	
30h - 3Eh	Vendor-specific pages	

8.3.3. Mode Parameters

This section defines the descriptors and pages for mode parameters used with direct-access devices.

The mode parameter list, including the mode parameter header and mode block descriptor are described in 7.3.3.

The medium-type code field is contained in the mode parameter header (see Table 7-61 and 7-62). Table 8-44 defines this field for direct-access devices.

Table 8-44: Direct-Access Medium-Type Codes

Code Value	Medium Type				
00h	Default medium type (currently mounted medium type)				
01h	Flexible disk, single-sided; unspecified medium				
02h	Flexible disk, double-sided; unspecified medium				
	Flexible Disks				
	Diameter mm (inch)	Bit Density Bits/Radian	Track Density /mm (/inch)	Number Of Sides	Reference Standard
05h	200 (8.0)	6631	1.9 (48)	1	ANSI X3.73-1980
06h	200 (8.0)	6631	1.9 (48)	2	(Note 1)
09h	200 (8.0)	13262	1.9 (48)	1	None
0Ah	200 (8.0)	13262	1.9 (48)	2	ANSI X3.121-1984
0Dh	130 (5.25)	3979	1.9 (48)	1	ANSI X3.82-1980
12h	130 (5.25)	7958	1.9 (48)	2	ANSI X3.125-1985
16h	130 (5.25)	7958	3.8 (96)	2	ANSI X3.126-1986
1Ah	130 (5.25)	13262	3.8 (96)	2	ISO DIS8630-1985
1Eh	90 (3.5)	7958	5.3 (135)	2	ANSI X3.137 (Note 2)
	Direct-Access Magnetic Tapes				
	Width mm (inch)	Tracks	Density ftpm (ftpi)	Reference Standard	
40h	6.3 (0.25)	12	394 (10000)	Note 1	
44h	6.3 (0.25)	24	394 (10000)	Note 1	
80h - FFh	Vendor-Specific				
All others	Reserved				

NOTES:

- (1) See Appendix D for additional standards information.
- (2) This standard is under development. Please contact the Secretariat for current information on status and availability.

The device specific parameter field (Table 8-45) is contained in the mode parameter header (see Table 7-61 and 7-62).

Table 8-45: Device Specific Parameter

Bit	7	6	5	4	3	2	1	0
	WP	Reserved		DPOFUA	Reserved			

A write protect (WP) bit of zero indicates that the medium is write enabled. A WP bit of one indicates that the medium is write protected.

When used with the MODE SELECT command, the DPOFUA bit is not used and the field is reserved.

When used with the MODE SENSE command, a DPOFUA bit of zero indicates that the target does not support the DPO and FUA bits. A DPOFUA bit of one indicates that the target supports the DPO and FUA bits (see 8.2.6).

The density code field is contained in the mode parameter block descriptor (see Table 7-63). This field is reserved for direct-access devices.

The mode page codes for direct-access devices are shown in Table 8-46.

Table 8-46: Mode Page Codes

Page Code	Description	Section
08h	Caching Page	8.3.3.1
0Ah	Control Mode Page	7.3.3.1
02h	Disconnect-Reconnect Page	7.3.3.2
05h	Flexible Disk Page	8.3.3.2
03h	Format Device Page	8.3.3.3
08h	Medium Types Supported Page	8.3.3.4
0Ch	Notch and Partition Page	8.3.3.5
09h	Peripheral Device Page	7.3.3.3
01h	Read-Write Error Recovery Page	8.3.3.6
04h	Rigid Disk Geometry Page	8.3.3.7
07h	Verify Error Recovery Page	8.3.3.8
00h	Vendor-Specific (does not require page format)	
06h	Reserved	
0Dh - 1Fh	Reserved	
3Fh	Return all pages (valid only for the MODE SENSE command)	
20h - 3Eh	Vendor-specific	

For direct-access devices, if the notch page is not supported, or if the active notch field in the notch page is zero then each page descriptor specifies mode parameters for the target to use for subsequent operations on the specified logical unit. If the notch page is supported, and the active notch is not zero, then each page descriptor specifies parameters for the target to use for subsequent operations on the disk notch specified by the current value of the active notch field on the specified logical unit.

8.3.3.1. Caching Page

Table 8-47: Caching Page

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved		Page Code (08h)				
1	Page Length (0Ah)							
2	Reserved					WCE	MF	RCD
3	Demand Read Retention Priority				Write Retention Priority			
4	(MSB)							
5	Disable Pre-fetch Transfer Length							
6	(MSB)							
7	Minimum Pre-fetch							
8	(MSB)							
9	Maximum Pre-fetch							
10	(MSB)							
11	Maximum Pre-fetch Ceiling							

The caching parameters page defines the parameters that affect the use of the cache.

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

A write cache enable (WCE) bit of zero specifies that the target shall return GOOD status for a WRITE command after successfully writing all of the data to the medium. A WCE bit of one specifies that the target may return GOOD status for a WRITE command after successfully receiving the data and prior to having successfully written it to the medium.

A multiplication factor (MF) bit of zero specifies that the target shall interpret the minimum and maximum pre-fetch fields in terms of the number of logical blocks for each of the respective types of pre-fetch. An MF bit of one specifies that the target shall interpret the minimum and maximum pre-fetch fields to be specified in terms of a scalar number which, when multiplied by the number of logical blocks to be transferred for the current command, yields the number of logical blocks for each of the respective types of pre-fetch.

A read cache disable (RCD) bit of zero specifies that the target may return data requested by a READ command by accessing either the cache or media. A RCD bit of one specifies that the target shall transfer all of the data requested by a READ command from the medium (i.e., data cannot be transferred from the cache).

The demand read retention priority field advises the target on the retention priority to assign data read into the cache that has also been transferred from the target to the initiator.

The write retention priority field advises the target on the retention priority to assign data written into the cache that has also been transferred from the cache memory to the medium.

Table 8-48: Demand Read and Write Retention Priority

Value	Description
0h	Indicates the target should not distinguish between retaining the indicated data and data placed into the cache memory by other means (e.g., pre-fetch).
1h	Demand Read Retention Priority: Data put into the cache via a READ command should be replaced sooner (has lower priority) than data placed into the cache by other means (e.g., pre-fetch). Write Retention Priority: Data put into the cache during a WRITE or WRITE AND VERIFY command should be replaced sooner (has lower priority) than data placed into the cache by other means (e.g., pre-fetch).
Fh	Demand Read Retention Priority: Data put into the cache via a READ command should not be replaced if there is other data in the cache that was placed into the cache by other means (e.g., pre-fetch) and it may be replaced (i.e., it is not locked). Write Retention Priority: Data put into the cache during a WRITE or WRITE AND VERIFY command should not be replaced if there is other data in the cache that was placed into the cache by other means (e.g., pre-fetch) and it may be replaced (i.e., it is not locked).
2h-Eh	Reserved

An anticipatory pre-fetch occurs when data is placed in the cache that has not been requested. This usually happens in conjunction with the reading of data that has been requested. The following parameters are advisory, and give an indication to the target how it should manage the cache based on the last READ command. An anticipatory pre-fetch may occur based on other information.

The disable pre-fetch transfer length field specifies the selective disabling of anticipatory pre-fetch on long transfer lengths. The value in this field is compared to the number of blocks requested by the current READ command. If the number of blocks is greater than the disable pre-fetch transfer length, then the target is advised that no anticipatory pre-fetch should be done for that command. Otherwise the target is advised that it should attempt an anticipatory pre-fetch. If the pre-fetch disable transfer length is set to zero, then all anticipatory pre-fetching is disabled for any request for data, including those for zero logical blocks.

The minimum pre-fetch field is an advisory parameter. It is interpreted as specifying either a number of blocks or a scalar multiplier of the transfer length, depending upon the setting of the MF bit. In either case, the resulting number of blocks is the number that the initiator advises the target to pre-fetch regardless of the delays it might cause in executing subsequent

commands.

The pre-fetching operation begins at the logical block immediately after the last logical block of the previous READ command. Pre-fetching shall always halt before the end of the media. Errors that occur during the pre-fetching operation shall not be reported to the initiator unless that target cannot, as a result of the error, execute subsequent commands correctly. In this case the error may be reported either immediately as an error for the current READ command, or as a deferred error, at the discretion of the target and according to the rules for reporting deferred errors.

By setting the minimum pre-fetch field to zero, the initiator is advising the target that pre-fetching should be terminated whenever another command is ready for executing. The initiator can advise the target to ignore this consideration by setting the minimum pre-fetch equal to the maximum pre-fetch.

The maximum pre-fetch field is an advisory parameter. It is interpreted as specifying either a number of blocks or a scalar multiplier of the transfer length, depending upon the setting of the MF bit. In either case, the resulting number of blocks is the number that the initiator advises the target to pre-fetch if there are no delays in executing subsequent commands.

The maximum pre-fetch field advises the target on the maximum amount of data to ever pre-fetch into the cache as a result of one READ command. It is used in conjunction with the disable pre-fetch transfer length and maximum pre-fetch ceiling parameters to advise the target on how to trade off pre-fetching new data from the media with displacing old data already stored in the cache.

The maximum pre-fetch ceiling field is a control parameter. It specifies an upper limit on the number of logical blocks computed as the maximum pre-fetch. If this number of blocks is greater than the maximum pre-fetch ceiling, then the number of logical blocks to maximally pre-fetch shall be truncated to the value stored in the ceiling field.

IMPLEMENTORS NOTE: If the MF bit is one the maximum pre-fetch ceiling field is useful in limiting the amount of data to be pre-fetched.

8.3.3.2. Flexible Disk Page

Table 8-49: Flexible Disk Page (Bytes 0-19)

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved	Page Code (05h)					
1	Page Length in Bytes (1Eh)							
2	(MSB)	Transfer Rate						(LSB)
3	---							
4	---	Number of Heads						---
5	Sectors per Track							
6	(MSB)	Data Bytes per Sector						(LSB)
7	---							
8	(MSB)	Number of Cylinders						(LSB)
9	---							
10	(MSB)	Starting Cylinder-Write Precompensation						(LSB)
11	---							
12	(MSB)	Starting Cylinder-Reduced Write Current						(LSB)
13	---							
14	(MSB)	Drive Step Rate						(LSB)
15	---							
16	---	Drive Step Pulse Width						---
17	(MSB)	Head Settle Delay						(LSB)
18	---							
19	Motor on Delay							

The definition of this page is continued in Table 8-50.

Table 8-50: Flexible Disk Page (Bytes 20-31)

Bit Byte	7	6	5	4	3	2	1	0
20	Motor off Delay							
21	TRDY	SSN	MO	Reserved				
22	Reserved				SPC			
23	Write Compensation							
24	Head Load Delay							
25	Head Unload Delay							
26	Pin 34				Pin 2			
27	Pin 4				Pin 1			
28	(MSB)							
29	Medium Rotation Rate							
30	Reserved							
31	Reserved							
	(LSB)							

The flexible disk page (Tables 8-49 and 8-50) contains parameters for control and reporting of flexible disk drive parameters.

This page is 32 bytes in length. Bytes 0-19 are defined in Table 8-49 and bytes 20-31 are defined in Table 8-50.

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

IMPLEMENTORS NOTE: This page is mainly intended for defining parameters of flexible disk drives, but may be used for other devices, if applicable.

The transfer rate indicates in kilobits per second the data rate of the peripheral device.

Common Transfer Rates

MSB	LSB	DESCRIPTION OF TRANSFER RATE
00h	FAh	250 kbit/second transfer rate
01h	2Ch	300 kbit/second transfer rate
01h	F4h	500 kbit/second transfer rate
03h	E8h	1 megabit/second transfer rate
07h	00h	2 megabit/second transfer rate
13h	88h	5 megabit/second transfer rate

The number of heads field specifies the number of heads used for reading and writing data on the medium. Heads used exclusively for servo information are excluded.

The sectors per track field specifies the number of sectors per revolution per head.

The data bytes per sector field specifies the number of bytes of data per sector that an initiator can read or write.

The number of cylinders field specifies the number of cylinders used for data storage.

The starting cylinder for write precompensation field specifies the cylinder at which write precompensation is to begin. Cylinders are numbered starting with zero. If the starting cylinder for write precompensation is equal to the value in the number of cylinders field, write precompensation shall be disabled by the target.

The starting cylinder for reduced write current field specifies cylinder at which write current is reduced. Cylinders are numbered starting with zero. If the starting cylinder for reduced write current is equal to the value in the number of cylinders field, reduced write current shall be disabled by the target.

The drive step rate field specifies the step rate in units of 100 microseconds. This value may be rounded as defined in 6.5.4. A value of zero requests the target to set its default value.

The drive step pulse width field specifies the width of the step pulse in microseconds. This value may be rounded as defined in 6.5.4. A value of zero requests the target to set its default value.

The head settle delay field specifies the head settle time in units of 100 microseconds. This value may be rounded as defined in 6.5.4. A value of zero requests the target to set its default value.

If a true ready signal is not available, the motor on delay field specifies in tenths of a second the time that the target shall wait before attempting to access the medium after the motor on signal is asserted. If a true ready signal is available, the motor on delay field specifies in tenths of a second the time that the target shall wait for drive ready status before aborting an attempt to access the medium. This value may be rounded as defined in 6.5.4.

The motor off delay field specifies in tenths of a second the time that the target shall wait before releasing the motor on signal after an idle condition exists. A value of FFh indicates that the motor on signal shall not be released. The START STOP UNIT command is not affected by this parameter. This value may be rounded as defined in 6.5.4.

A true ready (TRDY) bit of one specifies that a signal is provided that indicates the medium is ready to be accessed.

A start sector number (SSN) bit of one specifies that sectors are numbered starting with one. An SSN bit of zero specifies that sectors are numbered starting with zero.

A motor on (MO) bit of one specifies that pin 16 (motor on) shall remain released. This bit shall be set to one when using high capacity (192 tracks per inch) drives and their pre-formatted diskettes. An MO bit of zero indicates that pin 16 (motor on) shall be asserted.

The step pulse per cylinder (SPC) field is used to specify the number of additional step pulses required per cylinder. Non-zero values allow a drive to read a diskette formatted on a drive with a lower number of tracks per inch. For example, a value of one allows a 96 track-per-inch drive to access tracks on a diskette that was formatted for 48 tracks per inch.

The write compensation field is used to specify the amount of write compensation to be used starting at the cylinder specified in the starting cylinder for write precompensation field. The correlation of any values used in this field to actual write precompensation time values is vendor-specific. If a zero is specified in this field the target shall use its default write precompensation value. This value may be rounded as defined in 6.5.4.

The head load delay field specifies the head loading time in milliseconds. This value may be rounded as defined in 6.5.4. A value of zero requests the target to set its default value.

The head unload delay field specifies the head unloading time in milliseconds. This value may be rounded as defined in 6.5.4. A value of zero requests the target to set its default value.

The Pin 34 field defines the usage of pin 34 of the flexible disk drive interface. This use of this pin varies among vendors and drives. The following settings allow the initiator to select how pin 34 shall be used by the interface.

Bit	7	6	5	4	Pin 34 Use
P	0	0	0		Open
P	0	0	1		Ready
P	0	1	0		Disk Changed
P	0	1	1		Reserved
P	1	X	X		Reserved - X is 0 or 1

					Polarity Bit - "0" - Active Low "1" - Active High

The Pin 4 field defines the usage of pin 4 of the flexible disk drive interface. This use of this pin varies among drive vendors and drives. The following settings allow the initiator to specify how pin 4 shall be used by the interface.

BIT	7	6	5	4	Pin 4 Use
P	0	0	0		Open
P	0	0	1		In Use
P	0	1	0		Eject
P	0	1	1		Head Load
P	1	X	X		Reserved - X is "0" or "1"

					Polarity Bit - "0" - Active Low "1" - Active High

The Pin 1 field defines the usage of pin 1 of the flexible disk drive interface. This use of this pin varies among vendors and drives. The following settings allow the initiator to specify how pin 1 shall be used by the interface.

Bit	3	2	1	0	Pin 1 Use
P	0	0	0		Open (Pin 1 is not used)
P	0	0	1		Disk Change Reset
P	0	1	X		Reserved
P	1	X	X		Reserved

					Polarity Bit - "0" - Active Low "1" - Active High

The medium rotation rate field specifies the speed at which the medium rotates. The unit of measure is rotations per minute (e.g., 2400 rpm). This field cannot be changed by a MODE SELECT command.

8.3.3.3. Format Device Page

Table 8-51: Format Device Page

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved	Page Code (03h)					
1	Page Length (16h)							
2	(MSB)	Tracks per Zone						(LSB)
3								
4	(MSB)	Alternate Sectors per Zone						(LSB)
5								
6	(MSB)	Alternate Tracks per Zone						(LSB)
7								
8	(MSB)	Alternate Tracks per Logical Unit						(LSB)
9								
10	(MSB)	Sectors per Track						(LSB)
11								
12	(MSB)	Data Bytes per Physical Sector						(LSB)
13								
14	(MSB)	Interleave						(LSB)
15								
16	(MSB)	Track Skew Factor						(LSB)
17								
18	(MSB)	Cylinder Skew Factor						(LSB)
19								
20	SSEC	HSEC	RMB	SURF	Reserved			
21	Reserved							
23								

The format device page (Table 8-51) contains parameters which specify the medium format.

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

IMPLEMENTORS NOTE: If the initiator changes any of the current physical parameters defined below, the target may not be able to access the media until a FORMAT UNIT command has been successfully completed.

If the defect handling format parameters (tracks per zone, alternate sectors per zone, alternate tracks per zone and alternate tracks per logical unit) requested by the initiator are not supported by the target the target may round these fields to acceptable values as described in 6.5.4.

The tracks per zone field specifies the number of tracks per zone to use in dividing the capacity of the device for the purpose of allocating alternate sectors. A value of zero means that one zone is defined for the entire device. The last zone on the device might not contain the same number of tracks as the previous zone(s).

The alternate sectors per zone field specifies the number of sectors per zone the target shall reserve for defect handling. The target shall de-allocate these sectors from the initiator addressable blocks during the FORMAT UNIT command. If the notch page is implemented and the ND bit of the notch page is one and the active notch field of the notch page is zero, then a value of zero indicates that no alternate sectors shall be reserved. Otherwise, a value of zero indicates that the number of alternate sectors is target specific.

The alternate tracks per zone field specifies the number of tracks per zone the target shall reserve for defect handling. The target shall de-allocate these tracks from the initiator addressable blocks during the FORMAT UNIT command. If the notch page is implemented and the ND bit of the notch page is one and the active notch field of the notch page is zero, then a value of zero indicates that no alternate tracks shall be reserved. Otherwise, a value of zero indicates that the number of alternate tracks is target specific.

The alternate tracks per logical unit field specifies the number of tracks per logical unit the target shall reserve for defect handling. The target shall de-allocate these tracks from the initiator addressable blocks during the FORMAT UNIT command. If the notch page is implemented and the ND bit of the notch page is one and the active notch field of the notch page is zero, then a value of zero indicates that no alternate tracks shall be reserved. Otherwise, a value of zero indicates that the number of alternate tracks is target specific.

The sectors per track field specifies the number of physical sectors included within each track. This number includes any alternate sectors the target may allocate. A value of zero during MODE SELECT indicates that the target shall define the number of sectors per track. For devices with a variable number of sectors per track, the value in MODE SELECT shall be zero and the value reported in MODE SENSE for the number of sectors per track is

vendor-specific.

The data bytes per physical sector field specifies the number of data bytes per physical sector that the target shall use. This value may be different than the logical block size reported in the MODE SELECT data. The target shall return CHECK CONDITION status if it determines that the combination of this field and the sectors per track field exceed the capability of the medium. A value of zero indicates that the data bytes per physical sector is defined by the target.

For MODE SENSE the interleave field returns the same parameter passed in the FORMAT UNIT command. The target shall report this field as target defined in the corresponding MODE SENSE command. For MODE SELECT this field is ignored.

The track skew factor field specifies the number of physical sectors between the last logical block of one track and the first logical block on the next sequential track of the same cylinder.

The cylinder skew factor field specifies the number of physical sectors between the last logical block of one cylinder and the first logical block on the next sequential cylinder.

The SSEC bit set to one indicates that the target shall use soft sector formatting.

The HSEC bit set to one indicates that the target shall use hard sector formatting. The HSEC bit and the SSEC bit are mutually exclusive in MODE SELECT commands.

Possible combinations of MODE SENSE Report Default Values

SSEC	HSEC	Description
0	0	Target shall not return this combination
1	0	Target supports soft sector formatting only
0	1	Target supports hard sector formatting only
1	1	Target supports both soft and hard sector formatting

Possible combinations of MODE SENSE Report Changeable Values

SSEC	HSEC	Description
0	0	Sector formatting Not changeable
1	0	Target shall not return this combination
0	1	Target shall not return this combination
1	1	Target supports both Soft and Hard sector formatting

The removable (RMB) bit set to one indicates that the logical unit supports removable media. A RMB bit set to zero indicates that the logical unit does not support removable media. The status of this bit shall be reflected in the INQUIRY command removable media bit (RMB).

The surface (SURF) bit set to zero indicates that the target shall allocate progressive addresses to all logical blocks within a cylinder prior to allocating addresses on the next cylinder. A SURF bit set to one indicates that the target shall allocate progressive addresses to all logical blocks on a surface prior to allocating sector addresses on the the next surface.

IMPLEMENTORS NOTE: If the target supports savable parameters, all savable parameters for this initiator, including those in page codes 3, 4, and 5, will be saved to non-volatile memory when the save parameters bit (SP) in the command descriptor block is set to one. The savable parameters may also be saved to non-volatile memory during a FORMAT UNIT command (see 8.2.1).

8.3.3.4. Medium Types Supported Page

Table 8-52: Medium Types Supported Page

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved	Page Code (0Bh)					
1	Page Length (06h)							
2	Reserved							
3	Reserved							
4	Medium Type One Supported							
5	Medium Type Two Supported							
6	Medium Type Three Supported							
7	Medium Type Four Supported							

The medium types supported page (Table 8-52) contains a list of the medium types implemented by the target for logical units.

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

The code values for each medium type supported by the target (up to four maximum), as defined in the MODE SELECT command (see Table 8-34), are reported in ascending order. If only the default medium type is supported zero is reported. If less than four medium types are supported the unused entries shall be returned as zero.

8.3.3.5. Notch and Partition Page

Table 8-53: Notch Page

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved	Page Code (0Ch)					
1	Page Length (16h)							
2	ND	LPN	Reserved					
3	Reserved							
4	(MSB)	Maximum Number of Notches						---
5	---							(LSB)
6	(MSB)	Active Notch						---
7	---							(LSB)
8	(MSB)	Starting Boundary						--
11	--							(LSB)
12	(MSB)	Ending Boundary						--
15	--							(LSB)
16	(MSB)	Pages Notched						--
23	--							(LSB)

The notch page (Table 8-53) contains parameters for direct-access devices which implement a variable number of blocks per cylinder and support this page. Each section of the logical unit with a different number of blocks per cylinder is referred to as a notch.

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

A notched drive (ND) bit of zero indicates that the device is not notched and that all other parameters in this page shall be returned as zero by the target. A ND bit of one indicates that the device is notched. For each supported active notch value this page defines the starting and ending boundaries of the notch.

A logical or physical notch (LPN) bit of zero indicates that the boundaries are based on the physical parameters of the logical unit. The cylinder is considered most significant, the head least significant. A LPN bit of one indicates that the notch boundaries are based on logical blocks of the logical unit.

The maximum number of notches field indicates the maximum number of notches supported by the logical unit. This field shall be reported as unchangeable.

The active notch field indicates the notch that this and subsequent MODE SELECT and MODE SENSE commands shall refer to, until the active notch is changed by a later MODE SELECT command. The value of the active notch shall be greater than or equal to 0 and less than or equal to the maximum number of notches. An active notch value of zero indicates that this and subsequent MODE SELECT and MODE SENSE commands refer to the parameters that apply across all notches.

The starting boundary field indicates the beginning of the active notch or, if the active notch is zero, the beginning boundary of the logical unit. If the PLN bit is one, then the four bytes represent a logical block address. If the PLN bit is zero, then the three most significant bytes shall represent the cylinder number and the least significant byte shall represent the head number. This field shall be reported as unchangeable.

The ending boundary field indicates the ending of the active notch or, if the active notch is zero, the ending of the logical unit. If the PLN bit is one, then the four bytes represent logical block address. If the PLN bit is zero, then the three most significant bytes shall represent the cylinder number and the least significant byte shall represent the head number. This field shall be reported as unchangeable.

Each notch shall span a set of consecutive logical blocks on the logical unit, the notches shall not overlap, and no logical block shall be excluded from a notch.

The pages notched field is a bit map of the mode page codes that indicates which pages contain parameters that may be different for different notches. The most significant bit of this field corresponds to page code 3Fh and the least significant bit corresponds to page code 00h. If a bit is one, then the corresponding mode page contains parameters that may be different for different notches. If a bit is zero, then the corresponding mode page contains parameters that are constant for all notches. This field shall be reported as unchangeable.

8.3.3.6. Read-Write Error Recovery Page

Table 8-54: Read-Write Error Recovery Page

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved	Page Code (01h)					
1	Page Length (0Ah)							
2	AWRE	ARRE	TB	RC	EER	PER	DTE	DCR
3	Read Retry Count							
4	Correction Span							
5	Head Offset Count							
6	Data Strobe Offset Count							
7	Reserved							
8	Write Retry Count							
9	Reserved							
10	(MSB)	Recovery Time Limit						
11							(LSB)	

The read-write error recovery page (Table 8-54) specifies the error recovery parameters the target shall use during any command that performs a read or write operation to the medium (e.g., READ(6), READ(10), WRITE(6), WRITE(10), COPY, COMPARE, WRITE & VERIFY, etc.).

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

An automatic write reallocation enabled (AWRE) bit of one indicates that the target shall enable automatic reallocation to be performed during write operations. The automatic reallocation shall be performed only if the target has the valid data (e.g., original data in the buffer or recovered from the medium). The valid data shall be placed in the reallocated block. Error reporting as required by the error recovery bits (EER, PER, DTE, and DCR) shall be performed only after completion of the reallocation. The reallocation operation shall report any failures that occur. See the REASSIGN BLOCKS command (8.2.10) for error procedures.

An AWRE bit of zero indicates that the target shall not perform automatic reallocation of defective data blocks during write operations.

An automatic read reallocation enabled (ARRE) bit of one indicates that the target shall enable automatic reallocation of defective data blocks during read operations. All error recovery actions required by the error recovery bits (TB, EER, PER, DTE, and DCR) shall be executed. The automatic reallocation shall then be performed only if the target successfully recovers the data. The recovered data shall be placed in the reallocated block. Error posting as required by the error recovery bits shall be performed only after completion of the reallocation. The reallocation process shall present any failures that occur. See the REASSIGN BLOCKS command (8.2.10) for error procedures.

An ARRE bit of zero indicates that the target shall not perform automatic reallocation of defective data blocks during read operations.

A transfer block (TB) bit of one indicates that a data block that is not recovered within the recovery limits specified shall be transferred to the initiator before CHECK CONDITION status is returned. A TB bit of zero indicates that such a data block shall not be transferred to the initiator. The TB bit does ^{not} affect the action taken for recovered data.

A read continuous (RC) bit of one indicates the target shall transfer the entire requested length of data without adding delays to perform error recovery procedures. This implies that the target may send data that is erroneous or fabricated in order to maintain a continuous flow of data. The target shall assign priority to this bit over conflicting error control bits (EER, DCR, DTE, and PER) within this byte.

IMPLEMENTORS NOTE: Fabricated data may be data already in the buffer or any other target-specific data. This bit is typically used in image processing, audio, or video applications.

A RC bit of zero indicates that error recovery operations that cause delays are acceptable during the data transfer. Data shall not be fabricated.

The individual bit definitions for EER, PER, DTE and DCR are contained in Table 8-55. The combinations of these bits are explained in Tables 8-56 through 8-59.

Table 8-55: Error Recovery Bit Definitions

EER	PER	DTE	DCR	Description
1	-	-	-	An enable early recovery (EER) bit of one indicates that the target shall use of the most expedient form of error recovery first. This bit only applies to data error recovery and it does not affect positioning retries and the message system error recovery procedures.
0	-	-	-	An EER bit of zero indicates that the target shall use an error recovery procedure that minimizes the risk of mis-detection or mis-correction.
-	1	-	-	A post error (PER) bit of one indicates that the target shall report recovered errors.
-	0	-	-	A PER bit of zero indicates that the target shall not report recovered errors. Error recovery procedures shall be performed within the limits established by the error recovery parameters.
-	-	1	-	A disable transfer on error (DTE) bit of one indicates that the target shall terminate the data phase upon detection of a recovered error.
-	-	0	-	A DTE bit of zero indicates that the target shall not terminate the data phase upon detection of a recovered error.
-	-	-	1	A disable correction (DCR) bit of one indicates that error correction codes shall not be used for data error recovery.
-	-	-	0	A DCR bit of zero allows the use of error correction codes for data error recovery.

IMPLEMENTORS NOTE: An EER bit of one may imply an increase in the probability of mis-detection or mis-correction. An EER bit of zero typically means that the specified retry limit is exhausted prior to using error correction codes.

Table 8-56: Combined Error Recovery Parameter Descriptions(1)

EER	PER	DTE	DCR	Description
0	0	0	0	The full number of retries (specified in the read, write or verify retry count field) and error correction are attempted to recover the data (EER and DCR set to 0). A CHECK CONDITION is not reported at the completion of the command for recovered errors (PER set to 0). The command terminates with CHECK CONDITION status before the transfer count is exhausted only if an unrecoverable error is detected. If an unrecoverable data error occurred, the data in the block with the unrecoverable error may or may not be transferred to the initiator depending on the setting of the transfer block (TB) bit (read operation only).
0	0	0	1	Error correction is disabled (DCR set to one) so only the full number of retries (specified in the read, write or verify retry count field) are attempted to recover the data (EER set to 0). A CHECK CONDITION is not reported at the completion of the command for recoverable errors (PER set to 0). The command terminates with CHECK CONDITION status before the transfer count is exhausted only if an unrecoverable error is detected. If an unrecoverable data error occurred, the data in the block with the unrecoverable error may or may not be transferred to the initiator depending on the setting of the transfer block (TB) bit (read operation only).
0	0	1	0	Invalid mode (PER must be set to one if DTE is set to one). See the note at the end of this table.
0	0	1	1	Invalid mode (PER must be set to one if DTE is set to one). See the note at the end of this table.
0	1	0	0	The full number of retries (specified in the read, write or verify retry count field) and Error Correction are attempted to recover the data (EER and DCR set to 0). The command terminates with CHECK CONDITION status before the transfer count is exhausted only if an unrecoverable error is detected. If an unrecoverable data error occurred, the data in the block with the unrecoverable error may or may not be transferred to the initiator depending on the setting of the transfer block (TB) bit (read operation only). A CHECK CONDITION with a sense key of RECOVERED ERROR is reported at the completion of the command for any recoverable error that occurs (PER set to 1). The information field in the sense data shall contain the logical block address of the last recovered error which occurred during the transfer.

Table 8-57: Combined Error Recovery Parameter Descriptions(2)

EER	PER	DTE	DCR	Description
0	1	0	1	Error correction is disabled (DCR set to one) so only the full number of retries (specified in the read, write or verify retry count field) are attempted to recover the data (EER set to 0). The command terminates with CHECK CONDITION status before the transfer count is exhausted only if an unrecoverable error is detected. If an unrecoverable data error occurred, the data in the block with the unrecoverable error may or may not be transferred to the initiator depending on the setting of the transfer block (TB) bit (read operation only). A CHECK CONDITION with a sense key of RECOVERED ERROR is reported at the completion of the command for any recoverable error that occurs (PER set to 1). The information field in the sense data shall contain the logical block address of the last recovered error which occurred during the transfer.
0	1	1	0	The full number of retries (specified in the read, write or verify retry count field) and Error Correction are attempted to recover the data (EER and DCR set to 0). The command terminates with CHECK CONDITION status before the transfer count is exhausted if any error (recoverable or unrecoverable) is detected (DTE set to 1). The information field in the sense data shall contain the logical block address of the block in error. If an unrecoverable data error occurs the data in the block with the error may or may not be transferred to the initiator depending on the setting of the transfer block (TB) bit (read operation only).
0	1	1	1	Error correction is disabled (DCR set to one) so only the full number of retries (specified in the read, write or verify retry count field) are attempted to recover the data (EER set to 0). The command terminates with CHECK CONDITION status before the transfer count is exhausted if any error (recoverable or unrecoverable) is detected (DTE set to 1). The information field in the sense data shall contain the logical block address of the block in error. If an unrecoverable data error occurs the data in the block with the error may or may not be transferred to the initiator depending on the setting of the transfer block (TB) bit (read operation only).

Table 8-58: Combined Error Recovery Parameter Descriptions(3)

EER	PER	DTE	DCR	Description
1	0	0	0	The fewest possible retries and error correction are attempted to recover the data (EER set to one and DCR set to 0). A CHECK CONDITION is not reported at the completion of the command for recoverable errors (PER set to 0). The command terminates with CHECK CONDITION status before the transfer count is exhausted only if an unrecoverable error is detected. If an unrecoverable data error occurred, the data in the block with the unrecoverable error may or may not be transferred to the initiator depending on the setting of the transfer block (TB) bit (read operation only).
1	0	0	1	Invalid mode (DCR must be set to zero if EER is set to one). See the note at the end of this table.
1	0	1	0	Invalid mode (PER must be set to one if DTE is set to one). See the note at the end of this table.
1	0	1	1	Invalid mode (PER must be set to one if DTE is set to one). See the note at the end of this table.
1	1	0	0	The fewest possible retries and error correction are attempted to recover the data (EER set to one and DCR set to 0). The command terminates with CHECK CONDITION status before the transfer count is exhausted only if an unrecoverable error is detected. If an unrecoverable data error occurred, the data in the block with the unrecoverable error may or may not be transferred to the initiator depending on the setting of the transfer block (TB) bit (read operation only). A CHECK CONDITION with a sense key of RECOVERED ERROR is reported at the completion of the command for any recoverable error that occurs (PER set to 1). The information field in the sense data shall contain the logical block address of the last recovered error which occurred during the transfer.
1	1	0	1	Invalid mode (DCR must be set to zero if EER is set to one). See the note at the end of this table.

Table 8-59: Combined Error Recovery Parameter Descriptions(4)

EER	PER	DTE	DCR	Description
1	1	1	0	The fewest possible retries and error correction are attempted to recover the data (EER set to one and DCR set to 0). The command terminates with CHECK CONDITION status before the transfer count is exhausted if any error (recoverable or unrecoverable) is detected (DTE set to 1). The information field in the sense data shall contain the logical block address of the block in error. If an unrecoverable data error occurs the data in the block with the error may or may not be transferred to the initiator depending on the setting of the transfer block (TB) bit (read operation only).
1	1	1	1	Invalid mode (DCR must be set to zero if EER is set to one). See the note at the end of this table.

NOTE: If an invalid mode for the error recovery combination is sent by the initiator the target shall return CHECK CONDITION status with the sense key set to ILLEGAL REQUEST and the additional sense code set to INVALID FIELD IN PARAMETER LIST.

The read and write retry count fields specify the number of times that the target shall attempt its recovery algorithm during read and write operations, respectively. If the retry count field and the recovery time limit field are both specified in a MODE SELECT command, the field that requires the least time for data error recovery actions shall have priority.

The correction span field specifies the size, in bits, of the largest data error burst for which data error correction may be attempted. A correction span of zero specifies that the target shall use its default value or that this field is not supported.

The head offset count field specifies in two's-complement notation an incremental offset position from the track center to which the heads shall be moved. The effect of this field on write operations is unspecified. A head offset count of zero indicates that no offset is specified. A positive value indicates moving in the direction of increasing logical block addresses. A negative value indicates moving in the direction of decreasing logical block addresses. Any value specified in this field does not preclude the target from using positive or negative head offset during error recovery. However, after any error recovery is completed the target shall return the head offset to the value specified in this field.

IMPLEMENTORS NOTE: The degree of offset for each incremental value and the number of valid values are device specific. It is recommended that the number of valid values be equal for the positive and negative head offset counts.

The target shall return CHECK CONDITION status and shall set the sense key to ILLEGAL REQUEST if an unsupported head offset value is specified. The valid bit shall be set to one and the information field shall be set to the positive value of the maximum head offset count that is supported. The target shall set the valid bit to zero if it cannot determine the maximum head offset count supported. (Note: If the target does not support this field, it returns a zero value in the MODE SENSE command.)

The data strobe offset count field specifies in two's-complement notation an incremental position to which the recovered data strobe shall be adjusted from its nominal setting. The effect of this field on write operations is unspecified. A value of zero indicates that no data strobe offset is specified. A positive value indicates movement in a positive direction as defined by the target. A negative value indicates movement in the negative direction as defined by the target. Any value specified in this field does not preclude the target from using positive or negative data strobe offset during error recovery. However, after any error recovery is completed the target shall return the data strobe offset to the value specified in this field.

IMPLEMENTORS NOTE: The degree of offset for each incremental value and the number of valid values are device specific. It is recommended that the number of valid values be equal for the positive and negative data strobe offset counts.

The target shall return CHECK CONDITION status and shall set the sense key to ILLEGAL REQUEST if an unsupported data strobe offset count value is specified. The valid bit shall be set to one and the information field shall be set to the positive value of the maximum data strobe offset count that is supported. The target shall set the valid bit to zero if it cannot determine the maximum data strobe offset supported.

IMPLEMENTORS NOTE: If the target does not support the data strobe offset count field, it returns a zero value in the MODE SENSE command.

The recovery time limit field specifies in increments of one millisecond the maximum time duration that the target shall use for data error recovery procedures. The target may round this value as described in 6.5.4. The limits in this field specifies the maximum error recovery time allowed for any individual logical block. A recovery time limit of zero specifies that the target shall use its default value.

If both retry count and recovery time limit are specified, the field that specifies the recovery action of least duration shall have priority.

8.3.3.7. Rigid Disk Drive Geometry Page

Table 8-60: Rigid Disk Drive Geometry Page

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved	Page Code (04h)					
1	Page Length (16h)							
2	(MSB)	Number of Cylinders						(LSB)
4	-							
5	Number of Heads							
6	(MSB)	Starting Cylinder-Write Precompensation						(LSB)
8	-							
9	(MSB)	Starting Cylinder-Reduced Write Current						(LSB)
11	-							
12	(MSB)	Drive Step Rate						(LSB)
13	-							
14	(MSB)	Landing Zone Cylinder						(LSB)
16	-							
17	Reserved				RPL			
18	Rotational Offset							
19	Reserved							
20	(MSB)	Medium Rotation Rate						(LSB)
21	-							
22	Reserved							
23	Reserved							

The rigid disk drive geometry page (Table 8-60) specifies parameters for direct-access devices employing a rigid disk drive.

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

IMPLEMENTORS NOTE: This page is intended to define drive geometry parameters of rigid disk drives. It may be used for other devices if applicable.

The number of cylinders field defines the number of physical cylinders used for data storage.

The number of heads field defines the physical number of heads used for data storage. Heads used exclusively for servo information are excluded.

The starting cylinder for write precompensation field is the physical cylinder at which write precompensation is to begin. The first cylinder is number zero. If the starting cylinder for write precompensation is equal to the value in the number of cylinders field, write precompensation shall be disabled by the target.

The starting cylinder for reduced write current field is the physical cylinder at which write current is reduced. The first cylinder is number zero. If the starting cylinder for reduced write current is equal to the value in the number of cylinders field, reduced write current shall be disabled by the target.

The drive step rate field indicates the step rate in 100 nanoseconds increments. The target shall use the drive step rate, greater than or equal to the drive step rate specified. If the target rounds this field it shall terminate the command as described in 6.5.4. A value of zero requests the target to set its default value.

The landing zone cylinder field indicates two's complement location where the target shall position the disk heads. A negative value indicates that the heads are positioned below cylinder zero by that number of cylinders. A positive value greater than the number of cylinders indicates that the heads are positioned beyond the cylinders used for data storage at the cylinder location specified. A zero indicates that the default should be used.

The rotational position locking (RPL) field is used for spindle synchronization as defined in Table 8-61. See Appendix K of the Enhanced Small Device Interface (BSR X3.170-198X) for further information on synchronized spindles.

Table 8-61: Rotational Position Locking

RPL	Description
00b	Indicates that spindle synchronization is disabled or not supported.
01b	The target operates as a synchronized-spindle slave.
10b	The target operates as a synchronized-spindle master.
11b	The target operates as a synchronized-spindle master control.

IMPLEMENTORS NOTE: The signals and connectors used for rotational position locking are external to the SCSI bus and are not part of this standard.

If a target fails to achieve synchronization it shall create a unit attention condition to all initiators. The sense key shall be set to UNIT ATTENTION and the additional sense code set to RPL STATUS CHANGE.

If subsequent to achieving synchronization the target detects a change of synchronization:

(1) and, if the logical unit is not presently executing an I/O process for the initiator, the target shall create a unit attention condition. The sense key shall be set to UNIT ATTENTION and the additional sense code set to RPL STATUS CHANGE.

(2) and, if the logical unit is presently executing an I/O process and no other error occurs, then the target shall return CHECK CONDITION status. The sense key shall be set to RECOVERED ERROR if the target is able to complete the I/O process or HARDWARE ERROR if the target is unable to complete the I/O process. The additional sense code is set to RPL STATUS CHANGE.

The rotational offset indicates the amount of rotational skew that the target shall use when synchronized. The rotational skew is applied in the retarded direction (lagging the synchronized spindle master control). The value in the field is the numerator of a fractional multiplier that has 256 as its denominator (e.g., a value of 128 indicates a one-half revolution skew). A value of zero indicates that rotational offset shall not be used. This value may be rounded as defined in 6.5.4. The rotational offset is not used when a target is configured as synchronized-spindle master.

The medium rotation rate indicates the speed at which the medium rotates. The unit of measure is rotations per minute (e.g., 3600 rpm).

8.3.3.8. Verify Error Recovery Page

Table 8-62: Verify Error Recovery Page

Bit	7	6	5	4	3	2	1	0
Byte								
0	PS	Reserved	Page Code (07h)					
1	Parameter Length (0Ah)							
2	Reserved				EER	PER	DTE	DCR
3	Verify Retry Count							
4	Verify Correction Span							
5	Reserved							
6	Reserved							
7	Reserved							
8	Reserved							
9	Reserved							
10	(MSB)	Verify Recovery Time Limit						---
11	---							(LSB)

The verify error recovery page (Table 8-62) specifies the error recovery parameters the target shall use during the VERIFY command, the verify operation of the WRITE AND VERIFY command and the verify operation of the COPY AND VERIFY command.

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

The AWRE bit as defined in the read-write error recovery page (see Table 8-54) applies to the WRITE AND VERIFY command. The VERIFY and COMPARE commands shall not perform automatic reallocations. The COPY and COPY AND VERIFY commands might or might not perform automatic reallocations depending on the settings of the mode parameters.

The EER, PER, DTE and DCR bits are defined in Table 8-55. The combinations of these bits are defined in Tables 8-56 through 8-59.

The verify retry count field specifies the number of times that the target shall attempt its recovery algorithm during a verify operation. If the verify retry count and the verify recovery time limit are both specified, the one that requires the least time for data error recovery actions shall have priority.

The verify correction span field specifies the size, in bits, of the largest burst data error for which data error correction may be attempted. If the target does not implement this field, a value of zero is returned in MODE SENSE data.

The verify recovery time limit field specifies in increments of one millisecond the maximum time duration that the target shall use error recovery procedures to recover data for an individual logical block. The target may round this value as described in 6.5.4. If the verify retry count and the verify recovery time limit are both specified, the one that requires the least time for data error recovery actions shall have priority.

IMPLEMENTORS NOTE: To disable all types of correction and retries the initiator should set the EER bit to zero, the PER, DTE and DCR bits to one and the number of retries and recovery time limit to zero.

(This page is intentionally blank.)

8.4. Glossary for Direct-Access Devices

cache memory. A temporary (and often volatile) data storage area outside the user-accessible area. A cache memory is usually faster to access than the medium and thus has the effect of increasing data throughput by reducing the number of accesses to the medium.

non-volatile medium. Physical storage medium that retains data written to it for a subsequent read operation through a power-on condition. An example of this is a disk device that stores data as magnetic fields that do not require device power to exist.

notch. A notch refers to all or part of the medium having a consistent set of geometry parameters. Notches are used to increase storage capacity by optimizing the number of sectors per track between the inner and outer tracks.

read hit. When the logical block requested for a read operation is accessed from the cache memory, this is called a read hit.

user-accessible. The area of the medium that can be read from or written to by READ and WRITE commands.

volatile medium. Medium that may not retain data written to it for a subsequent read operation through a power-on condition. An example of this is a silicon memory device that loses data written to it if device power is lost.

9. Sequential-Access Devices

9.1. Sequential-Access Device Model

9.1.1. Physical Elements

Sequential-access devices (called devices below) optimize their use in storing or retrieving user data in a sequential manner. Since access is sequential, position changes typically take a long time, when compared to direct-access devices.

Sequential-access devices are usually tape devices. The remainder of this description is from the point of view of a tape device; however, other implementations are not precluded.

The recording medium for tape devices consists of various widths and lengths of a flexible substrate coated with a semi-permanent magnetic material. The recording medium may be wound onto single reels or encapsulated into cartridges containing both a supply reel and a take-up reel. Several American National Standards exist covering the construction of reels and cartridges for interchange as well as recording techniques for many of the format or density combinations.

A complete unit composed of the recording medium and its physical carrier (e.g., reel, cartridge, cassette) is called a volume. Volumes have an attribute of being mounted or demounted on a suitable transport mechanism.

Mounted is the state of a volume when the device is physically capable of executing commands that cause the the medium to be moved. A volume is demounted when it is being loaded, threaded, unloaded, unthreaded, or when not attached to the device.

Ready is the state of the device when medium access and non-medium access commands can be executed. The device is not ready when no volume is mounted or, from the initiator's perspective, whenever all medium access commands report CHECK CONDITION status and a NOT READY sense key. Some devices may have a separate switch function which places the device in a not ready state even when a volume is mounted.

The write enabled or write protected state determines when an initiator may write information on a volume. This attribute is usually controlled by the user of the volume through manual intervention (e.g., thumbwheel switch).

The recording medium has two physical attributes called beginning-of-medium (BOM) and end-of-medium (EOM). Beginning-of-medium is at the end of the medium that is attached to the take-up reel. End-of-medium is at the end of the medium that is attached to the supply reel. In some cases, the medium is permanently affixed to one or both of the reel hubs.

As shown in Figure 9-1, the entire physical length of medium is not usable for recording data. For most volumes, a length of the medium is reserved before the beginning-of-medium and after the end-of-medium position. This is done to provide sufficient tape wraps onto the reel hub(s) and to ensure that recording starts in an undamaged section of the medium.

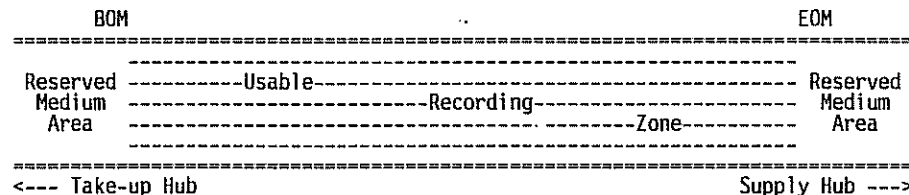


Figure 9-1: Typical Volume Layout

9.1.2. Data Storage Characteristics

The position on the medium where a pattern of recorded signal may be written by one write component is called a track (Figure 9-2). A device may write or read from one or more tracks at a time, depending on the format.

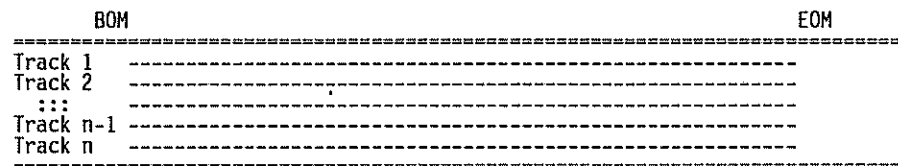


Figure 9-2: Typical Medium Track Layout

On a new volume, recording of one or more tracks begins after mounting the volume and moving from beginning-of-medium toward end-of-medium. The number of tracks written at one time is called a track group (TrkGrp). For recorded volumes, reading in the forward direction follows the same course of tracks as when writing.

If not all tracks are recorded at the same time, and the device reverses direction when approaching end-of-medium and begins writing on remaining tracks, the recording method is called serpentine. For serpentine devices that record only one track at a time, each physical track represents one track group (Figure 9-3).

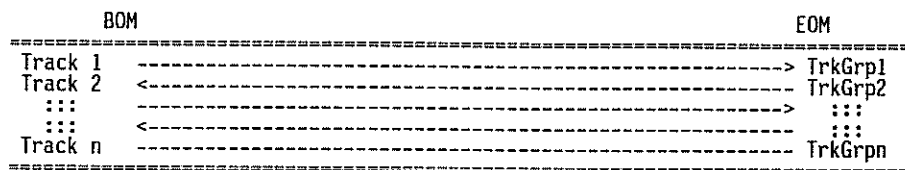


Figure 9-3: Serpentine Recording Example

Some multi-track devices have only one track group, using a parallel storage format that supports the simultaneous recording of all available tracks (Figure 9-4).

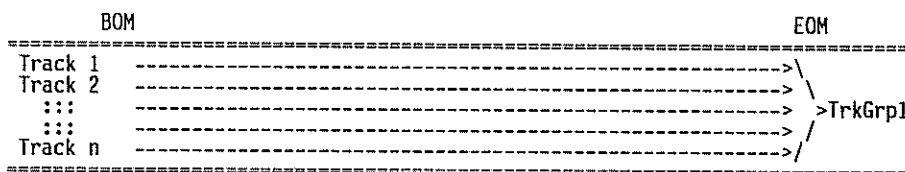


Figure 9-4: Parallel Recording Example

The serpentine and parallel recording formats shown in the previous examples define tracks as longitudinal patterns of recorded information. One other storage format used by some devices records tracks diagonally across the medium. This recording technique is known as helical scan (Figure 9-5).

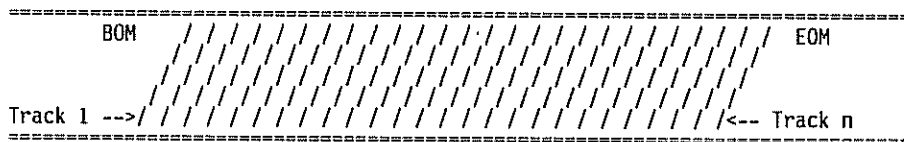


Figure 9-5: Helical Scan Recording Example

For most recording formats, an area at beginning-of-medium contains a format identification in the form of a tone burst or some other recognizable pattern. User data is not recorded in this area. The format identification is an attribute of a volume used for interchange purposes and is defined in applicable standards.

When writing, the initiator needs an indication that it is approaching the end of the permissible recording area. This position, called early-warning (EW), is typically reported to the initiator at a position early enough for the device to write any buffered data to the medium while still leaving enough room for additional recorded labels or filemarks. Some American National Standards include physical requirements for a marker placed on the medium to be detected by the device as early-warning (Figure 9-6).

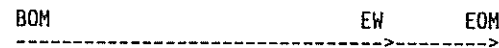


Figure 9-6: Early-Warning Example

For devices which implement large data buffers, the early-warning position defined by a physical marker may be too close to the end of the recording region to permit emptying the data buffer(s). For these devices, a logical concept of early-warning is often used to signal the initiator at an appropriate location prior to the physical marker.

9.1.3. Partitions within a Volume

Another attribute of a volume is called a partition. Partitions consist of one or more non-overlapped mini-volumes, each with its own beginning and ending points, occupying a single physical volume. Each partition (x) within a volume has a defined beginning-of-partition (BOPx), an early-warning position (EWx), and an end-of-partition (EOPx).

All volumes have a minimum of one partition called partition zero, the default data partition. For devices which support only one partition, the beginning-of-partition zero (BOP0) may be equivalent to the beginning-of-medium and the end-of-partition zero (EOP0) may be equivalent to the end-of-medium.

When a volume is mounted, it is logically positioned to beginning of the default data partition (BOP0). When a REWIND command is received in any partition (x), the device positions to the beginning-of-partition (BOPx).

Partitions on a volume do not need to be recorded in any defined order, nor do all partition numbers in a sequence need to be present on a volume. It is sufficient for a device to be able to locate a partition, given its code value, or determine that it does not exist on the volume. For interchange, information about which partitions are present on a volume may be stored on the volume in a device-defined area (possibly unavailable to the initiator) or the information may be an intrinsic attribute of the device implementation.

Figure 9-7 shows a possible partition implementation for a four-track serpentine recording device, assuming that each track group defines a partition.

	BOM	EOM
TrkGrp1	BOP0 ----->EWO----->	EOP0
TrkGrp2	EOP1 <-----EW1<----->	BOP1
TrkGrp3	BOP2 ----->EW2----->	EOP2
TrkGrp4	EOP3 <-----EW3<----->	BOP3

Figure 9-7: Partitioning Example - One Partition per Track Group

Another possible partition implementation for this four-track serpentine recording device is shown in Figure 9-8, using two track groups to define each partition.

	BOM	EOM
TrkGrp1	BOP0 ----->	
TrkGrp2	EOP0 <-----EW0<----->	
TrkGrp3	BOP1 ----->	
TrkGrp4	EOP1 <-----EW1<----->	

Figure 9-8: Partitioning Example - One Partition per Two Track Groups

The previous examples show the beginning and ending points for a partition aligned with physical bounds of the medium. This is not a mandatory requirement for partitioning; it is sufficient for a device to be able to locate to and stay in any partition bounded by a BOPx and EOPx. In this case, a recorded mark or some other device-recognizable attribute could be used to delineate the partitions. Figure 9-9 shows a possible two-partition implementation for a device with only one track group.

	BOM	EOM
	----->----->----->----->	
TrkGrp1	BOP0 ----->EWO----->EOP0/BOP1----->EW1----->	EOP1
	----->----->----->----->	

Figure 9-9: Partitioning Example - Two Partitions per Track Group

Three methods are defined in the MODE SENSE and MODE SELECT commands for managing partitions; each method is progressively more difficult to implement and manage:

- (1) device-defined fixed locations;
- (2) device-defined based on an initiator supplied number of partitions and a device specific allocation algorithm; and,
- (3) definition by name and capacity by an initiator.

9.1.4. Logical Elements within a Partition

The area between BOPx and EOPx on a typical recorded volume contains at least two types of initiator accessible elements, data blocks and tape marks. These elements are controlled and transferred between the initiator and the medium using READ, READ REVERSE, WRITE, and WRITE FILEMARKS commands.

A unit of data supplied or requested by an initiator is called a logical block. Logical blocks are stored according to the specifications of the format for the volume and may be recorded as one or more physical blocks on the medium. When the physical block and the logical block are not recorded in a one-to-one relationship, it is the responsibility of the device to perform all blocking, de-blocking, padding, stripping, splitting or rebuilding of the logical data block(s) sent by an initiator.

Filemarks are special recorded elements containing no user data. The filemark format is defined in some American National Standards. Initiators traditionally use filemarks to separate user data from labels and logical groupings of data from each other. Since some format standards do not define an explicit end-of-data (EOD), host software has often used conventions with filemarks to represent an EOD indication. At least one American National Standard specifically defines filemark use for this purpose. In some implementations, the device's EOD definition may be specified the initiator using the MODE SELECT command.

A setmark is another type of special recorded element containing no user data, providing a segmentation scheme hierarchically superior to filemarks. This level of segmentation is useful for some high capacity storage devices to provide concise addressability and fast positioning to specific sets of data within a partition. In some implementations, the detection and reporting of setmarks may be controlled by the initiator using the MODE SELECT command.

Inter-block gaps, the gaps between blocks, filemarks, and setmarks, are introduced on the medium at the time a block or mark is written without explicit action by the initiator. Minimum and maximum lengths for inter-block gaps are defined in some American National Standards. In some devices, the length of inter-block gaps may be preselected by the initiator using the MODE SELECT command while in other devices the gaps are fixed and non-changeable.

In addition to blocks, filemarks, and setmarks, erase gaps may be recorded on the medium through use of the ERASE command or device-initiated error recovery actions. Although explicitly recorded on the medium, there is normally no distinction between two contiguous erase gaps. An erase gap may be a length of erased medium or a recorded pattern not distinguishable as a block or mark. Minimum and maximum lengths for erase gaps are defined in some American National Standards while some devices may have no implementation of an erase gap.

After writing data from BOPx, the medium is considered to be a contiguous grouping of blocks, filemarks, setmarks, and gaps. Certain American National Standards define gap lengths which, if exceeded, are to be considered as having reached blank medium. Depending on the format, this blank medium may be treated as an end-of-data indication, an error recovery area, or an unrecoverable medium error causing an interchange error. Unrecorded volumes (new or erased) may exhibit blank medium characteristics if an attempt is made to read or space the volume before data has been written.

A sequential-access device may be capable of supporting fixed or variable length blocks. The concept of fixed or variable mode for writing and reading blocks only indicates the method by which the initiator specifies the size of a logical block for transfer and not the method of recording physical blocks on the medium. However, a device that supports only fixed-length physical blocks may only be capable of supporting logical blocks of the same length. The length of a logical block is always described in bytes. The length of a physical block may or may not be recorded as an exact byte count, depending on the format.

9.1.5. Data Buffering

A device may contain a temporary storage area capable of holding one or more logical blocks - a data buffer. A device data buffer may include any combination of blocks, filemarks, and setmarks in the process of being written to the medium, or it may contain read-ahead data blocks transferred from the medium.

A device with a data buffer may be capable of operating in either a buffered mode or an unbuffered mode. A device with no data buffer operates only in unbuffered mode. Either term is only applicable to the manner in which the device manages information to be written to the medium. Buffered mode is not applicable during read commands, regardless of whether read data passes through a data buffer.

A device operating in buffered mode may return GOOD status for write operations when all write data has been successfully transferred from the initiator into the device data buffer. For devices operating in unbuffered mode, GOOD status is not returned until all requested data, filemarks, or setmarks are successfully recorded on the medium.

When issuing a buffered WRITE FILEMARKS command with the immediate bit set to one, GOOD status is returned as soon as the command is validated. A WRITE FILEMARKS command with the immediate bit set to zero causes any buffered blocks, filemarks, and setmarks to be written to the medium. Upon successful completion of this process, which is called a synchronize operation, no blocks, filemarks, or setmarks remain in the data buffer which have not been written to the medium. A synchronize operation has no effect on a data buffer which contains only read-ahead data or write data which has already been successfully written to the medium.

Should an unrecoverable write error occur while in buffered mode, the device generates an error condition to the current active command. If no command is active, the error may be reported on the next applicable operation as a deferred error (see 7.2.14.2). For some implementations, asynchronous event notification (AEN) or extended contingent allegiance (ECA) may be required. Refer to 6.5.5 and 6.7 for descriptions of AEN and ECA protocol.

The READ POSITION command may be used to determine the number and storage space of buffered blocks not written before the unrecoverable error was encountered.

A device with read-ahead data blocks in the data buffer does not report an unrecovered read error until the data block in error is requested by an initiator.

9.1.6. Recorded Object Descriptors (Block Identifiers)

Some recording formats specify that recorded objects (blocks, filemarks, and setmarks) have identifiers included in the recorded information to help determine write sequence and also to help detect device positioning errors. The identifier values are unique within a partition and may be unique within a volume.

The use of the term block identifier may imply some arithmetic sequence applied to the assignment of recorded objects. The block identifier assignment algorithm may be defined in an applicable format standard.

For some pre-formatted volumes, the identifiers are associated with physical blocks. In variable-length implementations, the identifier can be associated with a physical block when the logical block and the physical block have a one-to-one relationship on the medium.

Some formats may carry both physical and logical block identifiers recorded on the medium. When a logical block is split over more than one physical block, or multiple logical blocks are concatenated to form a physical block, the logical block identifier and the physical block identifier are not the same. Filemarks and setmarks may or may not have recorded identifiers, but if identifiers are used in the format, then each mark is assigned a value even if it is not explicitly recorded.

The READ POSITION and LOCATE commands use four-byte fields to hold these format dependent identifiers. For some implementations, this value may correspond to a real physical location; however, it is sufficient for the device to map the identifier to a value representing the unique recorded object. With this capability, the READ POSITION command may be used to report a device-defined block identifier and the initiator may use this value with a LOCATE command to position to the same location at some future time (provided the volume has not been rewritten in the interim).

9.1.7. Direction and Position Definitions

For sequential-access devices, positioning has the connotation of logically being in, at, before, or after some defined place within a volume. This definition means the position is capable of being repeated under the same circumstances. The orientation of usage for the four words (in, at, before, or after) is in one direction, from BOPx toward EOPx. All positioning defined below is worded from this perspective. Devices without buffers have some physical position which relates to these logical positions. However, these definitions do not require the medium to have a physical position equivalent to the logical position unless explicitly stated.

The forward direction is defined as logically progressing from BOPx toward EOPx. The reverse direction is defined as logically progressing from EOPx toward BOPx. In serpentine devices, the logical forward or reverse direction has an alternating relationship to the physical motion of the medium.

The concept of being "in" some position means not being outside a defined region. The definition allows the position to be on the boundary of a defined region. When a volume is first mounted, the logical position is always at the beginning of the default data partition (BOP0). Whenever a volume is mounted and the medium motion is stopped, the position is in some partition. While moving between partitions, there is no stable position.

The concept of being "at" some position indicates being positioned to a logical or physical extremity of a partition. A sequential-access device may be positioned at beginning-of-medium, at BOPx, at end-of-data (EOD), at EOPx, or at end-of-medium (EOM), since these are stable positions at extremities of a partition.

The concept of being "before" some position indicates that there is some element (data block, filemark, setmark, or other defined point) which may be encountered when moving toward EOPx, if the proper commands are issued. Being positioned before a particular data block means that if the device receives a valid READ command, the data block is transferred to the initiator. This position may also be before EWx and EOPx, since these are defined points within any partition. However, if data has not been written to the end-of-partition, these points may not be accessible by the initiator.

The concept of being "after" some position indicates that there is some element (data block, filemark, setmark, or other defined point) on the BOPx side of the current position which may be encountered if the proper commands are issued. When a READ command for a single data block has been successfully executed, the logical position is after the transferred data block.

9.1.8. Error Reporting

If any of the following conditions occur during the execution of a command the target shall return CHECK CONDITION status. The appropriate sense key and additional sense code should be set. The following table illustrates some error conditions and the applicable sense keys. The table does not provide an exhaustive enumeration of all conditions that may cause the CHECK CONDITION status.

Condition	Sense Key
Invalid logical block address	ILLEGAL REQUEST
Unsupported option requested	ILLEGAL REQUEST
Target reset or medium change since last command from this initiator	UNIT ATTENTION
Self diagnostic failed	HARDWARE ERROR
Unrecovered read error	MEDIUM ERROR HARDWARE ERROR
Recovered read or write error	RECOVERED ERROR
Overrun or other error that might be resolved by repeating the command	ABORTED COMMAND
Attempt a WRITE, READ, READ REVERSE, VERIFY, or RECOVER BUFFERED DATA command with the fixed bit set to zero and variable block mode is not supported	ILLEGAL REQUEST
Attempt a WRITE, READ, READ REVERSE, VERIFY, or RECOVER BUFFERED DATA command with a fixed bit of zero and requested block length is not supported	ILLEGAL REQUEST
Attempt a WRITE, READ, READ REVERSE, VERIFY, or RECOVER BUFFERED DATA command with a fixed bit of one and MODE SENSE block length set to zero	ILLEGAL REQUEST
Attempt to execute an erase or write-type operation on write protected medium	DATA PROTECT
Deferred write error	MEDIUM ERROR VOLUME OVERFLOW

In the case of an unrecovered read or write error, if the read-write error recovery page (9.3.3.4) is implemented, the current values specify the target error recovery criteria. If this page is not implemented, the error recovery is vendor specific.

In the case of an unrecovered read error, if the fixed bit is one, the valid bit shall be set to one and the information field shall be set to the requested transfer length minus the actual number of blocks read (not including the unrecovered block). If the fixed bit is zero, the valid bit shall be set to one and the information field shall be set to the requested transfer length. Upon termination, the logical position shall be after the unrecovered block.

In the case of an unrecovered write error, if unbuffered mode is selected and the fixed bit is set to one, the valid bit shall be set to one and the information field shall be set to the requested transfer length minus the actual number of blocks written. If unbuffered mode is selected and the fixed bit is set to zero, the information field shall be set to the requested transfer length.

In the case of an unrecovered write error or a deferred write error, if buffered mode is selected and the fixed bit is one, the valid bit shall be set to one and the information field shall be set to the total number of blocks, filemarks, and setmarks not written (the number of blocks not transferred from the initiator plus the number of blocks, filemarks, and setmarks remaining in the target's buffer). If buffered mode is selected and the fixed bit is zero, the information field shall be set to the total number of bytes, filemarks, and setmarks not written (the number of bytes not transferred from the initiator plus the number of bytes, filemarks, and setmarks remaining in the target's buffer). In both cases, the value in the information field may exceed the transfer length.

In the case of an unrecovered write error or a deferred write error, if buffered mode 1h is selected, the error shall be reported to the first initiator issuing a command (other than INQUIRY or REQUEST SENSE) or the first initiator responding to asynchronous event notification. If buffered mode 2h is selected, the error shall be reported to the initiator with unwritten data in the buffer.

9.2. Command Descriptions for Sequential-Access Devices

The commands for sequential-access devices shall be as shown in Table 9-1.

Table 9-1: Commands for Sequential-Access Devices

Command Name	Operation Code	Type	Section
CHANGE DEFINITION	40h	0	7.2.1
COMPARE	39h	0	7.2.2
COPY	18h	0	7.2.3
COPY AND VERIFY	3Ah	0	7.2.4
ERASE	19h	M	9.2.1
INQUIRY	12h	M	7.2.5
LOAD UNLOAD	1Bh	0	9.2.2
LOCATE	2Bh	0	9.2.3
LOG SELECT	4Ch	0	7.2.6
LOG SENSE	4Dh	0	7.2.7
MODE SELECT(6)	15h	M	7.2.8
MODE SELECT(10)	55h	0	7.2.9
MODE SENSE(6)	1Ah	M	7.2.10
MODE SENSE(10)	5Ah	0	7.2.11
PREVENT ALLOW MEDIUM REMOVAL	1Eh	0	8.2.4
READ	08h	M	9.2.4
READ BLOCK LIMITS	05h	M	9.2.5
READ BUFFER	3Ch	0	7.2.12
READ POSITION	34h	0	9.2.6
READ REVERSE	0Fh	0	9.2.7
RECEIVE DIAGNOSTIC RESULTS	1Ch	0	7.2.13
RECOVER BUFFERED DATA	14h	0	9.2.8
RELEASE UNIT	17h	M	9.2.9
REQUEST SENSE	03h	M	7.2.14
RESERVE UNIT	16h	M	9.2.10
REWIND	01h	M	9.2.11
SEND DIAGNOSTIC	10h	M	7.2.15
SPACE	11h	M	9.2.12
TEST UNIT READY	00h	M	7.2.16
VERIFY	13h	0	9.2.13
WRITE	0Ah	M	9.2.14
WRITE BUFFER	3Bh	0	7.2.17
WRITE FILEMARKS	10h	M	9.2.15

Key: M = Command implementation is mandatory.
0 = Command implementation is optional.

The following command codes are vendor specific: 02h, 06h, 07h, 09h, 0Ch, 0Dh, 0Eh. All other command codes for sequential-access devices are reserved for future standardization.

9.2.1. ERASE Command

Table 9-2: ERASE Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (19h)							
1	Logical Unit Number			Reserved			Immed	Long
2	Reserved							
3	Reserved							
4	Reserved							
5	Control							

The ERASE command (Table 9-2) causes part or all of the medium to be erased beginning at the current position on the logical unit. As used here, "erased" means either the medium shall be erased or a pattern shall be written on the medium that appears as a gap to the target.

An immediate (Immed) bit of zero indicates that the target shall not return status until the erase operation has completed. An Immed bit of one indicates that the target shall return status as soon as the command has been validated. If CHECK CONDITION status is returned for an ERASE command with an Immed bit of one, the erase operation shall not be performed.

A long bit of one indicates that all remaining medium in the current partition shall be erased beginning at the current logical position. If the Immed bit is one, the target shall return status as soon as all buffered commands have completed execution and the command descriptor block of the ERASE command has been validated. The logical position following an ERASE command with a long bit of one is not specified by this standard.

IMPLEMENTORS NOTE: Some targets may reject an ERASE command with the long bit set to one if the logical unit is not at beginning-of-partition.

A long bit of zero specifies an erase gap defined by the gap size field in the device configuration page (9.3.3.1). If the gap size is zero or the field is not supported, a device defined erase gap operation shall be performed. If the Immed bit is one, the target shall return status as soon as the command descriptor block has been validated. Erase gaps may be used in host controlled error recovery or "update in place" applications.

If the logical unit encounters early-warning during an ERASE command, and any buffered data, filemarks, or setmarks remain to be written, the target action shall be as defined for the early-warning condition of the WRITE command (9.2.14). If the long bit is zero, the erase operation shall terminate with CHECK CONDITION status and return sense data as defined for the WRITE command. Any buffered erases are not reported as part of the information field.

9.2.2. LOAD UNLOAD Command

Table 9-3: LOAD UNLOAD Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (1Bh)							
1	Logical Unit Number			Reserved				Immed
2	Reserved							
3	Reserved							
4	Reserved				EOT		Re-Ten	Load
5	Control							

The LOAD UNLOAD command (Table 9-3) requests that the target enable or disable the logical unit for further operations. This command may also be used to request a re-tension function. Prior to performing the load or unload operation, the target shall ensure that all buffered data, filemarks, and setmarks have been transferred to the medium.

An immediate (Immed) bit of zero indicates that the target shall not return status until the load or unload operation has completed. An Immed bit of one indicates that the target shall return status as soon as all buffered commands have completed execution and the command descriptor block of the LOAD UNLOAD command has been validated. If CHECK CONDITION status is returned for a LOAD UNLOAD command with an Immed bit of one, the load or unload operation shall not be performed.

IMPLEMENTORS NOTE: For compatibility with devices implemented prior to this version of the standard, a WRITE FILEMARKS command with an Immed bit of zero should be used to ensure that all buffered data, filemarks, or setmarks have been transferred to the medium prior to issuing a LOAD UNLOAD command with an Immed bit of one.

An end-of-tape (EOT) bit of one indicates that an unload operation (load bit set to zero) shall position the medium at end-of-medium for removal from the device. An EOT bit of zero indicates that an unload operation shall position the medium at beginning-of-medium for removal from the device.

An EOT bit of one and a load bit of one shall cause the target to return CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST in the sense data.

A re-tension (Re-Ten) bit of one indicates that the medium on the logical unit shall be correctly tensioned. Implementation of the re-tensioning function is device specific.

If the load bit is set to one, the medium in the logical unit shall be loaded and positioned to the beginning-of-partition zero. If the load bit is zero, the medium in the logical unit shall be positioned for removal at the extreme position along the medium specified by the EOT bit. Following successful completion of an unload operation, the target shall return CHECK CONDITION status with the sense key set to NOT READY for all subsequent medium-access commands until a new volume is mounted or a load operation is successfully completed.

When operating in buffered mode 1h or 2h (see 9.3.3), the target shall discard any unwritten buffered data after the LOAD UNLOAD command is validated if the previous command was terminated with CHECK CONDITION status and the device is unable to continue successfully writing.

9.2.3. LOCATE Command

Table 9-4: LOCATE Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (2Bh)							
1	Logical Unit Number			Reserved		BT	CP	Immed
2	Reserved							
3	(MSB)							
4								
5	Block Address							
6								
7	(LSB)							
8	Reserved							
9	Partition							
	Control							

LOCATE command (Table 9-4) causes the target to position the logical unit to the specified block address in a specified partition. Upon completion, the logical position shall be before the specified location. Prior to performing the locate operation, the target shall ensure that all buffered data, filemarks, and setmarks have been transferred to the medium.

A block address type (BT) bit of one indicates the value in the block address field shall be interpreted as a device-specific value. A BT bit of zero indicates the value in the block address field shall be interpreted as an SCSI logical block address.

A change partition (CP) bit of one indicates that a change to the partition specified in the partition field is to occur prior to positioning to the block specified in the block address field. A CP bit of zero indicates no partition change is to be made and the partition field is to be ignored.

An immediate (Immed) bit of zero indicates that the target shall not return status until the locate operation has completed. An Immed bit of one indicates that the target shall return status as soon as all buffered commands have completed execution and the command descriptor block of the LOCATE command has been validated. If CHECK CONDITION status is returned for a LOCATE command with an Immed bit of one, the locate operation shall not be performed.

The block address field specifies the block address to which the target shall position the medium based on the current setting of the BT bit.

The partition field specifies which partition to select if the CP bit is one. Refer to the sequential-access device model (9.1.3) and the medium partition pages (9.3.3.2 and 9.3.3.3) for additional information about partitioning.

9.2.4. READ Command

Table 9-5: READ Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (08h)							
1	Logical Unit Number			Reserved			SILI	Fixed
2	(MSB)							
3	Transfer Length							
4								
5	Control							

The READ command (Table 9-6) requests that the target transfer one or more block(s) of data to the initiator beginning with the next block on the logical unit.

The fixed bit specifies whether fixed-length or variable-length blocks are to be transferred. Refer to the READ BLOCK LIMITS command (9.2.5) for additional information about fixed and variable block mode.

If the fixed bit is one, the transfer length specifies the number of fixed-length blocks to be transferred, using the current block length reported in the mode parameters block descriptor (7.3.3). If the fixed bit is zero, a variable-length block is requested with the transfer length specifying the maximum number of bytes allocated for the returned data.

A successful READ command with a fixed bit of one shall transfer the requested transfer length times the current block length in bytes to the initiator. A successful READ command with a fixed bit of zero shall transfer the requested transfer length in bytes to the initiator. Upon completion, the logical position shall be after the last block transferred (end-of-partition side).

If the suppress incorrect length indicator (SILI) bit is one and the fixed bit is zero, the target shall:

(1) report CHECK CONDITION status for an incorrect length condition only if the overlength condition exists (see 9.4) and the block length field in the mode parameter block descriptor is nonzero (see 7.3.3)

(2) not report CHECK CONDITION status if the only error is the underlength condition (see 9.4), or if the only error is the overlength condition and the block length field of the mode parameters block descriptor is zero.

IMPLEMENTORS NOTE: Since the residue information normally provided in the information field of the sense data may not be available when the SILI bit is set, other methods for determining the actual block length should be used (e.g., including length information in the data block).

If the SILI bit is one and the fixed bit is one, the target shall terminate the command with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST with an additional sense code of INVALID FIELD IN CDB.

If the SILI bit is zero and an incorrect length block is read, CHECK CONDITION status shall be returned and the ILI and valid bits shall be set to one in the sense data. Upon termination, the logical position shall be after the incorrect length block (end-of-partition side). If the fixed bit is one, the information field shall be set to the requested transfer length minus the actual number of blocks read (not including the incorrect length block). If the fixed bit is zero, the information field shall be set to the requested transfer length minus the actual block length. Targets that do not support negative values shall set the information field to zero if the overlength condition exists (see 9.4).

IMPLEMENTORS NOTE: In the above case with the fixed bit of one, only the position of the incorrect-length logical block can be determined from the sense data. The actual length of the incorrect block is not reported. Other means may be used to determine its actual length (e.g., read it again with the fixed bit set to zero).

A transfer length of zero indicates that no data shall be transferred. This condition shall not be considered an error and the logical position shall not be changed.

If the logical unit encounters a filemark during a READ command, CHECK CONDITION status shall be returned and the filemark and valid bits shall be set to one in the sense data. The sense key shall be set to NO SENSE or RECOVERED ERROR, as appropriate. Upon termination, the logical position shall be after the filemark (end-of-partition side). If the fixed bit is one, the information field shall be set to the requested transfer length minus the actual number of blocks read (not including the filemark). If the fixed bit is zero, the information field shall be set to the requested transfer length.

If the logical unit encounters a setmark during a READ command and the RSmk bit is set to one in the device configuration page (9.3.3.1), CHECK CONDITION status shall be returned and the filemark and valid bits shall be set to one in the sense data. The sense key shall be set to NO SENSE or RECOVERED ERROR, as appropriate, and the additional sense code shall be set to SETMARK DETECTED. Upon termination, the logical position shall be after the setmark (end-of-partition side). If the fixed bit is one, the information field shall be set to the requested transfer length minus the actual number of blocks read (not including the setmark). If the fixed bit is zero, the information field shall be set to the requested transfer length. The target shall not return CHECK CONDITION when a setmark is encountered if the RSmk bit is set to zero or if this option is not supported.

If the logical unit encounters early-warning during a READ command and the REW bit is set to one in the device configuration page (see 9.3.3.1), CHECK CONDITION status shall be returned upon completion of the current block. The sense key shall be set to NO SENSE or RECOVERED ERROR, as appropriate. The EOM and valid bits shall be set to one in the sense data. Upon termination, the logical position shall be after the last block transferred (end-of-partition side). If the fixed bit is one, the information field shall be set to the requested transfer length minus the actual number of blocks read. If the fixed bit is zero, the information field shall be set to the requested transfer length minus the actual block length. The target shall not return CHECK CONDITION status when early-warning is encountered if the REW bit is zero or if the REW option is not supported.

IMPLEMENTORS NOTE: A REW bit of one is not recommended for most system applications since read data may be present after early-warning.

If the logical unit encounters end-of-data during a READ command, CHECK CONDITION status shall be returned, the sense key shall be set to BLANK CHECK, and the valid bit shall be set to one in the sense data. If end-of-data is encountered at or after early-warning, the EOM bit shall also be set to one. Upon termination, the logical position shall be after the last recorded logical block (end-of-partition side). If the fixed bit is one, the information field shall be set to the requested transfer length minus the actual number of blocks read. If the fixed bit is zero, the information field shall be set to the requested transfer length.

If the logical unit encounters end-of-partition during a READ command, CHECK CONDITION status shall be returned, the sense key shall be set to MEDIUM ERROR, and the EOM and valid bits shall be set to one in the sense data. The medium position following this condition is not defined. If the fixed bit is one, the information field shall be set to the requested transfer length minus the actual number of blocks read. If the fixed bit is zero, the information field shall be set to the requested transfer length.

9.2.5. READ BLOCK LIMITS Command

Table 9-6: READ BLOCK LIMITS Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (05h)							
1	Logical Unit Number			Reserved				
2	Reserved							
3	Reserved							
4	Reserved							
5	Control							

The READ BLOCK LIMITS command (Table 9-6) requests that the logical units block length limits capability be returned. The READ BLOCK LIMITS data shown in Table 9-7 shall be returned during the DATA IN phase of the command.

Table 9-7: READ BLOCK LIMITS Data

Bit Byte	7	6	5	4	3	2	1	0
0	Reserved							
1	(MSB)							
2	Maximum Block Length Limit							
3	(LSB)							
4	(MSB)							
5	Minimum Block Length Limit							
	(LSB)							

If the maximum block length limit equals the minimum block length limit, only fixed-length blocks are supported. The size of the fixed-length blocks is maximum block length limit bytes in length.

If the maximum block length limit is not equal to the minimum block length limit, variable-length blocks are supported. If the maximum block length limit is zero, then a maximum block length is not specified.

9.2.6. READ POSITION Command

Table 9-8: READ POSITION Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (34h)							
1	Logical Unit Number			Reserved				BT
2	Reserved							
3	Reserved							
4	Reserved							
5	Reserved							
6	Reserved							
7	Reserved							
8	Reserved							
9	Control							

The READ POSITION command (Table 9-8) reports the current position of the logical unit and any data blocks in the buffer. No medium movement shall occur as a result of the command.

A block address type (BT) bit of one requests the target to return its current first block location and last block location as a device-specific value. A BT bit of zero requests the target to return the first block location and the last block location as an SCSI logical block address (relative to a partition).

If the logical unit is not capable of reporting block locations, the BPU bit shall be set to one in the returned data. Support for this option of the READ POSITION command is indicated by a BIS bit set to one in the device configuration page (9.3.3.1).

The READ POSITION data shown in Table 9-9 shall be returned during the DATA IN phase of the command.

Table 9-9: READ POSITION Data Format

Bit Byte	7	6	5	4	3	2	1	0
0	BOP	EOP	Reserved			BPU	Reserved	
1	Partition Number							
2	Reserved							
3	Reserved							
4	(MSB)	First Block Location						(LSB)
7	- -							
8	(MSB)	Last Block Location						(LSB)
11	- -							
12	Reserved							
13	(MSB)	Number of Blocks in Buffer						(LSB)
15	- -							
16	(MSB)	Number of Bytes in Buffer						(LSB)
19	- -							

A beginning of partition (BOP) bit of one indicates that the logical unit is at the beginning-of-partition in the current partition. A BOP bit of zero indicates that the current logical position is not at the beginning-of-partition.

An end of partition (EOP) bit of one indicates that the logical unit is positioned between early-warning and end-of-partition in the current partition. An EOP bit of zero indicates that the current logical position is not between early-warning and end-of-partition.

IMPLEMENTORS NOTE: The BOP and EOP indications are not necessarily a result of a physical tape marker (e.g., reflective marker).

A block position unknown (BPU) bit of one indicates that the first and last block locations are not known or cannot be obtained. A BPU bit of zero indicates that the first and last block location fields contain valid position information.

The partition number field reports the partition number for the current logical position. If the logical unit only supports one partition for the medium, this field shall be set to zero.

The first block location field indicates the block address associated with the current logical position. The value shall indicate the block address of the next data block to be transferred between the initiator and the target if a READ or WRITE command is issued.

The last block location field indicates the block address associated with the next block to be transferred from the buffer to the medium. The value shall indicate the block address of the next data block to be transferred between the buffer and the medium. If the buffer does not contain a whole block of data or is empty, the value reported for the last block location shall be equal to the value reported for the first block location.

IMPLEMENTORS NOTE: The information provided by the first and last block location fields can be used in conjunction with the LOCATE command to position the medium at the appropriate logical block on another device in the case of unrecoverable errors on the first device.

The number of blocks in buffer field indicates the number of data blocks in the target's buffer that have not been written to the medium.

The number of bytes in buffer field indicates the total number of data bytes in the target's buffer that have not been written to the medium.

9.2.7. READ REVERSE Command

Table 9-10: READ REVERSE Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (0Fh)							
1	Logical Unit Number			Reserved			SILI	Fixed
2	(MSB)							
3	Transfer Length							
4								
5	(LSB)							
	Control							

The READ REVERSE command (Table 9-10) requests that the target transfer one or more block(s) of data to the initiator beginning at the current position on the logical unit.

The execution of this command is similar to the READ command except that medium motion is in the reverse direction. All block(s), and the byte(s) within the block(s), are transferred in the reverse order. The order of bits within each byte shall not be changed. Upon completion of a READ REVERSE command, the logical position shall be before the last block transferred (beginning-of-partition side).

Refer to the READ command (9.2.4) for a description of the fixed bit, the SILI bit, the transfer length field, and any associated error conditions.

Filemarks, setmarks, incorrect length blocks, and unrecovered read errors are handled the same as in the READ command, except that upon termination the logical position shall be before the filemark, setmark, incorrect length block, or unrecovered block (beginning-of-partition side).

If the logical unit encounters beginning-of-partition during a READ REVERSE command, CHECK CONDITION status shall be returned and the EOM and valid bits shall be set to one in the sense data. The sense key shall be set to NO SENSE or RECOVERED ERROR, as appropriate. If the fixed bit is one, the information field shall be set to the requested transfer length minus the actual number of blocks transferred. If the fixed bit is zero, the information field shall be set to the requested transfer length.

9.2.8. RECOVER BUFFERED DATA Command

Table 9-11: RECOVER BUFFERED DATA Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (14h)							
1	Logical Unit Number			Reserved			SILI	Fixed
2	(MSB)							
3	Transfer Length							
4								
5	(LSB)							
	Control							

The RECOVER BUFFERED DATA command (Table 9-11) is used to recover data that has been transferred to the target's buffer but has not been successfully written to the medium. It is normally used to recover from error or exception conditions that make it impossible to write the buffered data to the medium. One or more RECOVER BUFFERED DATA commands may be required to recover all unwritten buffered data.

The execution of this command is similar to the READ command except that the data is transferred from the target's buffer instead of the medium. The order in which block(s) are transferred is defined by the RBO bit in the device configuration page (9.3.3.1). If the RBO bit is not implemented, block(s) are transferred in the same order they would have been transferred to the medium.

Refer to the READ command (9.2.4) for a description of the fixed bit, the SILI bit, the transfer length field, and any associated error conditions.

If the fixed bit is zero, no more than the requested transfer length shall be transferred to the initiator. If the requested transfer length is smaller than the actual length of the logical block to be recovered, only the requested transfer length shall be transferred to the initiator and the remaining data for the current logical block shall be discarded.

IMPLEMENTORS NOTE: During recovery operations involving unknown block sizes, the initiator should select the maximum block length supported by the target to ensure that all buffered data will be transferred.

If a buffered filemark is encountered during a RECOVER BUFFERED DATA command, CHECK CONDITION status shall be returned, the sense key shall be set to NO SENSE, and the filemark and valid bits shall be set to one in the sense data. Upon termination, the logical position shall be after the filemark. If the fixed bit is one, the information field shall be set to the requested transfer length minus the actual number of blocks transferred (not including the filemark). If the fixed bit is zero, the information field shall be set to the requested transfer length.

If a buffered setmark is encountered during a RECOVER BUFFERED DATA command and the RSmk bit is set to one in the device configuration page (9.3.3.1), CHECK CONDITION status shall be returned and the filemark and valid bits shall be set to one in the sense data. The sense key shall be set to NO SENSE and the additional sense code shall be set to SETMARK DETECTED. Upon termination, the logical position shall be after the setmark. If the fixed bit is one, the information field shall be set to the requested transfer length minus the actual number of blocks transferred (not including the setmark). If the fixed bit is zero, the information field shall be set to the requested transfer length. The target shall not return CHECK CONDITION when a setmark is encountered if the RSmk bit is zero or if this option is not supported.

If an attempt is made to recover more logical blocks of data than are contained in the target's buffer, CHECK CONDITION status shall be returned, the sense key shall be set to NO SENSE, and the EOM and valid bits shall be set to one in the sense data. If the fixed bit is one, the information field shall be set to the requested transfer length minus the actual number of blocks transferred. If the fixed bit is zero, the information field shall be set to the requested transfer length.

9.2.9. RELEASE UNIT Command

Table 9-12: RELEASE UNIT Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (17h)							
1	Logical Unit Number			3rdPty	Third Party Device ID			Reserved
2	Reserved							
3	Reserved							
4	Reserved							
5	Control							

The RESERVE UNIT and RELEASE UNIT commands provide the basic mechanism for contention resolution in multiple-initiator systems. The RELEASE UNIT command (Table 9-12) is used to release previously reserved logical units for the requesting initiator, or if it is a third-party release, to another specified SCSI device.

If a valid reservation exists for the I_I_L nexus, the target shall release the reservation and return GOOD status.

It is not an error to attempt to release a reservation that is not currently valid. In this case, the target shall return GOOD status without altering any other reservation.

9.2.9.1. Third-Party Release

Third-party release allows an initiator to release a logical unit that was previously reserved using a third-party reservation (see 9.2.10.1).

If the third party (3rdPty) bit is zero, then a third-party release is not requested. If the 3rdPty bit is one, and if the reservation was made using a third-party reservation by the initiator that is requesting the release for the same SCSI device as specified in the third-party device ID, then the target shall release the reservation.

If the 3rdPty bit is one, the target shall not modify the mode parameters for commands received from the third-party device even if the target implements the transfer of mode parameters with a third-party RESERVE UNIT command.

IMPLEMENTORS NOTE: When a target implements independent storage of mode parameters for each initiator, a third-party RESERVE UNIT command effects a transfer of the current mode parameters. Those set up by the initiator of the RESERVE UNIT are to be set as the mode parameters used for commands from the third-party device (usually a copy master device). A unit attention condition notifies the third-party device of the changed mode parameters. A successful third-party RELEASE UNIT command leaves the transferred parameters intact. The third-party device can issue MODE SENSE and MODE SELECT commands to query and modify the mode parameters.

9.2.10. RESERVE UNIT Command

Table 9-13: RESERVE UNIT Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (16h)							
1	Logical Unit Number			3rdPty	Third Party Device ID			Reserved
2	Reserved							
3	Reserved							
4	Reserved							
5	Control							

The RESERVE UNIT and RELEASE UNIT commands provide the basic mechanism for contention resolution in multiple-initiator systems. The RESERVE UNIT command (Table 9-13) is used to reserve logical units for the exclusive use of the requesting initiator, or if it is a third-party reservation, to another specified SCSI device.

This command requests that the entire logical unit be reserved for the exclusive use of the initiator until the reservation is superseded by another valid RESERVE UNIT command from the initiator that made the reservation or until released by a RELEASE UNIT command from the same initiator that made the reservation, by a BUS DEVICE RESET message from any initiator, by a hard reset condition, or by a power on cycle. The reservation shall not be granted if the logical unit is reserved by another initiator. It shall be permissible for an initiator to reserve a logical unit that is currently reserved by that initiator.

If the logical unit is reserved for another initiator, the target shall return RESERVATION CONFLICT status.

If, after honoring the reservation, any other initiator attempts to perform any command on the reserved logical unit other than an INQUIRY command, a REQUEST SENSE command, or a RELEASE UNIT command, which shall be ignored, then the command shall be rejected with RESERVATION CONFLICT status.

9.2.10.1. Third-Party Reservation

Third-party reservation allows an initiator to reserve a logical unit for another SCSI device. This is intended for use in multiple-initiator systems that use the COPY command.

If the third-party (3rdPty) bit is zero, a third-party reservation is not requested. If the 3rdPty bit is one the target shall reserve the logical unit for the SCSI device specified in the third-party device ID field. The target shall preserve the reservation until it is superseded by another valid RESERVE UNIT command from the initiator that made the reservation or until it is released by the same initiator, by a BUS DEVICE RESET message from any initiator, or a hard reset condition. The target shall ignore any attempt to release the reservation made by any other initiator.

If independent sets of parameters are implemented, a third party reservation shall cause the target to transfer the set of parameters in effect for the initiator of the RESERVE command to the parameters used for commands from the third party device. Any subsequent command issued by the third-party device is executed according to the mode parameters in effect for the initiator that sent the RESERVE command.

If independent sets of parameters are implemented, a third party reservation shall cause the target to transfer the set of parameters in effect for the initiator of the RESERVE command to the parameters used for commands from the third party device. Any subsequent command issued by the third-party device is executed according to the mode parameters in effect for the initiator that sent the RESERVE command.

IMPLEMENTORS NOTE: This transfer of the mode parameters is applicable to target devices which store mode information independently for different initiators. This mechanism allows an initiator to set the mode parameters of a target for the use of a copy master (i.e., the third-party device). The third-party copy master may subsequently issue a MODE SELECT command to modify the mode parameters.

9.2.10.2. Superseding Reservations

An initiator that currently has a logical unit reserved may modify the current reservation by issuing another RESERVE UNIT command to the same logical unit. The superseding reservation shall release the current reservation if the superseding reservation request is granted. The current reservation shall not be modified if the superseding reservation request cannot be granted. If the superseding reservation cannot be granted because of conflicts with a previous reservation (other than the current reservation), then the target shall return RESERVATION CONFLICT status.

IMPLEMENTORS NOTE: Superseding reservations allow the SCSI device ID to be changed on a reservation using the third-party reservation option. This capability is necessary for certain situations when using COMPARE, COPY, and COPY AND VERIFY commands.

9.2.11. REWIND Command

Table 9-14: REWIND Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (01h)							
1	Logical Unit Number			Reserved				Immed
2	Reserved							
3	Reserved							
4	Reserved							
5	Control							

The REWIND command (Table 9-14) causes the target to position the logical unit to the beginning-of-partition in the current partition. Prior to performing the rewind operation, the target shall ensure that all buffered data, filemarks, and setmarks have been transferred to the medium.

An immediate (Immed) bit of zero indicates that the target shall not return status until the rewind operation has completed. An Immed bit of one indicates that the target shall return status as soon as all buffered commands have completed execution and the command descriptor block of the REWIND command has been validated. If CHECK CONDITION status is returned for a REWIND command with an Immed bit of one, the rewind operation shall not be performed.

IMPLEMENTORS NOTE: For compatibility with devices implemented prior to this standard, it is suggested that a WRITE FILEMARKS command with an Immed bit of zero be used to ensure that all buffered data, filemarks, or setmarks have been transferred to the medium before issuing a REWIND command with an Immed bit of one.

When operating in buffered mode 1h or 2h (see 9.3.3), the target shall discard any unwritten buffered data after the REWIND command is validated if the previous command was terminated with CHECK CONDITION status and the device is unable to continue successfully writing.

9.2.12. SPACE Command

Table 9-15: SPACE Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation (11h)							
1	Logical Unit Number			Reserved		Code		
2	(MSB)							
3	Count							
4	(LSB)							
5	Control							

The SPACE command (Table 9-15) provides a variety of positioning functions that are determined by the code and count. Both forward and reverse positioning are provided, although some targets may only support a subset of this command. If an initiator requests an unsupported function, the command shall be terminated with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST.

The code is defined as follows:

Code	Description	Support
000b	Blocks	Mandatory
001b	Filemarks	Mandatory
010b	Sequential Filemarks	Optional
011b	End-of-Data	Optional
100b	Setmarks	Optional
101b	Sequential Setmarks	Optional
110b-111b	Reserved	

When spacing over blocks, filemarks, or setmarks, the count field specifies the number of blocks, filemarks, or setmarks to be spaced over in the current partition. A positive value N in the count field shall cause forward positioning (toward end-of-partition) over N blocks, filemarks, or setmarks ending on the end-of-partition side of the last block, filemark, or setmark. A zero value in the count field shall cause no change of logical position. A negative value -N (two's complement notation) in the count field shall cause reverse positioning (toward beginning-of-partition) over N blocks, filemarks, or setmarks ending on the beginning-of-partition side of the last block, filemark, or setmark. Support of spacing in the reverse direction is optional.

If a filemark is encountered while spacing over blocks, the command shall be terminated. The logical position shall be on the end-of-partition side of the filemark if movement was in the forward direction and on the beginning-of-partition side of the filemark if movement was in the reverse direction. CHECK CONDITION status shall be returned to the initiator, the sense key shall be set to NO SENSE, and the filemark and valid bits shall be set to one in the sense data. The information field shall be set to the requested count minus the actual number of blocks spaced over (not including the filemark).

If a setmark is encountered while spacing over blocks or filemarks and the RSmk bit is set to one in the device configuration page (9.3.3.1), the command shall be terminated, CHECK CONDITION status shall be returned to the initiator, and the filemark and valid bits shall be set to one in the sense data. The sense key shall be set to NO SENSE and the additional sense code shall be set to SETMARK DETECTED. The information field shall be set to the requested count minus the actual number of blocks or filemarks spaced over (not including the setmark). The logical position shall be on the end-of-partition side of the setmark if movement was in the forward direction and on the beginning-of-partition side of the setmark if movement was in the reverse direction. The target shall not return CHECK CONDITION status when a setmark is encountered if the RSmk bit is set to zero or if this option is not supported.

If early-warning is encountered while spacing over blocks, filemarks, or setmarks and the REW bit is set to one in the device configuration page (9.3.3.1), CHECK CONDITION status shall be returned, the sense key shall be set to NO SENSE, and the EOM and valid bits shall be set to one in the sense data. The information field shall be set to the requested count minus the actual number of blocks, filemarks, or setmarks spaced over as defined by the code value. If the REW bit is zero or the option is not supported by the target, the target shall not report CHECK CONDITION status at the early-warning point.

IMPLEMENTORS NOTE: Setting the REW bit is not recommended for most system applications since data may be present after early-warning.

If end-of-data is encountered while spacing over blocks, filemarks, or setmarks, CHECK CONDITION status shall be returned, the sense key shall be set to BLANK CHECK, and the valid bit shall be set to one in the sense data. Additionally, the EOM bit shall be set to one if end-of-data is encountered at or after early-warning. The information field shall be set to the requested count minus the actual number of blocks, filemarks, or setmarks spaced over as defined by the code value.

If the end-of-partition is encountered while spacing forward over blocks, filemarks, or setmarks, CHECK CONDITION status shall be returned, the sense key shall be set to MEDIUM ERROR, the EOM bit shall be set to one, and the valid bit shall be set to one. The information field shall be set to the requested count minus the actual number of blocks, filemarks, or setmarks spaced over as defined by the code value.

If beginning-of-partition is encountered while spacing over blocks, filemarks, or setmarks in the reverse direction, the target shall return CHECK CONDITION status and shall set the sense key to NO SENSE. The EOM and valid bits shall be set to one, and the information field set to the total number of blocks, filemarks, or setmarks not spaced over (the requested number of blocks, filemarks, or setmarks minus the actual number of blocks, filemarks, or setmarks spaced over).

When spacing over sequential filemarks (or setmarks), the count field is interpreted as follows:

- (1) A positive value N shall cause forward movement to the first occurrence of N or more consecutive filemarks (or setmarks) being logically positioned after the Nth filemark (or setmark).
- (2) A zero value shall cause no change in the logical position.
- (3) A negative value -N (2's complement notation) shall cause reverse movement to the first occurrence of N or more consecutive filemarks (or setmarks) being logically positioned on the beginning-of-partition side of the Nth filemark (or setmark).

If a setmark is encountered while spacing to sequential filemarks and the RSmk bit is set to one in the device configuration page (9.3.3.1), CHECK CONDITION status shall be returned, the filemark bit shall be set to one, and the valid bit shall be set to zero. The sense key shall be set to NO SENSE and the additional sense code shall be set to SETMARK DETECTED. The target shall not return CHECK CONDITION status when a setmark is encountered if the RSmk bit is set to zero or if this option is not supported.

If end-of-partition is encountered while spacing to sequential filemarks or setmarks, CHECK CONDITION status shall be returned, the sense key shall be set to MEDIUM ERROR, the EOM bit shall be set to one, and the valid bit shall be set to zero.

If end-of-data is encountered while spacing to sequential filemarks or setmarks, CHECK CONDITION status shall be returned, the sense key shall be set to BLANK CHECK, and the valid bit shall be set to zero. Additionally, the EOM bit shall be set to one if end-of-data is encountered at or after early-warning.

When spacing to end-of-data, the count field is ignored. Upon successful completion, the medium shall be positioned such that a subsequent write operation would append to the last logically recorded information.

If end-of-partition is encountered while spacing to end-of-data, CHECK CONDITION status shall be returned, the sense key shall be set to MEDIUM ERROR, the EOM bit shall be set to one, and the valid bit shall be set to zero.

9.2.13. VERIFY Command

Table 9-16: VERIFY Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (13h)							
1	Logical Unit Number			Reserved		Immed	BytCmp	Fixed
2	(MSB)							
3	Verification Length							---
4								---
	(LSB)							
5	Control							

The VERIFY command (Table 9-16) requests that the target verify one or more block(s) beginning with the next block on the logical unit.

An immediate (Immed) bit of zero indicates that the command shall not return status until the verify operation has completed. An Immed bit of one indicates that status shall be returned as soon as the command descriptor block has been validated (but after the data transfer from the initiator, if the BytCmp bit is one).

IMPLEMENTORS NOTE: In order to ensure that no errors are lost, the initiator should set the Immed bit to zero on the last VERIFY command when issuing a series of VERIFY commands.

A byte compare (BytCmp) bit of zero indicates that the verification shall be simply a medium verification (e.g., CRC, ECC). No data shall be transferred between the initiator and target.

A BytCmp bit of one indicates that the target shall perform a byte-by-byte compare of the data on the medium and the data transferred from the initiator during the DATA OUT phase of the VERIFY command. Data shall be transferred from the initiator to the target as in a WRITE command. If the BytCmp bit is one and the byte compare option is not supported, the target shall terminate the command with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST.

The verification length specifies the amount of data to verify, in blocks or bytes, as indicated by the fixed bit. Refer to the READ command (9.2.4) for a description of the fixed bit and any associated error conditions. When the verification length is zero, no data shall be verified and the current logical position shall not be changed. This condition shall not be considered as an error.

The VERIFY command shall terminate when the verification length has been satisfied, when an incorrect length block is encountered, when a filemark is encountered, when a setmark is encountered (if the RSmk bit is one in the device configuration page, 9.3.3.1), when end-of-data is encountered, when the end-of-partition is encountered, when early-warning is encountered (if the REW bit is one in the device configuration page, 9.3.3.1), or when an unrecoverable read error is encountered. The status and sense data for each of these conditions are handled in the same manner as in the READ command (9.2.4). Upon successful completion of a VERIFY command, the logical position shall be after the last block verified.

If the data does not compare (BytCmp bit of one), the command shall terminate with CHECK CONDITION status, the valid bit shall be set to one, and the sense key shall be set to MISCOMPARE. If the fixed bit is one, the information field shall be set to the requested verification length minus the actual number of blocks successfully verified. If the fixed bit is zero, the information field shall be set to the requested verification length minus the actual number of bytes successfully verified. This number may be larger than the requested verification length if the error occurred on a previous VERIFY command with an Immed bit of one. Upon termination, the medium shall be positioned after the block containing the miscompare (end-of-partition side).

9.2.14. WRITE Command

Table 9-17: WRITE Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (0Ah)							
1	Logical Unit Number			Reserved				Fixed
2	(MSB)							
3	Transfer Length							
4								(LSB)
5	Control							

The WRITE command (Table 9-17) requests that the target write the data that is transferred from the initiator to the current position on the logical unit.

The fixed bit specifies whether fixed-length or variable-length blocks are to be transferred. See the READ BLOCK LIMITS command (9.2.5) for additional information about fixed and variable block mode.

If the fixed bit is one, the transfer length specifies the number of fixed-length blocks to be transferred, using the current block length reported in the mode parameter block descriptor (7.3.3). If the fixed bit is zero, a single block is transferred with the transfer length specifying the block length in bytes.

If the transfer length is zero, no data shall be transferred and the current position on the logical unit shall not be changed. This condition shall not be considered an error.

A WRITE command may be buffered or unbuffered, as indicated by the buffered mode field of the mode parameter header (9.3.3). For unbuffered operation (buffered mode 0h), the target shall not return GOOD status until all data block(s) are successfully written to the medium. For buffered operation (buffered mode 1h or 2h), the target may return GOOD status as soon as all data block(s) are successfully transferred to the target's buffer.

IMPLEMENTORS NOTE: For compatibility with devices implemented prior to this version of the standard, a WRITE FILEMARKS command with the Immed bit set to zero should be issued when completing a buffered write operation to ensure that all buffered data, filemarks, and setmarks are written to the medium.

If the logical unit encounters early-warning during a WRITE command, an attempt to finish writing any data may be made, as determined by the current settings of the REW and SEW bits in the device configuration page (9.3.3.1). The command shall terminate with CHECK CONDITION status and the EOM and valid bits shall be set to one in the sense data. If all data that is to be written is successfully transferred to the medium, the sense key shall be set to NO SENSE or RECOVERED ERROR, as appropriate. If any data that is to be written cannot be transferred to the medium when early-warning is encountered, the sense key shall be set to VOLUME OVERFLOW.

The information field shall be defined as follows:

(1) If unbuffered mode is reported in the mode parameter header and the fixed bit is set to one, the information field shall be set to the requested transfer length minus the actual number of blocks written.

(2) If unbuffered mode is reported and the fixed bit is set to zero, the information field shall be set to the requested transfer length.

(3) If buffered mode is reported in the mode parameter header and the fixed bit is set to one, the information field shall be set to the total number of blocks, filemarks, and setmarks not written (the number of blocks not transferred from the initiator plus the number of blocks, filemarks, and setmarks remaining in the target's buffer). Note that the value in the information field may exceed the transfer length.

(4) If buffered mode is reported and the fixed bit is set to zero, the information field shall be set to the total number of bytes, filemarks, and setmarks not written (the number of bytes not transferred from the initiator plus the number of bytes, filemarks, and setmarks remaining in the target's buffer). Note that the value in the information field may exceed the transfer length.

IMPLEMENTORS NOTE: The target should ensure that some additional data can be written to the medium (e.g., labels, filemarks, or setmarks) after the first early-warning indication has been returned to the initiator.

If a WRITE command is received while the logical unit is positioned between early-warning and end-of-partition, the target shall return CHECK CONDITION status after attempting to perform the command. The EOM and valid bits shall be set to one in the sense data. If all data that is to be written is successfully transferred to the medium, the information field shall be set to zero. If any data that is to be written is not transferred to the medium prior to encountering end-of-partition, the sense key shall be set to VOLUME OVERFLOW and the information field shall be defined as follows:

(1) If the fixed bit is one, the information field shall be set to the requested transfer length minus the actual number of blocks written to the medium.

(2) If the fixed bit is zero, the information field shall be set to the requested transfer length.

9.2.15. WRITE FILEMARKS Command

Table 9-18: WRITE FILEMARKS Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (10h)							
1	Logical Unit Number			Reserved			WSmk	Immed
2	(MSB)							
3	Transfer Length							
4								
5	Control							

The WRITE FILEMARKS command (Table 9-18) requests that the target write the specified number of filemarks or setmarks to the current position on the logical unit.

If the Write Setmark (WSmk) bit is one, the transfer length specifies the number of setmarks to be written. If the WSmk bit is zero, the transfer length specifies the number of filemarks to be written.

An immediate (Immed) bit of one indicates that the target shall return status as soon as the command descriptor block has been validated. An Immed bit of one is only valid if buffered mode is reported in the mode parameter header (9.3.3).

An Immed bit of zero indicates that the target shall not return status until the write operation has completed. Any buffered data, filemarks, and setmarks shall be written to the medium prior to completing the command.

IMPLEMENTORS NOTE: Upon completion of any buffered write operation, the initiator can issue a WRITE FILEMARKS command with the Immed bit set to zero and the transfer length field set to zero to ensure that all buffered data, filemarks, and setmarks are successfully written to the medium.

If the logical unit encounters early-warning during a WRITE FILEMARKS command, an attempt to finish writing any buffered data, filemarks, or setmarks may be made, as determined by the current settings of the REW and SEW bits in the device configuration page (9.3.3.1). The command shall terminate with CHECK CONDITION status and the EOM and valid bits shall be set to one in the sense data. If all buffered data, filemarks, and setmarks are successfully transferred to the medium, the sense key shall be set to NO SENSE or RECOVERED ERROR, as appropriate. If any buffered data, filemarks, or setmarks to be written are not transferred to the medium when early-warning is encountered, the sense key shall be set to VOLUME OVERFLOW.

The information field shall be defined as follows:

- (1) If unbuffered mode is reported in the mode parameter header, the information field shall be set to the requested transfer length minus the actual number of filemarks or setmarks written.
- (2) If buffered mode is reported in the mode parameter header and the buffered data was written in variable block mode (see 9.5.14), the information field shall be set to the total number of bytes, filemarks, and setmarks not written (the number of filemarks or setmarks not transferred from the initiator plus the number of bytes, filemarks and setmarks remaining in the target's buffer). Note that it is possible for the value in the information field to exceed the transfer length.
- (3) If buffered mode is reported and the buffered data was written in fixed - block mode (see 9.5.14), the information field shall be set to the total number of blocks, filemarks, and setmarks not written (the number filemarks or setmarks not transferred from the initiator plus the number of blocks, filemarks, and setmarks remaining in the target's buffer). Note that it is possible for the value in the information field to exceed the transfer length.

IMPLEMENTORS NOTE: The target should ensure that some additional data can be written to the medium (e.g., labels, filemarks, or setmarks) after the first early-warning indication has been returned to the initiator.

If a WRITE FILEMARKS command is received while the logical unit is positioned between early-warning and end-of-partition, the target shall return CHECK CONDITION status after attempting to perform the command. The EOM and valid bits shall be set to one in the sense data. If all filemarks or setmarks to be written are successfully transferred to the medium, the information field shall be set to zero. If any filemarks or setmarks to be written are not transferred to the medium prior to encountering end-of-partition, the sense key shall be set to VOLUME OVERFLOW and the information field shall be set to the requested transfer length minus the actual number of filemarks or setmarks written to the medium.

9.3. Parameters for Sequential-Access Devices

9.3.1. Diagnostic Parameters

This section defines the descriptors and pages for diagnostic parameters used with sequential-access devices.

The diagnostic page codes for sequential-access devices are defined in Table 9-19.

Table 9-19: Diagnostic Page Codes

Page Code	Description	Section
00h	Supported Diagnostic Pages	7.3.1.1
01h - 3Fh	Reserved (for all device type pages)	
40h - 7Fh	Reserved	
80h - FFh	Vendor-specific pages	

9.3.2. Log Parameters

This section defines the descriptors and pages for log parameters used with sequential-access devices.

The log page codes for sequential-access devices are defined in Table 9-20.

Table 9-20: Log Page Codes

Page Code	Description	Section
01h	Buffer Over-Run/Under-Run Page	7.3.2.1
02h	Error Counter Page (Write) Page	7.3.2.2
03h	Error Counter Page (Read) Page	7.3.2.2
04h	Error Counter Page (Read Reverse) Page	7.3.2.2
05h	Error Counter Page (Verify) Page	7.3.2.2
07h	Last n Error Events Page	7.3.2.3
06h	Non-Media Error Page	7.3.2.4
00h	Supported Log Pages	7.3.2.5
08h - 2Fh	Reserved	
3Fh	Reserved	
30h - 3Eh	Vendor Specific	

9.3.3. Mode Parameters

This section defines the descriptors and pages for mode parameters used with sequential-access devices.

The mode parameter list, including the mode parameter header and mode block descriptor, are described in 7.3.3.

The medium-type code field in the mode parameter header is reserved for sequential-access devices.

The device-specific parameter byte of the mode parameter header (see Tables 7-61 and 7-62) is defined in Table 9-21 for sequential-access devices.

Table 9-21: Device-Specific Parameter

Bit	7	6	5	4	3	2	1	0
	WP	Buffered Mode			Speed			

For the MODE SENSE command, a write protect (WP) bit of zero indicates that the medium is write enabled. A WP bit of one indicates that the medium is write protected. For the MODE SELECT command, this field is ignored.

Values for the buffered mode field are defined as follows:

- 0h The target shall not report GOOD status on write commands until the data blocks are actually written on the medium.
- 1h The target may report GOOD status on write commands as soon as all the data specified in the write command has been transferred to the target's buffer. One or more blocks may be buffered prior to writing the block(s) to the medium.
- 2h The target may report GOOD status on write commands as soon as:
 - (1) All the data specified in the write command has been successfully transferred to the target's buffer, and
 - (2) All buffered data from different initiators has been successfully written to the medium.

3h - 7h Reserved

Values for the speed field shall be assigned as follows:

- 0h Default (Use the peripheral device's default speed).
- 1h Use the peripheral device's lowest speed.
- 2h - Fh Use increasing peripheral device speeds.

For the MODE SELECT command, the density code field of the sequential-access device block descriptor (7.3.3) indicates the density selected by the initiator for use in subsequent read and write operations. For devices capable of automatic density recognition, the density code selected by the initiator may be overridden by the target for a subsequent read operation if the selected value does not match the current recorded density of the medium.

For the MODE SENSE command, the density code field reflects the current operating density of the device. For some devices, the default density code value returned in response to a MODE SENSE command may change dynamically to match the most recently selected density. The density code value returned in response to a MODE SENSE command shall be as described below:

- (1) Following a UNIT ATTENTION condition for a power on or hard reset condition, while not ready, the target shall report the default density.
- (2) Following a UNIT ATTENTION condition for a not-ready-to-ready transition, the target shall:
 - (a) report the default density if no attempt has been made by the target to determine the density.
 - (b) report the default density if the target cannot automatically determine the density from the medium.
 - (c) report the current recorded density if the target can automatically determine the density from the medium.
- (3) Following a successful read operation at or after beginning-of-medium, the target shall report a density code value reflecting the recorded density of the medium. For some implementations, the target may automatically determine this value from the medium. For devices not capable of automatic density determination, the default density is reported if the density code value is not provided by the preceding MODE SELECT command.
- (4) Following an unsuccessful read operation or a successful write operation, while at beginning-of-partition, the target shall:
 - (a) report a density code value as described for item (2) if a previous MODE SELECT command has not established a density code for the currently mounted volume.
 - (b) report a density code value as provided by the last successful MODE SELECT command for the currently mounted volume.
- (5) Following a successful unload operation the target shall report the most recent density code value as determined by items (2) through (4) above.

Table 9-22 lists the sequential-access device density codes

Table 9-22: Sequential-Access Density Codes

Code Value	Density							
00h	Default (target or peripheral device's default density)							
	Magnetic Tapes							
	Width	Tracks	Density		Code	Type	Reference	Note
	mm (Inch)		bpm	bpi				
01h	12.7 (0.5)	9	32	(800)	NRZI	R	X3.22-1983	2
02h	12.7 (0.5)	9	63	(1600)	PE	R	X3.39-1986	2
03h	12.7 (0.5)	9	246	(6250)	GCR	R	X3.54-1986	2
04h	6.3 (0.25)	4/9	315	(8000)	GCR	C	X3.136-1986	1,3
05h	6.3 (0.25)	4/9	315	(8000)	GCR	C	X3.136-1986	1
06h	12.7 (0.5)	9	126	(3200)	PE	R	X3.157-1987	2,4
07h	6.3 (0.25)	4	252	(6400)	IMFM	C	X3.116-1986	1
08h	3.81 (0.15)	4	315	(8000)	GCR	CS	X3.158-1987	1,4
09h	12.7 (0.5)	18	1491	(37871)	GCR	C	X3B5/87-099	2,4
0Ah	12.7 (0.5)	22	262	(6667)	MFM	C	X3B5/86-199	1,4
0Bh	6.3 (0.25)	4	63	(1600)	PE	C	X3.56-1986	1
0Ch	12.7 (0.5)	24	500	(12690)	GCR	C	HI-TC1	1,6
0Dh	12.7 (0.5)	24	999	(25380)	GCR	C	HI-TC2	1,6
0Eh	Reserved for ECMA							
0Fh	6.3 (0.25)	15	394	(10000)	GCR	C	QIC-120	1,6
10h	6.3 (0.25)	18	394	(10000)	GCR	C	QIC-150	1,6
11h	6.3 (0.25)	26	630	(16000)	GCR	C	QIC-320	1,6
12h	6.3 (0.25)	30	2034	(51667)	RLL	C	QIC-1350	1,6
13h	3.81 (0.15)	1	2400	(61000)	DDS	CS	X3B5/88-185A	5
14h	8.00 (0.315)	1	2126	(54000)		CS	X3B5/88-036	5
15h - 7Eh	Reserved							
7Fh	No change from previous density (NO-OP)							7
80h - FFh	Vendor unique							

NOTES:

- (1) Serial Recorded.
- (2) Parallel Recorded.
- (3) Old format known as QIC-11.
- (4) See Appendix D for additional standards information.
- (5) Helical Scan
- (6) This is not an American National Standard. The reference is based on an industry standard definition of the media format.
- (7) This density code value is defined for the MODE SELECT command and shall not be returned by the MODE SENSE command.

The mode page codes for sequential-access devices are defined in Table 9-23.

Table 9-23: Mode Page Codes

Page Code	Description	Section
0Ah	Control Mode Page	7.3.3.1
10h	Device Configuration Page	9.3.3.1
02h	Disconnect-Reconnect Page	7.3.3.2
11h	Medium Partition Page(1)	9.3.3.2
12h	Medium Partition Page(2)	9.3.3.3
13h	Medium Partition Page(3)	9.3.3.3
14h	Medium Partition Page(4)	9.3.3.3
09h	Peripheral Device Page	7.3.3.3
01h	Read-Write Error Recovery Page	9.3.3.4
00h	Vendor specific (does not require page format)	7.3.3
03h - 08h	Reserved	
08h - 10h	Reserved	
15h - 3Eh	Vendor Specific (page format required)	
3Fh	Returns all pages (valid only for the MODE SENSE command)	

Key:

	Code
NRZI	Non Return to Zero, change on ones
GCR	Group Code Recording
PE	Phase Encoded
IMFM	Inverted Modified Frequency Modulation
MFM	Modified Frequency Modulation
DDS	DAT Data Storage
RLL	Run Length Limited

	Type
R	Reel-to-Reel
C	Cartridge
CS	Cassette

9.3.3.1. Device Configuration Page

Table 9-24: Device Configuration Page

Bit Byte	7	6	5	4	3	2	1	0					
0	PS	Reserved	Page Code (10h)										
1	Page Length (0Eh)												
2	Reserved	CAP	CAF	Active Format									
3	Active Partition												
4	Write Buffer Full Ratio												
5	Read Buffer Empty Ratio												
6	(MSB)	Write Delay Time						(LSB)					
7													
8	DBR	BIS	RSmk	AVC	SOCF		RBO	REW					
9	Gap Size												
10	EOD Defined			EEG	SEW	Reserved							
11	(MSB)							(LSB)					
12	Buffer Size at Early Warning												
13													
14	Select Data Compression Algorithm												
15	Reserved												

This page (Table 9-24) is used to specify the appropriate sequential-access device configuration.

A change active partition (CAP) bit of one indicates that the logical partition is to be changed to the one specified by the active partition field. A CAP bit of zero indicates no partition change is specified.

A change active format (CAF) bit of one indicates that the active format is to be changed to the value specified in the active format field. A CAF bit of zero indicates no active format change is specified. For some devices, the format may only be changed when the logical unit is at beginning-of-partition.

The active format field indicates which recording format is in use for the selected density code when reading or writing data on a logical unit. The value of the active format field is vendor specific.

The active partition field indicates the current logical partition number in use on the medium.

The write buffer full ratio, on WRITE commands, indicates to the target how full the buffer shall be before writing data to the medium. A value of zero indicates that the value is not specified.

The read buffer empty ratio, on READ commands, indicates to the target how empty the buffer shall be before retrieving additional data from the medium. A value of zero indicates that the value is not specified.

The write delay time indicates the maximum time, in 100 millisecond increments, that the target should wait before any buffered data that is to be written, is forced to the medium after the last buffered WRITE command that did not cause the buffer to exceed the buffer full ratio. A value of zero indicates that the target shall never force buffered data to the medium under these conditions.

A data buffer recovery (DBR) bit of one indicates that the target supports data buffer recovery using the RECOVER BUFFERED DATA command. A DBR bit of zero indicates that the target does not support data buffer recovery. This bit is target-defined.

A block identifiers supported (BIS) bit of zero indicates that block IDs are not supported in the format written on the medium. A BIS bit of one indicates that the format on the medium has recorded information about the block IDs relative to a partition. This bit is target-defined.

A report setmarks (RSmk) bit of one indicates that the target shall recognize and report setmarks during appropriate read or space operations. A RSmk bit of zero indicates that the target shall not report setmarks.

The automatic velocity control (AVC) bit of one, indicates that the device shall select the speed (if the device supports more than one speed) based on the data transfer rate that should optimize streaming activity and minimize medium repositioning. An AVC bit of zero indicates the speed chosen should be the device's default speed.

A stop on consecutive filemarks (SOCF) field of 00b indicates that the device shall pre-read data from the medium in buffered mode to the limits of the buffer capacity without regard for filemarks. This implies that the device can differentiate between data blocks and filemarks in the buffer. Values 01b, 10b, and 11b specify that the device shall terminate the pre-read operation if 1, 2, or 3 consecutive filemarks are detected, respectively. If the RSmk bit is one, the target shall interpret this field as stop on consecutive setmarks.

A recover buffer order (RBO) bit of one indicates that data blocks shall be returned from the target's buffer on a RECOVERED BUFFERED DATA command in LIFO order (last-in-first-out) from which they were written to the buffer. A RBO bit of zero indicates data blocks shall be returned in FIFO (first-in-first-out) order.

A report early-warning (REW) bit of zero indicates that the target shall not report the early-warning condition for read operations and it shall report early-warning at or before any medium-defined early-warning position during write operations.

A REW bit of one indicates that the target shall return CHECK CONDITION status with the EOM bit set in the sense data when the early-warning position is encountered during read and write operations. If the REW bit is one and the SEW bit is zero, the target shall return CHECK CONDITION status with the sense key set to VOLUME OVERFLOW when early-warning is encountered during write operations.

IMPLEMENTORS NOTE: A REW bit of one is intended for compatibility with those systems using old tape formats that require an early-warning indication during read operations. Other systems should set this bit to zero to avoid potential data loss when interchanging tapes between peripheral devices.

The GapSize field value determines the size of the inter-block gap when writing data. A value of 00h specifies the device's defined gap size. A value of 01h specifies a device defined gap size sufficiently long to support update-in-place. Values of 02h through 0fh are multipliers on the device's defined gap size. Values 10h through 7fh are reserved. Values 80h through FFh are vendor specific.

The end-of-data (EOD) defined field indicates which format type the logical unit shall use to detect and generate the EOD area. The types of EOD formats are specified below:

Value	Description
000b	Logical unit's default EOD definition
001b	Format-defined erased area of medium
010b	As specified in the SOCF field
011b	EOD recognition and generation is not supported
100b-111b	Reserved

An enable EOD generation (EEG) bit set to one indicates that the logical unit shall generate the appropriate EOD area, as determined by the EOD defined field. A value of zero indicates that EOD generation is disabled.

IMPLEMENTORS NOTE: Some logical units may not generate EOD at the completion of any write-type operation.

The synchronize at early-warning (SEW) bit set to one indicates that the target shall cause any buffered write data, filemarks, or setmarks to be transferred to the medium when early-warning is encountered. A value of zero indicates that the target shall retain any unwritten buffered data, filemarks, or setmarks in the buffer when early-warning is encountered (see the WRITE command, 9.2.14, and the WRITE FILEMARKS command, 9.2.15).

The buffer size at early-warning field indicates the value, in bytes, to which the target shall reduce its logical buffer size when writing. The target should reduce the buffer size only when the logical unit is positioned between its early-warning and end-of-partition. A value of zero indicates that the implementation of this function is device specific.

NOTE: The intent is to prevent the loss of data by limiting the size of the buffer when near the end-of-partition.

The select data compression algorithm field set to 00h indicates that the target shall not use a compression algorithm on any data sent to it prior to writing the data to the medium. A value of 01h indicates that the data to be written shall be compressed using the target's default compression algorithm. Values 02h through 7fh are reserved. Values 80h through FFh are vendor specific.

9.3.3.2. Medium Partition Page(1)

Table 9-25: Medium Partition Page(1)

Bit Byte	7	6	5	4	3	2	1	0
0	PS		Reserved		Page Code (11h)			
1	Page Length (n-1)							
2	Maximum Additional Partitions							
3	Additional Partitions Defined							
4	FDP	SDP	IDP	PSUM		Reserved		
5	Medium Format Recognition							
6	Reserved							
7	Reserved							
8 - n	Partition Size Descriptor(s)							
0	(MSB)							
1	Partition Size							
	(LSB)							

This page (Table 9-25) is used to specify the first group of medium partitions. Additional groups are specified in medium partition pages(2-4).

The maximum additional partitions field is a target-defined value indicating the maximum number of additional partitions supported by the logical unit. A value of zero indicates that no value is specified.

The additional partitions defined field specifies the number of additional partitions to be defined for a volume when the SDP or IDP bit is set to one. The maximum value allowed is the value returned in the maximum additional partitions field.

A fixed data partitions (FDP) bit of one indicates that the device assigns partitions based on its fixed definition of partitions. Setting this bit to one may only be valid at beginning-of-partition and is mutually exclusive with the SDP and IDP bits.

A select data partitions (SDP) bit of one indicates that the device is to partition the medium into the number of partitions as specified by the additional partitions defined field using partition sizes defined by the device. Setting this bit to one may only be valid at beginning-of-partition and it is mutually exclusive with the FDP and IDP fields.

An initiator-defined partitions (IDP) bit of one indicates that the initiator is defining the number and size of the data partitions using the additional partitions defined field and the partition size descriptors. Setting this bit to one may only be valid at beginning-of-partition and is mutually exclusive with the FDP and SDP fields.

IMPLEMENTORS NOTE: Since defining partitions may require re-formatting the medium for some implementations, an implicit write to the medium may occur as a result of a MODE SELECT command that supplies any of these parameters.

Partition size unit of measure (PSUM) field defines the units in which the partition size descriptors select the partition size. The values 00b, 01b, and 10b define the units as bytes, kilobytes, and megabytes, respectively. The value 11b is reserved.

The medium format recognition field is a target-defined value indicating the device's capability to automatically identify the medium format and partition information when reading an unknown volume.

Values for the medium format recognition field shall be assigned as follows:

00h Logical unit is incapable of format or partition recognition.
 01h Logical unit is capable of format recognition only.
 02h Logical unit is capable of partition recognition only.
 03h Logical unit is capable of format and partition recognition.
 04h-FFh Reserved.

IMPLEMENTORS NOTE: If a target indicates that it is not capable of medium format recognition, the initiator must supply all necessary parameters for the device to identify the specific format. The value in this field may be different following a medium change.

Partition size descriptors define the size of the respective partitions in the units specified in the PSUM field. Up to 64 partitions may be defined using this page. If additional partitions need to be defined, up to 64 partitions can be defined in each of the medium partition pages(2-4) defined in 9.3.3.3.

9.3.3.3. Medium Partition Page(2-4)

Table 9-26: Medium Partition Page(2-4)

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved	Page Code (p)					
1	Page Length (n-1)							
2 - n	Partition Size Descriptor(s)							
0	(MSB)	Partition Size						---
1	---							(LSB)

The medium partition pages (Table 9-26) are used to specify additional groups of medium partitions. The first group is specified in the medium partition page(1) (see 9.3.3.2).

The additional page codes (p) defined for partition size definition are 12h, 13h, and 14h. Up to a maximum of 64 partitions can be defined in each of these pages. The partition size is defined by the value of the partition size field. The units of size used by the partition size field is specified in the PSUM field of the medium partition page(1) (9.3.3.2).

IMPLEMENTORS NOTE: Since defining partitions may require re-formatting the medium for some implementations, an implicit write to the medium may occur as a result of a MODE SELECT command that supplies these parameters.

9.3.3.4. Read-Write Error Recovery Page

Table 9-27: Read-Write Error Recovery Page

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved	Page Code (01h)					
1	Page Length (0Ah)							
2	Reserved	Reserved	TB	Reserved	EER	PER	DTE	DCR
3	Read Retry Count							
4	Reserved							
5	Reserved							
6	Reserved							
7	Reserved							
8	Write Retry Count							
9	Reserved							
10	Reserved							
11	Reserved							

The read-write error recovery page (Table 9-27) specifies the error recovery and reporting parameters that the target shall use when transferring data between the initiator and the medium. These parameters only apply to read-write errors and do not affect message system retries or positioning error recovery procedures.

NOTE: The parameters in this page also apply to verify operations.

A transfer block (TB) bit of one indicates that a data block that is not recoverable within the specified recovery limits shall be transferred to the initiator before CHECK CONDITION status is returned. A TB bit of zero indicates that the unrecoverable data block shall not be transferred to the initiator. Data blocks that can be recovered within the recovery limits are always transferred, regardless of the value of the TB bit.

An enable early recovery (EER) bit of one indicates that the target shall use the most expedient error recovery algorithm (e.g., attempt error correction prior to retries). An EER bit of zero indicates that the target shall use the most deliberate error recovery algorithm, within the limits established by the other error recovery parameters (e.g., attempt to recover the block error-free prior to using error correction).

A post error (PER) bit of one indicates that the target shall return CHECK CONDITION status to report recovered errors. A PER bit of zero indicates that the target shall not report errors recovered within the limits established by the error recovery parameters. If this bit is zero, the DTE bit must also be set to zero.

A disable transfer on error (DTE) bit of one indicates that the target shall terminate the data transfer after a recovered read-write error occurs. All data from the recovered block shall be transferred prior to terminating the read-write operation. A DTE bit of zero indicates that the target shall not terminate the transfer for errors recovered within the limits established by the read-write error recovery parameters.

A disable correction (DCR) bit of one indicates that the target shall not use error correction codes during error recovery. A DCR bit of zero allows the use of error correction codes for error recovery.

The read retry count specifies the number of times that the target should attempt its recovery algorithm during a read operation before an unrecoverable error is reported. A read retry count of zero indicates that the target shall not use its recovery algorithm during read operations.

The write retry count specifies the number of times that the target should attempt its recovery algorithm during a write operation before an unrecoverable error is reported. A write retry count of zero indicates that the target shall not use its recovery algorithm during write operations.

9.4. Glossary for Sequential Access Devices

beginning-of-partition. The position at the beginning of the permissible recording region of a partition. If only one partition is defined, this position is typically equivalent to the beginning-of-medium.

beginning-of-medium. The extreme position along the medium in the direction away from the supply reel which can be accessed by the device.

early-warning. A physical mark or device computed position near but logically before the end-of-partition (independent of physical direction). See the REW bit in the device configuration page in 9.3.3.1.

end-of-data. End of data in a partition is denoted in format-specific manner. See the EOD defined field in the device configuration page in 9.3.3.1.

end-of-medium. The extreme position along the medium in the direction away from the take-up reel which can be accessed by the device. This position may be accessed by devices that support the LOAD UNLOAD command with the EOT bit set to one.

end-of-partition. The position at the end of the permissible recording region of a partition.

logical block. A unit of data supplied or requested by an initiator.

overlength. The incorrect length condition that exists after executing a read command when the the length of the actual block read exceeds the requested transfer length in the command descriptor block.

partition. The entire usable region of recording and reading paths in a volume or in a portion of a volume, defined in a device-specific manner. If there is more than one partition, they shall be numbered starting with zero (i.e., beginning-of-partition-zero).

setmark. A special recorded element within a partition, containing no user data, which provides a segmentation scheme hierarchically superior to filemarks for use in addressing or fast positioning on high capacity storage devices.

underlength. The incorrect length condition that exists after executing a read command when the requested transfer length in the command descriptor block exceeds the length of the actual block read.

volume. A recording medium together with its physical carrier.

10. Commands for Printer Devices

10.1. Model for Printer Devices

This command set includes capability for the printer-controlling device, which is an SCSI target to be functionally separate from the physical printer device (see Figure 10-1) as well as integrated with it. The physical printer device is connected to the SCSI target via one of several common device-level interfaces. There may be more than one physical printer device attached to the printer controlling device. In such a case, each physical printer device is assigned a separate logical unit number, beginning with zero. The printer-controlling device, printer device-level interface, and the physical printer device are collectively referred to as the printer device.

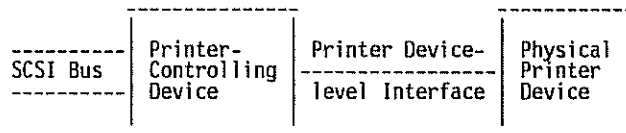


Figure 10-1: SCSI Printer Model

Specific control mechanisms are defined in mode pages for two industry-standard interfaces known as the line printer interface (e.g., the Data Products interface or equivalent) and the EIA RS-232C interface. These mode pages are used to control optional features of these interfaces. No mode page was defined for the popular industry-standard parallel interface because the options requiring controls are embedded in the data.

The printer-controlling device may be physically integrated within the printer device and it is not required to use one of the industry-standard interfaces referenced above.

The printer device commands are structured on the assumption that specific printer control codes may be embedded in the data transferred by the FORMAT, PRINT, and SLEW AND PRINT commands. The transparent control codes may take the form of "escape code sequences". Commands for the operation of the target function and some printer controls, which are not convenient to handle in a transparent way, are specified in 10.2.

This standard does not specify which character set is used by the printer device nor does it specify the meaning of the "escape code sequences" that may be used.

10.2. Commands for Printer Devices

The commands for printer devices shall be as shown in Table 10-1.

Table 10-1: Commands for Printer Devices

Command Name	Operation		Section	Page
	Code	Type		
CHANGE DEFINITION	40h	O	7.2.1	
COMPARE	39h	O	7.2.2	
COPY	18h	O	7.2.3	
COPY AND VERIFY	3Ah	O	7.2.4	
FORMAT	04h	O	10.2.1	
INQUIRY	12h	M	7.2.5	
LOG SELECT	4Ch	O	7.2.6	
LOG SENSE	4Dh	O	7.2.7	
MODE SELECT(6)	15h	O	7.2.8	
MODE SELECT(10)	55h	O	7.2.9	
MODE SENSE(6)	1Ah	O	7.2.10	
MODE SENSE(10)	5Ah	O	7.2.11	
PRINT	0Ah	M	10.2.2	
READ BUFFER	3Ch	O	7.2.12	
RECEIVE DIAGNOSTIC RESULTS	1Ch	O	7.2.13	
RECOVER BUFFERED DATA	14h	O	10.2.3	
RELEASE UNIT	17h	M	9.2.9	
REQUEST SENSE	03h	M	7.2.14	
RESERVE UNIT	16h	M	9.2.10	
SEND DIAGNOSTIC	1Dh	M	7.2.15	
SLEW AND PRINT	08h	O	10.2.4	
STOP PRINT	18h	O	10.2.5	
SYNCHRONIZE BUFFER	10h	O	10.2.6	
TEST UNIT READY	00h	M	7.2.16	
WRITE BUFFER	3Bh	O	7.2.17	

Key: M = Command implementation is mandatory.
O = Command implementation is optional.

The following operation codes are vendor specific: 01h, 02h, 05h, 06h, 07h, 08h, 09h, 0Ch, 0Dh, 0Eh, 0Fh, 11h, 13h, 19h, and C0h through FFh. All remaining operation codes for printer devices are reserved for future standardization.

10.2.1. FORMAT Command

Table 10-2: FORMAT Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (04h)							
1	Logical Unit Number			Reserved			Format Type	
2	(MSB)							
3	Transfer Length							---
4	---							---
5	(LSB)							
5	Control							

The FORMAT command (Table 10-2) provides a means for the initiator to specify forms or fonts to printers that support programmable forms or fonts. The format information is peripheral-device specific.

The format type field specifies the type of format information to be transferred from the initiator to the target. This field is defined as follows:

DB(1)	DB(0)	Format Type
0	0	Set Form
0	1	Set Font
1	0	Vendor Unique
1	1	Reserved

The transfer length specifies the length in bytes of format information that shall be transferred during the DATA OUT phase. A transfer length of zero indicates that no format information shall be transferred. This condition shall not be considered as an error.

10.2.2. PRINT Command

Table 10-3: PRINT Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (0Ah)							
1	Logical Unit Number				Reserved			
2	(MSB)							
3	Transfer Length							
4	(LSB)							
5	Control							

The PRINT command (Table 10-3) transfers the specified number of bytes from the initiator to the target to be printed.

The transfer length specifies the length in bytes of data that shall be transferred during the DATA OUT phase. A transfer length of zero indicates that no data shall be transferred. This condition shall not be considered as an error.

10.2.3. RECOVER BUFFERED DATA Command

Table 10-4: RECOVER BUFFERED DATA Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (14h)							
1	Logical Unit Number				Reserved			
2	(MSB)							
3	Transfer Length							
4	(LSB)							
5	Control							

The RECOVER BUFFERED DATA command (Table 10-4) returns to the initiator data that has been previously transferred to the target, but not yet printed.

This command is normally used only to recover from error or exception conditions that make it impossible to print the buffered data. The order in which data are transferred from the target to the initiator is the same as when the data were previously transferred using the PRINT command or SLEW AND PRINT command. Data transferred by this command are deleted from the target data buffer. One or more RECOVER BUFFERED DATA commands may be used to return the buffered data that is not printed.

If an attempt is made to recover more data than is contained in the buffer, the command shall be terminated with CHECK CONDITION status and the sense key shall be set to NO SENSE. In addition, the EOM, valid, and ILI bits shall be set to one. The information bytes shall be set to the difference (residue) between the transfer length and the actual number of bytes returned.

The transfer length specifies the maximum length in bytes of data that shall be transferred during the DATA IN phase. A transfer length of zero indicates that no data shall be transferred. This condition shall not be considered as an error.

10.2.4. SLEW AND PRINT Command

Table 10-5: SLEW AND PRINT Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (0Bh)							
1	Logical Unit Number				Reserved			Channel
2	Slew Value							
3	(MSB)							
4	Transfer Length							
5	(LSB)							
6	Control							

The SLEW AND PRINT command (Table 10-5) transfers the specified number of bytes from the initiator to the target to be printed. This command is provided for printer devices that do not support forms control information embedded within the print data.

The transfer length specifies the length in bytes of data that shall be transferred during the DATA OUT phase. A transfer length of zero indicates that no data shall be transferred. This condition shall not be considered as an error.

If the channel bit is zero, the slew value specifies the number of lines the form shall be advanced before printing. A value of 255 indicates that the form shall be advanced to the first line of the next form before printing. If the channel bit is one, the slew value specifies the forms control channel number to which the form shall be advanced prior to printing the data. If implemented, the printer options page in the MODE SELECT command provides additional control over the usage of the slew value field.

If the channel bit is one, and the channel option is not implemented, the command shall be terminated with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST.

10.2.5. STOP PRINT Command

Table 10-6: STOP PRINT Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (18h)							
1	Logical Unit Number				Reserved			Retain
2	Vendor Unique							
3	Reserved							
4	Reserved							
5	Control							

The STOP PRINT command (Table 10-6) requests the target to halt printing on buffered devices in an orderly fashion.

A retain bit of zero requests that the target data buffer be discarded; otherwise, the data not printed shall be retained. The data not printed may be recovered using the RECOVER BUFFERED DATA command, if implemented by the target. A subsequent SYNCHRONIZE BUFFER command, PRINT command or SLEW AND PRINT command shall cause the remaining data not already printed and to be printed followed by the data transferred by the subsequent command, if any. The point at which printing is suspended by this command is vendor specific.

10.2.6. SYNCHRONIZE BUFFER Command

Table 10-7: SYNCHRONIZE BUFFER Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (10h)							
1	Logical Unit Number			Reserved				
2				Reserved				
3				Reserved				
4				Reserved				
5				Control				

The SYNCHRONIZE BUFFER command (Table 10-7) provides a means for an initiator to ensure that the data have been successfully printed prior to releasing the peripheral device. This is useful for applications that wish to handle any error or exception conditions (e.g., end-of-medium) prior to termination of the application.

When all buffered data are actually printed the command shall be terminated with a GOOD status. If it is not possible to finish printing all of the buffered data (due to an error or exception condition on the peripheral device), then this command shall be terminated with a CHECK CONDITION status and the appropriate sense key.

The printer options page in the MODE SELECT command, if implemented, provides additional control over termination sequences when using this command.

10.3. Parameters for Printer Devices

10.3.1. Diagnostic Parameters

This section defines the descriptors and pages for diagnostic parameters used with printer devices.

The diagnostic page codes for printer devices are defined in Table 10-8.

Table 10-8: Diagnostic Page Codes

Page Code	Description	Section
00h	Supported Diagnostic Pages	7.3.1.1
01h-3Fh	Reserved (for all device type pages)	
40h-7Fh	Reserved	
80h-FFh	Vendor-specific pages	

10.3.2. Log Parameters

This section defines the descriptors and pages for log parameters used with printer devices.

The log page codes for printer devices are defined in Table 10-9.

Table 10-9: Log Page Codes

Page Code	Description	Section
01h	Buffer Over-Run/Under-Run Page	7.3.2.1
07h	Last n Error Events Page	7.3.2.3
06h	Non-Medium Error Page	7.3.2.4
00h	Supported Log Pages	7.3.2.5
02h - 05h	Reserved	
08h - 2Fh	Reserved	
3Fh	Reserved	
30h - 3Eh	Vendor-specific pages	

10.3.3. Mode Parameters

This section defines the descriptors and pages for mode parameters used with printer devices.

The mode parameter list, including the mode parameter header and mode block descriptor are described in 7.3.3.

The medium-type code field is contained in the mode parameter header (see Table 7-61 and 7-62). This field is reserved for printer devices.

The device specific parameter field is contained in the mode parameter header (see Table 7-61 and 7-62). Table 10-10 defines the device specific parameter used for printer devices.

Table 10-10: Printer Device Specific Parameter

Bit	7	6	5	4	3	2	1	0
	Reserved	Buffered Mode				Reserved		

When using the MODE SELECT command, a buffered mode of zero indicates that the target shall not report a GOOD status on PRINT commands or SLEW AND PRINT commands until the data are actually printed. A buffered mode of one indicates that the target may report a GOOD status on PRINT commands or SLEW AND PRINT commands as soon as the data have been transferred to the SCSI device buffer. The data from one or more commands may be buffered prior to printing. Buffered modes of 2h through 7h are reserved.

When using the MODE SENSE command, the buffered mode field returns the current value of this parameter.

The density code field is contained in the mode parameter block descriptor (see Table 7-63). This field is reserved for printer devices.

The mode page codes for printer devices are shown in Table 10-11.

Table 10-11: Mode Page Codes

Page Code	Description	Section
00h	Vendor Specific (does not require page format)	
0Ah	Control Mode Page	7.3.3.1
02h	Disconnect-Reconnect Page	7.3.3.2
03h	Parallel Printer Interface Parameters	10.3.3.1
09h	Peripheral Device Page	7.3.3.3
05h	Printer Options	10.3.3.2
04h	Serial Printer Interface Parameters	10.3.3.3
01h	Reserved	
06h - 08h	Reserved	
0Bh - 1Fh	Reserved	
20h - 3Eh	Vendor Specific (page format required)	
3Fh	Returns all pages (valid only for the MODE SENSE command)	

10.3.3.1. Parallel Printer Interface Parameters Page

This page (Table 10-12) is intended to support printer devices which use the industry-standard line printer interface.

Table 10-12: Parallel Printer Interface Parameters

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved	Page Code (03h)					
1	Parameter Length (03h)							
2	Parity Select	PIPC	Reserved	VCBP	VCBS	VES	Autofd	
3	Reserved							

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

The parity select field specifies parity generation on the printer interface and is defined as follows:

Code	Parity Select
00b	No parity generation
01b	Even Parity
10b	Odd Parity
11b	Reserved

A paper instruction parity check (PIPC) bit of one indicates that the printer interface "Paper Instruction" signal is included in parity generation on the printer interface by the target. A PIPC bit of zero indicates that the "Paper Instruction" signal is not included in parity generation on the printer interface by the target.

A VFU control bit polarity (VCBP) bit of one indicates the state of the VFU control bit (C) is true for a VFU channel command and false for a line skip command. A VCBP bit of zero indicates the opposite polarity.

A VFU control bit select (VCBS) bit of one indicates that the control bit (C) is in bit position 6 and the EE bits are in bit positions 5 and 4. A VCBS bit of zero indicates that the control bit (C) is in bit position 4 and the EE bits are in bit positions 6 and 5.

A VFU extended skip (VES) bit of one indicates that the EE bits are supported for line skip operations. A VES bit of zero indicates that the EE bits are not supported.

For MODE SELECT an automatic line feed (Autofd) bit of one specifies that the target assert the printer interface "Auto Line Feed" signal. An Autofd bit of zero specifies that the target negate the printer interface "Auto Line Feed" signal. For MODE SENSE the target shall return the current value of this bit.

The format of the VFU control byte is specified by the VCBP, VCBS, and VES bits. The VFU control byte is not part of this standard, however a discussion of it is included here for a better understanding. This VFU control byte (see table, below) includes a control bit (C) to select whether to slew over a number of lines or to skip to a vertical forms unit (VFU) channel number. This bit may be located in two different bit positions and may have either polarity. Additionally, the number of lines to slew over may be coded as an unsigned four-bit number (NNNN) or an unsigned six-bit number (EENNNN). The upper two bits of the six-bit number (EE) may be positioned contiguously with the lower four bits (NNNN) or they may be separated from the lower four bits by the control bit (C). Furthermore, the upper two bits may or may not be supported.

Pictorially, the control byte is as follows:

VFU Control Byte	Description
7 6 5 4 3 2 1 0	
0 E E C N N N N	Advance (EE)NNNN lines (C = line skip)
0 0 0 C N N N N	Skip to VFU channel NNNN (C = VFU channel)
0 C E E N N N N	Advance (EE)NNNN lines (C = line skip)
0 C 0 0 N N N N	Skip to VFU channel NNNN (C = VFU channel)

10.3.3.2. Printer Options Page

The printer options page (Table 10-13) supports control and reporting of various target functions and features.

Table 10-13: Printer Options

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved	Page Code (05h)					
1	Parameter Length (0Ah)							
2	EVFU	Font Identification						
3	Reserved		Slew Mode		Reserved		SCTE	AFC
4	(MSB)							
5	Maximum Line Length							
6	(LSB)							
7	EVFU Format Start Character							
8	EVFU Format Stop Character							
9	Line Slew options				Form Slew Options			
10	Data Termination Options				Reserved			
11	Reserved							

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

An electronic vertical forms unit (EVFU) bit of one indicates that a tape or electronic vertical forms unit is present in the printer device. An EVFU bit of zero indicates that the electronic vertical forms unit should not be used or that the unit is not present in the printer device.

For the MODE SELECT command, the font identification field specifies which font is to be used. For the MODE SENSE command, this field identifies the currently selected font. The font identification codes are defined as follows:

Code	Font Identification
00h	Default Font
01h - 3Fh	Reserved
40h - 7Fh	Vendor specific Font

For the MODE SELECT command, the slew mode field controls the target's behavior when a SLEW AND PRINT command is received with a channel bit of zero (see 10.2.4). For the MODE SENSE command, this field reports the current slew mode. This slew mode is defined as follows:

Code	Slew Mode
00b	SLEW AND PRINT commands with a channel bit of zero are supported without any required setup.
01b	SLEW AND PRINT commands with a channel bit of zero are only accepted after the initiator issues a FORMAT command with a format type of "set form" (to initialize the electronic vertical forms unit).
10b	SLEW AND PRINT commands with a channel bit of zero are always accepted. The target shall use the electronic vertical forms unit if it has previously been initialized or it shall initialize the electronic vertical forms unit for line skipping by sending the EVFU start format character immediately followed by the EVFU stop format character (both are defined below) prior to executing the SLEW AND PRINT command.
11b	Reserved

For the MODE SELECT command, a step count truncate enable (SCTE) bit of one indicates that the target shall assume the printer truncates slew values in SLEW AND PRINT commands so that the data is printed on the first line of the next form. A SCTE bit of zero indicates that the target shall assume the printer continues slewing over forms boundaries in SLEW AND PRINT commands if the slew value exceeds the number of remaining lines on the current form. The target shall return the current value of this field in the MODE SENSE data.

An ASCII forms control (AFC) bit of one indicates that the printer supports ASCII forms control characters. An AFC bit of zero indicates that the printer does not support ASCII forms control characters.

The maximum line length field specifies the maximum transfer length (maximum number of bytes per line) to be accepted in the SLEW AND PRINT command. A value of 0000h in MODE SELECT specifies that the target shall use its default value. In any case, the target shall report its actual value in the MODE SENSE data (not 0000h).

The EVFU format start character field specifies the character code to be used by the target to start the initialization of the electronic vertical forms unit if slew mode option 10b is selected.

The EVFU format stop character field specifies the character code to be used by the target to stop the initialization of the electronic vertical forms unit if slew mode option 10b is selected.

The line slew options field specifies the implementation of line slewing (using ASCII forms control characters) in the SLEW AND PRINT command. Code values in this field are defined as follows:

Code	Line Slew Option
0h	Not implemented. (SLEW AND PRINT commands cause CHECK CONDITION status with ILLEGAL REQUEST sense key.)
1h	The target shall insert an ASCII carriage return character (0Dh) for each line slewed over.
2h	The target shall insert an ASCII line feed character (0Ah) for each line slewed over.
3h	The target shall insert an ASCII carriage return character (0Dh) and line feed character (0Ah) for each line slewed over.
4h-7h	Reserved
8h-Fh	Vendor Unique

The form slew options field specifies the implementation of form slewing in the SLEW AND PRINT command. Code values in this field are defined as follows:

Code	Form Slew Option
0h	Not implemented. (SLEW AND PRINT commands cause CHECK CONDITION status with ILLEGAL REQUEST sense key.)
1h	The target shall insert an ASCII form feed character (0Ch) to move to the beginning of the next form.
2h	The target shall insert an ASCII carriage return character (0Dh) and form feed character (0Ch) to move to the beginning of the next form.
3h-7h	Reserved
8h-Fh	Vendor Unique

The data termination options field specifies the termination sequence to be issued to the printer device when a SYNCHRONIZE BUFFER command is received. Code values for this field are defined as follows:

Code	Data Termination Option
0h	Selects the target default implementation (MODE SELECT).
1h	No termination sequence. (The target sends any remaining data in its buffer to the printer device with no termination sequence.)
2h	The target shall print any buffered data followed by an ASCII carriage return character (0Dh).
3h	The target shall print any buffered data followed by an ASCII line feed character (0Ah).
4h	The target shall print any buffered data followed by an ASCII carriage return, line feed character sequence (0Dh, 0Ah).
5h	The target shall send any buffered data followed by an ASCII form feed character (0Ch).
6h	The target shall print any buffered data followed by an ASCII carriage return, form feed character sequence (0Dh, 0Ch).
7h	The target shall issue a zero line slew command to the printer device.
8h-Bh	Reserved
Ch-Fh	Vendor Unique

10.3.3.3. Serial Printer Interface Parameters Page

This page (Table 10-14) is intended to support printer devices which use the industry-standard serial interface usually referred to as EIA RS-232C.

Table 10-14: Serial Printer Interface Parameters

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved	Page Code (04h)					
1	Parameter Length (06h)							
2	Reserved		Stop Bit Length					
3	Parity Selection			Reserved	Bits per Character			
4	RTS	CTS	Reserved		Pacing Protocol			
5	(MSB)							
6	Baud Rate							
7	(LSB)							

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

The stop bit length field specifies the length of the stop bit(s) in units of 1/16 bit. The target may round this field as described in 6.5.4. A stop bit length of zero in the MODE SELECT command indicates that the target shall use its default value.

The parity selection field specifies parity generation and checking as defined in the following table:

Code	Parity Selection
----	-----
000b	None
001b	Mark
010b	Space
011b	Odd
100b	Even
101b	Reserved
110b	Reserved
111b	Reserved

The bits per character field specifies the number of bits in each character. A value of zero in the MODE SELECT command indicates that the target shall use its default character size.

For the MODE SELECT command, a request to send (RTS) bit of one specifies that the target shall insure the "Request to Send" (RTS) signal in the serial interface follows the line state of the "Data Terminal Ready" (DTR) signal in the serial interface. A RTS bit of zero specifies that the target shall set the "Request to Send" (RTS) signal to "high" whenever the target power is on. For the MODE SENSE command, the target shall report the current value of this field.

For the MODE SELECT command, a clear to send (CTS) bit of one specifies that the target shall delay data transmission to the printer device until the printer device asserts the "Clear to Send" (CTS) signal in the serial interface. A CTS bit of zero indicates that the target shall ignore the "Clear to Send" signal. For the MODE SENSE command, the target shall report the current value of this field.

The code values for the pacing protocol field are defined as follows:

Code	Pacing Protocol
----	-----
0h	None
1h	XON/XOFF
2h	ETX/ACK
3h	DTR
4h-7h	Reserved
8h-Fh	Vendor specific

The baud rate field specifies the baud rate in bits per second. The target may round this value as described in 6.5.4. A baud rate of zero in the MODE SELECT command specifies that the target shall use its default baud rate.

11. Processor Devices

11.1. Model for Processor Devices

The SCSI processor device is a target with the characteristics of a primary computing device, typically a personal computer, minicomputer, mainframe computer, or auxiliary computing device or server. Such a primary computing device is often called a host. The processor device receives or provides packets of data as requested by the initiator.

In the SCSI processor device, the target accepts and provides the data packets transferred according to the commands of the initiator. The initiator and the processor device are both assumed to know the rules by which information is exchanged between them, how the information is interpreted by the processor device, and when it is allowable to exchange the information. These rules are not specified by this standard.

The initiator requests that the processor device accept a packet of data by transmitting a SEND command. The initiator requests that the processor device return a packet of data by transmitting a RECEIVE command. A COPY command can also be transmitted to the processor device to request that it serve as a copy manager. The actual data flow can be between the processor device and another SCSI device or can be between two SCSI devices under control of the processor device acting as a copy manager.

If a processor device temporarily has no resource available to manage a data packet from the initiator, has no data packet to provide to the initiator, or has no resources assigned to perform the operation, the device may then choose one of the following responses:

- 1) Terminate the command with CHECK CONDITION status and the sense key NOT READY with the appropriate additional sense code for the condition. This is the appropriate response to a TEST UNIT READY command.
- 2) Disconnect until the necessary resource or data packet becomes available, and then reconnect to the initiator and resume the operation.
- 3) Terminate the command with BUSY status.

More than one Logical Unit can be implemented by a processor device. Logical Units can serve as additional paths to a single resource, and/or each logical unit can serve as a path to different resources within the device. A single logical unit may also serve as a path to multiple resources if the processor device can interpret information within the data packet and route the packet to the appropriate resource. If the logical unit addressed by the initiator does not have an available resource or data packet associated with it, the processor device may choose to treat the logical unit as an invalid logical unit (see 6.5.3) or respond as described in the previous paragraph.

If the processor device determines that an error or unusual condition has occurred while performing an operation specified by the contents of a data packet, the information describing the condition is normally returned as a part of a data packet. If the processor device determines that an error or unusual condition has occurred while executing the SCSI command from the initiator, the command is terminated with a CHECK CONDITION and the failures are identified through a REQUEST SENSE command.

The SCSI processor device is distinguished from a SCSI communications device by the fact that the primary destination of the data packets is within the target device. A SCSI communications device, in contrast, passes the data on to an ultimate destination outside the target through a network. Many types of devices may find it convenient to function as processor devices if no other suitable SCSI device type exists and if the packet exchange protocol dictated by the processor device model meets their functional requirements. Devices requiring totally incompatible protocols and command sets should be examined carefully to be sure that the incompatibilities are based on functional requirements. If they are, they should be treated as vendor unique device types.

Several examples of processor device implementations are provided to clarify the range of utility of the SCSI processor device.

11.1.1. Host to Host Communication, SEND Only

A host system, (Host A), takes the initiator role and selects a processor device (Host B), transmitting a packet to Host B using the SEND command. The SEND command contains an operating system call that requests data from Host B's local storage devices. After performing the functions required by the data packet, Host B assumes the initiator mode and selects Host A as a processor device and uses a SEND command to transmit the requested data back to Host A. Host A thus acts as a primary computer and Host B as a specialized data server computer. Note that the SEND command is sufficient to perform a complete transaction if both Host A and Host B can assume the initiator mode. This provides the function of a Very Local Area Network (VLAN) for very high bandwidth inter-communication among nearby host processors.

11.1.2. Host to Host Communication, SEND and RECEIVE

A host system (Host A) takes the initiator role and selects a processor device (Host B), transmitting a packet using the SEND command to Host B containing instructions about an operation to be performed. Host A again takes the initiator role and sends the data to be used by Host B in the operation. Host A then assumes that a result will be obtained consistent with rules understood by both devices. Host A generates a RECEIVE command to obtain the result from Host B. If the result is not yet ready, Host B may disconnect until the calculation is complete and the requested data packet can be returned to Host A. This method also provides the function of the VLAN between nearby hosts. Note that Host A need not support target mode and Host B need not support initiator mode to successfully complete an exchange between the two devices.

11.1.3. Host to Special Output Peripheral

A special co-processor device which can use the processor device command set is a high performance graphics display terminal. The initiator sends control and data packets to the display terminal that contain the image to be displayed. Only the SEND command would be required. A peripheral failure would be indicated through the normal CHECK CONDITION / REQUEST SENSE protocol.

11.1.4. Host to Special Input Peripheral

A second special co-processor device which can use the processor device command set is a data acquisition subsystem. Such subsystems may multiplex and compact streams of data from many sources. A host could control the data acquisition modes and the selection of data streams by transmitting control packets to the processor device using the SEND command. The host could then obtain the acquired data by executing a series of RECEIVE commands. The data acquisition device could also serve as an initiator, selecting peripheral storage devices and storing the compacted acquired data there for later access directly by the host or through Host to Host communication protocols. A peripheral failure would be indicated through the normal CHECK CONDITION / REQUEST SENSE protocol.

11.2. Commands for Processor Devices

The commands for processor devices shall be as shown in Table 11-1.

Table 11-1: Commands for Processor Devices

Command Name	Operation Code	Type	Section	Page
CHANGE DEFINITION	40h	O	7.2.1	
COMPARE	39h	O	7.2.2	
COPY	18h	O	7.2.3	
COPY AND VERIFY	3Ah	O	7.2.4	
INQUIRY	12h	M	7.2.5	
LOG SELECT	4Ch	O	7.2.6	
LOG SENSE	4Dh	O	7.2.7	
READ BUFFER	3Ch	O	7.2.12	
RECEIVE	08h	O	11.2.1	
RECEIVE DIAGNOSTIC RESULTS	1Ch	O	7.2.13	
REQUEST SENSE	03h	M	7.2.14	
SEND	0Ah	M	11.2.2	
SEND DIAGNOSTIC	1Dh	M	7.2.15	
TEST UNIT READY	00h	M	7.2.16	
WRITE BUFFER	3Bh	O	7.2.17	

Key: M = Command implementation is mandatory.
O = Command implementation is optional.

The following operation codes are vendor specific: 02h, 05h, 06h, 09h, 0Ch, 0Dh, 0Eh, 0Fh, 10h, 11h, 13h, 14h, 19h, C0h through FFh. All remaining operation codes for processor devices are reserved for future standardization.

11.2.1. RECEIVE Command

Table 11-2: RECEIVE Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (08h)							
1	Logical Unit Number				Reserved			
2	(MSB)							
3	Allocation Length							
4								
5	(LSB)							
	Control							

The RECEIVE command (Table 11-2) requests the target transfer data to the initiator. The contents of the data are not defined by this standard.

11.2.2. SEND Command

Table 11-3: SEND Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (0Ah)							
1	Logical Unit Number				Reserved			AEN
2	(MSB)							
3	Transfer Length							
4								
5	(LSB)							
	Control							

The SEND command (Table 11-3) requests the target transfer data from the initiator.

An asynchronous event notification (AEN) bit of one indicates that the data to be transferred conforms to AEN data format as defined in Table 11-4. A SEND command with an AEN bit of one shall only be issued to logical unit zero.

An AEN bit of zero indicates that the data to be transferred are vendor-specific.

The transfer length specifies the length in bytes of data that shall be sent during the DATA OUT phase. A transfer length of zero indicates that no data shall be sent. This condition shall not be considered as an error.

Table 11-4: SEND Command - AEN Data Format

Bit Byte	7	6	5	4	3	2	1	0
0	Reserved		LUNTAR	Reserved		LUNTRN		
1	Reserved							
2	Reserved							
3	Reserved							
4 to	Sense Data Byte (0)							
n+4	Sense Data Byte (n)							

A logical unit target (LUNTAR) bit of zero specifies that the asynchronous event occurred on a logical unit. A LUNTAR bit of one specifies that the asynchronous event occurred on a target routine.

If the LUNTAR bit is zero, the logical unit number target routine number (LUNTRN) field specifies on which logical unit the asynchronous event occurred. If the LUNTAR bit is one, the LUNTRN field specifies on which target routine the asynchronous event occurred.

The sense data bytes are defined in Table 7-35.

11.3. Parameters for Processor Devices

11.3.1. Diagnostic Parameters

This section defines the descriptors and pages for diagnostic parameters used with processor devices.

The diagnostic page codes for processor devices are defined in Table 11-5.

Table 11-5: Diagnostic Page Codes

Page Code	Description	Section
00h	List of Supported Parameters Page	7.3.1.1
01h - 3Fh	Reserved (for all device type pages)	
40h - 7Fh	Reserved	
80h - FFh	Vendor-specific pages	

11.3.2. Log Parameters

This section defines the descriptors and pages for log parameters used with processor devices.

The log page codes for processor devices are defined in Table 11-6.

Table 11-6: Log Page Codes

Page Code	Description	Section
01h	Buffer Over-Run/Under-Run Page	7.3.2.1
07h	Last n Error Events Page	7.3.2.3
00h	List of Supported Parameters Page	7.3.2.5
06h	Non-Medium Error Page	7.3.2.4
02h - 05h	Reserved	
08h - 2Fh	Reserved	
30h - 3Eh	Vendor-specific pages	
3Fh	Reserved	

11.4. Glossary for Processor Devices

host. A device with the characteristics of a primary computing device, typically a personal computer, workstation, minicomputer, mainframe computer, or auxiliary computing device or server.

data packet. The data transferred during the DATA IN phase of a RECEIVE command, or during the DATA OUT phase of a SEND command. A data packet often contains information at the beginning or end of the packet that describes the contents of the packet. A data packet might contain control or status information for the destination device.

resource. A part of the device required to operate on or store the data packet.

12. Write-Once Devices

12.1. Model for Write-Once Devices

The model for the write-once device is a variation on the direct access device type. Most of the aspects of the direct access device are retained by the write-once device, such as removability, logical block addressing, and data cache. The differences are summarized below.

12.1.1. Logical Blocks

Data may be written to a logical block only once. A subsequent write to a logical block already written may or may not be corrupted, depending on the implementation. Write-once physical medium is non-volatile.

The "updating" of blocks (see 15.1) is discouraged. SCSI write-once devices are intended to be "archival" in nature. Data at a logical block address is not expected to change once written. The "update" commands are not defined for this device type. Devices are not precluded from implementing the "update" function by this standard, though devices that require the update function should use the optical memory device type.

Devices may be able to determine the state of a logical block prior to access. These devices can determine whether a block is blank or written. This is useful in detecting previously written blocks, preventing a destructive overwrite. This is also useful in finding blank areas for later writing. The MEDIUM SCAN command can be used to find blank and written areas prior to WRITE and READ access.

12.1.2. Initialization

The FORMAT UNIT command is not defined for write-once devices. Write-once devices are shipped pre-formatted by the manufacturer and are ready for use when mounted.

12.1.3. Physical Medium Defects

The raw defect rate for optical medium is higher than for magnetic medium used in direct access devices. Data is usually recovered through the use of sophisticated error correction algorithms. The level of error correction used for data recovery is often selectable. However, many write-once devices have a minimum level that is always used and cannot be deactivated through the error recovery mode parameter. Control of the error correction algorithms and level of correction is vendor specific.

12.1.4. Error Reporting

If any of the following conditions occur during the execution of a command the target shall return CHECK CONDITION status. The appropriate sense key and additional sense code should be set. The following table illustrates some error conditions and the applicable sense keys. The table does not provide an exhaustive enumeration of all conditions that may cause the CHECK CONDITION status.

Condition	Sense Key
Invalid logical block address	ILLEGAL REQUEST
Unsupported option requested	ILLEGAL REQUEST
Target reset or medium change since last command from this initiator	UNIT ATTENTION
Self diagnostic failed	HARDWARE ERROR
Unrecovered read error	MEDIUM ERROR HARDWARE ERROR
Recovered read error	RECOVERED ERROR
Overrun or other error that might be resolved by repeating the command	ABORTED COMMAND
Attempt to write on write protected medium	DATA PROTECT
Attempt to read a blank or previously unwritten block	BLANK CHECK
Attempt to write a previously written block and blank checking is enabled	BLANK CHECK

In the case of an invalid logical block address, the sense data information field shall be set to the logical block address of the first invalid address.

In the case of an attempt to read a blank or previously unwritten block, the sense data information field shall be set to the logical block address of the first blank block encountered. The data read up to that block shall be transferred.

In the case of an attempt to write a previously written block and blank checking is enabled, the sense information field shall be set to the logical block address of the first non-blank block encountered.

12.2. Commands for Write-Once Devices

The commands for write-once devices shall be as shown in Table 12-1 and Table 12-2.

Table 12-1: Commands for Write-Once Devices(1)

Command Name	Code	Type	Section
CHANGE DEFINITION	40h	O	7.2.1
COMPARE	39h	O	7.2.2
COPY	18h	O	7.2.3
COPY AND VERIFY	3Ah	O	7.2.4
INQUIRY	12h	M	7.2.5
LOCK UNLOCK CACHE	36h	O	8.2.2
LOG SELECT	4Ch	O	7.2.6
LOG SENSE	4Dh	O	7.2.7
MEDIUM SCAN	38h	O	15.2.3
MODE SELECT(6)	15h	O	7.2.8
MODE SELECT(10)	55h	O	7.2.9
MODE SENSE(6)	1Ah	O	7.2.10
MODE SENSE(10)	5Ah	O	7.2.11
PRE-FETCH	34h	O	8.2.3
PREVENT ALLOW MEDIUM REMOVAL	1Eh	O	8.2.4
READ(6)	08h	O	8.2.5
READ(10)	28h	M	8.2.6
READ(12)	A8h	O	15.2.4
READ BUFFER	3Ch	O	7.2.12
READ CAPACITY	25h	M	8.2.7
READ LONG	3Eh	O	8.2.9
REASSIGN BLOCKS	07h	O	8.2.10
RECEIVE DIAGNOSTIC RESULTS	1Ch	O	7.2.13
RELEASE	17h	M	8.2.11
REQUEST SENSE	03h	M	7.2.14
RESERVE	16h	M	8.2.12
REZERO UNIT	01h	O	8.2.13

Table 12-2: Commands for Write-Once Devices(2)

Command Name	Code	Type	Section
SEARCH DATA EQUAL(10)	31h	O	8.2.14.1
SEARCH DATA EQUAL(12)	B1h	O	15.2.8
SEARCH DATA HIGH(10)	30h	O	8.2.14.2
SEARCH DATA HIGH(12)	80h	O	15.2.8
SEARCH DATA LOW(10)	32h	O	8.2.14.3
SEARCH DATA LOW(12)	B2h	O	15.2.8
SEEK(6)	0Bh	O	8.2.15
SEEK(10)	2Bh	O	8.2.15
SEND DIAGNOSTIC	1Dh	M	7.2.15
SET LIMITS(10)	33h	O	8.2.16
SET LIMITS(12)	B3h	O	15.2.9
START STOP UNIT	1Bh	O	8.2.17
SYNCHRONIZE CACHE	35h	O	8.2.18
TEST UNIT READY	00h	M	7.2.16
VERIFY(10)	2Fh	O	15.2.11
VERIFY(12)	AFh	O	15.2.12
WRITE(6)	0Ah	O	8.2.20
WRITE(10)	2Ah	M	8.2.21
WRITE(12)	AAh	O	15.2.14
WRITE AND VERIFY(10)	2Eh	O	8.2.22
WRITE AND VERIFY(12)	AEh	O	15.2.16
WRITE BUFFER	3Bh	O	7.2.17
WRITE LONG	3Fh	O	8.2.23

Key: M = Command implementation is mandatory.
O = Command implementation is optional.

The following command codes are vendor specific: 02h, 05h, 06h, 09h, 0Ch, 0Dh, 0Eh, 0Fh, 10h, 11h, 13h, 14h, 19h, 20h, 21h, 22h, 23h, 24h, 26h, 27h, 29h, and C0h through FFh. All remaining command codes for write-once devices are reserved for future standardization.

12.3. Parameters for Write-Once Devices

Refer to the parameters for optical memory devices (see 15.3).

12.4. Glossary for Write-Once Devices

Refer to the glossary for optical memory devices (see 15.4).

13. CD-ROM Devices

13.1. Model for CD-ROM Devices

CD-ROM devices permit reading data from a rotating media. Writing the media is not supported. Data transfer can begin with any of the consecutively numbered logical blocks. Data on CD-ROM devices is addressed the same as usual (magnetic) direct-access devices. Some CD-ROM devices support a separate information stream (e.g. audio and/or video but referred to as audio in this section) transmitted via a connection other than the SCSI Bus. This device type defines commands for controlling these other information streams.

CD-ROM drives are designed to work with any disc that meets the "Red Book" or "Yellow Book" (IEC 908) documents at the current revision level. Many new drives read CD-ROM data discs, digital audio discs, and audio-combined discs (i.e. some tracks are audio, some tracks are data).

13.1.1. CD-ROM Media Organization

The formats written on the CD-ROM and CD-DA (Digital Audio) media require special interfacing considerations.

NOTE: This section contains a number of terms which have special meanings peculiar to CD-ROM technology or that may be unfamiliar to many readers of this standard. The glossary (see 13.4.) defines these terms.

These discs may contain either audio, data or a mixture of the two. Table 13-2 shows the format of a audio-combined disc to illustrate the relationship between the logical block addresses reported in SCSI and the MSF address encoded on the media.

NOTE: The term frame is used in two different ways in the CD-ROM media standard. The intended meaning can only be determined from the context. Whenever possible, this description replaces the larger data unit with the more familiar term "sector". The primary exception to this policy is the use of frame when referring to the MSF address. In the MSF context, one frame (F Field unit) equals one sector. On a typical two channel CD-DA media, each frame (F Field unit) is played in 1/75th of a second.

The physical format defined by the CD-ROM media standards provides 2352 bytes per sector. For usual computer data applications, 2048 bytes are used for user data, 12 bytes for a synchronization field, 4 bytes for a sector address tag field and 288 bytes - the auxiliary field - for L-EC (CD-ROM Data Mode 1). In less critical applications, the auxiliary field may also be used for user data (CD-ROM Data Mode 2). A CD-ROM physical sector size is 2048, 2336 or 2340 bytes per sector. These values correspond to user data field only, user data plus auxiliary data, the 4 byte address tag plus user data plus auxiliary data.

Selection among these options is made by setting the density code field of the MODE SELECT Block Descriptor. The usual selection is 2048 bytes per physical sector via a Density Code setting of 1. (See Table 13-32 CD-ROM Density Code Values.)

This same area of the CD-ROM or CD audio media may store 1/75th of a second of 2 channel audio information formatted according to the CD-DA specification. (These audio channels are usually the Left and Right components of a stereo pair.) An audio only density code value can be used to declare an area of the media to be invalid for data operations.

Data sectors are addressed by their logical block address relative to the first addressable block on the medium which is located at MSF address 00/02/00. This MSF address is the logical beginning of medium, and is designated logical block zero.

Logical addressing of CD-ROM information may use any logical block length. When the specified logical block length is an exact divisor or integral multiple of the selected number of bytes per CD-ROM sector, the device shall map one to one the bytes transferred from CD-ROM sectors to the bytes of Logical Blocks. For instance, if 2048 bytes are transferred from each CD-ROM sector (specified by the CD-ROM Density Code value, see Table 13-32), and the logical block length is 512 bytes, then each CD-ROM sector shall map to exactly four logical blocks. This standard does not define the mapping of Logical Block lengths which do not evenly divide or are not exact multiples of the selected number of bytes per CD-ROM sector.

A track may be viewed as a partition of the CD-ROM address space. A CD-ROM media contains from one to ninety-nine tracks. All information sectors of a track are required to be of the same type (audio or data) and mode. Each change in the type of information on the disc requires a change in track number. A disc containing both audio and data would have at least two tracks, one for audio and one for data.

The tracks of a CD media are numbered consecutively with values between 1 and 99. However, the first information track may have a number greater than 1. Tracks have a minimum length of 300 sectors including any transition area that is part of a track.

The CD-ROM media standards require transition areas between tracks encoded with different types of information. In addition, transition areas may be used at the beginning or end of any track. For audio tracks the transition areas are called pause areas. For data tracks transition areas are called pre-gap and post-gap areas. See Table 13-2 for an example. The "Red Book" and "Yellow Book" documents specify minimum time durations for these areas. Maximum time durations are not specified.

Transition areas are formatted and the logical address continues to increment through transition areas. Some media (i.e. discs with only one track) may not have transition areas. The means to determine of the location of the transition areas is vendor or application specific and is addressed by other standards (e.g. ISO 9660).

CD-ROM is a unique SCSI device in the respect that some logical blocks on a disc may not be accessible by all commands. SEEK commands may be issued to any logical block address within the reported capacity of the disc. READ commands cannot be issued to logical blocks that occur in some transition areas, or to logical blocks within an audio track. PLAY commands cannot be issued to logical blocks within a data track.

The blank verify (BlkVfy) option of the Verify command (15.2.11), for CD-ROM devices, shall verify that the addressed blocks are transition area blocks.

CD-ROM media have lead-in and lead-out areas. These areas are outside of the user accessible area as reported in the READ CD-ROM CAPACITY data. The lead-in area of the media is designated track zero. The lead-out area is designated track 0AAh. The sub-channel Q in the lead-in track contains a table of contents (TOC) of the disc.

IMPLEMENTORS NOTE: The READ CD-ROM CAPACITY command returns the logical block address of the last block prior to the lead-out area. This location may be in a transition area and therefore not a valid address for read operations.

The table of contents gives the absolute MSF location of the first information sector of each track. Control information (audio/data, method of audio encoding, etc.) for each track is also given in the TOC. However, the TOC does not distinguish between the different modes of data tracks (i.e. CD-ROM Data Mode one vs. CD-ROM Data Mode two).

The MSF locations of the beginning of data tracks in the TOC are required to be accurate, however, the TOC values for audio tracks have a tolerance of plus or minus 75 sectors. Information from the TOC can be used to reply to a READ CD-ROM CAPACITY command. When this is done the drive implementor should consider the possible tolerances and return a value which allows access to all information sectors.

An index is a partition of a track. Pre-gap areas are encoded with an index value of zero. Pause areas at the beginning of audio tracks are also encoded with an index value of zero. The first information sector of a track has index value of one. Consecutive values up to 99 are permitted. Index information is not contained in the TOC. Not all sectors are encoded with the index value in the Q-sub-channel data (the requirement is 9 out of 10). A sector without an index value is presumed to have the same index as the preceding sector.

Tracks and indexes are not defined to be any particular length. (Except for a minimum track length of 300 sectors.) A CD disc may be created with a single information track with a single index or with 99 information tracks each with 99 indexes.

The sub-channel information which is part of each sector includes a track relative MSF location value giving the distance from the first information sector of the track. On the media this value decreases during the pre-gap area (sectors with index values of 0) and increases for the rest of the track. The data returned by the READ SUB-CHANNEL command with MSF bit set to zero, converts this to a track relative logical block address (TRLBA). The TRLBA is continually increasing over the whole track and pre-gap areas shall return negative values. When the MSF bit in the Read Sub-channel command is set to one, the MSF track relative location value from the media is reported without change.

13.1.2. CD-ROM Physical Data Format

The physical format of CD-ROM and CD-DA media uses a smaller unit of synchronization than the more familiar magnetic or optical recording systems. The basic data stream synchronization unit is a small frame. This is not the same large frame (sector) as referred to in the MSF unit. Each small frame consists of 588 bits. A sector on CD-ROM media consists of 98 small frames.

A CD-ROM small frame consists of:

1 Synchronization pattern	(24+3 bits)	
1 Byte of Sub-channel Data	(14+3 bits)	
24 Bytes of Data	(24 x (14+3) bits)	
8 Bytes of CIRC code	(8 x (14+3) bits)	Total: 588 bits.

Data, sub-channel and CIRC bytes are encoded with an eight-to-fourteen bit code then three merging bits are added. The merging bits are chosen to provide minimum low frequency signal content and optimize phase lock loop performance.

13.1.2.1. Frame Format for Audio

Each small frame of an audio track on a two channel CD-DA or CD-ROM media consists of six digitized 16 bit samples of each audio channel. These 24 bytes of data are combined with a synchronization pattern, CIRC bytes and a sub-channel byte to make a frame. Each frame takes approximately 136.05 μ s to play. This gives a sampling rate of 44.1 KHz for each channel. The sub-channel information creates the higher level sector grouping for audio tracks.

13.1.2.2. Sector Format for Data

The data bytes of 98 small frames comprise the physical unit of data which is referred to as a sector. (98 small frames times 24 bytes per small frame equal 2352 bytes of data per sector.)

A sector which contains CD-ROM Data Mode one data has the following format:

12 bytes	Synchronization Field
4 bytes	CD-ROM data header
	Absolute M Field in bcd format
	Absolute S Field in bcd format
	Absolute F Field in bcd format
	CD-ROM Data Mode Field
2048 bytes	User Data Field
4 bytes	Error Detection Code
8 bytes	zero
276 bytes	Layered Error Correction Code

A sector which contains CD-ROM Data Mode two data has the following format:

12 bytes	Synchronization Field
4 bytes	CD-ROM Data Header
	Absolute M Field in bcd format
	Absolute S Field in bcd format
	Absolute F Field in bcd format
	CD-ROM Data Mode Field
2336 bytes	User Data Field (2048 bytes of Mode 1 data plus 288 bytes of Auxiliary data)

IMPLEMENTORS NOTE: Many drives are capable of returning CD-ROM Data Mode one data in a CD-ROM Data Mode two format. This allows the user to investigate the error detection and error correction codes. However data encoded as CD-ROM Data Mode two cannot be read as CD-ROM Data Mode one data.

13.1.2.3. Sub-channel Information Formats

The Sub-channel byte of each frame is assigned one bit to each of the 8 sub-channels, designated P, Q, R, S, T, U, V, W. Only sub-channels P and Q are defined for CD-ROM media. The other sub-channel bits are zero.

Sub-channel P is a simple flag bit that may be used for audio muting control and track boundary determination.

Sub-channel Q has a higher level of structure. All the sub-channel Q bits of a sector define the sub-channel Q information block. (For audio tracks, decoding the Q sub-channel is the only way to distinguish sector boundaries.)

The sub-channel Q block consists of 98 bits, one bit from each small frame in a sector. Three formats are defined for the sub-channel Q information block. The first format provides location information and is defined as follows:

2 bits	Sub-channel Synchronization Field
4 bits	ADR Field - defines the format
4 bits	Control field - defines the type of information in this sector.
8 bits	Track Number (in BCD)
8 bits	Index Number (in BCD)
24 bits	Track Relative MSF address
8 bits	reserved (0)
24 bits	Absolute MSF address
16 bits	CRC error detection code.

This format is required to exist in at least nine out of ten consecutive sectors.

The second and third formats are optional. If used they must exist in at least one out of 100 consecutive sectors. They include the absolute Frame byte of the MSF address to provide location information continuity.

The second format gives the catalog number of the disc (UPC/EAN Bar Code number). This information is constant over the whole media.

The third format gives the International-Standard-Recording-Code (ISRC) for each track. The ISRC is defined in DIN-31-621. This format is not present on lead-in or lead-out tracks and may change only after the track number changes.

13.1.3. CD Audio Error Reporting

PLAY AUDIO commands with the immediate bit set in the audio control mode return status as soon as the command has been validated (which may involve a seek to the starting address). The playback operation continues and may complete without notification to the initiator. Error termination of audio operations shall be reported to the initiator by returning immediate CHECK CONDITION status to the next command (except for REQUEST SENSE and INQUIRY.) The deferred error sense data error code, 71h (see Table 7-35 and 7.2.14.2.) is used to indicate that the error is not due to the current command.

The status of the play operation may be determined by issuing a REQUEST SENSE command. The sense key is set to NO SENSE and the play operation status (see Table 13-20) is reported in the additional sense code qualifier field.

13.1.4. CD-ROM Ready Condition/Not Ready Condition

The ready condition occurs after a cartridge is inserted and the drive has performed its initialization tasks. These may include reading the table of contents from the media.

A not ready condition may occur for the following reasons:

- (1) There is no cartridge inserted.
- (2) The drive is unable to load or unload the cartridge.
- (3) The drive is unable to recover the table of contents.
- (4) The controller cannot select the drive.

13.1.5. CD-ROM Address Reporting Formats (MSF Bit)

Several CD-ROM specific commands can report addresses either in logical or in MSF format (Table 13-1). The READ HEADER, READ SUB-CHANNEL and READ TABLE OF CONTENTS commands have this feature.

An MSF bit of zero requests that the logical block address format be used for the CD-ROM absolute address field or for the offset from the beginning of the current track expressed as a number of logical blocks in a CD-ROM track relative address field. This track relative logical block address (TRLBA) value is reported as a negative value in twos-complement notation for transition areas that have decreasing MSF encoded relative addresses.

An MSF bit of one requests that the MSF format be used for these fields. In certain transition areas the relative MSF addresses are decreasing positive values. The absolute MSF addresses are always increasing positive values (see entry in 13.4).

Table 13-1: MSF Address Format

Bit Byte	7	6	5	4	3	2	1	0
0	Reserved							
1	M Field							
2	S Field							
3	F Field							

The M, S, and F fields are expressed as binary numbers. The values match those on the media except for the encoding. The ratios of M field units to S field units and S field units to F field units is reported in the mode parameters page (see 13.3.3).

13.1.6. Sensing Support for CD-Audio Commands.

If any commands related to audio operations are implemented then the PLAY AUDIO command shall be implemented to allow a method for the initiator to determine if audio operations are supported. A target responding to a PLAY AUDIO command which has a transfer length of zero with CHECK-CONDITION status and setting the sense key to ILLEGAL REQUEST does not support audio play operations.

NOTE: The PLAY AUDIO command may be implemented without support for any of the other audio operations commands.

Table 13-2 gives an example of a disc layout.

Table 13-2: Example Mixed Mode CD-ROM Disc Layout

The logical addresses given below assume a block length of 2048, 2336 or 2340 bytes. For block lengths of 1024, 512 or 256 multiply all values given in the SCSI Address columns by 2, 4 or 8 respectively.

Block Description	SCSI Address		/--- Sub-channel Information ---\					CD-ROM Data Mode (2)
	Logical Address (decimal)	Track Relative Logical Address	Absolute MSF Address (1)	Track and Index	Track Relative MSF Address	Sector is Info or Pause	Mode or Data Audio	
Lead-In Area (3)	---	---	---	0/-	---	---	Audio	---
Pre-gap (3)	---	---	00/00/00	1/0	00/02/00	Pause	Data	Null
1st Trk Data	0000 (5)	0	00/02/00 (4)	1/1	00/00/00	Info	Data	L-EC
2nd Trk Data	6000 (5)	0	01/22/00 (4)	2/1	00/00/00	Info	Data	L-EC
	7500	1500	01/42/00	2/2	00/20/00	Info	Data	L-EC
Post-gap	9000	3000	02/02/00	2/3	00/40/00	Pause	Data	Null
Pause- Silence	9150	-150 (9)	02/04/00	3/0	00/02/00 (8)	Pause	Audio	---
3rd Trk Audio	9300 (7)	0	02/04/00 (6)	3/1	00/00/00	Info	Audio	---
	11400	2250	02/34/00	3/2	00/30/00	Info	Audio	---
Track 4 Audio	21975 (7)	0	04/53/00 (6)	4/1	00/00/00	Info	Audio	---
Pre-gap Part 1	30000	-225 (9)	06/40/00	5/0	00/03/00	Pause	Audio	---
Pre-gap Part 2	30075	-150	06/41/00	5/0	00/02/00	Pause	Data	Null
Track 5 Data	30225	0	06/43/00	5/1	00/00/00	Info	Data	L-EC
Last Info	263999 (10)	233774	58/39/74	5/1	51/56/74	Info	Data	L-EC
Post-gap	---	233775	58/40/00	5/2	51/57/00	Pause	Data	Null
Lead-out track	264000 (11)	0	58/42/00 (12)	AA/1 (13)	00/00/00	Pause	Audio	---

Notes for Table 13-2:

- (1) Absolute MSF address repeated in the header field of data blocks.
- (2) The CD-ROM Data Mode is stored in the header of data tracks. This indicates that the block is part of a data pre-gap or post-gap (null), that this is a data block using the auxiliary field for L-EC symbols (ECC - CD-ROM Data Mode one), or that this is a data block using the auxiliary field for user data (CD-ROM Data Mode two).
- (3) Table of contents information is stored in the sub-channel of lead-in area. The lead-in area is coded as track zero. Track zero and the initial 150 sector pre-gap (or audio pause) are not accessible with logical addressing.
- (4) Value stored in table of contents with zero tolerance.
- (5) Exact value returned by READ TOC command.
- (6) Value stored in table of contents plus or minus 75 sectors.
- (7) Value returned by READ TOC command plus or minus 75 blocks.
- (8) Track relative MSF value decreases to 0 in pre-gap areas.
- (9) Track relative logical addresses are negative in the pre-gap areas. Pre-gap areas have index values of zero.
- (10) Minimum value returned by READ CD-ROM CAPACITY - Exact value depends on encoding of this track and the lead out track and whether this is derived from the TOC data.
- (11) Value returned by READ TOC command - exact if lead-out track is encoded as data, or plus or minus 75 blocks if encoded as audio.
- (12) Value stored in table of contents - exact if lead-out track is coded as data, or plus or minus 75 blocks if coded as audio.
- (13) Lead-out track number is defined as 0AAh.

13.1.7. Error Reporting

If any of the following conditions occur during the execution of a command the target shall return CHECK CONDITION status. The appropriate sense key and additional sense code should be set. The following table illustrates some error conditions and the applicable sense keys. The table does not provide an exhaustive enumeration of all conditions that may cause the CHECK CONDITION status.

Condition	Sense Key
Invalid logical block address	ILLEGAL REQUEST
Unsupported option requested	ILLEGAL REQUEST
Target reset or medium change since last command from this initiator	UNIT ATTENTION
Self diagnostic failed	HARDWARE ERROR
Unrecovered read error	MEDIUM ERROR HARDWARE ERROR
Recovered read error	RECOVERED ERROR
Overrun or other error that might be resolved by repeating the command	ABORTED COMMAND
Attempt to read a blank block.	BLANK CHECK
Attempt to play a data block as audio.	BLANK CHECK

In the case of an invalid logical block address, the sense data information field shall be set to the logical block address of the first invalid address.

In the case of an attempt to read a blank or previously unwritten block, the sense data information field shall be set to the logical block address of the first blank block encountered. The data read up to that block shall be transferred.

There are other special error situations for CD-ROM devices. In the following cases the sense key shall be set to BLANK CHECK and the additional sense code set to END OF USER AREA ENCOUNTERED ON THIS TRACK.

- (1) A post-gap area is encountered. (i.e. a block with CD-ROM Data Mode 0.)
- (2) A pre-gap area is encountered. (i.e. a block with index equal to 0.)
- (3) The information type (data vs. audio) changes.

When not performing audio playback, if the logical block address requested is not within a data track the command shall be terminated with CHECK CONDITION status. The sense key shall be set to BLANK CHECK and the additional sense code set to ILLEGAL MODE FOR THIS TRACK. This applies to audio-combined and audio media.

13.2. Commands for CD-ROM Devices

The commands for CD-ROM devices shall be as shown in Table 13-3, and 13-4.

Table 13-3: Commands for CD-ROM Devices(1)

Command Name	Operation Code	Type	Section
CHANGE DEFINITION	40h	0	7.2.1
COMPARE	39h	0	7.2.2
COPY	18h	0	7.2.3
COPY AND VERIFY	3Ah	0	7.2.4
INQUIRY	12h	M	7.2.5
LOCK/UNLOCK CACHE	36h	0	8.2.2
LOG SELECT	4Ch	0	7.2.6
LOG SENSE	40h	0	7.2.7
MODE SELECT(6)	15h	0	7.2.8
MODE SELECT(10)	55h	0	7.2.9
MODE SENSE(6)	1Ah	0	7.2.10
MODE SENSE(10)	5Ah	0	7.2.11
PAUSE/RESUME	48h	0	13.2.1
PLAY AUDIO(10)	45h	0*	13.2.2
PLAY AUDIO(12)	A5h	0*	13.2.3
PLAY AUDIO MSF	47h	0*	13.2.4
PLAY AUDIO TRACK/INDEX	48h	0*	13.2.5
PLAY TRACK RELATIVE(10)	49h	0*	13.2.6
PLAY TRACK RELATIVE(12)	A9h	0*	13.2.7
PRE-FETCH	34h	0	8.2.3
PREVENT/ALLOW MEDIUM REMOVAL	1Eh	0	8.2.4
READ(6)	08h	0	8.2.5
READ(10)	28h	M	8.2.6
READ(12)	A8h	0	15.2.4
READ BUFFER	3Ch	0	7.2.12
READ CD-ROM CAPACITY	25h	M	13.2.8
READ HEADER	44h	0	13.2.9
READ LONG	3Eh	0	8.2.9
READ SUB-CHANNEL	42h	0	13.2.10
READ TOC	43h	0	13.2.11
RECEIVE DIAGNOSTIC RESULTS	1Ch	0	7.2.13
RELEASE	17h	M	8.2.11
REQUEST SENSE	03h	M	7.2.14
RESERVE	16h	M	8.2.12
REZERO UNIT	01h	0	8.2.13

Key: M = Command implementation is mandatory.
 0 = Command implementation is optional.
 * = If any of the PLAY AUDIO commands (as indicated by an * in the 'Type' column in Table 13-3) are implemented, the PLAY AUDIO command shall be implemented by the target.

Table 13-4: Commands for CD-ROM Devices(2)

Command Name	Operation Code	Type	Section
SEARCH DATA EQUAL(10)	31h	0	8.2.14.1
SEARCH DATA EQUAL(12)	81h	0	15.2.8
SEARCH DATA HIGH(10)	30h	0	8.2.14.2
SEARCH DATA HIGH(12)	80h	0	15.2.8
SEARCH DATA LOW(10)	32h	0	8.2.14.3
SEARCH DATA LOW(12)	82h	0	15.2.8
SEEK(6)	0Bh	0	8.2.15
SEEK(10)	2Bh	0	8.2.15
SEND DIAGNOSTIC	1Dh	M	7.2.15
SET LIMITS(10)	33h	0	8.2.16
SET LIMITS(12)	83h	0	15.2.9
START STOP UNIT	1Bh	0	8.2.17
SYNCHRONIZE CACHE	35h	0	8.2.18
TEST UNIT READY	00h	M	7.2.16
VERIFY(10)	2Fh	0	15.2.11
VERIFY(12)	AFh	0	15.2.12
WRITE BUFFER	3Bh	0	7.2.17

Key: M = Command implementation is mandatory.
 0 = Command implementation is optional.

The following command codes are vendor specific: 02h, 05h, 06h, 09h, 0Ch, 0Dh, 0Eh, 0Fh, 10h, 11h, 13h, 14h, 19h, 20h, 21h, 22h, 23h, 24h, 26h, 27h, 29h, and C0h through FFh. All remaining command codes for CD-ROM devices are reserved for future standardization.

13.2.1. PAUSE RESUME Command

Table 13-5: PAUSE RESUME Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (48h)							
1	Logical Unit Number			Reserved				
2	Reserved							
3	Reserved							
4	Reserved							
5	Reserved							
6	Reserved							
7	Reserved							
8	Reserved							Resume
9	Control							

The PAUSE RESUME command (Table 13-5) requests that the device stop or start an audio play operation. This command is used with PLAY AUDIO commands issued while the immediate bit is set to one.

A resume bit of zero causes the drive to enter the hold track state with the audio output muted after the current block is played. A resume bit of one causes the drive to release the pause and begin play at the block following the last block played.

If an audio play operation cannot be resumed and the resume bit is one, the command is terminated with CHECK CONDITION status. If the resume bit is zero and an audio play operation cannot be paused, (no audio play operation has been requested, or the requested audio play operation has been completed) the command is terminated with CHECK CONDITION status.

It shall not be considered an error to request a pause when a pause is already in effect or to request a resume when a play operation is in progress.

13.2.2. PLAY AUDIO(10) Command

Table 13-6: PLAY AUDIO(10) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (45h)							
1	Logical Unit Number			Reserved				RelAdr
2	(MSB)							
3								
4	Starting Logical Block Address							
5								
6	Reserved							
7	(MSB)							
8	Transfer Length							
9	(LSB)							
	Control							

The PLAY AUDIO command (Table 13-6) requests the target to begin an audio playback operation. The command function (Immed and S0TC bits) and the output of audio signals shall be as specified by the settings of the mode parameters (see 13.3.3.1.).

If any commands related to audio operations are implemented then the PLAY AUDIO command shall be implemented to allow a method for the initiator to determine if audio operations are supported. A target responding to a PLAY AUDIO command which has a transfer length of zero with CHECK CONDITION status and setting the sense key to ILLEGAL REQUEST does not support audio play operations.

The logical block address field specifies the logical block at which the audio playback operation shall begin.

The transfer length field specifies the number of contiguous logical blocks that shall be played. A transfer length field of zero indicates that no audio operation shall occur. This condition shall not be considered as an error.

If the logical block length is not equal to the sector size the target may adjust the starting logical block address and the transfer length. In such case, it is recommended that the target start the audio play operation with the beginning of a sector whenever the starting logical address falls within that sector (MSF unit). If the requested transfer length causes the end of an audio play operation to fall within a sector the target may continue the play operation through the end of that sector.

If the starting address is not found, if the address is not within an audio track, or if a not ready condition exists, the command shall be terminated with CHECK CONDITION status.

If the CD-ROM information type (data vs. audio) changes the sense key shall be set to BLANK CHECK and the additional sense code set to END OF USER AREA ENCOUNTERED ON THIS TRACK.

If the logical block address requested is not within an audio track the command shall be terminated with CHECK CONDITION status. The sense key shall be set to BLANK CHECK and the additional sense code set to ILLEGAL MODE FOR THIS TRACK.

13.2.3. PLAY AUDIO(12) Command

Table 13-7: PLAY AUDIO(12) Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (A5h)							
1	Logical Unit Number				Reserved			RelAdr
2	(MSB)							
3								
4	Logical Block Address							
5								
6	(MSB)							
7								
8	Transfer Length							
9								
10	Reserved							
11	Control							

The PLAY AUDIO(12) command (Table 13-7) requests the target to begin an audio playback operation. The command function (Immed and SOTC bits) and the output of audio signals shall be as specified by the settings of the mode parameters (see 13.3.3.1.). See the PLAY AUDIO(10) command for a description of the fields in this command.

13.2.4. PLAY AUDIO MSF Command

Table 13-8: PLAY AUDIO MSF Command

Byte	7	6	5	4	3	2	1	0
0	Operation Code (47h)							
1	Logical Unit Number				Reserved			
2	Reserved							
3	Starting M Field							
4	Starting S Field							
5	Starting F Field							
6	Ending M Field							
7	Ending S Field							
8	Ending F Field							
9	Control							

The PLAY AUDIO MSF command (Table 13-8) requests the target to begin an audio playback operation. The command function (Immed and SOTC bits) and the output of audio signals shall be as specified by the settings of the mode parameters (see 13.3.3.1.).

The starting M field, the starting S field, and the starting F field specify the absolute MSF address at which the audio play operation shall begin. The ending M field, the ending S field, and the ending F field specify the absolute MSF address at which the audio play operation shall end. All contiguous audio sectors between the starting and the ending MSF address shall be played.

A starting MSF address equal to an ending MSF address causes no audio play operation to occur. This shall not be considered an error. If the starting MSF address is less than the ending MSF address, the command shall be terminated with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST.

If the starting address is not found, or if the address is not within an audio track, or if a not ready condition exists, the command shall be terminated with CHECK CONDITION status. See 13.1.7 for a description of error reporting information.

13.2.5. PLAY AUDIO TRACK INDEX Command

Table 13-9: PLAY AUDIO TRACK INDEX Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (48h)							
1	Logical Unit Number			Reserved				
2	Reserved							
3	Reserved							
4	Starting Track							
5	Starting Index							
6	Reserved							
7	Ending Track							
8	Ending Index							
9	Control							

The PLAY AUDIO TRACK INDEX command (Table 13-9) requests the target to begin an audio play operation. The command function (Immed and SOTC bits) and the output of audio signals shall be as specified by the settings of the mode parameters (see 13.3.3.1.).

The starting track field specifies the track number of the starting audio track. The starting index field specifies the index number within the track at which the audio play operation shall begin.

The ending track field specifies the track number of the ending audio track. The ending index field specifies the index number within the track after which the audio play operation shall stop. The audio play operation shall terminate at the last block with an index number equal to the ending index. All contiguous audio sectors between the starting and the ending address shall be played.

If the starting address is less than the ending address, the command shall be terminated with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST.

If the starting address is not found, or if the address is not within an audio track, or if a not ready condition exists, the command shall be terminated with CHECK CONDITION status. See 13.1.7 for a description of error reporting information.

IMPLEMENTORS NOTE: Valid values for the track and index fields are 1 to 99. A starting index value of one specifies that playback is to start with the first audio sector of the track following the (optional) pause. A last index value of 99 specifies that playback continues through the last sector of the track.

If the ending track is greater than the last information track on the media, the playback shall continue until the last track is complete. If the ending index is greater than the largest index value on the ending track, the playback shall continue until this track is complete then terminate. These conditions shall not be considered errors.

If the starting index is greater than the largest index value on the starting track, and the stop on track crossing (SOTC) bit of the audio control MODE SELECT parameters page (see 13.3.3.1) is zero, the playback operation shall start at the beginning of the next track. This situation is not an error.

If the starting index is greater than the largest index value on the starting track, and the stop on track crossing (SOTC) bit of the audio control MODE SELECT parameters page (see 13.3.3.1) is one, the playback shall not begin. The target shall return CHECK CONDITION, and the sense key shall be set to ILLEGAL REQUEST.

IMPLEMENTORS NOTE: The operation of the SOTC bit described above comes about because the user may not be able to determine the largest index value on a track, either from the table of contents or by other means. The SOTC bit one case allows the user to determine the largest index. The SOTC bit zero case allows the user to set up play operations without complete knowledge of the media layout.

13.2.6. PLAY AUDIO TRACK RELATIVE(10) Command

Table 13-10: PLAY AUDIO TRACK RELATIVE(10) Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (49h)							
1	Logical Unit Number				Reserved			
2	(MSB)							
3								
4	Track Relative Logical Block Address							
5								
6	(LSB)							
7								
8	Starting Track							
9								
10	(MSB)							
11	Transfer Length							
12								
13	(LSB)							
14								
15	Control							

The PLAY AUDIO TRACK RELATIVE command (Table 13-10) requests that the device begin an audio playback operation. The starting address is specified as a track relative logical block address within the specified starting track. The command function (Immed and SOTC bits) and the output of audio signals shall be as specified by the settings of the mode parameters (see 13.3.3.1.).

The starting track field specifies the track number of the starting audio track.

The track relative logical block address (TRLBA) field specifies the two's complement starting logical block address relative to the beginning of the first sector on the track with an index value of one. Negative values indicate a starting location within the audio pause area at the beginning of the requested track.

The transfer length field specifies the number of contiguous logical blocks that shall be output as audio data. A transfer length field of zero indicates that no audio playback operation shall occur. This condition shall not be considered as an error. Any other value indicates the number of logical blocks that shall be output.

If the logical block length is not equal to the sector size the target may adjust the starting logical block address and the transfer length. In such case, it is recommended that the target start the audio play operation with the beginning of a sector whenever the starting logical address falls within that sector (MSF unit). If the requested transfer length causes the end of an audio play operation to fall within a sector the target may continue the play operation through the end of that sector.

If the starting address is not found, or if the address is not within an audio track, or if a not ready condition exists, the command is terminated with CHECK CONDITION status. See 13.1.7 for a description of error reporting information.

13.2.7. PLAY AUDIO TRACK RELATIVE(12) Command

Table 13-11: PLAY AUDIO TRACK RELATIVE(12) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (A9h)							
1	Logical Unit Number				Reserved			
2	(MSB)							
3								
4	Track Relative Logical Block Address							
5								
6	(MSB)							
7								
8	Transfer Length							
9								
10	Starting Track							
11	Control							

The PLAY AUDIO TRACK RELATIVE(12) command (Table 13-11) requests that the device begin an audio playback operation. The command function (Immed and SOTC bits) and the output of audio signals shall be as specified by the settings of the mode parameters (see 13.3.3.1.). See the PLAY AUDIO TRACK RELATIVE(10) command for a description of the fields in this command.

13.2.8. READ CD-ROM CAPACITY Command

Table 13-12: READ CD-ROM CAPACITY Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (25h)							
1	Logical Unit Number			Reserved				RelAdr
2	(MSB)							
3								
4	Logical Block Address							
5								
6								
7	Reserved							
8	Reserved							
9	Reserved							
10	Control							
11								

The READ CD-ROM CAPACITY command (Table 13-12) provides a means for the initiator to request information regarding the capacity of the logical unit.

IMPLEMENTORS NOTE: This command has the same operation code (25h) as the READ CAPACITY command (see 8.2.7). The general function is the same but the exact definitions of the returned logical block address is modified to allow returning a possibly inexact value (but one with a known error bound) based on CD-ROM table of contents data. For many CD-ROM devices, this implementation allows a quicker response.

A partial medium indicator (PMI) bit of zero indicates that the information returned in the READ CD-ROM CAPACITY data shall be the logical block address and block length (in bytes) of the last valid logical block of the logical unit for seek operations. The logical block address returned shall be greater than or equal to the last readable or playable block. If greater, this address may be in a transition area beyond the last valid logical block for read or audio play operations. The value returned shall not be more than 75 sectors (MSF units) greater than the logical block address of the last readable or playable block. (This value arises because the CD-ROM Table of Contents lead-out track location has a +/- 75 sector tolerance when the lead-out track is encoded as an audio track.) The logical block address in the command descriptor block shall be set to zero for this option.

A PMI bit of one indicates that the information returned shall be the logical block address and block length (in bytes) of the last logical block address after which a substantial delay in data transfer will be encountered. This logical block address shall be greater than or equal to the logical block address specified in the command descriptor block. On CD-ROM media, this is interpreted to be the last readable or playable logical block of the information area containing or immediately following the specified logical block address.

IMPLEMENTORS NOTE: This option may take several seconds to complete on CD-ROM media.

See 8.1.10 for a description of the relative address bit (RelAdr).

Eight bytes of READ CD-ROM CAPACITY data (Table 13-13) shall be sent during the DATA IN phase of the command.

Table 13-13: READ CAPACITY Data Format

Byte	Description
0	(MSB)
3	Logical Block Address (LSB)
4	(MSB)
7	Block Length (LSB)

13.2.9. READ HEADER Command

Table 13-14: READ HEADER Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (44h)							
1	Logical Unit Number			Reserved			MSF	Reserved
2	(MSB)							
3								
4	Logical Block Address							
5								
6	Reserved							
7	(MSB)							
8	Allocation Length							
9	(LSB)							
	Control							

The READ HEADER command (Table 13-14) requests that the device return the CD-ROM data block address header (Table 13-15) of the requested logical block.

See 13.1.5. for a description of the MSF bit.

The logical block address field specifies the logical block at which the read header operation shall begin.

See the READ command for exception handling. If the logical block size is other than the physical block size, it shall be mapped into the appropriate physical block from which the data would have been read.

Table 13-15: READ HEADER Data Format

Bit	7	6	5	4	3	2	1	0
Byte								
0	CD-ROM Data Mode							
1	Reserved							
2	Reserved							
3	Reserved							
4	(MSB)							
7	Absolute CD-ROM Address						(LSB)	

The CD-ROM data mode field specifies the CD-ROM Data Mode of the logical blocks in this sector of data. The values in this field are defined in Table 13-16.

Table 13-16: CD-ROM Data Mode Codes

CD-ROM Mode	User Data Field Contents (2048 bytes)	Auxiliary Field Contents (288 bytes)
00h	All bytes zero	All bytes zero
01h	User Data	L-EC symbols
02h	User Data	User Data
03h - FFh	Reserved	Reserved

If the MSF bit is zero, the absolute address field gives the logical block address of the first logical block in the physical sector where the data for the requested logical block address is found. If the MSF bit is one, the absolute address field gives the MSF address of the sector where the data for the requested logical block address is found. (See 13.1.5.)

13.2.10. READ SUB-CHANNEL Command

Table 13-17: READ SUB-CHANNEL Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (42h)							
1	Logical Unit Number			Reserved			MSF	Reserved
2	Reserved	SubQ	Reserved					
3	Sub-channel Data Format							
4	Reserved							
5	Reserved							
6	Track Number							
7	(MSB)	Allocation Length						(LSB)
8								
9	Control							

The READ SUB-CHANNEL command (Table 13-17) requests that the target return the requested sub-channel data plus the state of audio play operations.

IMPLEMENTORS NOTE: Sub-channel data returned by this command may be from the last appropriate sector encountered by a current or previous media accessing operation. When there is no current audio play operation, the target may access the media to read the sub-channel data. The target is responsible that the data returned are current and consistent. For example with sub-channel data format 0, the International Standard Recording Code (ISRC) data reported must have been read from the same track as the reported current position data.

See 13.1.5. for a description of the MSF bit.

The Sub Q bit set to one requests the target return the Q sub-channel data. The Sub Q bit set to zero requests that no sub-channel data be returned. This shall not be considered an error.

NOTE: The other bits in this byte are reserved for future standardization when they may be defined to request other sub-channel data.

The Sub-channel Data Format field specifies the returned Sub Channel Data. If this field is 00h, Sub-Q Channel data is returned. If this field is 01h, 02h or 03h, the requested Sub-Q data item is returned.

Table 13-18: Sub-channel Data Format Codes

Format Code	Returned data
00h	Sub-Q Channel data
01h	CD-ROM Current Position
02h	Media Catalog Number (UPC/Bar Code)
03h	Track International-Standard-Recording-Code (ISRC)
04h - EFh	Reserved
F0h - FFh	Vendor Specific

The Track Number field specifies the track from which ISRC data is read. This field must have a value between 01h and 63h (99bcd), and is valid only when the Sub-channel Data Format field is 03h. In this case, the target returns ISRC data for this track.

The READ SUB-CHANNEL command data formats (Tables 13-19, 13-23, 13-24, and 13-25) consist of a four-byte header followed by a sub-channel data block. The header contains the audio status byte and the sub-channel data length field. If the SUB-Q bit is zero, the target shall not return the sub-channel data block, in this case the sub-channel data length is 0.

Table 13-19: Sub-Q Channel Data Format

Bit	7	6	5	4	3	2	1	0
Byte								
	Sub-channel Data Header							
0	Reserved							
1	Audio Status							
2	(MSB)	Sub-channel Data Length						(LSB)
3								
	Sub-Q Channel Data Block							
4	Sub Channel Data Format code (00h)							
5	ADR				Control			
6	Track Number							
7	Index Number							
8	(MSB)	Absolute CD-ROM Address						(LSB)
11								
12	(MSB)	Track Relative CD-ROM Address						(LSB)
15								
16	MCVal	Reserved						
17	(MSB)	Media Catalog Number (UPC/Bar Code)						(LSB)
31								
32	TCVal	Reserved						
33	(MSB)	Track International-Standard-Recording-Code (ISRC)						(LSB)
47								

The audio status field indicates the status of audio play operations. The audio status values are defined in Table 13-20. Audio status values greater than zero are returned only to the initiator that requested the last audio play operation. Audio status values 13h and 14h return information on previous audio operations; they are returned only once after the condition has occurred. If another audio play operation is not requested, the audio status returned for subsequent READ SUB-CHANNEL commands is 15h.

Table 13-20: Audio Status Codes

Status	Description
00h	Audio status byte not supported or not valid
11h	Audio play operation in progress.
12h	Audio play operation paused.
13h	Audio play operation successfully completed.
14h	Audio play operation stopped due to error.
15h	No current audio status to return

The Sub-Channel data length specifies the length in bytes of the following sub-channel data block. A sub-channel data length of zero indicates that no sub-channel data block is included in the returned data.

IMPLEMENTORS NOTE: Usual values for Sub-channel data length are 0, 12, 20, 28 and 44 bytes. Sub-channel data length does not include the Sub Channel Header.

The Sub-Q Channel data block consists of control data (bytes 4-5), current position data (bytes 6 - 15) and identification data (bytes 16 - 47). The control data and current position data is obtained from the Q sub-channel information of the current block. Identification data may be reported that was obtained from a previous block. If identification data is reported, the data shall be valid for the sector addressed by the current position data.

(1) If an audio play operation is proceeding in the background, position data for the last sector played shall be reported.

(2) In other cases, for instance after a READ command, the target may either report position data for the last sector processed for that operation or may report position data from the sector at the current read head position.

IMPLEMENTORS NOTE: When the type of information encoded in the Q sub-channel of the current sector is the media catalog number or ISRC; the track, index, and address fields should be extrapolated from the previous sector.

The ADR field gives the type of information encoded in the Q sub-channel of this block, as shown in Table 13-21.

Table 13-21: ADR Sub-channel Q Field

ADR Code	Description
0h	Sub-channel Q mode information not supplied.
1h	Sub-channel Q encodes current position data. (i.e. track, index, absolute address, relative address.)
2h	Sub-channel Q encodes media catalog number.
3h	Sub-channel Q encodes ISRC.
4h - Fh	Reserved

The control bits are defined in Table 13-22.

Table 13-22: Sub-channel Q Control Bits

Bit	equals zero	equals one
0	Audio without pre-emphasis	Audio with pre-emphasis
1	Digital copy prohibited	Digital copy permitted
2	Audio track	Data track
3	Two channel audio	Four channel audio

The track number specifies the current track number.

The index number specifies the index number in the current track.

The absolute CD-ROM address field gives the current location relative to the logical beginning of the media. If the MSF bit is zero, this field is a logical block address. If the MSF bit is one, this field is an absolute MSF address. (See 13.1.5.)

The track relative CD-ROM address field gives the current location relative to the logical beginning of the current track. If the MSF bit is zero, this field is a track relative logical block address. (If the current block is in the pre-gap area of a track, this will be a negative value, expressed as a two's complement number. See 13.1.5). If the MSF bit is one, this field is the relative MSF address from the Q sub-channel.

A media catalog valid (MCVal) bit of one indicates that the media catalog number field is valid. A MCVal bit of zero indicates that the media catalog number field is not valid.

The media catalog number field contains the identifying number of this media according to the uniform product code values (UPC/EAN Bar Coding) expressed in ASCII. Non-zero values in this field are controlled by the Uniform Product Code Council and the European Article Number Council (see Section 2). A value in this field of all ASCII zeros indicates that the media catalog number is not supplied.

The track code valid (TCVal) bit of one indicates that the track ISRC field is valid. A TCVal bit of zero indicates that the track International-Standard-Recording-Code (ISRC) field is not valid.

The track ISRC field contains the identifying number of this media according to the ISRC standards (DIN-31-621) expressed in ASCII.

Table 13-23: CD-ROM Current Position Data Format

Bit Byte	7	6	5	4	3	2	1	0
	Sub-channel Data Header							
0	Reserved							
1	Audio Status							
2	(MSB)	Sub-Channel Data Length						(LSB)
3								
	CD-ROM Current Position Data Block							
4	Sub Channel Data Format code (01h)							
5	ADR				Control			
6	Track Number							
7	Index Number							
8	(MSB)	Absolute CD-ROM Address						(LSB)
11								
12	(MSB)	Track Relative CD-ROM Address						(LSB)
15								

Table 13-24: Media Catalog Number Data Format

Bit Byte	7	6	5	4	3	2	1	0
	Sub-channel Data Header							
0	Reserved							
1	Audio Status							
2	(MSB)	Sub-Channel Data Length						(LSB)
3								
	Media Catalog Number Data Block							
4	Sub Channel Data Format code (02h)							
5	Reserved							
6	Reserved							
7	Reserved							
8	MCVal	Reserved						
9	(MSB)	Media Catalog Number (UPC/Bar Code)						(LSB)
23								

If Media Catalog Number data is found, the MCVal bit is set to one. If MCN data is not detected, the MCVal bit is set to zero to indicate the Media Catalog Number field is invalid.

IMPLEMENTORS NOTE: Media Catalog Number data returned by this command with Sub-channel Data Format field code 02h may be from any block which has UPC/Bar Code Q sub-channel data. (This code is constant anywhere in every applicable disc.)

Table 13-25: Track International Standard Recording Code Data Format

Bit	7	6	5	4	3	2	1	0
Byte								
Sub-channel Data Header								
0	Reserved							
1	Audio Status							
2	(MSB)							
3	Sub-Channel Data Length							
	(LSB)							
Track ISRC Data Block								
4	Sub Channel Data Format code (03h)							
5	ADR				Control			
6	Track Number							
7	Reserved							
8	TCVa1		Reserved					
9	(MSB)							
23	Track International-Standard-Recording-Code (ISRC)							
	(LSB)							

If ISRC data is detected, the TCVal bit is set to one. If ISRC data is not detected, the TCVal bit is set to zero to indicate the ISRC field is invalid.

IMPLEMENTORS NOTE: Track ISRC data returned by this command with Sub-Channel Data Format field 03h may be from any block in the specified track which has ISRC data.

13.2.11. READ TOC Command

Table 13-26: READ TOC Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (43h)							
1	Logical Unit Number			Reserved			MSF	Reserved
2	Reserved							
3	Reserved							
4	Reserved							
5	Reserved							
6	Starting Track							
7	(MSB)							
8	Allocation Length							
9	(LSB)							
	Control							

The READ TOC command (Table 13-26) requests the target transfers data from the table of contents to the initiator. The format of the data returned is specified in Table 13-27.

See 13.1.5. for a description of the MSF bit.

The starting track field specifies the starting track number for which the data shall be returned. If this value is zero, the table of contents data shall begin with the first track on the medium. The data is returned in contiguous ascending track number order.

If the starting track field is not valid for the currently installed medium, the command shall be terminated with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST and the additional sense code set to INVALID FIELD IN CDB.

IMPLEMENTORS NOTE: The maximum TOC data length possible on currently available CD-ROM media is 804 bytes, or 100 TOC track descriptors.

Table 13-27: READ TOC Data Format

Bit Byte	7	6	5	4	3	2	1	0
0	(MSB)							
1	TOC Data Length							(LSB)
2	First Track Number							
3	Last Track Number							
	TOC Track Descriptor(s)							
0	Reserved							
1	ADR					Control		
2	Track Number							
3	Reserved							
4	(MSB)							
7	Absolute CD-ROM Address							(LSB)

The TOC data block contains a four byte header followed by zero or more TOC track descriptors.

The TOC data length specifies the length in bytes of the following TOC data that is available to be transferred during the DATA IN phase. The TOC data length value does not include the TOC data length field itself.

The first track number field indicates the first track number in the table of contents.

The last track number field indicates the last track number in the table of contents before the lead-out track number.

IMPLEMENTORS NOTE: The first track number is not required to be one. A disc may start at any valid track number. The track numbers between the first track number and the last track number are required to be in contiguous ascending order, except for the lead-out track.

The ADR field gives the type of information encoded in the Q Sub-channel of the block where this TOC entry was found. The possible ADR values are defined in Table 13-21.

The control field indicates the attributes of the track. The possible control field values are defined in Table 13-22.

The track number field indicates the track number for which the data in the TOC track descriptor is valid. A track number of 0AAh indicates the track descriptor is for the start of the lead-out area.

The absolute CD-ROM address contains the address of the first block with user information for that track number as read from the table of contents. An MSF bit of zero indicates that the absolute CD-ROM address field contains a logical block address. An MSF bit of one indicates the absolute CD-ROM address field contains an MSF address (see 13.1.5).

IMPLEMENTORS NOTE: The starting logical block address value recovered from the TOC has a tolerance of zero for data tracks and plus or minus 75 CD sectors for audio tracks. This tolerance is multiplied by a factor dependent on the logical block Length.

13.3. Parameters for CD-ROM Devices

13.3.1. Diagnostic Parameters

This section defines the descriptors and pages for diagnostic parameters used with CD-ROM devices.

The diagnostic page codes for CD-ROM devices are defined in Table 13-28.

Table 13-28: Diagnostic Page Codes

Page Code	Description	Section
00h	Supported Diagnostic Pages	7.3.1.1
01h - 3Fh	Reserved (for all device type pages)	
40h - 7Fh	Reserved	
80h - FFh	Vendor-specific pages	

13.3.2. Log Parameters

This section defines the descriptors and pages for log parameters used with CD-ROM devices.

The log page codes for CD-ROM devices are defined in Table 13-29.

Table 13-29: Log Page Codes

Page Code	Description	Section
01h	Buffer Over-Run/Under-Run Page	7.3.2.1
03h	Error Counter Page (Read) Page	7.3.2.2
07h	Last n Error Events Page	7.3.2.3
06h	Non-Medium Error Page	7.3.2.4
02h, 04h, 05h, 08h - 2Fh, 3Fh	Reserved	
00h	Supported Log Pages	7.3.2.5
30h - 3Eh	Vendor-specific pages	

13.3.3. Mode Parameters

This section defines the descriptors and pages for mode parameters used with CD-ROM devices.

The mode parameter list, including the mode parameter header and mode block descriptor are defined in 7.3.3.

The medium-type code field is contained in the mode parameter header (see Table 7-61 and 7-62). Table 13-30 defines the medium type values for CD-ROM devices.

Table 13-30: CD-ROM Medium Type Codes

Code Value	Medium Type
00h	Default (only one type supported)
01h	120 mm CD-ROM data only
02h	120 mm CD-DA audio only
03h	120 mm CD-ROM data and audio combined
04h	Reserved
05h	80 mm CD-ROM data only
06h	80 mm CD-DA audio only
07h	80 mm CD-ROM data and audio combined
08h - 7Fh	Reserved
80h - FFh	Vendor unique

The device specific parameter field is contained in the mode parameter header (see Table 7-61, (byte 2) and 7-62 (byte 3)). Table 13-31 defines the device specific parameter field for CD-ROM devices.

Table 13-31: CD-ROM Device Specific Parameter

Bit	7	6	5	4	3	2	1	0
Byte								
	Reserved			DPOFUA	Reserved			EBC

When used with the MODE SELECT command the DPOFUA bit is reserved.

When used with the MODE SENSE command, a DPOFUA bit of zero indicates that the target does not support the DPO and FUA bits. A DPOFUA bit of one indicates that the target supports the DPO and FUA bits (see 8.2.6).

The enable blank check (EBC) bit is reserved.

The density code field is contained in the mode parameter block descriptor (see Table 7-63). Table 13-32 defines the density code values for CD-ROM devices.

Table 13-32: CD-ROM Density Codes

Code Value	Data types to be transferred
00h	Default density code
01h	User Data Only
02h	- 2048 Bytes per physical sector User Data plus Auxiliary Data field
03h	- 2336 bytes per sector 4 Byte Tag Field, User Data plus Auxiliary Data
04h	- 2340 bytes per sector Audio Information Only
05h - 7Fh	- (1/75th of a second per logical block) Reserved
80h - FFh	Vendor specific

IMPLEMENTORS NOTE: The number of bytes per sector specified by this parameter is used with the block Length to map CD-ROM sectors to logical block addresses.

The mode page codes for CD-ROM devices are defined in Table 13-33.

Table 13-33: Mode Page Codes

Page Code	Description	Section
08h	Caching Page	8.3.3.1
0Eh	CD-ROM Audio Control Page	13.3.3.1
0Dh	CD-ROM Page	13.3.3.2
0Ah	Control Mode Page	7.3.3.1
02h	Disconnect-Reconnect Page	7.3.3.2
08h	Medium Types Supported Page	8.3.3.4
09h	Peripheral Device Page	7.3.3.3
01h	Read Error Recovery Page	13.3.3.3
03h - 06h, 0Ch, 0Fh - 1Fh	Reserved	
3Fh	Returns all pages (valid only for the MODE SENSE command)	
20h - 3Eh	Vendor Specific (page format required)	
07h	Verify Error Recovery Page	13.3.3.4

13.3.3.1. CD-ROM Audio Control Parameters

Table 13-34: CD-ROM Audio Control Parameters Page

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved	Page Code (0Eh)					
1	Parameter Length (0Eh)							
2	Reserved				Immed		SOTC	Reserved
3	Reserved							
4	Reserved							
5	APRVa1	Reserved			Format of LBAs / Sec.			
6	(MSB)							
7	Logical Blocks per Second of Audio Playback							
	(LSB)							
8	Reserved				Output Port 0 Channel Selection			
9	Output Port 0 Volume							
10	Reserved				Output Port 1 Channel Selection			
11	Output Port 1 Volume							
12	Reserved				Output Port 2 Channel Selection			
13	Output Port 2 Volume							
14	Reserved				Output Port 3 Channel Selection			
15	Output Port 3 Volume							

The CD-ROM audio control parameters page (Table 13-34) sets the playback modes and output controls for subsequent PLAY AUDIO commands and any current audio playback operation.

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

An immediate (Immed) bit of zero indicates the target shall not send completion status until the audio playback operation is terminated.

An Immed bit of one indicates the target shall send completion status as soon as the playback operation has been started.

IMPLEMENTORS NOTE: It is recommended that a Logical Unit type RESERVE be issued prior to starting audio play operations with an Immed bit of one in any multiple initiator environment.

A stop on track crossing (SOTC) bit of zero indicates the target shall terminate the audio playback operation when the transfer length is satisfied. Multiple tracks shall be played as necessary. Periods of time encoded as audio pause/silence at the beginning of tracks, (index 0,) shall also be played.

A stop on track crossing (SOTC) bit of one indicates the target shall terminate the audio playback operation when the beginning of a following track is encountered.

The audio playback rate valid (APRVAl) bit value of one indicates that the format of logical blocks per second field and the logical blocks per second of audio playback field are valid.

The format of logical blocks addresses per second field gives the multiplier to be used with the logical blocks per second of audio playback, as defined in the following table:

Format of LBAs / Sec value	Multiplier for LBAs / Sec field
0h	1
1h - 7h	Reserved
8h	1/256
9h - Fh	Reserved

IMPLEMENTORS NOTE: This field is provided as a means to return fractional (i.e., non-integral) values in the logical block addresses per second of audio playback. This shall occur when logical block sizes that are not even multiples or divisions of the physical block size are used.

The logical blocks per Second of Audio Playback field gives the relationship between time and the duration of play per logical block address. The value in this field is to be multiplied by the value in Format of LBAs per Second field.

IMPLEMENTORS NOTE: The logical blocks per second of audio playback field and its companion format of LBAs per second field may not be supported by most current CD-ROM devices as a modifiable mode select parameter.

The output port channel selection specifies the audio channels from the disc to which this output port should be connected. See Table 13-35. More than one output port may be connected to an audio channel. More than one audio channel may be connected to an output port.

Table 13-35: Output Port Channel Selection

0000	output port muted
0001	connect audio channel 0 to this output port
0010	connect audio channel 1 to this output port
0100	connect audio channel 2 to this output port
1000	connect audio channel 3 to this output port

The channel volume control indicates the relative volume level for this audio output port. A value of zero indicates the output is muted, a value of FFh indicates maximum volume level.

IMPLEMENTORS NOTE: If volume controls are implemented, the default volume level should be no more than 25% of the maximum level as a personal safety consideration.

13.3.3.2. CD-ROM Device Parameters

Table 13-36: CD-ROM Parameters Page

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved		Page Code (0Dh)				
1	Parameter Length (06h)							
2	Reserved							
3	Reserved				Inactivity Timer Multiplier			
4								
5	- - -	Number of MSF - S Units per MSF - M Unit						- - -
6								
7	- - -	Number of MSF - F Units per MSF - S Unit						- - -

The CD-ROM parameters page (Table 13-36) specifies parameters that affect all CD-ROM data types.

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

The inactivity timer multiplier specifies the length of time that the drive shall remain in the hold track state after completion of a seek or read operation (see Table 13-37).

IMPLEMENTORS NOTE: Higher values in this parameter may have an adverse effect on the drive MTBF, in some implementations.

Table 13-37: Inactivity Timer Multiplier Values

Inactivity Timer Multiplier	Minimum Time in Hold Track State	Inactivity Timer Multiplier	Minimum Time in Hold Track State
0h	Vendor specific	8h	16 Seconds
1h	125 milli Seconds	9h	32 Seconds
2h	250 milli Seconds	Ah	1 Minute
3h	500 milli Seconds	Bh	2 Minutes
4h	1 Second	Ch	4 Minutes
5h	2 Seconds	Dh	8 Minutes
6h	4 Seconds	Eh	16 Minutes
7h	8 Seconds	Fh	32 Minutes

The number of S Units per M Unit field gives the ratio of these MSF address values. For media conforming to the CD-ROM and CD-DA standard, this value is 60.

The number of F Units per S Unit field gives the ratio of these MSF address values. For media conforming to the CD-ROM and CD-DA standard, this value is 75.

13.3.3.3. Read Error Recovery Parameters

Table 13-38: Read Error Recovery Parameters Page

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved	Page Code (01h)					
1	Parameter Length (06h)							
2	Error Recovery Parameter							
3	Read Retry Count							
4	Reserved							
5	Reserved							
6	Reserved							
7	Reserved							

The read error recovery parameters page (Table 13-38) specifies the error recovery parameters the target shall use during any command that performs a data read operation to the media (e.g., READ, READ TOC, etc.).

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

The Error Recovery Parameter bits are defined in 8.3.3.6.

IMPLEMENTORS NOTE: The implementation of error recovery procedures for CD-ROM devices is markedly different from those used for magnetic medium disk drives. At least one level of error correction (i.e. CIRC) is required to transfer the data stream. Therefore the performance of the drive may differ substantially from what would be expected by sending the same error recovery parameters to a magnetic medium device.

The correlation of the error recovery parameter and the bit settings defined for CD-ROM devices is given in Table 13-39. The interpretation of these bit settings for CD-ROM devices is given in Tables 13-40 through 13-45. If the error recovery parameter is set to any other value, the command shall be terminated with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST and the additional sense code set to INVALID FIELD IN PARAMETER LIST.

The read retry count field specifies the number of times that the controller shall attempt its read recovery algorithm.

Table 13-39: Error Recovery Parameter Bit Settings

Bit Byte 2 Value	7	6	5 TB	4 RC	3	2 PER	1 DTE	0 DCR
00			0	0		0	0	0
01			0	0		0	0	1
04			0	0		1	0	0
05			0	0		1	0	1
06			0	0		1	1	0
07			0	0		1	1	1
10			0	1		0	0	0
11			0	1		0	0	1
14			0	1		1	0	0
15			0	1		1	0	1
20			1	0		0	0	0
21			1	0		0	0	1
24			1	0		1	0	0
25			1	0		1	0	1
26			1	0		1	1	0
27			1	0		1	1	1
30			1	1		0	0	0
31			1	1		0	0	1
34			1	1		1	0	0
35			1	1		1	0	1

NOTE: Reserved bits are not shown and shall be zero.

A CIRC recovered data error is defined as a block for which the CIRC based error correction algorithm was unsuccessful for a read attempt, but on a subsequent read operation no error was reported. The number of subsequent read operations is limited to the read retry count. Layered error correction was not used.

A CIRC unrecovered data error is defined as a block for which the CIRC based error correction algorithm was unsuccessful on all read attempts up to the read retry count. Layered error correction was not used.

An L-EC recovered data error is defined as a block for which the CIRC based error correction algorithm was unsuccessful but the layered error correction was able to correct the block within the read retry count.

An L-EC uncorrectable data error is defined as a block which could not be corrected by layered error correction within the read retry count.

Table 13-40: CD-ROM Error Recovery Parameters 00h - 04h

===== Error Recovery Description =====

- 00h The maximum error recovery procedures available are used. If an error occurs which is uncorrectable with the error correction codes (ECC) on the media, data transfer is terminated with CHECK CONDITION status. The block with the error is not transferred. The sense key is set to MEDIUM ERROR. The information bytes give the address of the block where the unrecovered error was detected. Recovered errors are not reported.
- 01h Only retries of the read operation and CIRC are used (layered error correction is not used). Only CIRC unrecovered data errors are reported. If an CIRC unrecovered data error occurs, data transfer is terminated with CHECK CONDITION status. The block with the error is not transferred. The sense key is set to MEDIUM ERROR. The information bytes give the address of the block where the unrecovered error was detected. Recovered errors are not reported.
- 04h The maximum error recovery procedures available are used. Recovered data errors are reported. If a recovered data error occurs, data transfer is not terminated. However, when the data transfer has completed CHECK CONDITION status is reported. The sense key is set to RECOVERED ERROR. The information bytes give the address of the last block where a recovered data error was detected.
- If a data error occurs which is uncorrectable with the ECC information available on the media, data transfer is terminated and CHECK CONDITION status is reported. The block with the error is not transferred. The sense key is set to MEDIUM ERROR. The information bytes give the address of the block where the uncorrectable error was detected.

Table 13-41: CD-ROM Error Recovery Procedures 05h - 07h

===== Error Recovery Description =====

- 05h Only retries of the read operation and CIRC are used (layered error correction is not used). Recovered data errors are reported. If a recovered data error occurs, data transfer is not terminated. However, when the data transfer has completed CHECK CONDITION status is reported. The sense key is set to RECOVERED ERROR. The information bytes give the address of the last block where a CIRC recovered data error was detected.
- If an unrecovered data error occurs data transfer is terminated and CHECK CONDITION status is reported. The block with the error is not transferred. The sense key is set to MEDIUM ERROR. The information bytes give the address of the block where the unrecovered error was detected.
- 06h The maximum error recovery procedures are used. Recovered data errors are reported. If a recovered data error occurs data transfer is terminated and CHECK CONDITION status is reported. The block with the recovered error is not transferred. The sense key is set to RECOVERED ERROR. The information bytes give the address of the block where the recovered data error was detected.
- If a data error occurs which is uncorrectable with the ECC information on the medium, data transfer is terminated with CHECK CONDITION status. The block with the error is not transferred. The sense key is set to MEDIUM ERROR. The information bytes give the address of the block where the uncorrectable error was detected.
- 07h Only retries of the read operation are used (layer error correction is not used). CIRC recovered data errors are reported. If a CIRC recovered data error occurs, data transfer is terminated with CHECK CONDITION status. The block with the recovered error is not transferred. The sense key is set to RECOVERED ERROR. The information bytes give the address of the block where the recovered data error was detected.
- If an CIRC unrecovered data error occurs, data transfer is terminated with CHECK CONDITION status. The block with the error is not transferred. The sense key is set to MEDIUM ERROR. The information bytes give the address of the block where the unrecovered error was detected.

Table 13-42: CD-ROM Error Recovery Procedures 10h - 15h

Error Recovery Description

10h If data transfer can be maintained, the maximum error recovery procedures available are used. (RC = 1.) If an error occurs which is uncorrectable with the error correction codes (ECC) on the media, or is uncorrectable in time to maintain data transfer, the data transfer is not terminated. However, when the data transfer has completed CHECK CONDITION status is reported. The sense key is set to MEDIUM ERROR. The information bytes give the address of the block where the first unrecovered error was detected. Recovered errors are not reported.

11h If data transfer can be maintained, retries of the read operation and CIRC are used (layered error correction is not used). (RC = 1.) Only CIRC unrecovered data errors are reported. If an CIRC unrecovered data error occurs, data transfer is not terminated. However, when the data transfer has completed CHECK CONDITION status is reported. The sense key is set to MEDIUM ERROR. The information bytes give the address of the block where the first unrecovered error was detected. Recovered errors are not reported.

14h If data transfer can be maintained, the maximum error recovery procedures available are used. (RC = 1.) Recovered data errors are reported. If a recovered data error occurs, data transfer is not terminated. However, when the data transfer has completed CHECK CONDITION status is reported. The sense key is set to RECOVERED ERROR. The information bytes give the address of the block where a recovered data error was detected.

If an data error occurs which is uncorrectable with the ECC information available on the media, or is uncorrectable in time to maintain data transfer, the data transfer is not terminated. However, when the data transfer has completed CHECK CONDITION status is reported. The sense key is set to MEDIUM ERROR. The information bytes give the address of the block where the first uncorrectable error was detected. Reporting unrecovered errors takes precedence over reporting recovered errors.

15h If data transfer can be maintained, retries of the read operation and CIRC are used (layered error correction is not used). (RC = 1.) Recovered data errors are reported. If a recovered data error occurs, data transfer is not terminated. However, when the data transfer has completed CHECK CONDITION status is reported. The sense key is set to RECOVERED ERROR. The information bytes give the address of the block where a CIRC recovered data error was detected.

If an unrecovered data error occurs, data transfer is not terminated. However, when the data transfer has completed CHECK CONDITION status is reported. The sense key is set to MEDIUM ERROR. The information bytes give the address of the block where the first unrecovered error was detected. Reporting unrecovered errors takes precedence over reporting recovered errors.

Table 13-43: CD-ROM Error Recovery Procedures 20h - 24h

Error Recovery Description

20h The maximum error recovery procedures available are used. If an error occurs which is uncorrectable with the error correction codes (ECC) on the media, data transfer is terminated with CHECK CONDITION status. The block with the error is transferred. The sense key is set to MEDIUM ERROR. The information bytes give the address of the block where the unrecovered error was detected. Recovered errors are not reported.

21h Only retries of the read operation and CIRC are used (layered error correction is not used). Only CIRC unrecovered data errors are reported. If an CIRC unrecovered data error occurs data transfer is terminated with CHECK CONDITION status. The block with the error is transferred. The sense key is set to MEDIUM ERROR. The information bytes give the address of the block where the unrecovered error was detected. Recovered errors are not reported.

24h The maximum error recovery procedures available are used. Recovered data errors are reported. If a recovered data error occurs data transfer is not terminated. However, when the data transfer has completed CHECK CONDITION status is reported. The sense key is set to RECOVERED ERROR. The information bytes give the address of the last block where a recovered data error was detected.

If a data error occurs which is uncorrectable with the ECC information available on the media data transfer is terminated and CHECK CONDITION status is reported. The block with the error is transferred. The sense key is set to MEDIUM ERROR. The information bytes give the address of the block where the uncorrectable error was detected.

Table 13-44: CD-ROM Error Recovery Procedures 25h - 27h

Error Recovery Description

25h Only retries of the read operation and CIRC are used (layered error correction is not used). Recovered data errors are reported. If a recovered data error occurs, data transfer is not terminated. However, when the data transfer has completed CHECK CONDITION status is reported. The sense key is set to RECOVERED ERROR. The information bytes give the address of the last block where an CIRC recovered data error was detected.

If an unrecovered data error occurs, data transfer is terminated and CHECK CONDITION status is reported. The block with the error is transferred. The sense key is set to MEDIUM ERROR. The information bytes give the address of the block where the unrecovered error was detected.

26h The maximum error recovery procedures are used. Recovered data errors are reported. If a recovered data error occurs, data transfer is terminated and CHECK CONDITION status is reported. The block with the recovered error is transferred. The sense key is set to RECOVERED ERROR. The information bytes give the address of the block where the recovered data error was detected.

If a data error occurs which is uncorrectable with the ECC information on the media, data transfer is terminated with CHECK CONDITION status. The block with the error is transferred. The sense key is set to MEDIUM ERROR. The information bytes give the address of the block where the uncorrectable error was detected.

27h Only retries of the read operation are used (layer error correction is not used). CIRC recovered data errors are reported. If a CIRC recovered data error occurs, data transfer is terminated with CHECK CONDITION status. The block with the recovered error is transferred. The sense key is set to RECOVERED ERROR. The information bytes give the address of the block where the recovered data error was detected.

If an CIRC unrecovered data error occurs, data transfer is terminated with CHECK CONDITION status. The block with the error is transferred. The sense key is set to MEDIUM ERROR. The information bytes give the address of the block where the unrecovered error was detected.

Table 13-45: CD-ROM Error Recovery Procedures 30h - 35h

Error Recovery Description

30h Same as code 10h.

31h Same as code 11h.

34h Same as code 14h.

35h Same as code 15h.

13.3.3.4. Verify Error Recovery Parameters

Table 13-46: Verify Error Recovery Parameters Page

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved		Page Code (07h)				
1	Parameter Length (06h)							
2	Error Recovery Parameter							
3	Verify Retry Count							
4	Reserved							
5	Reserved							
6	Reserved							
7	Reserved							

The verify error recovery parameters page (Table 13-46) specifies the error recovery parameter the target shall use during verify operations.

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

The error recovery parameters for verify operations are as defined by the read error recovery parameters (see 13.3.3.3).

13.4. Glossary for CD-ROM Devices

absolute F field. The sub-division of the S Field part of the absolute physical address from the beginning of the media. One physical sector. Written on CD-ROM media in binary coded decimal notation. The value is from 00bcd to 74bcd. (See MSF Address, Frame.)

absolute M field. The most significant part of the absolute physical address from the beginning of the media. Written on CD-ROM media in binary coded decimal notation. The value is from 00bcd to 74bcd. (See MSF Address.)

absolute S field. The intermediate significant part of the absolute physical address from the beginning of the media. Written on CD-ROM media in binary coded decimal notation. The value is from 00bcd to 59bcd. (See MSF Address.)

bcd. Binary Coded Decimal is the number system used on the physical CD-ROM and CD-DA media. Numbers that use this notation have the 'bcd' suffix attached. A byte has two four-bit values each of which can have a value from 0 to 9. The maximum value is 99bcd (99 decimal). Hence the maximum number of tracks is 99bcd.

Example: 00 01 02 03 ... 08 09 10 11 ... 19 20 21 ... 98 99.

block. The term block refers to the data in one logical block - the number of bytes defined by the logical block length in the mode block descriptor.

blocks per sector. The number of logical blocks read from each CD-ROM physical sector. The value depends on the logical block length as defined in the mode block descriptor. The value for blocks per sector is:

- 1 if the logical block length is 2048,
- 2 if the logical block length is 1024,
- 4 if the logical block length is 512, and
- 8 if the logical block length is 256 bytes.

(Assuming a density code of 1 - 2048 bytes of data per sector.)

CD-DA. Compact Disc - Digital Audio. The standardized media for recording digital audio information. The 'Red Book' defines CD-DA media.

CD-ROM. Compact Disc - Read Only Memory. CD-ROM is a standardized medium for recording digitized audio and digital data. CD-ROM is used to describe media with digital data rather than discs which encode audio only. The 'Yellow Book' defines CD-ROM media.

CD-ROM control field. A four bit field in the Q sub-channel data indicating the type of information encoded on the current track. Indicates audio versus data and the type of audio encoding, etc. The control field is also found in the Table of Contents entries. See Table 13-22.

CD-ROM data mode. A byte in the header of CD-ROM data sectors. This indicates if data is present and if layered error correction information is present. See Table 13-16.

CIRC. Cross Interleaved Reed-Solomon Code. The error detection and correction technique used within small frames of audio or data. The CIRC bytes are present in all CD-ROM Data Modes. The error correction procedure which uses the CIRC bytes is referred to as the CIRC based algorithm. In most CD-ROM drives, this function is implemented in hardware.

frame. A physical sector on CD-ROM media. Also the F field unit of a MSF CD-ROM address. The term frame is also used in the CD-ROM model to describe the amount of data received between synchronization patterns. 98 of these frames make a sector. This sort of frame is referred to as 'small frames' where the meaning is not clear from context.

hold track state. When a CD-ROM device enters the hold track state the optical pick-up is maintained at an approximately constant position on the media. This allows a paused operation to be resumed without latency due to seeking. Rotational latency may be incurred, however.

index. Index Number. An index is a subdivision of a CD-ROM track. A track can have from 1 to 99 index numbers. Index numbers within a track are sequential starting with 1.

lead-in area. The area on a CD-ROM disc prior to track one. The main channel in the lead-in area contains audio or data null information. This area is coded as track zero but is not addressable via the SCSI command set. The Q sub-channel in this area is coded with the table of contents information.

lead-out area. The area on a CD-ROM disc beyond the last information track. The main channel in the lead-out area contains audio or data null information. This area is coded as track AAbcd but is not addressable via the SCSI command set. The READ CD-ROM CAPACITY data is the first logical block address of this area minus one.

L-EC. Layered Error Correction. The error correction technique used with CD-ROM Data Mode one sectors. In such sectors, 276 bytes of the auxiliary data field contain L-EC bytes. The algorithm which uses these bytes to perform error correction is often implemented in the drive firmware.

MSF Address. M Field / S Field / F Field. The physical address written on CD-ROM discs. Expressed as a sector count relative to either the beginning of the medium (absolute) or to the beginning of the current track (relative). As defined by the CD-DA and CD-ROM standards, each F field unit is one physical sector; each S field unit is 75 F field units; each M field unit is 60 S field units. Valid contents of F fields are binary values from 0 through 74. Valid contents of S fields are binary values from 0 through 59. Valid contents of M fields are binary values from 0 through 74.

output port. A means for connecting the audio signal being played to equipment outside the CD-ROM device.

pause area. A transition area at the beginning or end of an audio track encoded with audio silence. This transition area is required where the audio track immediately precedes a data track.

post-gap area. A transition area at the end of a data track encoded with null information. This transition area is required where the data track immediately precedes an audio track.

pre-gap area. A transition area at the beginning of a data track encoded with null information. This transition area is required where the data track immediately follows an audio track.

Red Book. Compact disc digital audio (CD-DA) media standard. Defines the media formats and facilities to support the recording of digitized audio information.

relative F field. The sub-division of the S field part of the relative physical address from the beginning of the current track. One physical sector. Written on CD-ROM medium in binary coded decimal notation. The value is from 00bcd to 74bcd. (See MSF Address, Frame.)

relative M field. The most significant part of the relative physical address from the beginning of the current track. Written on CD-ROM media in binary coded decimal notation. The value is from 00bcd to 74bcd. (See MSF Address.)

relative S field. The intermediate significant part of the relative physical address from the beginning of the current track. Written on CD-ROM medium in binary coded decimal notation. The value is from 00bcd to 59bcd. (See MSF Address.)

sector. Refers to the data contained during one frame time. (On CD-DA medium with two channels of audio this is 1/75th of a second). In the CD-ROM standard document the term block is used for this unit. There may be more than one logical block per sector. Similarly, a single logical block may map to multiple sectors.

sub-channel. CD-ROM and CD-DA media have a main channel and a sub-channel. The sub-channel area has eight parts called P, Q, R, S, T, U, V, and W. The Q-sub-channel contains information useful to the controller and drive, such as the control field and MSF addresses. The data rate of each sub-channel (P, Q, etc.) is 1/192nd of that of the main channel.

TOC. The table of contents has information on the type of disc and the starting address of the tracks. This information is encoded in the Q sub-channel.

track. Track Number. A track is a sub-division of the CD-ROM media. A disc has from one to ninety-nine tracks. The data within a track is always of the same type. A track can contain be either CD-ROM or CD-Audio. A disc can start at any track number.

track relative logical address. An offset from the beginning of the recorded information on a track expressed in units of the logical block length. The value is used to address logical blocks relative to the beginning of a track using the relative MSF address encoded in the CD-ROM Q sub-channel.

transition area. Sectors at the beginning or end of tracks coded with null information. Where required by the media standards, these areas have minimum lengths of 1 or 2 seconds. The maximum lengths are not specified. Transition

areas at the beginning of a track are encoded with Index zero. Addresses within transition areas may not be read.

Yellow Book. A standard (IEC 108) which defines CD-ROM media. The Yellow Book defines the media formats and facilities that support the recording of digital data on Compact Disc media.

14. Scanner Devices

14.1. Model for Scanner Devices

Scanner devices generate a digital representation of two or three dimensional objects (e.g., a page of text, a photograph, or a piece of art). This is accomplished by sensing the amount of light reflected from the object and generating the digital data. The digital data can then be sent across the SCSI bus to an initiator for further processing.

There are two types of scanners in use; in one type, the operations and functions of the scanner are fixed; in the other type, the operations and functions are programmable and need to be set up prior to being used for scanning objects.

The scanner device generates the data and transfers it in accordance with the commands received from the initiator. The contents of the data is vendor specific, therefore the initiator and the scanner must know how to use the contents of the data.

Scanners generate a digital image of an object in a two dimensional plane. The x-axis dimension is along the cross-scan direction which is perpendicular to the direction in which a scan occurs. The y-axis dimension is along the scan direction and is parallel to the direction in which a scan occurs. The coordinates are measured from the upper left hand corner of the two dimensional plane. The x-axis measurement increases in a positive manner going from left to right. The left side of the two dimensional plane (i.e., where x equals zero) is called the base element line. The y-axis measurement increases in a positive manner going from top to bottom. The top side of the two dimensional plane (i.e. where y equals zero) is called the base line. The scanning range encompasses the area in which the scanner can operate. It is the area from the scan line and base line to the maximum x and y position. These conventions are adopted to aid in understanding the fields within the command descriptor blocks and parameters used for scanner devices. As such this is a conceptual model and may not accurately reflect the physical device.

The displacements used for positioning windows is independent of the resolution with which a window is scanned. The measurement of displacements is controlled by the scan measurement mode parameters.

In the event of a scanner automatic creation of sub-windows within a defined window (i.e., the auto bit in the DEFINE WINDOW parameters is one), one of the following responses is appropriate:

1) the initiator may issue a GET WINDOW PARAMETERS command prior to any READ commands

2) if the initiator issues a READ command before issuing a GET WINDOW PARAMETERS command, the target shall return CHECK CONDITION status. The ILI and valid bits in the sense data shall be set to one. The initiator should then issue a GET WINDOW PARAMETERS command. This feature is useful when the scanner has the ability to distinguish between image and text data and to define windows accordingly.

It may occur that a scanner device temporarily does not have resources available to manage a data transfer from the initiator or does not have data available to transfer to the initiator. One of the following responses is appropriate in such a case:

1) A CHECK CONDITION status is returned and the sense key is set to NOT READY with the appropriate additional sense code. This response is applicable to a TEST UNIT READY command.

2) The target disconnects until the resource or data is available, and then re-connects to resume the operation.

3) A BUSY status is returned.

If the scanner device determines that an error or exception condition has occurred while executing the SCSI command from the initiator, a CHECK CONDITION status is returned. A REQUEST SENSE command can then be used to determine additional information regarding the error or exception condition.

14.2. Commands for Scanner Devices

The commands for scanner devices shall be as shown in Table 14-1.

Table 14-1: Commands for Scanner Devices

Command Name	Operation Code	Type	Section	Page
CHANGE DEFINITION	40h	O	7.2.1	
COMPARE	39h	O	7.2.2	
COPY	18h	O	7.2.3	
COPY AND VERIFY	3Ah	O	7.2.4	
GET DATA BUFFER STATUS	34h	O	14.2.1	
GET WINDOW	25h	O	14.2.2	
INQUIRY	12h	M	7.2.5	
LOG SELECT	4Ch	O	7.2.6	
LOG SENSE	4Dh	O	7.2.7	
MODE SELECT(6)	15h	O	7.2.8	
MODE SELECT(10)	55h	O	7.2.9	
MODE SENSE(6)	1Ah	O	7.2.10	
MODE SENSE(10)	5Ah	O	7.2.11	
OBJECT POSITION	31h	O	14.2.3	
READ	28h	M	14.2.4	
READ BUFFER	3Ch	O	7.2.12	
RECEIVE DIAGNOSTIC RESULTS	1Ch	O	7.2.13	
RELEASE UNIT	17h	M	9.2.9	
REQUEST SENSE	03h	M	7.2.14	
RESERVE UNIT	16h	M	9.2.10	
SCAN	1Bh	O	14.2.5	
SET WINDOW	24h	M	14.2.6	
SEND	2Ah	O	14.2.7	
SEND DIAGNOSTIC	1Dh	M	7.2.15	
TEST UNIT READY	00h	M	7.2.16	
WRITE BUFFER	3Bh	O	7.2.17	

Key: M = Command implementation is mandatory.
 O = Command implementation is optional.

All other operation codes for scanner devices are reserved for future standardization.

14.2.1. GET DATA BUFFER STATUS Command

Table 14-2: GET DATA BUFFER STATUS Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (34h)							
1	Logical Unit Number			Reserved				Wait
2	Reserved							
3	Reserved							
4	Reserved							
5	Reserved							
6	Reserved							
7	(MSB)							
8	Allocation Length						(LSB)	
9	Control							

The GET DATA BUFFER STATUS command (Table 14-2) provides a means for the initiator to get information about the data buffer. Information is returned only for window identifiers for which a SCAN command has been received (see 14.1.9).

A wait bit of zero indicates that the target shall respond immediately. A wait bit of one indicates that the target shall wait for image data to be available before returning scan status data.

The data buffer status data format is defined in Table 14-3.

Table 14-3: Data Buffer Status Format

Bit Byte	7	6	5	4	3	2	1	0	
0	(MSB)								
1	Data Buffer Status Length								
2									(LSB)
3	Reserved								Block
	Data Buffer Status Descriptor(s)								
0	Window Identifier								
1	Reserved								
2	(MSB)								
3	Available Data Buffer								
4									(LSB)
5	(MSB)								
6	Filled Data Buffer								
7									(LSB)

The data buffer status length indicates the length in bytes of the following scan status data that is available to be transferred during the DATA IN phase. The data buffer status length does not include itself. The data buffer status data transferred to the initiator includes zero or more data buffer status descriptors. Each descriptor returns information for the window specified by the window identifier.

The block bit specifies the buffering capabilities of the scanner. A block bit of one indicates that the data buffer is full and all image data must be transferred to the initiator before the scan operation resumes. A block bit of zero indicates that data buffer is not full and scan operations can continue with the available data buffer space.

The available data buffer field indicates in bytes, the amount of buffer available for transfers from the initiator. This field is valid only in scanners with the ability to accept data from an initiator for processing.

The filled data buffer field indicates the amount of image data in bytes available for transfer to the initiator.

14.2.2. GET WINDOW Command

Table 14-4: GET WINDOW Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (25h)							
1	Logical Unit Number			Reserved				Single
2	Reserved							
3	Reserved							
4	Reserved							
5	Window Identifier							
6	(MSB)							
7	Transfer Length							
8	(LSB)							
9	Control							

The GET WINDOW command (Table 14-4) provides a means for the initiator to get information about previously defined windows.

A single bit of one specifies that a single window descriptor be returned for the specified window identifier. A single bit of zero specifies that window descriptors be returned for all window identifiers that were defined by a SET WINDOWS command or by the target if the automatic bit was set to one.

The get window data shall consist of a header followed by one or more window descriptors (see Tables 14-5, 14-6, and 14-7). Each window descriptor specifies the location, size and scanning method used for a window.

Table 14-5: Get Window Data Header

Bit Byte	7	6	5	4	3	2	1	0
0	(MSB) Window Data Length							(LSB)
1								
2	Reserved							
3	Reserved							
4	Reserved							
5	Reserved							
6	(MSB) Window Descriptor Length							(LSB)
7								

The window data length specifies the length in bytes of the following data that is available to be transferred. The window data length does not include itself. If the allocation length is not sufficient to return all the get window data the window data length shall not be adjusted to reflect the truncation data.

The window descriptor length specifies the length in bytes of a single window descriptor. Each descriptor shall be of equal length. The first forty-eight bytes are defined in this standard and the remaining bytes in each descriptor are vendor specific.

Table 14-6: Window Descriptor Bytes (1-23)

Bit Byte	7	6	5	4	3	2	1	0
0	Window Identifier							
1	Reserved							Auto
2	(MSB)							
3	X-Axis Resolution							(LSB)
4	(MSB)							
5	Y-Axis Resolution							(LSB)
6	(MSB)							
9	X-Axis Upper Left							(LSB)
10	(MSB)							
13	Y-Axis Upper Left							(LSB)
14	(MSB)							
17	Window Width							(LSB)
18	(MSB)							
21	Window Length							(LSB)
22	Brightness							
23	Threshold							

Table 14-7: Window Descriptor Bytes (24 - 39)

Bit Byte	7	6	5	4	3	2	1	0
24	Contrast							
25	Image Composition							
26	Bits Per Pixel							
27	(MSB)	Halftone Pattern						(LSB)
28	+-							
29	RIF	Reserved				Padding Type		
30	(MSB)	Bit Ordering						(LSB)
31	--							
32	+-							
33	Compression Type							
33	Compression Argument							
34 to 39	-- --	Reserved						-- --
40 to n	-- --	Vendor Specific Parameter Byte(s)						-- --

The window identifier specifies the window defined by the window descriptor. A window is referenced by the window identifier during data transfers and parameter updates.

When used with the SET WINDOW command, an automatic (auto) bit of one indicates that the target is allowed to create sub-windows within the window specified. An auto bit of zero indicates that target is not allowed to create sub-windows.

When used with the GET WINDOW command, an auto bit of zero indicates that the window was defined directly by the SET WINDOW command. A value of one indicates that the window was defined by the target. This is a sub-window within a window defined by a SET WINDOW command.

The window identifiers assigned by the target shall be unique and not currently in use. The parameters for the sub-windows may be retrieved using the GET WINDOW command.

The x-axis resolution field specifies the resolution in the scan line direction. The unit of measure is picture elements (pixels) per inch. A value of zero specifies the default resolution.

The y-axis resolution field specifies the resolution in the base line direction. The unit of measure is scan lines per inch. A value of zero specifies the default resolution.

The x-axis upper left field specifies the x-axis coordinate of the upper left corner of the window. This coordinate is measured from the scan line using the target's current measurement unit divisor (see 14.3.3.1).

The y-axis upper left field specifies the y-axis coordinate of the upper left corner of the window. This coordinate is measured from the base line using target's current measurement unit divisor (see 14.3.3.1).

The window width field specifies the width of window in the scan line direction. The window width is measured using the target's current measurement unit divisor (see 14.3.3.1).

The window length field specifies the length of the window in the base line direction. The window length is measured using the target's current measurement unit divisor (see 14.3.3.1).

The brightness field specifies the level of brightness used to scan the object. A value of zero specifies the default brightness or automatic brightness control if it is supported. Any other value indicates a relative brightness setting with 255 being the highest setting, one being the lowest setting, and 128 being the nominal setting.

The threshold field specifies the threshold at which scan data is converted to binary data. A value of zero specifies the the default threshold or automatic threshold control if it is supported. Any other value indicates relative threshold setting with 255 being the highest setting, one being the lowest setting, and 128 being the nominal setting.

The contrast field specifies the level of contrast used to scan the object. A value of zero specifies the default contrast or automatic contrast control if it is supported. Any other value indicates a relative contrast setting with 255 being the highest setting, one being the lowest setting, and 128 being the nominal setting.

The halftone field specifies the level of halftone at which the scan data is converted to binary data. The values in this field are vendor specific. The halftone field is used in conjunction with the image composition field.

The image composition field specifies the type of scan operation requested. The image composition is defined as shown in Table 14-8.

Table 14-8: Image Composition Codes

Code	Description
00h	Bi-level black & white
01h	Dithered/halftone black & white
02h	Multi-level black & white (gray scale)
03h	Bi-level RGB Color
04h	Dithered/halftone RGB Color
05h	Multi-level RGB Color
06h - FFh	Reserved

The bits per pixel field specifies the number of bits used to represent the intensity of a single color.

A reverse image format (RIF) bit of zero indicates that white pixels are to be indicated by zeros and black pixels are to be indicated by ones. A RIF bit of one indicates that white pixels are to be indicated by ones and black pixels are to be indicated by zeros. The RIF bit is applicable only for images represented by one bit per pixel.

The padding type field specifies how the target shall pad the image data transmitted to the initiator if it is not an integral number of bytes. The padding type is defined in Table 14-9.

Table 14-9: Padding Types

Code	Description
00h	No padding
01h	Pad with 0's to byte boundary
02h	Pad with 1's to byte boundary
03h	Truncate to byte boundary
04h - 07h	Reserved

The bit ordering field specifies the order in which data is transferred to the host from the window. The bit ordering specifies the direction of pixels in a scan line, the direction of scan lines within a window and the image data packing within a byte. The values in this field are vendor specific.

The compression type and compression argument fields specify the compression technique to be applied to the image data (see Table 14-10).

Table 14-10: Compression Types and Arguments

Compression Code	Description	Compression Argument
00h	No compression	Reserved
01h	CCITT Group III, 1 dimensional	Reserved
02h	CCITT Group III, 2 dimensional	K factor
03h	CCITT Group IV, 2 dimensional	Reserved
04h - 0Fh	Reserved	Reserved
10h	Optical Character Recognition (OCR)	Vendor Specific
11h - 7Fh	Reserved	Reserved
80h - FFh	Vendor Specific	Vendor Specific

14.2.3. OBJECT POSITION Command

Table 14-11: OBJECT POSITION Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (31h)							
1	Logical Unit Number			Reserved		Position Function		
2	(MSB)							
3	Count							
4	(LSB)							
5	Reserved							
6	Reserved							
7	Reserved							
8	Reserved							
9	Control							

The OBJECT POSITION command (Table 14-11) provides positioning functions. Absolute as well as relative positioning is provided. A target shall return CHECK CONDITION status and set the sense key to ILLEGAL REQUEST if a positioning function is requested that is not supported.

The position function field specifies the requested function (see Table 14-12).

The count field is used with the position function to specify the displacement of the object.

Table 14-12: Position Function

Position Function	Description
000b	Unload Object
001b	Load Object
010b	Absolute Positioning
011b	Relative Positioning
100b	Rotate Object
101b	Reserved
110b	Reserved
111b	Reserved

Unload Object. This position function specifies that the object shall be positioned for removal. If upon receipt of this command there is no object loaded, the target shall return a GOOD status. This condition shall not be considered as an error. If the target is unable to unload the object (i.e., paper jam or misfeed condition), the target shall return CHECK CONDITION status and set the sense key sense to MEDIUM ERROR.

Load object. This position function specifies that the object is to be loaded and positioned to the base line. If upon receipt of this command there is a object already loaded, the target shall return GOOD status. This condition shall not be considered as an error. If an object is not loaded and the target is unable to load an object, the target shall return CHECK CONDITION status and set the EOM bit to one and the sense key to MEDIUM ERROR.

Absolute Positioning. This position function specifies that the object is to be positioned at a y-axis displacement from the base line. The y-axis displacement is determined using the count field and the target's current measurement unit divisor (see 14.3.3.1). A count field of zero positions the object at the base line.

Any other value in the count field shall cause the target to position the object that number of units in the scan line direction. If there is no object loaded or if the specified y-axis displacement is not achieved the target shall return CHECK CONDITION status and set the EOM bit to one and the sense key to MEDIUM ERROR.

Relative Positioning. This position function specifies that the object is to be positioned at a y-axis displacement relative to the current position. The y-axis displacement is determined using the count field and the target's current measurement unit divisor (see 14.3.3.1). A count field of zero causes no change in position of the object.

A positive value in the count field shall cause the target to position the object that number of units in the scan line direction. If the scan range is

exceeded the target shall return CHECK CONDITION status. The EOM bit is set to one, the ILI bit is set to one, and sense key is set to MEDIUM ERROR. The valid bit is set to one and the information bytes are set to the requested count minus the actual number of units moved.

A negative value (e.g., two's complement notation) in the count field shall cause the target to position the object that number of units toward the base line. If there is no object loaded or if the specified y-axis displacement is not achieved the target shall return CHECK CONDITION status and set the EOM bit to one and the sense key to MEDIUM ERROR. If the base line is encountered the target shall position the object at the base line and return CHECK CONDITION status. The ILI bit is set to one and the sense key is set to MEDIUM ERROR. The valid bit is set to one and the information bytes are set to the requested count minus the actual number of units moved.

Rotate object. This position function specifies that the object is to be rotated in a counter-clockwise direction in units of one one-thousandth of a degree. The count field specifies the number of units that the object is to be moved.

14.2.4. READ Command

Table 14-13: READ Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (28h)							
1	Logical Unit Number			Reserved				
2	Data Type Code							
3	Reserved							
4	(MSB)				Data Type Qualifier			
5								
6	(MSB)				Transfer Length			
7								
8								
9	Control							

The READ command (Table 14-13) requests that the target transfer data to the initiator.

The transfer data type distinguishes between the different types of data

which may be transferred between the initiator and the target. The types of transfers are specified in Table 14-14.

Table 14-14: Data Type Codes

Code	Description
00h	Image
01h	Vendor Specific
02h	Halftone Mask
03h	Gamma Function
04h - 7Fh	Reserved
80h - FFh	Vendor Specific

The data type qualifier field provides a means to differentiate data transfers of the same data type code. The values used in this field are vendor specific.

The transfer length specifies the number of blocks the target shall transfer to the initiator during the DATA IN phase. The block size is the current block size in the mode parameters block descriptor (see 7.3.3). A transfer length of zero, is not considered an error and no data shall be transferred.

If the target transfers less than transfer length blocks a CHECK CONDITION status shall be returned. The ILI bit is set to one, the valid bit is set to one, and the information bytes are set to the requested transfer length minus the actual number of blocks transferred.

This command shall be terminated with a status of RESERVATION CONFLICT if any reservation access conflict (see 9.2.10) exists and no data shall be transferred.

14.2.5. SCAN Command

Table 14-15: SCAN Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (1Bh)							
1	Logical Unit Number			Reserved				
2	Reserved							
3	Reserved							
4	Transfer Length							
5	Control							

The SCAN command (Table 14-15) requests the target begin a scan operation.

The transfer length specifies the length in bytes of the window identifier list that shall be sent during the DATA OUT phase. A transfer length of zero indicates that no data shall be transferred. This condition shall not be considered as an error.

The window identifier list consists of zero or more window identifiers each of which specifies a window to be scanned.

14.2.6. SEND Command

Table 14-16: SEND Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (2Ah)							
1	Logical Unit Number				Reserved			
2	Data Type Code							
3	Reserved							
4	(MSB)							
5	Data Type Qualifier							
6	(MSB)							
7	Transfer Length							
8	(LSB)							
9	Control							

The SEND command (Table 14-16) transfers data from the initiator to the target.

The data type code and data type qualifier are defined in the READ command (see 14.2.5).

The transfer length specifies the number of blocks the target shall transfer from the initiator during the DATA OUT phase. The block size is the current block size in the mode parameters block descriptor (see 7.3.3). A transfer length of zero, is not considered an error and no data shall be transferred.

This command shall be terminated with a status of RESERVATION CONFLICT if any reservation access conflict (see 9.2.10) exists and no data shall be transferred.

14.2.7. SET WINDOW Command

Table 14-17: SET WINDOW Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (24H)							
1	Logical Unit Number			Reserved				
2	Reserved							
3	Reserved							
4	Reserved							
5	Reserved							
6	(MSB)							
7	Transfer Length							
8	(LSB)							
9	Control							

The SET WINDOW command (Table 14-17) provides a means for the initiator to specify one or more windows within the scanning range of the device.

The transfer length specifies the length in bytes of the data that shall be sent during the DATA OUT phase. A transfer length of zero indicates that no window parameters data shall be transferred. This condition shall not be considered as an error.

The window parameters data shall consist of a header followed by one or more window descriptors (see Tables 14-18, 14-6, and 14-7). Each window descriptor specifies the location, size and scanning method used for a window.

Table 14-18: Set Window Data Header

Bit Byte	7	6	5	4	3	2	1	0
0	Reserved							
1	Reserved							
2	Reserved							
3	Reserved							
4	Reserved							
5	Reserved							
6	(MSB)	Window Descriptor Length						
7							(LSB)	

The window descriptor length specifies the length in bytes of a single window descriptor. Each descriptor shall be of equal length. The first forty-eight bytes are defined in this standard and the remaining bytes in each descriptor are vendor specific.

See Tables 14-6 and 14-7 for the definition of a window descriptor.

14.3. Parameters for Scanner Devices

14.3.1. Diagnostic Parameters

This section defines the descriptors and pages for diagnostic parameters used with scanner devices.

The diagnostic page codes for scanner devices are defined in Table 14-19.

Table 14-19: Diagnostic Page Codes

Page Code	Description	Section
00h	Supported Diagnostic Pages	7.3.1.1
01h - 7Fh	Reserved	
80h - FFh	Vendor-specific pages	

14.3.2. Log Parameters

This section defines the descriptors and pages for log parameters used with scanner devices.

The log page codes for scanner devices are defined in Table 14-20.

Table 14-20: Log Page Codes

Page Code	Description	Section
01h	Buffer Over-Run/Under-Run Page	7.3.2.1
03h	Error Counter Page (Read) Page	7.3.2.2
07h	Last n Error Events Page	7.3.2.3
06h	Non-Medium Error Page	7.3.2.4
00h	Supported Log Pages	7.3.2.5
08h - 2Fh	Reserved	
3Fh	Reserved	
30h - 3Eh	Vendor-specific pages	

14.3.3. Mode Parameters

This section defines the descriptors and pages for mode parameters used with scanner devices.

The mode parameter list, including the mode parameter header and mode block descriptor is defined in 7.3.3.

The medium-type code field is contained in the mode parameter header (see Tables 7-61 and 7-62). This field is reserved for scanners devices.

The device specific parameter field is contained in the mode parameter header (see Tables 7-61 and 7-62). This field is reserved for scanners devices.

The density code field is contained in the mode parameter block descriptor (see Table 7-63). This field is reserved for scanners devices.

The mode page codes for scanner devices are defined in Table 14-21.

Table 14-21: Mode Page Codes

Page Code	Description	Section
0Ah	Control Mode Page	7.3.3.1
02h	Disconnect-Reconnect Page	7.3.3.2
03h	Measurement Units Page	14.3.3.1
09h	Peripheral Device Page	7.3.3.3
01h	Reserved	
03h - 08h	Reserved	
08h - 1Fh	Reserved	
3Fh	Returns all pages (valid only for the MODE SENSE command)	
20h - 3Eh	Vendor-specific pages (page format required)	
00h	Vendor Specific (does not require page format)	

14.3.3.1. Measurement Units Page

Table 14-22: Measurement Units Page

Bit	7	6	5	4	3	2	1	0						
Byte	0	PS	Reserved	Page Code (03h)										
1	Parameter Length (06h)													
2	Basic Measurement Unit													
3	Reserved													
4	(MSB)	Measurement Unit Divisor						(LSB)						
5														
6	Reserved													
7	Reserved													

The measurement units page (Table 14-22) specifies the units of measurement used for calculating the displacement of window and for positioning an object. The measurement units are independent of the horizontal and vertical scan resolutions.

The parameter savable (PS) bit is only used with the MODE SENSE command. This bit is reserved for the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

The basic measurement unit field is defined in Table 14-23. Targets shall use inches as the default basic measurement unit.

Table 14-23: Basic Measurement Units

Code	Description
00h	Inch
01h	Millimeter
02h	Point
03h - FFh	Reserved

The measurement unit divisor specifies the number of units needed to equal one basic measurement unit. Targets shall use 1200 as the default measurement unit divisor. If a value of zero is specified the target shall return CHECK CONDITION status and set the sense key to ILLEGAL REQUEST.

IMPLEMENTORS NOTE: A target which does not implement this page or only supports default values uses one twelve-hundredths (1/1200) of an inch as the unit of measure.

14.4. Glossary for Scanner Devices

base element line. A x-axis displacement equal to zero.

base line. A y-axis displacement equal to zero.

beginning-of-medium. A x-axis and y-axis of zero displacement. Alternatively this is being positioned at the intersection of the base and scan lines.

end-of-medium. The maximum x-axis and y-axis displacement.

image. The digital result of a scan.

object. The original or item being scanned.

pixel. Picture-element, is the smallest photo sight in the array.

platen. The surface in which the target is positioned.

scan. An operation that generates a digital image from the reflected light of an object.

scan line. A y-axis displacement from the base line.

scanning range. The total area that a scanner can generate an image from. For two dimensional objects this may correspond to the largest object that can be scanned.

window. All or part of the scanning range of a scanner. A window defines the part of the object scanned.

15. Optical Memory Devices

15.1. Model for Optical Memory Devices

An optical memory device is a is one that can potentially support a variety of optical media, (e.g., read-only, write-once, erasable, or reversible). In several respects an optical memory device is similar to a direct-access device. However optical memory devices typically offer features that are not available with other devices, including very large capacity removable media. These devices often require the functions which are not found in direct-access devices such as logical block update, pre-erasure before writing, or scanning for blank medium and twelve-byte command descriptor blocks.

This standard has specific device types for write-once and CD-ROM devices which also use optical media, but are not capable of supporting a several types of optical media. It is possible to implement devices using write-once media as optical memory devices. It is possible to implement devices using read-only media as optical memory devices, however devices using CD-ROM media have certain unique characteristics and should not be implemented as optical memory devices.

A model of optical memory devices is complicated by the nature of one of its potential advantages, that it can support media which has different characteristics. There are three types of optical media in general use, read-only, write-once and reversible. Read-only media is typically used as for publishing applications requiring dissemination of large amounts of data, since the data can be replicated on a disk at low cost. Write-once media is used in applications that have large backup or archiving requirements. It is also used in applications that have a need for large amounts of on-line reference information. Reversible media is typically used in applications that need large amounts of temporary storage (e.g., graphics workstations) and can take advantage of removable media. In some applications reversible media drives are used in place of direct-access devices.

Reversible media usually needs to be reversed (erased, blanked) before new data can be written. In such cases an erase operation is required before data can be written. Some optical memory devices perform this erase operation implicitly with each write operation which typically will impact the data throughput. Some devices can perform the erase as a separately. The ERASE command may be used to erase areas of the medium with a corresponding increase in data throughput on subsequent write operations. Products using optical media should not be implemented as direct-access devices due to the overhead penalty on performance from the emulation and lack of support in direct-access devices to take advantage of the sophisticated features available in with optical memory devices.

The type of medium supported by the device and the type of medium currently loaded can be determined by examining the MODE SENSE data. One unique feature of optical memory devices is support of media with mixed types (e.g., media with read-only and write-once areas). The INQUIRY command informs the initiator that the target is an optical memory device; the initiator should then determine the medium type from the MODE SENSE data. The initiator needs to be cognizant of medium type since the devices characteristics can change when the media is changed.

Write-once media can have valid data written to a logical block once. This is an important feature where audit trails and permanent archives are needed. Many optical memory devices supporting write-once media have the ability to update a logical block, preserving the previous generation of data. These devices usually provide a means to recover the previous data through use of commands that allow read access to the different generations of data that are stored at the same logical block address.

Defect management can also be performed by use of logical block can be updating in a manner similar to that used by direct-access devices with the REASSIGN BLOCKS command. The advantage of using the updating (which is not supported by direct-access devices) is access to the previous data.

The update operation assigns an alternate physical block to the logical block while simultaneously writing the data to the block. Commands are provided to allow the recovery of previous generations of updated blocks.

An important requirement in dealing with optical media is determining the which logical blocks are contain written data and which are blank. A blank logical block is one that is properly initialized so that data written to it can be recovered. The logical blocks usually have a flag associated with each that indicates whether they have been written or not.

Many of the strategies used to manage write once and erasable media depend on being able to determine the boundary between written and blank areas of the medium. The MEDIUM SCAN command is useful in finding blank areas for subsequent write operations.

15.1.1. Error Reporting

If any of the following conditions occur during the execution of a command the target shall return CHECK CONDITION status. The appropriate sense key and additional sense code should be set. The following table illustrates some error conditions and the applicable sense keys. The table does not provide an exhaustive enumeration of all conditions that may cause the CHECK CONDITION status.

Condition	Sense Key
Invalid logical block address	ILLEGAL REQUEST
Unsupported option requested	ILLEGAL REQUEST
Target reset or medium change since last command from this initiator	UNIT ATTENTION
Self diagnostic failed	HARDWARE ERROR
Unrecovered read error	MEDIUM ERROR HARDWARE ERROR
Recovered read error	RECOVERED ERROR
Overrun or other error that might be resolved by repeating the command	ABORTED COMMAND
Attempt to write on write protected medium	DATA PROTECT
Attempt to read a blank or previously unwritten block.	BLANK CHECK
Attempt to write a previously written block and blank checking is enabled.	BLANK CHECK
Attempt to write on read-only medium.	DATA PROTECT

In the case of an invalid logical block address, the sense data information field shall be set to the logical block address of the first invalid address.

In the case of an attempt to read a blank or previously unwritten block, the sense data information field shall be set to the logical block address of the first blank block encountered. The data read up to that block shall be transferred.

In the case of an attempt to write a previously written block and blank checking is enabled, the sense information field shall be set to the logical block address of the first non-blank block encountered.

15.2. Commands for Optical Memory Devices

The commands for optical memory devices shall be as shown in Table 15-1 and 15-2.

Table 15-1: Commands for Optical Memory Devices(1)

Command Name	Operation Code	Type	Section	Page
CHANGE DEFINITION	40h	O	7.2.1	
COMPARE	39h	O	7.2.2	
COPY	18h	O	7.2.3	
COPY AND VERIFY	3Ah	O	7.2.4	
ERASE(10)	2Ch	O	15.2.1	
ERASE(12)	ACh	O	15.2.2	
FORMAT UNIT	04h	O	8.2.1	
INQUIRY	12h	M	7.2.5	
LOCK UNLOCK CACHE	36h	O	8.2.2	
LOG SELECT	4Ch	O	7.2.6	
LOG SENSE	4Dh	O	7.2.7	
MEDIUM SCAN	38h	O	15.2.3	
MODE SELECT(6)	15h	O	7.2.8	
MODE SELECT(10)	55h	O	7.2.9	
MODE SENSE(6)	1Ah	O	7.2.10	
MODE SENSE(10)	5Ah	O	7.2.11	
PRE-FETCH	34h	O	8.2.3	
PREVENT ALLOW MEDIUM REMOVAL	1Eh	O	8.2.4	
READ(6)	08h	O	8.2.5	
READ(10)	28h	M	8.2.6	
READ(12)	A8h	O	15.2.4	
READ BUFFER	3Ch	O	7.2.12	
READ CAPACITY	25h	M	8.2.7	
READ DEFECT DATA(10)	37h	O	8.2.8	
READ DEFECT DATA(12)	B7h	O	15.2.5	
READ GENERATION	29h	O	15.2.6	
READ LONG	3Eh	O	8.2.9	
READ UPDATED BLOCK	2Dh	O	15.2.7	
REASSIGN BLOCKS	07h	O	8.2.10	
RECEIVE DIAGNOSTIC RESULTS	1Ch	O	7.2.13	
RELEASE	17h	M	8.2.11	
REQUEST SENSE	03h	M	7.2.14	
RESERVE	16h	M	8.2.12	
REZERO UNIT	01h	O	8.2.13	

Table 15-2: Commands for Optical Memory Devices(2)

Command Name	Operation Code	Type	Section	Page
SEARCH DATA EQUAL(10)	31h	0	8.2.14.1	
SEARCH DATA EQUAL(12)	B1h	0	15.2.8	
SEARCH DATA HIGH(10)	30h	0	8.2.14.2	
SEARCH DATA HIGH(12)	B0h	0	15.2.8	
SEARCH DATA LOW(10)	32h	0	8.2.14.3	
SEARCH DATA LOW(12)	B2h	0	15.2.8	
SEEK(6)	0Bh	0	8.2.15	
SEEK(10)	2Bh	0	8.2.15	
SEND DIAGNOSTIC	1Dh	M	7.2.15	
SET LIMITS(10)	33h	0	8.2.16	
SET LIMITS(12)	B3h	0	15.2.9	
START STOP UNIT	1Bh	0	8.2.17	
SYNCHRONIZE CACHE	35h	0	8.2.18	
TEST UNIT READY	00h	M	7.2.16	
UPDATE BLOCK	3Dh	0	15.2.10	
VERIFY(10)	2Fh	0	15.2.11	
VERIFY(12)	AFh	0	15.2.12	
WRITE(6)	0Ah	0	8.2.20	
WRITE(10)	2Ah	M	15.2.13	
WRITE(12)	AAh	0	15.2.14	
WRITE AND VERIFY(10)	2Eh	0	15.2.15	
WRITE AND VERIFY(12)	AEh	0	15.2.16	
WRITE BUFFER	3Bh	0	7.2.17	
WRITE LONG	3Fh	0	8.2.23	

Key: M = Command implementation is mandatory.
 0 = Command implementation is optional.

The following codes are vendor specific: 20h, 21h, 22h, 23h, and C0h through FFh. All remaining codes for optical memory devices are reserved for future standardization.

15.2.1. ERASE(10) Command

Table 15-3: ERASE(10) Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (2Ch)							
1	Logical Unit Number			Reserved		ERA	Reserved	RelAdr
2	(MSB)							
3								
4	Logical Block Address							
5								
6	(LSB)							
7	Reserved							
8	(MSB)							
9	Transfer Length							
	(LSB)							
	Control							

The ERASE(10) command (Table 15-3) requests that the target erase the specified number of blocks starting at the specified logical block address on the medium. As used here, "erased" means either the medium shall be erased, or a pattern shall be written on the medium that appears to the target as no data present. The blocks erased shall be considered blank for purposes of blank checking (see 15.1). The previous data recorded on the medium, if any, shall not be recoverable.

The erase all (ERA) bit set to one indicates that all remaining blocks on the medium shall be erased. If the ERA bit is set to one and if the number of blocks is not zero, the target shall return CHECK CONDITION, and the sense key shall be set to ILLEGAL REQUEST, with an additional sense code of INVALID FIELD IN CDB.

See 8.2.2 for a description of the RelAdr bit and logical block address field.

The transfer length specifies the number of contiguous logical blocks that shall be erased when the ERA bit is zero. If the ERA bit is zero a transfer length of zero indicates that no blocks shall be erased. This condition shall not be considered as an error and no data shall be erased. Any other value indicates the number of logical blocks that shall be erased.

This command shall be terminated with a status of RESERVATION CONFLICT if any reservation access conflict (see 8.2.12) exists and no data shall be erased.

IMPLEMENTORS NOTE: This command allows the user to separate the erase and write operations. This may increase system performance in certain applications.

15.2.2. ERASE(12) Command

Table 15-4: ERASE(12) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (ACh)							
1	Logical Unit Number		Reserved		ERA	Reserved	RelAdr	
2	(MSB)							
3								
4	Logical Block Address							
5								
6	(MSB)							
7								
8	Transfer Length							
9								
10	Reserved							
11	Control							

The ERASE(12) command (Table 15-4) requests that the target erase the specified number of blocks starting at the specified logical block address on the medium.

See the ERASE(10) command (15.2.1) for a description of the fields in this command.

15.2.3. MEDIUM SCAN Command

Table 15-5: MEDIUM SCAN Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (38h)							
1	Logical Unit Number			WBS	ASA	RSD	PRA	RelAdr
2	(MSB)							
3								
4	Logical Block Address							
5	(LSB)							
6	Reserved							
7	Reserved							
8	Parameter List Length							
9	Control							

The MEDIUM SCAN command (Table 15-5) requests the target to scan the medium for a contiguous set of written or blank logical blocks.

A written block search (WBS) bit of zero indicates that the scan is for blank blocks. A WBS bit of one indicates that the scan is for written blocks.

An advanced scan algorithm (ASA) bit of zero indicates that the scan area is scanned in sequential order (as selected by the RSD bit). An ASA bit of one indicates to the target that the written and blank areas within the scan area form contiguous extents (as opposed to scattered blocks). This indication is advisory to the target.

IMPLEMENTORS NOTE: The purpose of this bit is to allow the target to use a more advanced algorithm (such as a binary search) to locate the requested blocks.

A reverse scan direction (RSD) bit of zero indicates the scan shall begin with the first logical block of the scan area. A RSD bit of one indicates the scan shall begin with the last logical block of the scan area.

A partial results acceptable (PRA) bit of zero indicates that the scan shall not be considered satisfied until a contiguous set of blocks is found within the scan area which is greater than or equal in size to the number of blocks requested, and meets the other criteria specified in the command descriptor block. A PRA bit of one indicates that the scan may be satisfied by a contiguous set of blocks within the scan area which is less than the number of blocks requested, and meets the other criteria specified in the command descriptor block. In this case, the target shall determine the largest contiguous set of blocks in the scan area for reporting to the initiator.

See 8.2.2 for a description of the RelAdr bit and logical block address field.

The parameter list length specifies the length in bytes of the parameter list that shall be transferred during the DATA OUT phase. A parameter list length of zero indicates that the number of blocks requested field has a value of one, and the number of blocks to scan field has a value of zero. This condition shall not be considered an error. The contents of the parameter list are specified in Table 15-6.

A link bit of zero indicates a non-linked command and if the scan is satisfied, the command shall be terminated with a CONDITION MET status. A REQUEST SENSE command can then be issued to determine the starting logical block address of the area that meets the request. If the scan is not satisfied and no error occurs, the command shall be terminated with GOOD status.

A link bit of one indicates a command is linked to the MEDIUM SCAN command and if the search is satisfied, CONDITION MET status is returned and the next command is executed. If the RelAdr bit in the next command is one, the logical block address of the next command is used as a displacement from the logical block address at which the search was satisfied. If a linked scan is not satisfied, the command is terminated with a CHECK CONDITION status. A REQUEST SENSE command may then be issued.

A REQUEST SENSE command following a satisfied MEDIUM SCAN command shall:

(1) Return a sense key of EQUAL if the scan was satisfied by a contiguous set of blocks equal in size to the number of blocks requested. If the scan was satisfied by a contiguous set of blocks not equal in size to the number of blocks requested, then a sense key of NO SENSE shall be returned.

(2) Return the valid bit set to one.

(3) Return the logical block address of the first logical block of the contiguous set of blocks that satisfied the scan criteria in the information bytes.

(4) Return the number of contiguous logical blocks meeting the scan criteria in the command specific information bytes.

A REQUEST SENSE command following an unsatisfied MEDIUM SCAN command shall:

(1) Return a sense key of NO SENSE if no errors occurred during the command execution.

(2) Return the valid bit set to zero.

Table 15-6: MEDIUM SCAN Parameter List

Bit	7	6	5	4	3	2	1	0
Byte								
0	(MSB)							
1								---
2								---
3								---
4	(MSB)							
5								---
6								---
7								(LSB)

The number of blocks requested field specifies the number of blocks that meet the specified requirements. The number of blocks requested field, if set to zero, indicates that the scan shall not take place. This shall not be considered an error condition.

The number of blocks to scan field specifies the length in blocks of the area to be scanned on the medium. The number of blocks to scan field, if set to zero, indicates that the scan shall continue for all remaining blocks on the medium or until the scan is satisfied. See 15.1.3 for a description of error handling.

15.2.4. READ(12) Command

Table 15-7: READ(12) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (A8h)							
1	Logical Unit Number		DPO	FUA	Reserved		RelAdr	
2	(MSB)							
3								
4	Logical Block Address							
5	(LSB)							
6	(MSB)							
7								
8	Transfer Length							
9	(LSB)							
10	Reserved							
11	Control							

The READ(12) command (Table 15-7) requests that the target transfer data to the initiator from the medium. See the READ(10) command (8.2.6) for a complete description of the fields in this command.

15.2.5. READ DEFECT DATA(12) Command

Table 15-8: READ DEFECT DATA(12) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (B7h)							
1	Logical Unit Number			Plist	GList	Defect List Format		
2	Reserved							
3	Reserved							
4	Reserved							
5	Reserved							
6	(MSB)							
7								
8	Allocation Length							
9	(LSB)							
10	Reserved							
11	Control							

The READ DEFECT DATA(12) command (Table 15-8) requests that the target transfer the medium defect data to the initiator. See the READ DEFECT DATA(10) command (8.2.8) for a description of the fields in this command.

The READ DEFECT DATA(12) list header (Table 15-9) contains an eight byte header, followed by zero or more defect descriptors.

Table 15-9: READ DEFECT DATA(12) LIST HEADER

Bit Byte	7	6	5	4	3	2	1	0
0	Reserved							
1	Reserved			PList	GList	Defect List Format		
2	Reserved							
3	Reserved							
4	(MSB)							
5								
6	Defect List Length							
7	(LSB)							
Defect Descriptors								
0-n								

See the description of the READ DEFECT DATA defect list (Table 8-19) for a description of the fields in this header.

15.2.6. READ GENERATION Command

Table 15-10: READ GENERATION Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (29h)							
1	Logical Unit Number			Reserved				RelAdr
2	(MSB)							
3								
4	Logical Block Address							
5								(LSB)
6	Reserved							
7	Reserved							
8	Allocation Length							
9	Control							

The READ GENERATION command (Table 15-10) requests that the target transfer to the initiator the maximum generation address for the logical block specified.

See 8.2.2 for a description of the RelAdr bit and logical block address field.

The READ GENERATION data is defined in Table 15-11.

Table 15-11: Maximum Generation Data Block

Bit Byte	7	6	5	4	3	2	1	0
0	(MSB)							
1	Maximum Generation Address							(LSB)
2	Reserved							
3	Reserved							

The maximum generation address field defines the maximum generation address available for the logical block address specified.

15.2.7. READ UPDATED BLOCK(10) Command

Table 15-12: READ UPDATED BLOCK(10) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (2Dh)							
1	Logical Unit Number		DPO	FUA	Reserved		RelAdr	
2	(MSB)							
3								
4	Logical Block Address							
5								
6	(LSB)							
7	Latest (MSB)		Generation Address				(LSB)	
8	Reserved							
9	Control							

The READ UPDATED BLOCK(10) command (Table 15-12) requests that the target transfer data to the initiator from the specified generation and logical block.

See 8.2.6 for a description of the cache control bits (DPO and FUA). See 8.2.2 for a description of the RelAdr bit and logical block address field.

One block of data is transferred during the DATA IN phase.

The latest bit determines the meaning of the generation address field. A latest bit of zero indicates that the generation address is specified relative to the first generation of the block; generation address zero specifies the first generation. Increasing generation addresses specify later generations.

A latest bit of one indicates that the generation address is specified relative to the latest generation of the block; generation address zero specifies the most recent generation. Increasing generation addresses specify earlier generations.

If the requested generation does not exist, the command shall be terminated

with CHECK CONDITION status. The sense key shall be set to BLANK CHECK, and the additional sense code shall be set to GENERATION DOES NOT EXIST.

This command shall be terminated with a status of RESERVATION CONFLICT if any reservation access conflict (see 8.2.12) exists and no data shall be transferred.

15.2.8. SEARCH DATA(12) Commands

Table 15-13: SEARCH DATA(12) Commands

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (B0h, B1h, B2h)							
1	Logical Unit Number		Invert	Reserved		SpnDat	RelAdr	
2	(MSB)							
3								
4	Logical Block Address							
5	(LSB)							
6	(MSB)							
7								
8	Transfer Length							
9	(LSB)							
10	Reserved							
11	Control							

The SEARCH DATA(12) commands (Table 15-13) search one or more logical blocks for equality or inequality to a data pattern. See the SEARCH DATA(10) commands (8.2.14) for a description of the fields in this command.

15.2.9. SET LIMITS(12) Command

Table 15-14: SET LIMITS(12) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (B3h)							
1	Logical Unit Number			Reserved			RdInh	WrInh
2	(MSB)							
3								
4	Logical Block Address							
5	(LSB)							
6	(MSB)							
7								
8	Number of Blocks							
9	(LSB)							
10	Reserved							
11	Control							

The SET LIMITS(12) command (Table 15-14) defines the range within which subsequent linked commands may operate. See the SET LIMITS(10) command (8.2.16) for a description of the fields in this command.

15.2.10. UPDATE BLOCK Command

Table 15-15: UPDATE BLOCK Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (3Dh)							
1	Logical Unit Number			Reserved				RelAdr
2	(MSB)							
3								
4	Logical Block Address							
5								
6	(LSB)							
7	Reserved							
8	Reserved							
9	Reserved							
	Control							

The UPDATE BLOCK command (Table 15-15) requests that the target logically replace data on the medium with the data sent during the DATA OUT phase.

See 8.2.2 for a description of the RelAdr bit and logical block address field.

One block of data is transferred during the DATA OUT phase.

IMPLEMENTORS NOTE: See the MODE Parameters (15.3) for a description of the behavior of the UPDATE BLOCK command relative to the enable blank check (EBC) bit.

This standard does not define the result of a WRITE command issued to a block previously updated by an UPDATE BLOCK command when blank checking is disabled.

IMPLEMENTORS NOTE: It is recommended that the target inhibit this behavior.

A logical block may be updated until the alternate block area is exhausted. The alternate blocks used for the update operation shall not be reported in the READ CAPACITY data. If the alternate block area is exhausted the command shall be terminated with CHECK CONDITION and the sense key shall be set to MEDIUM ERROR, and the additional sense code set to NO DEFECT SPARE LOCATION AVAILABLE.

If the report updated block read parameter specifies posting of recovered errors for a read operation of a logical block that has had a successful update operation performed occurs the command shall terminate with a CHECK CONDITION status. The sense key shall be set to RECOVERED ERROR and the additional sense code set to UPDATED BLOCK READ.

15.2.11. VERIFY(10) Command

Table 15-16: VERIFY Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (2Fh)							
1	Logical Unit Number		DPO	Reserved	BlkVfy	BytChk	RelAdr	
2	(MSB)							
3								
4	Logical Block Address							
5								
6	Reserved							
7	(MSB)							
8	Verification Length							
9	(LSB)							
	Control							

The VERIFY command (Table 15-16) requests that the target verify the data on the medium.

If the MODE SELECT command is implemented, and the Verify Error Recovery Parameters page is also implemented, then the current settings in that page define the verification criteria. If the Verify Error Recovery Parameters page is not implemented, then the verification criteria is vendor unique.

A byte check (BytChk) bit of zero causes a medium verification to be performed with no data comparison. A BytChk bit of one causes a byte-by-byte compare of the data written on the medium and the data transferred from the initiator. The data shall be transferred as it would be for a WRITE command. If the compare is unsuccessful, for any reason, the target shall return CHECK CONDITION status and the sense key shall be set to MISCOMPARE.

A blank verify (BlkVfy) bit of one causes a verification that the blocks are blank.

If the BytChk is one when the BlkVfy bit is one, this shall be considered an error. The target shall return CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST, and the additional sense code set to INVALID FIELD IN CDB.

See 8.2.6 for a description of the cache control bits (DPO). See 8.2.2 for a description of the RelAdr bit and logical block address field.

The verification length specifies the number of contiguous logical blocks of data or blanks that shall be verified. A verification length of zero indicates that no logical blocks shall be verified. This condition shall not be considered as an error. Any other value indicates the number of logical blocks that shall be verified.

15.2.12. VERIFY(12) Command

Table 15-17: VERIFY(12) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (AFh)							
1	Logical Unit Number		DPO	Reserved	BlkVfy	BytChk	RelAdr	
2	(MSB)							
3								
4	Logical Block Address							
5								
6	(MSB)							
7								
8	Verification Length							
9								
10	Reserved							
11	Control							

The VERIFY(12) command (Table 15-17) requests that the target verify the data on the medium. See the VERIFY(10) command (15.2.11) for a description of the fields in this command.

15.2.13. WRITE(10) Command

Table 15-18: WRITE(10) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (2Ah)							
1	Logical Unit Number			DPO	FUA	EBP	Reserved	RelAdr
2	(MSB)							
3								
4	Logical Block Address							
5								
6	Reserved							
7	(MSB)							
8	Transfer Length							
9	(LSB)							
	Control							

The WRITE(10) command (Table 15-18) requests that the target write the data transferred from the initiator to the medium.

An erase by-pass (EBP) bit of zero indicates that the device will default to the normal write operation. An EBP bit of one indicates that the target is allowed to by-pass the erase operation prior to writing the data. When accessing write-once media, the EBP bit shall be considered reserved.

See 8.2.6 for a description of the cache control bits (DPO and FUA). See 8.2.2 for a description of the RelAdr bit and logical block address field.

The transfer length specifies the number of contiguous logical blocks of data that shall be transferred. A transfer length of zero indicates that no data shall be transferred. This condition shall not be considered as an error and no data shall be written. Any other value indicates the number of logical blocks that shall be transferred.

This command shall be terminated with a status of RESERVATION CONFLICT if any reservation access conflict (see 8.2.12) exists and no data shall be written.

15.2.14. WRITE(12) Command

Table 15-19: WRITE(12) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (AAh)							
1	Logical Unit Number		DPO	FUA	EBP	Reserved	RelAdr	
2	(MSB)							
3								
4	Logical Block Address							
5								
6	(MSB)							
7								
8	Transfer Length							
9								
10	(LSB)							
11	Reserved							
12	Control							

The WRITE(12) command (Table 15-19) requests that the target write the data transferred from the initiator to the medium. See the WRITE(10) command (15.2.13) for a description of the fields in this command.

15.2.15. WRITE AND VERIFY(10) Command

Table 15-20: WRITE AND VERIFY(10) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (2Eh)							
1	Logical Unit Number			DPO	Reserved	EBP	BytChk	RelAdr
2	(MSB)							
3								
4	Logical Block Address							
5								
6	(LSB)							
7	Reserved							
8	(MSB)							
9	Transfer Length							
	(LSB)							
	Control							

The WRITE AND VERIFY(10) command (Table 15-20) requests that the target write the data transferred from the initiator to the medium and then verify that the data is correctly written.

If the MODE SELECT command is implemented, and the Verify Error Recovery Parameters page is also implemented, then the current settings in that page define the verification criteria. If the Verify Error Recovery Parameters page is not implemented, then the verification criteria is vendor unique.

A byte check (BytChk) bit of zero causes a medium verification to be performed with no data comparison. A BytChk bit of one causes a byte-by-byte compare of the data written on the medium and the data transferred from the initiator. The data shall be transferred as it would be for a WRITE command. If the compare is unsuccessful, for any reason, the target shall return CHECK CONDITION status and the sense key shall be set to MISCMPARE.

An erase by-pass (EBP) bit of zero indicates that the device will default to the normal write operation. An EBP bit of one indicates that the target is allowed to by-pass the erase operation prior to writing the data. When accessing write-once media, the EBP bit shall be considered reserved.

See 8.2.6 for a description of the cache control bits (DPO). See 8.2.2 for a description of the RelAdr bit and logical block address field.

The transfer length specifies the number of contiguous logical blocks of data that shall be transferred. A transfer length of zero indicates that no logical blocks shall be transferred. This condition shall not be considered as an error and no data shall be written. Any other value indicates the number of logical blocks that shall be transferred.

15.2.16. WRITE AND VERIFY(12) Command

Table 15-21: WRITE AND VERIFY(12) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (AEh)							
1	Logical Unit Number		DPO	Reserved	EBP	BytChk	RelAdr	
2	(MSB)							
3								
4	Logical Block Address							
5								
6	(MSB)							
7								
8	Transfer Length							
9								
10	Reserved							
11	Control							

The WRITE AND VERIFY(12) command (Table 15-21) requests that the target write the data transferred from the initiator to the medium and then verify that the data is correctly written. See the WRITE AND VERIFY(10) command (15.2.15) for a description of the bits in this command.

15.3. Parameters for Optical Memory Devices

15.3.1. Diagnostic Parameters

This section defines the descriptors and pages for diagnostic parameters used with optical memory devices.

The diagnostic page codes for optical memory devices are defined in Table 15-22.

Table 15-22: Diagnostic Page Codes

Page Code	Description	Section
00h	List of Supported Parameters Page	7.3.1.1
01h - 3Fh	Reserved (for all device type pages)	
40h	Translate Address Page	8.3.1
41h - 7Fh	Reserved	
80h - FFh	Vendor specific pages	

15.3.2. Log Parameters

This section defines the descriptors and pages for log parameters used with optical memory devices.

The log page codes for optical memory devices are defined in Table 15-23.

Table 15-23: Log Page Codes

Page Code	Description	Section
01h	Buffer Over-Run/Under-Run Page	7.3.2.1
03h	Error Counter Page (Read) Page	7.3.2.2
05h	Error Counter Page (Verify) Page	7.3.2.2
02h	Error Counter Page (Write) Page	7.3.2.2
07h	Last n Error Events Page	7.3.2.3
00h	List of Supported Parameters Page	7.3.2.5
06h	Non-Media Error Page	7.3.2.4
04h	Reserved	
08h - 2Fh	Reserved	
30h - 3Eh	Vendor Specific	
3Fh	Reserved	

15.3.3. Mode Parameters

This section defines the descriptors and pages for mode parameters used with optical memory devices.

The mode parameter list, including the mode parameter header and mode block descriptor are defined in 7.3.3.

The medium-type code field is contained in the mode parameter header (see Table 7-61 and 7-62). Table 15-24 defines the medium-type code values used for optical memory devices.

Table 15-24: Optical Memory Medium-Type Codes

Code Value	Medium Type
00h	Default (only one medium type supported)
01h	Optical Read Only medium
02h	Optical Write Once medium
03h	Optical Reversible or Erasable medium
04h	Combination of Read only and Write Once medium
05h	Combination of Read only and Reversible or Erasable medium
06h	Combination of Write Once and Reversible or Erasable medium
07h - 7Fh	RESERVED
80h - FFh	Vendor unique

The device specific parameter field is contained in the mode parameter header (see Table 7-61 and 7-62). Table 15-25 defines the device specific parameter values used for optical memory devices.

Table 15-25: Optical Memory Device Specific Parameter

Bit	7	6	5	4	3	2	1	0
	WP	Reserved	DPOFUA	Reserved	Reserved	Reserved	Reserved	EBC

A write protected (WP) bit of zero indicates that the medium is write enabled. A WP bit of one indicates that the medium is write protected.

The DPOFUA bit is reserved for the MODE SELECT command.

For the MODE SENSE command, a DPOFUA bit of zero indicates that the target does not support the DPO and FUA bits. A DPOFUA bit of one indicates that the target supports the DPO and FUA bits (see 8.2.6).

For the MODE SELECT command, an enable blank check (EBC) bit of zero advises the target to disable the blank checking operation of the medium during write operations or during an UPDATE BLOCK command. An EBC bit of one enables blank checking. If a non-blank block is found during a write operation, the command shall be terminated with a CHECK CONDITION status and the sense key shall be set to BLANK CHECK. If a blank block is found during an UPDATE BLOCK command, the command shall be terminated with a CHECK CONDITION status and the sense key shall be set to BLANK CHECK. For read-only media the EBC bit is reserved.

For the MODE SENSE command, an enable blank check (EBC) bit of zero indicates that blank checking of the medium during write operations is disabled. An EBC bit of one indicates that blank checking during write and update operations is enabled. For read-only media the EBC bit is reserved.

The density code field is contained in the mode parameter block descriptor (see Table 7-63). Table 15-26 defines the density code values used for optical memory devices.

Table 15-26: Optical Memory Density Codes

Density Code	Optical Media						
00h	Default density (currently mounted density)						
	Diameter mm (inch)	Type	Sector Size	Tracks	Sides	Servo	Reference Standard
01h	86 (3.5)	R/W	512/1024	12500	1		660-D
02h	89 (3.5)	R/W	512/1024	12500	2		581-D
03h	130 (5.25)	R/W	512/1024	18750	2	CS	607-I
04h	130 (5.25)	W-O	512/1024	30000	2	SS	655-D
05h	130 (5.25)	W-O	512/1024	20000	2	SS	659-D
06h	130 (5.25)	W-O	512/1024	18750	2	CS	457-D
07h	200 (8.0)						407-D
08h	300 (12.0)		1024		2		408-D
09h	356 (14.0)		1024	56350	2		456-D
80h - FFh	Vendor unique						
All others	Reserved						

Key:

	Type	Servo
R/W	Erasable	CS Continuous Servo
W-O	Write Once	SS Sampled Servo
R/O	Read Only	

NOTES:

- (1) The reference standard refers to an ANSI X3B11 project under development. The project may have several documents describing different aspects of the media. Please contact the Secretariat for current information on status and availability.
- (2) RZ modulation.
- (3) 4/15 modulation.
- (4) This project is not presently active.

The mode page codes for optical memory devices are shown in Table 15-27.

Table 15-27: Mode Page Codes

Page Code	Description	Section
08h	Caching Page	8.3.3.1
0Ah	Control Mode Page	7.5.3.1
02h	Disconnect-Reconnect Page	7.3.3.2
0Bh	Medium Types Supported Page	8.3.3.4
06h	Optical Memory Page	15.3.3.1
09h	Peripheral Device Page	7.3.3.3
01h	Read-Write Error Recovery Page	8.3.3.6
07h	Verify Error Recovery Page	8.3.3.8
03h - 05h	Reserved	
0Ch - 1Fh	Reserved	
3Fh	Returns all pages (valid only for the MODE SENSE command)	
00h	Vendor specific (does not require page format)	
20h - 3Eh	Vendor specific	

15.3.3.1. Optical Memory Page

Table 15-28: Optical Memory Page

Bit	7	6	5	4	3	2	1	0
Byte	0	PS	Reserved	Page Code (06h)				
1	Parameter Length (02h)							
2	Reserved							RUBR
3	Reserved							

The optical memory page (Table 15-28) defines parameters for control of optical memory devices.

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

A report updated block read (RUBR) bit of zero indicates the target shall not report an error when a command performs a successful read of a block that has been updated. A RUBR bit of one indicates the target shall terminate a command that performs a read of a block that has been updated with CHECK CONDITION status, the sense key set to RECOVERED ERROR, the additional sense code set to UPDATED BLOCK READ, and the data shall be transferred to the initiator. The default state of the RUBR bit for write-once devices (as reported in the INQUIRY command) shall be one.

15.4. Glossary for Optical Memory Devices

blank. The logical block contains no information detectable by the target, or is written with a pattern that appears to the target as no data present. The logical block is considered ready for a write operation.

generation. Indicates a relative revision level of a logical block that has been updated via the UPDATE BLOCK command. A logical block that has never been updated has only one generation associated with it.

read-only medium. This is medium that may not be written by the initiator. The media contains data prepared in a manner not defined by this standard.

update. To write new data to a logical block without destroying the previous data. After a block has been updated, a normal read returns the most recent generation of the data. Earlier generations are still available after the update.

write-once medium. This is medium that may be written only once by the initiator. Logical blocks on write-once media that have not been written are considered blank. Logical blocks on write-once media that have been written may not be written again.

16. Medium Changer Devices

16.1. Medium Changer Device Model

Medium changer devices mechanize the movement of media to and from primary devices (such as disk or tape drives) and other locations within the range of the medium changer. The medium changer command set is based on a physical model of this functionality.

This command set supports varied physical implementations of the medium changer function. Most of these variations are hidden from the initiator by the high level of function provided by the MOVE MEDIUM and EXCHANGE MEDIUM commands and by the generalized nature of the element addressing scheme. However, initiators may need to be aware of the capabilities of the particular medium changer device. These characteristics and capabilities are conveyed via MODE SENSE pages.

16.1.1. Medium Changer Elements

The medium changer command set uses as its address space the set of physical locations and mechanisms within the scope of a medium changer device. The term element is used throughout this section to refer to one member of the medium changer address space. Each element is a discrete physical entity that may hold zero or one physical unit of media - one disk cartridge, one spool of tape, etc. Element addresses do not extend across multiple physical units of media. Likewise, element addresses are independent of any logical partitioning that the primary device may support within a physical unit of media.

A medium changer is viewed as a set of addressable elements, each of which may contain a unit of media or be used to move a unit of media. Each medium changer element has a unique 16 bit element address. Each element is an instance of one of four classes or element types.

1. MEDIUM TRANSPORT ELEMENT
2. STORAGE ELEMENT
3. IMPORT EXPORT ELEMENT
4. DATA TRANSFER ELEMENT

Units of media (cartridges, cassettes, caddies, etc.) are referred to only indirectly by this device model. Units of media can be moved to or from any of the elements of the medium changer device. The presence of media at the various elements in a medium changer can be sensed. In order to ensure exclusive access to a unit of media, the element where the unit of media is located (the element address) must be reserved.

Elements of the medium transport, import export and (rarely) data transport types may not provide independent storage for medium. The capabilities of a particular medium changer in this respect can be sensed via the device capabilities page of the Mode Sense data. The following hypothetical medium changer implementation illustrates one case of an element not providing independent storage for medium. Consider a medium changer which has a carousel style storage for medium. The import export function could be provided by a port which allows operator access to one of the storage elements. In such a device, the MOVE ELEMENT command from storage element to

import export element would rotate the carousel to align the addressed storage element to the import export position. In this case the import export element does not provide independent storage but rather access to one of the storage elements.

16.1.1.1. Medium Transport Elements

Medium transport elements address the functions of the medium changer device that perform the movement of units of media. Where a medium transport element can serve (even temporarily) as a storage location for medium, the location of each unit of media must have a separate element address.

In larger medium changer devices, the medium movement functions may be performed by multiple independent robotic subsystems. Each of these may have a number of medium transport element addresses. The element addresses within each subsystem shall be contiguous. Any of the element addresses within a subsystem may be used interchangeably in the medium transport element address field of MOVE MEDIUM and EXCHANGE MEDIUM commands. An initiator may determine the capabilities of the medium movement facilities of a medium changer device via the transport geometry MODE SENSE page, 16.3.3.3.

Element address zero is reserved for use in the medium transport element address field of MOVE MEDIUM and EXCHANGE MEDIUM commands to direct the medium changer to use a default or medium changer selected medium transport element.

In some implementations, medium transport elements may be source and/or destination addresses in MOVE MEDIUM and EXCHANGE MEDIUM commands. They may or may not provide independent storage of a unit of media. See the device capabilities MODE SENSE page, 16.3.3.

16.1.1.2. Storage Elements

Storage elements are locations of units of media while not in some other element type. Medium in storage elements is available for access by medium transport elements.

Storage elements may be source and/or destination addresses in MOVE MEDIUM and EXCHANGE MEDIUM commands.

16.1.1.3. Import Export Elements

Import export elements are locations of units of media which are being inserted into or withdrawn from the medium changer device. Medium in these elements is accessible by both medium transport elements, by the operator, or by another independent medium changer device.

Import export elements may be source and/or destination addresses in MOVE MEDIUM and EXCHANGE MEDIUM commands. They may or may not provide independent storage of a unit of media, see the device capabilities MODE SENSE page, 16.3.3.

Particular import export elements may be capable of either import actions, export actions, both or neither (if an element is not present).

16.1.1.4. Data Transfer Element

Data transfer elements are locations of the primary devices which are capable of reading or writing the medium. Data transfer elements may also be viewed as medium changer element addresses of units of media loaded in or available for loading in or removal from primary devices such as disk or tape drives. Note that the medium changer function specified in this section does not control the primary device. That is the responsibility of the system.

Data transfer elements may be source and/or destination addresses in MOVE MEDIUM and EXCHANGE MEDIUM commands. They may or may not provide independent storage of a unit of media, see the device capabilities MODE SENSE page, 16.3.3.

16.1.2. SCSI Addressing of Medium Changer Devices

Medium Changer devices respond to a SCSI ID / LUN address different from those used by the data transfer element devices. If the data transfer elements served by the medium changer device are SCSI devices, they may be addressed though the same SCSI ID as the changer device but with a different logical unit number or through a different SCSI ID. There is no requirement that communication with the data transfer elements be done with the same SCSI bus as the medium changer device or by an SCSI interface at all. This command set does not support addressing both the changer and the data transfer elements with the same SCSI ID / LUN because the device at a particular SCSI ID / LUN address must have a single peripheral device type.

The READ ELEMENT STATUS response page for data transfer elements provides the SCSI ID / LUN address of the primary devices serviced by a medium changer device if known to the medium changer device.

16.1.3. Data Access Operations Using a Medium Changer Device

The medium changer device moves units of media among the several elements accessible to it. It is able to report the full/empty status of any of these elements. Data transfer elements represent the interface between the medium changer device and some primary device such as a removable media optical disk drive or tape drive.

In order to access data on a unit of media currently located in a medium changer storage element, a system will have to issue commands to both the medium changer device and to the primary device. The commands to the medium changer may include MOVE MEDIUM or EXCHANGE MEDIUM plus READ ELEMENT STATUS commands. The commands to the primary device may include TEST UNIT READY, INQUIRY, START STOP and LOAD UNLOAD commands in addition to READ and WRITE commands.

16.1.4. Element Status Maintenance Requirements

If the medium changer device chooses to implement the READ ELEMENT STATUS command, the medium changer device must be capable of reporting the various data (i.e. full, error, etc.) required by each page type. The medium changer may maintain this information at all times or regenerate it in response to the READ ELEMENT STATUS command. The INITIALIZE ELEMENT STATUS command can be used to force regeneration of this information.

16.1.5. Volume Tags

The Read Element Status descriptor format for all element types includes two sets of fields that contain volume tag information. These optional fields are used to report media identification information that the medium changer has acquired either by reading an external label (e.g. bar code labels), by a SEND VOLUME TAG command or by other means which may be vendor unique. The same volume tag information shall be available to all initiators whether assigned by that initiator, by some other initiator or by the media changer itself.

Volume tag information provides a means to confirm the identity of a unit of media that resides in a medium changer element. This command set does not define any direct addressing of units of media based on these fields. However, commands are defined that provide translation between volume tag information and the element address where that unit of media currently resides.

The medium changer command set definition does not impose the requirement that volume tag information be unique over the units of media within the scope of the changer device.

The following commands support the optional volume tag functionality:

SEND VOLUME TAG - either as a translation request or to associate a Volume Tag with the unit of media currently residing at an element address.

REQUEST VOLUME ELEMENT ADDRESS - return the element address currently associated with the volume tag information transferred with the last Send Volume Tag command.

RED ELEMENT STATUS - optionally reports volume tag information for all element types.

MOVE MEDIUM and EXCHANGE MEDIUM commands - if volume tags are implemented, these commands must retain the association between volume tag information and units of media as they are moved from element to element.

16.1.5.1. Volume Tag Format

Volume tag information consists of a volume identifier field of 32 bytes plus a volume sequence number field of 2 bytes. The volume identifier shall consist of a left justified sequence of ASCII characters. Unused positions shall be blank (20h) filled. In order for the SEND VOLUME TAG translate with template to work the characters '*' and '?' (2Ah and 3Fh) must not appear in volume identification data and there must be no blanks (20h) within the significant part of the volume identifier. If volume tag information for a particular element is undefined, the volume identifier field shall be zero.

The volume sequence number is a 2 byte integer field. If the volume sequence number is not used this field shall be zero. The volume tag contents are independent of the volume identification information recorded on the media.

IMPLEMENTORS NOTE: For compatibility with the volume identifier defined by volume and file structure standards, it is recommended that the characters in the significant non-blank portion of the volume identifier field be restricted to the set: '0'..'9', 'A'..'Z', and ' ' (30h .. 39h, 41h .. 5Ah, 5Fh). Specific systems may have differing requirements that may take precedence over this recommendation.

Table 16-1 defines the fields within the 36 byte primary and alternate volume tag information fields found in READ ELEMENT STATUS descriptors and in the data format for the SEND VOLUME TAG command.

Table 16-1: Volume Tag Information Format

Bit Byte	7	6	5	4	3	2	1	0
0	Volume Identification Field							
31								
32								
33								
34	Reserved							
35								
34	(MSB)							
35								
	Volume Sequence Number							
	(LSB)							

16.1.5.2. Primary and Alternate Volume Tag Information

Element status descriptors as reported by the READ ELEMENT STATUS command define a primary volume tag and an alternate volume tag. Alternate volume tag information provides a means for a system to use different volume identification information for each side of double sided media. In such a system, the primary volume tag information refers to the logical medium accessible via a MOVE MEDIUM command without the invert bit set. The alternate volume tag information refers to the other side of the media, i.e. the side that would be accessed via a MOVE MEDIUM command with the invert bit set.

16.2. Commands for Medium Changer Devices

The commands for medium changer devices shall be as shown in Table 16-2.

Table 16-2: Commands for Medium Changer Devices

Command Name	Operation Code	Type	Section
CHANGE DEFINITION	40h	O	7.2.1
EXCHANGE MEDIUM	A6h	O	16.2.1
INITIALIZE ELEMENT STATUS	07h	O	16.2.2
INQUIRY	12h	M	7.2.5
LOG SELECT	4Ch	O	7.2.6
LOG SENSE	4Dh	O	7.2.7
MODE SELECT(6)	15h	O	7.2.8
MODE SELECT(10)	55h	O	7.2.9
MODE SENSE(6)	1Ah	O	7.2.10
MODE SENSE(10)	5Ah	O	7.2.11
MOVE MEDIUM	A5h	M	16.2.3
POSITION TO ELEMENT	2Bh	O	16.2.4
PREVENT ALLOW MEDIUM REMOVAL	1Eh	O	8.2.4
READ BUFFER	3Ch	O	7.2.12
READ ELEMENT STATUS	B8h	O	16.2.5
RECEIVE DIAGNOSTIC RESULTS	1Ch	O	7.2.13
RELEASE	17h	O	16.2.6
REQUEST VOLUME ELEMENT ADDRESS	B5h	O	16.2.7
REQUEST SENSE	03h	M	7.2.14
RESERVE	16h	O	16.2.8
REZERO UNIT	01h	O	8.2.13
SEND DIAGNOSTIC	1Dh	M	7.2.15
SEND VOLUME TAG	B6h	O	16.2.9
TEST UNIT READY	00h	M	7.2.16
WRITE BUFFER	3Bh	O	7.2.17

Key: M = Command implementation is mandatory.
O = Command implementation is optional.

Operation codes C0h through FFh are vendor specific. All other operation codes are reserved.

16.2.1. EXCHANGE MEDIUM Command

Table 16-3: EXCHANGE MEDIUM Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (A6h)							
1	Logical Unit Number				Reserved			
2	(MSB)							
3	Transport Element Address							
4	(MSB)							
5	Source Address							
6	(MSB)							
7	First Destination Address							
8	(MSB)							
9	Second Destination Address							
10	Reserved				Inv2		Inv1	
11	Control							

The EXCHANGE MEDIUM command (Table 16-3) provides a means to exchange the medium in the source element with the medium located at a destination element.

IMPLEMENTORS NOTE: Support of this command requires that the device have the capability of handling two units of media at the same time or that it emulate this capability.

The medium in the source element is moved to the destination 1 element and the medium which previously occupied the destination 1 element is moved to the destination 2 element. The destination 2 element may or may not be the same element as the source element. In the case of a simple exchange, the source element and the destination 2 element are the same.

If this command is received and the source element is empty or the destination 1 element is empty or the destination 2 element (if different from the source element) is full or the destination 2 element (if the same as the source element) is empty, the target shall return CHECK CONDITION status and set the sense key to ILLEGAL REQUEST.

The transport element address specifies the medium transport element that is

to be used in executing this command. The default medium transport element address of zero may be used if this functionality is supported by the medium changer device. If the address specified has not been assigned or has been assigned to an element other than a medium transport element, the target shall return CHECK CONDITION status and set the sense key to ILLEGAL REQUEST.

The source address, the destination 1 address, and the destination 2 address may represent a storage element, an import export element, a data transfer element, or a medium transport element. If the address specified has not been assigned to a specific element of the medium changer, the target shall return CHECK CONDITION status and set the sense key to ILLEGAL REQUEST.

The device capabilities MODE SENSE page provides a matrix with the supported source element / destination 1 element combinations for EXCHANGE MEDIUM commands with source element the same as destination 2 element. See 16.3.3.

An Inv1 bit of one specifies that the medium shall be inverted prior to depositing the medium into the destination 1 element.

An Inv2 bit of one specifies that the medium shall be inverted prior to depositing the medium into the destination 2 element.

If the medium changer does not support medium rotation for handling double sided media, the Inv1 and Inv2 bits should be set to zero. If either of these bits is nonzero, a target which is not capable of medium rotation shall return CHECK CONDITION status and set the sense key to ILLEGAL REQUEST.

16.2.2. INITIALIZE ELEMENT STATUS Command

Table 16-4: INITIALIZE ELEMENT STATUS Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (07h)							
1	Logical Unit Number			Reserved				
2	Reserved							
3	Reserved							
4	Reserved							
5	Control							

The INITIALIZE ELEMENT STATUS command (Table 16-4) will cause the Medium Changer to check all elements for medium and any other status relevant to that element. The intent of this command is to enable the initiator to get a quick response from a following READ ELEMENT STATUS command. It may be useful to issue this command after a power failure, or if medium has been changed by an operator, or if configurations have been changed.

16.2.3. MOVE MEDIUM Command

Table 16-5: MOVE MEDIUM Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (A5h)							
1	Logical Unit Number			Reserved				
2	(MSB)							
3	Transport Element Address							
4	(MSB)							
5	Source Address							
6	(MSB)							
7	Destination Address							
8	Reserved							
9	Reserved							
10	Reserved							
11	Control							

The MOVE MEDIUM command (Table 16-5) requests that the target move a unit of media from a source element to a destination element.

The source address specifies the location that the medium is taken from, and the destination address specifies the location that the medium is moved to.

If this command is received and the source element is empty or the destination element (if different from the source element) is full, the target shall return CHECK CONDITION status and set the sense key to ILLEGAL REQUEST.

The transport element address specifies the medium transport element that is to be used in executing this command. The default medium transport element address of zero may be used if this functionality is supported by the medium changer device. If the address specified has not been assigned or has been assigned to an element other than a medium transport element, the target shall return CHECK CONDITION status and set the sense key to ILLEGAL REQUEST.

The source address and the destination address may represent a storage element, an import export element, a data transfer element, or a medium transport element. If the address specified has not been assigned to a

specific element of the medium changer, the target shall return CHECK CONDITION status and set the sense key to ILLEGAL REQUEST.

The device capabilities MODE SENSE page provides a matrix with the supported source element / destination element combinations for the MOVE MEDIUM command. See 16.3.3.

An invert bit of one specifies that the medium shall be inverted or rotated prior to depositing the medium into the destination element. If the medium changer does not support medium rotation for handling double sided media, the invert bit should be set to zero. If this bit is nonzero, a target which is not capable of medium rotation shall return CHECK CONDITION status and set the sense key to ILLEGAL REQUEST.

16.2.4. POSITION TO ELEMENT Command

Table 16-6: POSITION TO ELEMENT Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (2Bh)							
1	Logical Unit Number				Reserved			
2	(MSB)							
3	Transport Element Address (LSB)							
4	(MSB)							
5	Destination Element Address (LSB)							
6	Reserved							
7	Reserved							
8	Reserved							Invert
9	Control							

The POSITION TO ELEMENT command (Table 16-6) will position the transport element specified in front of the destination element specified.

IMPLEMENTORS NOTE: In this context, "in front of" is defined to mean -- positioned so that no further motion of the medium transport element is necessary to execute an appropriate MOVE MEDIUM command between the two elements.

An invert bit value of one requests that the transport element be inverted

or rotated before positioning in front of the destination element. A value of zero requests that the orientation be unchanged before positioning in front of the destination element.

16.2.5. READ ELEMENT STATUS Command

Table 16-7: READ ELEMENT STATUS Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (B8h)							
1	Logical Unit Number			VolTag	Element Type Code			
2	(MSB)							
3	Starting Element Address							(LSB)
4	(MSB)							
5	Number of Elements							(LSB)
6	Reserved							
7	(MSB)							
8	Allocation Length							(LSB)
9	(LSB)							
10	Reserved							
11	Control							

The READ ELEMENT STATUS command (Table 16-7) requests the target report the status of its internal elements to the initiator.

A volume tag (VolTag) bit of one indicates that the target shall report volume tag information if this feature is supported. A value of zero indicates that volume tag information shall not be reported. If the volume tag feature is not supported this field shall be treated as reserved.

The element type code field specifies the particular element type(s) selected for reporting by this command. A value of zero specifies that status for all element types shall be reported. The element type codes are defined in Table 16-8.

Table 16-8: Element Type Code

Code	Description
0h	All element types reported, (valid in CDB only)
1h	Medium Transport Element
2h	Storage Element
3h	Import Export Element
4h	Data Transfer Element
5h-Fh	Reserved

The starting element address specifies the minimum element address to report. Only elements with an element type code permitted by the element type code specification and an element address greater than or equal to the starting element address shall be reported. Element descriptor blocks are not generated for undefined element addresses.

The number of elements specifies the maximum number of element descriptors to be created by the target for this command. The value specified by this field is not the range of element addresses to be considered for reporting but rather the number of defined elements to report. If the allocation length is not sufficient to transfer all the element descriptors, the target shall transfer all those descriptors that can be completely transferred and this shall not be considered an error.

16.2.5.1. Element Status Data

The data returned by the READ ELEMENT STATUS command shall be as shown in Tables 16-9 through 16-14 and Figure 16-1. Element status data consists of an eight-byte header, (Table 16-9) followed by one or more element status pages.

Figure 16-1: Illustration of Element Status Data Structure

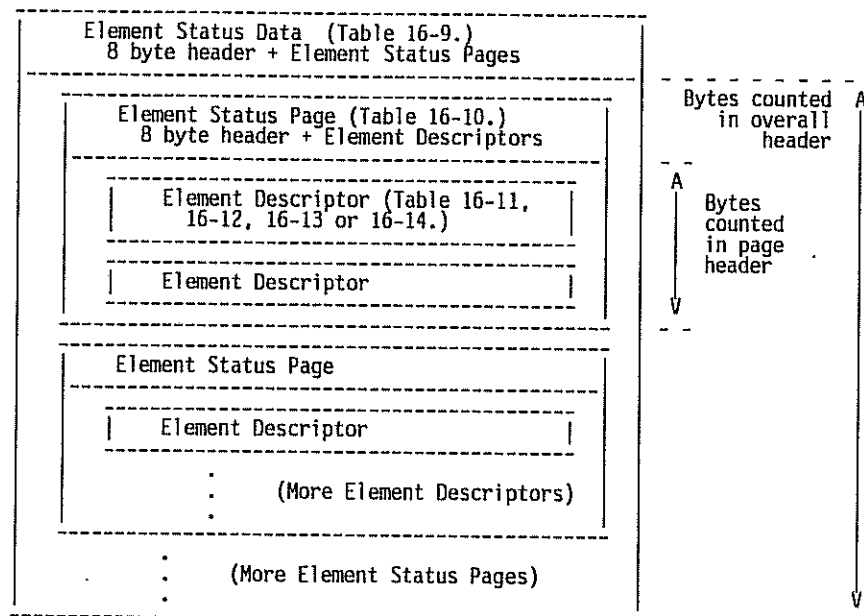


Table 16-9: Element Status Data

Bit Byte	7	6	5	4	3	2	1	0
0	(MSB)	First Element Address Reported						(LSB)
1								
2	(MSB)	Number of Elements Reported						(LSB)
3								
4		Reserved						
5	(MSB)							
6		Byte Count of Report Available						
7		(all pages, x - 7)						(LSB)
8 to								
x		Element Status Page(s)						

The first element address reported field indicates the element address of the element with the smallest element address found to meet the CDB request.

The number of elements field indicates the number of elements meeting the request in the command descriptor block. The status for these elements is returned if sufficient allocation length was specified.

The byte count of report available field indicates the number of bytes of element status page data available for all elements meeting the request in the command descriptor block. This value shall not be adjusted to match the allocation length available.

IMPLEMENTORS NOTE: The READ ELEMENT STATUS command can be issued with an allocation length of eight bytes in order to determine the allocation length required to transfer all the element status data specified by the command.

16.2.5.2. Element Status Page

Table 16-10: Element Status Page

Bit Byte	7	6	5	4	3	2	1	0
0	Element Type Code							
1	PVolTag	AVolTag	Reserved					
2	(MSB)	Element Descriptor Length						(LSB)
3		(z, from tables 16-11 - 16-14.)						
4		Reserved						
5	(MSB)							
6		Byte Count of Descriptor Data Available						
7		(this page, y - 7)						(LSB)
8 to								
y		Element Descriptor(s)						

The element status page is defined in Table 16-9. Each element status page includes an eight-byte header followed by one or more element descriptor blocks. The page header (Table 16-10) includes the element type code, the length of each descriptor block and the number of bytes of element descriptor information which follow the header for this element type.

The element type code field indicates the element type reported by this page.

A primary volume tag (PVolTag) bit of one indicates that the primary volume tag information field is present in each of the following element descriptor blocks. A value of zero indicates that these bytes are omitted from the element descriptors that follow.

An alternate volume tag (AVolTag) bit of one indicates that the alternate volume tag information field is present in each of the following element descriptor blocks. A value of zero indicates that these bytes are omitted from the element descriptors that follow.

The element descriptor length field indicates the number of bytes in each element descriptor.

The byte count of descriptor data available field indicates the number of bytes of element descriptor data available for elements of this element type meeting the request in the CDB. This value shall not be adjusted to match the

allocation length available.

Each element descriptor includes the element address, status flags and may contain sense code information as well as other information depending on the element type. (Tables 16-11 - 16-14.)

16.2.5.3. Medium Transport Element Descriptor

Table 16-11: Medium Transport Element Descriptor

Bit Byte	7	6	5	4	3	2	1	0
0	(MSB)							
1	Element Address (LSB)							
2	Reserved				Except		Reserved	Full
3	Reserved							
4	Additional Sense Code							
5	Additional Sense Code Qualifier							
6	Reserved							
8								
9	SValid	Invert	Reserved					
10	(MSB)							
11	Source Storage Element Address (LSB)							
12	Primary Volume Tag Information (Field omitted if PVolTag = 0)							
47								
48	Alternate Volume Tag Information (Field omitted if AVolTag = 0)							
83								
84	Reserved (Field moved up if Volume Tag Information field(s) are omitted.)							
87								
88	Vendor Unique (Field moved up if Volume Tag Information field(s) are omitted.)							
z-1								

The element address field gives the address of the medium changer element whose status is reported by this element descriptor block.

An exception (Except) bit of one indicates the element is in an abnormal state. An exception bit of zero indicates the element is in a normal state. If this bit is one, information on the abnormal state may be available in the additional sense code and additional sense code qualifier bytes.

A full bit value of one indicates that the element contains a unit of media. A value of zero indicates that the element does not contain a unit of media.

The additional sense code field may provide specific information on an abnormal element state. The values in this field are as defined for the additional sense code of the REQUEST SENSE data (see Table 7-42).

The additional sense code qualifier field may provide more detailed information on an abnormal element state. The values in this field are as defined for the additional sense code qualifier of the REQUEST SENSE data (see Table 7-42).

A source valid (SValid) bit value of one indicates that the source storage element address field and the invert bit information is valid. A value of zero indicates that the values in these fields are not valid.

An invert bit value of one indicates that the unit of media now in this element was inverted by MOVE MEDIUM or EXCHANGE MEDIUM operations since it was last in the source storage element. A value of zero indicates that no inversion occurred during the operation.

The source storage element address field provides the address of the last storage element this unit of media was moved from. This field is valid only if the SValid bit is one.

The primary and alternate volume tag information fields provide for identifying the unit of media residing in this element (see 16.1.5.). Either or both of these fields may be omitted for all the element descriptor blocks that comprise an element status page as indicated by the PVolTag and AVolTag bits in the element status page header (Table 16-10).

16.2.5.4. Storage Element Descriptor

Table 16-12: Storage Element Descriptor

Bit Byte	7	6	5	4	3	2	1	0
0	(MSB)							
1	Element Address (LSB)							
2	Reserved				Access	Except	Reserved	Full
3	Reserved							
4	Additional Sense Code							
5	Additional Sense Code Qualifier							
6								
8	Reserved							
9	SValid	Invert	Reserved					
10	(MSB)							
11	Source Element Address (LSB)							
12								
47	Primary Volume Tag Information (Field omitted if PVolTag = 0)							
48								
83	Alternate Volume Tag Information (Field omitted if PVolTag = 0)							
84								
87	Reserved (Field moved up if Volume Tag Information field(s) are omitted.)							
88								
z-1	Vendor Unique (Field moved up if Volume Tag Information field(s) are omitted.)							

An access bit value of one indicates access to the element by a medium transport element is allowed. An access bit of zero indicates access to the element by the medium transport element is denied.

The source storage element address field provides the address of the last storage element this unit of media was moved from. This element address value may or may not be the same as this element. This field is valid only if the

SValid bit is one.

(For fields not defined in this section, see 16.2.5.3.)

16.2.5.5. Import Export Element Descriptor

Table 16-13: Import Export Element Descriptor

Bit Byte	7	6	5	4	3	2	1	0
0	(MSB) Element Address (LSB)							
1								
2	Reserved		InEnab	ExEnab	Access	Except	ImpExp	Full
3	Reserved							
4	Additional Sense Code							
5	Additional Sense Code Qualifier							
6								
8	Reserved							
9	SValid	Invert	Reserved					
10	(MSB) Source Storage Element Address (LSB)							
11								
12								
47	Primary Volume Tag Information (Field omitted if PVolTag = 0)							
48								
83	Alternate Volume Tag Information (Field omitted if PVolTag = 0)							
84								
87	Reserved (Field moved up if Volume Tag Information field(s) are omitted.)							
88								
z-1	Vendor Unique (Field moved up if Volume Tag Information field(s) are omitted.)							

An import enable (InEnab) bit of one indicates that the import export element supports movement of media into the scope of the medium changer device. An InEnab bit of zero indicates that this element does not support import actions.

An export enable (ExEnab) bit of one indicates that the import export element supports movement of media out of the scope of the medium changer device. An ExEnab bit of zero indicates that this element does not support export actions.

An access bit of one indicates that access to the import export element by a medium transport element is allowed. An access bit of zero indicates access to the import export element by medium transport elements is denied.

IMPLEMENTORS NOTE: An example of when access would be denied is when the operator has exclusive access to the import export element.

An import export (ImpExp) bit of one indicates the unit of media in the import export element was placed there by an operator. An ImpExp bit of zero indicates the unit of media in the import export element was placed there by the medium transport element.

(For fields not defined in this section, see 16.2.5.3.)

16.2.5.6. Data Transfer Element Descriptor

Table 16-14: Data Transfer Element Descriptor

Bit Byte	7	6	5	4	3	2	1	0
0	(MSB)							
1	Element Address (LSB)							
2	Reserved				Access	Except	Reserved	Full
3	Reserved							
4	Additional Sense Code							
5	Additional Sense Code Qualifier							
6	Not Bus	Reserved	ID Valid	LU Valid	Reserved	Logical Unit Number		
7	SCSI Bus Address							
8	Reserved							
9	SValid	Invert	Reserved					
10	(MSB)							
11	Source Storage Element Address (LSB)							
12								
47	Primary Volume Tag Information (Field omitted if PVolTag = 0)							
48								
83	Alternate Volume Tag Information (Field omitted if PVolTag = 0)							
84								
87	Reserved (Field moved up if Volume Tag Information field(s) are omitted.)							
88								
z-1	Vendor Unique (Field moved up if Volume Tag Information field(s) are omitted.)							

An access bit value of one indicates access to the data transfer element by the medium transport element is allowed. A value of zero indicates access to the data transfer element by a medium transport element is denied.

IMPLEMENTORS NOTE: Access to the data transfer element by medium transport elements might be denied if a data transfer operation were under way. Note that a one value in this bit may not be sufficient to ensure a successful operation. This bit can only reflect the best information available to the medium changer device, which may not accurately reflect the state of the primary (data transfer) device.

A not this bus (Not Bus) bit value of one indicates that the SCSI bus address and logical unit number values are not valid for the SCSI bus used to select the medium changer device. A Not Bus bit value of zero indicates that the SCSI address and logical unit values, if valid, are on the same bus as the medium changer device.

An ID Valid bit value of one indicates that the SCSI bus address field contains valid information. An LU Valid bit value of one indicates that the logical unit number field contains valid information.

The SCSI bus address field, if valid, provides the SCSI address (binary representation) of the primary device served by the medium changer at this element address.

The logical unit number field, if valid, provides the logical unit number within the SCSI bus device of the primary device served by the medium changer at this element address.

(For fields not defined in this section, see 16.2.5.3.)

16.2.6. REQUEST VOLUME ELEMENT ADDRESS Command

Table 16-15: REQUEST VOLUME ELEMENT ADDRESS Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (B5h)							
1	Logical Unit Number			VolTag	Element Type Code			
2	(MSB)							
3	Element Address							(LSB)
4	(MSB)							
5	Number of Elements							(LSB)
6	Reserved							
7	(MSB)							
8	Allocation Length							(LSB)
9	Reserved							
10	Control							

The REQUEST VOLUME ELEMENT ADDRESS command (Table 16-15) is used to transfer the results of a SEND VOLUME TAG command. Multiple REQUEST VOLUME ELEMENT ADDRESS commands may be used to recover the results of a single SEND VOLUME TAG command with the translate option.

For each SEND VOLUME TAG command, the target shall be able to report multiple elements that match a volume tag template in element address order. Once information for a given element address has been reported, only higher element addresses will be reported by subsequent REQUEST VOLUME ELEMENT ADDRESS commands.

IMPLEMENTORS NOTE: In order to ensure the successful completion of a SEND VOLUME TAG, REQUEST VOLUME ELEMENT ADDRESS sequence in a multi-initiator environment it may be necessary to reserve the medium changer device. To do this, issue a unit RESERVE command prior to the SEND VOLUME TAG command and delay issuing the RELEASE until after the last REQUEST VOLUME ELEMENT ADDRESS.

A volume tag (VolTag) bit of one indicates that the target shall report volume tag information. A value of zero indicates that volume tag information

shall not be reported. This is provided for compatibility with the READ ELEMENT STATUS command.

The element type code field specifies the particular element type(s) selected for reporting by this command. A value of zero specifies that status for all element types shall be reported. The codes for this field are defined in Table 16-8. This acts as a qualification of the element(s) determined by SEND VOLUME TAG to match the template specified.

The Element Address field gives a medium changer element address whose interpretation depends on the Send Action Code received with the last SEND VOLUME TAG command. This code is echoed in the volume element address return format transferred by this command. When the last Send Action Code was a translate, (code values 0h .. 7h), the Element Address field gives the minimum element address to be reported by this command. When the Send Action Code is other than translate, (i.e. assert, replace, and undefine, code values 8h .. Dh), the Element Address field gives the particular element whose volume tag information was modified.

The number of elements specifies the maximum number of elements and volume tags to be reported by the target for this command. The value specified by this field is not the range of element addresses to be considered for reporting but rather the number of elements to report of those that match the last Send Volume Tag translate template. If the allocation length is not sufficient to transfer all the element descriptors, the target shall transfer all those descriptors that can be completely transferred and this shall not be considered an error.

The data returned by this command consists of a header as defined by Table 16-16, plus one or more element type specific pages in the same format as defined by the READ ELEMENT STATUS command. (See Tables 16-10 - 16-14.)

Table 16-16: Volume Element Address Header Format

Bit	7	6	5	4	3	2	1	0
Byte								
0	(MSB)							
1	First Element Address Reported							(LSB)
2	(MSB)							
3	Number of Elements Reported							(LSB)
4	Reserved			Send Action Code				
5	(MSB)							
6	Byte Count of Report Available							(LSB)
7	(all pages, x - 7)							
8 to	Element Status Page(s)							
x								

The send action code field gives the function performed by the last SEND VOLUME TAG command as listed in Table 16-21. This command reports the results of this send action code operation.

(For fields not defined in this section, see 16.2.5.1.)

16.2.7. RELEASE Command

Table 16-17: RELEASE Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (17h)							
1	Logical Unit Number			3rdPty	Third Party Device ID			Element
2	Reservation Identification							
3	Reserved							
4	Reserved							
5	Control							

The RELEASE command (Table 16-17) is used to release previously reserved units, or, if the element release option is implemented, previously reserved elements within units. It is not an error for an initiator to attempt to release a reservation that is not currently active. In this case, the target shall return GOOD status without altering any other reservation.

16.2.7.1. Logical Unit Release.

If the element bit is zero, this command shall cause the target to terminate all unit and element reservations that are active from the initiator to the specified logical unit.

16.2.7.2. Element Release (Optional).

If the element bit is one and the element release option is not implemented, the RELEASE command shall be terminated with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST. This option shall be implemented if the element reservation option (see 16.2.8.1) is implemented.

If the element bit is one and the element release option is implemented, this command shall cause any reservation from the requesting initiator with a matching reservation identification to be terminated. Other reservations from the requesting initiator shall remain in effect.

16.2.7.3. Third Party Release

The third-party release option of the RELEASE command allows an initiator to release a unit or elements within a unit that were previously reserved using third-party reservation (see 16.2.8.3). Third-party release shall be implemented if third-party reserve is implemented.

If the third-party (3rdPty) bit is zero, the third-party release option is not requested. If the 3rdPty bit is one and the target implements the third-party release option, the target shall release the specified unit or elements, but only if the reservation was made using the third-party reservation option by the initiator that is requesting the release for the same SCSI device as

specified in the third-party device ID field.

If the 3rdPty bit is one and the target does not implement the third-party release option, the target shall terminate the command with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST.

16.2.8. RESERVE Command

Table 16-18: RESERVE Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (16h)							
1	Logical Unit Number			3rdPty	Third Party Device ID		Element	
2	Reservation Identification							
3	(MSB)							
4	Element List Length							(LSB)
5	Control							

The RESERVE command (Table 16-18) is used to reserve units or, if the element reservation option is implemented, elements within units for the use of the initiator. The third-party reservation allows units or elements to be reserved for another specified SCSI device. The RESERVE and RELEASE commands provide the basic mechanism for contention resolution in multiple-initiator systems.

16.2.8.1. Logical Unit Reservation

If the element bit is zero, this command shall request that the entire unit be reserved for the exclusive use of the initiator until the reservation is superseded by another valid RESERVE command from the initiator that made the reservation or until released by a RELEASE command from the same initiator that made the reservation, by a BUS DEVICE RESET message from any initiator, by a hard RESET condition, or by a power off/on cycle. A unit reservation shall not be granted if the unit or any element is reserved by another initiator. It shall be permissible for an initiator to reserve a unit that is currently reserved by that initiator. If the element bit is zero, the reservation identification and the element list length shall be ignored.

If the unit, or any element within the unit is reserved for another initiator, the target shall respond by returning RESERVATION CONFLICT status.

If, after honoring the reservation, any other initiator then subsequently attempts to perform any command on the reserved unit other than an INQUIRY command, a REQUEST SENSE command, a RESERVE command, or a RELEASE command,

which shall be ignored, the command shall be rejected with RESERVATION CONFLICT status.

16.2.8.2. Element Reservation (Optional)

The reservation identification provides a means for an initiator to identify each element reservation. This allows an initiator in a multiple-tasking environment, to have multiple reservations outstanding. The reservation identification is used in the RELEASE command to specify which reservation is to be released. It is also used in superseding RESERVE commands to specify which reservation is to be superseded.

If the element reservation option is implemented, the element release option (see 16.2.7.2) shall also be implemented. These options permit multiple elements within the unit to be reserved, each with a separate reservation identification.

If the element bit is one, and the element reservation option is implemented, the target shall process the reservation request as follows:

(1) The element list shall be checked for valid element addresses. If any element address is invalid for this unit, the command shall be terminated with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST. The element list shall be checked for invalid element overlaps with other element descriptors in the element list and if invalid overlaps are found, the command shall be terminated with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST.

(2) If the requested reservation does not conflict with any active or previously requested reservation, the elements specified shall be reserved until superseded by another valid RESERVE command from the initiator that made the reservation or until released by a RELEASE command from the same initiator, by a BUS DEVICE RESET message from any initiator, or by a hard RESET condition. If either of the last two conditions occur, the next command from each initiator shall be terminated with CHECK CONDITION status and the sense key shall be set to UNIT ATTENTION.

(3) If the reservation request conflicts with a reservation already active the target shall return RESERVATION CONFLICT status.

If the element bit is one, and the element reservation option is not implemented, the RESERVE command shall be terminated with CHECK CONDITION status. The sense key shall be set to ILLEGAL REQUEST.

Table 16-19: Data Format of Element List Descriptors

Bit	7	6	5	4	3	2	1	0
Byte								
0	Reserved							
1								
2	(MSB)	Number of Elements						
3								(LSB)
4	(MSB)	Element Address						
5								(LSB)

The size of the element list shall be defined by the element list length parameter. The element list shall consist of zero or more descriptors as shown in Table 16-19. Each element list descriptor defines a series of elements beginning at the specified element address for the specified number of elements. If the number of elements is zero, the element list shall begin at the specified element address and continue through the last element address on the unit.

If an initiator issues a command to an element that has been reserved by a different initiator, the command shall not be performed and the command shall be terminated with a RESERVATION CONFLICT status. If a reservation conflict precludes any part of the command, none of the command shall be performed.

16.2.8.3. Third Party Reservation

The third-party reservation option of the RESERVE command allows an initiator to reserve a unit or elements within a unit for another SCSI device.

If the third-party (3rdPty) bit is zero, the third-party reservation option is not requested. If the 3rdPty bit is one and the third-party reservation option is implemented, the RESERVE command shall reserve the specified unit or elements for the SCSI device specified in the third-party device ID field. The target shall preserve the reservation until it is superseded by another valid RESERVE command from the initiator that made the reservation or until it is released by the same initiator, by a BUS DEVICE RESET message from any initiator, or a hard RESET condition. The target shall ignore any attempt to release the reservation made by any other initiator.

If the 3rdPty bit is one and the third-party reservation option is not implemented, the target shall reject the RESERVE command with CHECK CONDITION status and the sense key shall be set to ILLEGAL REQUEST.

16.2.8.4. Superseding Reservations

An initiator that holds a current reservation may modify that reservation by issuing another RESERVE command to the same unit and, if the element bit is one, using the same reservation identification. The superseding RESERVE command shall release the previous reservation in place when the new reservation request is granted. The current reservation shall not be modified if the new reservation request cannot be granted. If the superseding reservation cannot be granted because of conflicts with a previous active reservation (other than the reservation being superseded), the target shall return RESERVATION CONFLICT status.

IMPLEMENTORS NOTE: Superseding reservations are principally intended to allow the SCSI device ID to be changed on a reservation using the third-party reservation option.

16.2.9. SEND VOLUME TAG Command

Table 16-20: SEND VOLUME TAG Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (B6h)							
1	Logical Unit Number			Reserved		Element Type Code		
2	(MSB)							
3	Element Address							(LSB)
4	Reserved							
5	Reserved			Send Action Code				
6	Reserved							
7	Reserved							
8	(MSB)							
9	Parameter List Length							(LSB)
10	Reserved							
11	Control							

The SEND VOLUME TAG command (Table 16-20) transfers a volume tag template to be searched for or new volume tag information for one or more media changer elements. The function of the command is conveyed by the Send Action Code field. The REQUEST VOLUME ELEMENT ADDRESS command is used to transfer the results of a translate search operation.

The element type code field specifies an element type specification as defined by Table 16-8. If the send action code indicates a translate operation this field indicates the element types to be searched. If this is a zero, all element types are candidates for a translate operation. If the Send action code does not indicate a translate, this field shall be treated as reserved.

The send action code field gives the function to be performed by this command as listed in Table 16-21.

Table 16-21: Send Volume Tag Action Codes

Code	Description
0h	Translate - search all defined volume tags.
1h	Translate - search only primary volume tags.
2h	Translate - search only alternate volume tags.
3h	Reserved
4h	Translate - search all defined tags - ignore sequence numbers.
5h	Translate - search primary tags - ignore sequence numbers.
6h	Translate - search alternate tags - ignore sequence numbers.
7h	Reserved
8h	Assert - as the primary volume tag - if tag now undefined.
9h	Assert - as the alternate volume tag - if tag now undefined.
Ah	Replace - the primary volume tag - current tag ignored.
Bh	Replace - the alternate volume tag - current tag ignored.
Ch	Undefine - the primary volume tag - current tag ignored.
Dh	Undefine - the alternate volume tag - current tag ignored.
Eh-1Bh	Reserved
1Ch-1Fh	Vendor Specific

Translate operations request that the target search the volume tag information available to it for elements whose volume tag information matches the template given by the command parameters. The resulting information is reported via the REQUEST VOLUME ELEMENT ADDRESS command.

Assert operations define volume tag information for a single element. A CHECK CONDITION with ILLEGAL REQUEST is reported if the subject element already has defined volume tag information, in this case the original volume tag information shall remain unaltered.

Replace operations define volume tag information for a single element. Any previously defined volume tag information is overwritten.

Undefine operations cause previously defined volume tag information for the specified element to be cleared. It shall not be considered an error to undefine volume tag information that was not previously defined.

A medium changer device that implements volume tag information may choose to not implement the facilities that modify the volume tag information. With such an implementation a request for the assert, replace or undefine functions would cause the SEND VOLUME TAG command to be terminated with CHECK CONDITION and a sense key of ILLEGAL REQUEST.

The element address field gives a medium changer element address whose interpretation depends on the send action code received with this command. When the send action code is a translate, (code values 0h .. 7h), the Element Address field gives the starting element to be examined for satisfaction of the search criteria. When the send action code is other than translate, (i.e. assert, replace, and undefine, code values 8h .. Dh), the element address field gives the particular element whose volume tag information is to be modified.

The volume tag information used for translate, assert and replace functions is defined in Table 16-22. No parameters are necessary for undefine operations.

Table 16-22: Send Volume Tag Parameters Format

Bit	7	6	5	4	3	2	1	0						
Byte														
0														
31	Volume Identification Template Field													
32														
33	Reserved													
34	(MSB)	Minimum Volume Sequence Number												
35	(LSB)													
36														
37	Reserved													
38	(MSB)	Maximum Volume Sequence Number												
39	(LSB)													

The volume identification template field specifies a search template for translate functions and the exact value of the new volume identification information for other SEND VOLUME TAG command functions.

As a search template, this field may contain the wildcard characters '*' and '?' (2Ah and 3Fh). '*' will match any string of characters, when it appears in a template the remainder of the template is not used. '?' will match any

single character.

For a function other than translate, this field may not contain the '*' or '?' wildcard characters.

The minimum volume sequence number field specifies the new sequence number for the assert and replace functions. For a translate, this specifies the least value that will meet the search specification.

The maximum volume sequence number field specifies the greatest value that will meet the search specification. This field is ignored on functions other than translate.

16.3. Parameters for Medium Changer Devices

16.3.1. Diagnostic Parameters

This section defines the descriptors and pages for diagnostic parameters used with medium changer devices.

The diagnostic page codes for medium changer devices are defined in Table 16-23.

Table 16-23: Diagnostic Page Codes

Page Code	Description	Section
00h	Supported Diagnostic Pages	7.3.1.1
01h - 3Fh	Reserved (for all device type pages)	
40h - 7Fh	Reserved	
80h - FFh	Vendor-specific pages	

16.3.2. Log Parameters

This section defines the descriptors and pages for log parameters used with medium changer devices.

The log page codes for medium changer devices are defined in Table 16-24.

Table 16-24: Log Page Codes

Page Code	Description	Section
00h	Supported Log Pages	7.3.2.5
06h	Non-Medium Error Page	7.3.2.4
07h	Last n Error Events Page	7.3.2.3
01h - 05h, 08h - 2Fh, 3Fh	Reserved	
30h - 3Eh	Vendor-specific pages	

16.3.3. Mode Parameters

This section defines the descriptors and pages for mode parameters used with medium changer devices.

The mode parameter list, including the mode parameter header and mode block descriptor are defined in 7.3.3.

The medium-type code field is contained in the mode parameter header (see Table 7-61 and 7-62). This field is reserved for medium changer devices.

The device specific parameter field is contained in the mode parameter header (see Table 7-61 and 7-62). This field is reserved for medium changer devices.

The density code field is contained in the mode parameter block descriptor (see Table 7-63). This field is reserved for medium changer devices.

The mode page codes for medium changer devices are shown in Table 16-25.

Table 16-25: Mode Page Codes

Page Code	Description	Section
1Fh	Device Capabilities	16.3.3.1
1Dh	Element Address Assignment	16.3.3.2
1Eh	Transport Geometry Parameters	16.3.3.3
01h - 1Dh	Reserved	
00h	Vendor Specific (does not require page format)	
20h - 3Eh	Vendor-specific Pages	
3Fh	Returns all pages (valid only for the MODE SENSE command)	

16.3.3.1. Device Capabilities Page

Table 16-26: Device Capabilities Page

Bit Byte	7	6	5	4	3	2	1	0
0	PS	Reserved	Page Code (1Fh)					
1	Parameter Length (12h)							
2	Reserved				StorDT	StorI/E	StorST	StorMT
3	Reserved							
4	Reserved				MT->DT	MT->I/E	MT->ST	MT->MT
5	Reserved				ST->DT	ST->I/E	ST->ST	ST->MT
6	Reserved				I/E->DT	I/E->I/E	I/E->ST	I/E->MT
7	Reserved				DT->DT	DT->I/E	DT->ST	DT->MT
8 to 11	- - -	Reserved						- - -
12	Reserved				MT<>DT	MT<>I/E	MT<>ST	MT<>MT
13	Reserved				ST<>DT	ST<>I/E	ST<>ST	ST<>MT
14	Reserved				I/E<>DT	I/E<>I/E	I/E<>ST	I/E<>MT
15	Reserved				DT<>DT	DT<>I/E	DT<>ST	DT<>MT

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates

that the target is capable of saving the page in a non-volatile vendor-specific location.

The device capabilities page (Table 16-26) defines characteristics of the element types on this medium changer implementation. This information may be employed by the initiator to determine functions permitted by the MOVE MEDIUM and EXCHANGE MEDIUM commands.

In the field names on this page, the following element type abbreviations are used:

MT a medium transport element,
ST a storage element,
I/E an import export element, and
DT a data transfer element.

In the descriptions, XX and YY are any of the element type abbreviations.

A StorXX bit value of one indicated that the defined elements of type XX may provide independent storage for a unit of media. A value of zero indicates that elements of type XX provide virtual sources or destinations, that the location of the unit of media is provided by an element of some other type. The value of StorST is one by the definition of that type. See 16.1.1.

An XX->YY bit value of one indicates that the medium changer device supports MOVE MEDIUM commands where the source is element type XX and the destination is element type YY. An XX->YY bit value of zero indicates that these MOVE MEDIUM commands will be rejected with ILLEGAL REQUEST.

An XX<->YY bit value of one indicates that the medium changer device supports EXCHANGE MEDIUM commands where the source is element type XX, destination 1 is element type YY, and destination 2 is the same type as the source element type. These bits indicate the support for the possible types of simple exchange operations. An XX<->YY bit value of zero indicates that these EXCHANGE MEDIUM commands will be rejected with ILLEGAL REQUEST.

16.3.3.2. Element Address Assignment Page

Table 16-27: Element Address Assignment Page

Bit Byte	7	6	5	4	3	2	1	0
0	PS	[Reserved]	Page Code (10h)					
1	Parameter Length (12h)							
2	(MSB)	Medium Transport Element Address						(LSB)
3								
4	(MSB)	Number of Medium Transport Elements						(LSB)
5								
6	(MSB)	First Storage Element Address						(LSB)
7								
8	(MSB)	Number of Storage Elements						(LSB)
9								
10	(MSB)	First Import Export Element Address						(LSB)
11								
12	(MSB)	Number of Import Export Elements						(LSB)
13								
14	(MSB)	First Data Transfer Element Address						(LSB)
15								
16	(MSB)	Number of Data Transfer Elements						(LSB)
17								
18	Reserved							
19								

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

The element address assignment page (Table 16-27) is used to assign addresses to the elements of the medium changer (MODE SELECT) and to report those assignments (MODE SENSE). This page also defines the number of each type of element present.

The first medium transport element address field identifies the first medium transport element contained in the medium changer (other than the default medium transport address of zero). The number of medium transport elements field defines the total number of medium transport elements contained in the medium changer. If the number of medium transport elements field in a MODE SELECT command is greater than the default value returned in the MODE SENSE parameter data, the target shall return CHECK CONDITION status and set the sense key to ILLEGAL REQUEST.

The first storage element address field identifies the first medium storage element contained in the medium changer. The number of storage elements field defines the total number of medium storage elements contained in the medium changer. If the number of medium storage elements field in a MODE SELECT command is greater than the default value returned in the MODE SENSE parameter data, the target shall return CHECK CONDITION status and set the sense key to ILLEGAL REQUEST.

The first import export element address field identifies the first medium portal that is accessible both by the medium transport devices and also by an operator from outside the medium changer. The number of import export elements field defines the total number of import export elements contained in the medium changer and accessible to the medium transport elements. If the number of import export elements field in a MODE SELECT command is greater than the default value returned in the MODE SENSE parameter data, the target shall return CHECK CONDITION status and set the sense key to ILLEGAL REQUEST.

IMPLEMENTORS NOTE: The number of import export elements may be zero.

The first data transfer element address field identifies the first data transfer element contained in the medium changer. The data transfer elements may be either read/write or read-only devices. The number of data transfer field defines the total number of data transfer elements contained within the medium changer and accessible to the medium transport elements. If the number of data transfer elements field in a MODE SELECT command is greater than the default value returned in the MODE SENSE parameter data, the target shall return CHECK CONDITION status and set the sense key to ILLEGAL REQUEST.

Each element in the medium changer must have a unique address. If the address ranges defined for any of the element types overlap, the target shall return CHECK CONDITION status and set the sense key to ILLEGAL REQUEST.

16.3.3.3. Transport Geometry Parameters Page

Table 16-28: Transport Geometry Parameters Page

Bit	7	6	5	4	3	2	1	0
Byte								
0	PS	Reserved	Page Code (1Eh)					
1	Parameter Length							
	Transport Geometry Descriptor(s)							
0	Reserved							Rotate
1	Member Number in Transport Element Set							

The parameters savable (PS) bit is only used with the MODE SENSE command. This bit is reserved with the MODE SELECT command. A PS bit of one indicates that the target is capable of saving the page in a non-volatile vendor-specific location.

The transport geometry parameters page (Table 16-28) defines whether each medium transport element is a member of a set of elements that share a common robotics subsystem and whether the element is capable of media rotation. One transport geometry descriptor is transferred for each medium transport element, beginning with the first medium transport element. (Other than the default transport element address of zero.)

The parameter length specifies the number of bytes of transport geometry descriptors that follow. The geometry of each medium transport element is defined using a two-byte field as defined below.

A rotate bit of one indicates that the medium transport element supports media rotation for handling double-sided media. A rotate bit of zero indicates that the medium transport element does not support media rotation.

The member number in set indicates the position of this element in a set of medium transport elements that share a common robotics subsystem. The first element in a set has a member number of zero.

16.4. Glossary for Medium Changer Devices

data transfer element. A component of a medium changer used to access the data stored on a unit of media. The address in medium changer element space of a primary device.

element. An addressable physical component of a medium changer device that can serve as the location of a unit of media.

import export element. A location within a medium changer device that can be accessed by both the medium transport elements and by the operator.

medium transport element. A component of a medium changer device that is used to move units of media.

media rotation. The process of changing the orientation of a unit of media. In particular this refers to inverting a two sided media cartridge so that a data transport element which can access only side at a time can access data on the other side.

primary device. A device for reading or writing data on medium. These devices would typically implement the direct access, sequential access, optical memory device or CD-ROM command sets in this standard. Examples are magnetic disk drives, cartridge tape drives, optical disk drives and CD-ROM drives. Use in a medium changer environment implies that the device supports removable media.

storage element. A component of a medium changer device used for inactive storage of a unit of media.

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17. Communications Devices

17.1. Communications Device Model

A communications device provides a facility to send and receive information over some medium (usually an electrical or fiber-optic cable) using a defined protocol to one or more other systems which support the same protocol. The media and protocols are often specified in national or international standards, although some are proprietary. Within this section, the protocol used on the external medium is called the "external protocol".

The SCSI communications device model assumes that information to control the external protocol is embedded within the data transferred by the SEND and GET MESSAGE commands. The structure of the data transferred by these commands is not defined by this standard. Vendor-unique MODE pages may be defined to provide additional control over the external protocol.

The SCSI communications device is a target with the characteristics of a network access device, typically attached to a local area network (LAN), public telecommunications network, private telecommunications network, or packet switching network, etc., with no theoretical limits. Such a communications device is often called a network node. The communications device transmits or receives packets of data as requested by the initiator. The contents and meaning of the data packets is not defined by this standard. The bytes in the data packets may or may not contain addressing, path selection, an path control information. Auxiliary path selection information is optionally provided by the path selection identifier found in ten and twelve byte SEND MESSAGE and GET MESSAGE commands.

In the SCSI communications device, the target accepts and provides the data packets transferred in accordance with commands received from the initiator. There is an assumption that the initiator and the communications device know the rules by which information is exchanged between them, how the information is to be interpreted by the communications device, and when it is allowable to exchange the information. These rules are not specified by this standard.

The initiator requests that the communications device accept a packet of data by transmitting a SEND MESSAGE command. The initiator requests that the communications device return a packet of data by transmitting a GET MESSAGE command. The communications device also supports commands related to configuration of the network access device and network management of the particular environment. MODE SENSE and MODE SELECT are examples of these configuration commands.

If a communications device temporarily has no resource available to handle a data packet from the initiator, or has no data packet to provide to the initiator, or has no resources assigned to perform the operation, the device may then choose one of the following responses:

- (1) Terminate the command with CHECK CONDITION status. The sense key shall be set to NOT READY and the appropriate additional sense code for the condition.
- (2) Disconnect until the necessary resource or data packet becomes available, and then reconnect to the initiator and resume the operation.
- (3) Terminate the command with BUSY status.

More than one logical unit can be implemented by a communications device. Logical units can serve as multiple paths to a single network access device, and/or each logical unit can serve as a discrete path to different resources within the device. A single logical unit may also serve as a path to multiple resources if the communications device can interpret information within the data packet and route the packet to the appropriate resource. If the logical unit addressed by the initiator does not have an available resource or data packet associated with it, the communications device may choose to treat the logical unit as an invalid logical unit (see section 6.5.3) or respond as described in the previous paragraph.

If the communications device determines that an error or unusual condition has occurred while performing an operation specified by the contents of a data packet, the information describing the condition is normally returned as a part of a data packet from another network device. If the communications device determines that an error or unusual condition has occurred while either executing the SCSI command from the initiator, or during a network medium access transaction, the command is terminated with a CHECK CONDITION. The failures are identified through a REQUEST SENSE command.

The SCSI communications device is distinguished from a SCSI processor device by the fact that the primary destination of the data packets transferred to the communications device is not the target device itself, but another network node. A SCSI communications device passes the data on to an ultimate destination outside the target through a network. In contrast, the SCSI processor device is itself the primary destination of the data packets. Devices requiring protocols and command sets that are totally incompatible with the communications device protocols should be examined carefully to be sure that the incompatibilities are based on functional requirements.

17.1.1. Implementation Examples

Several examples of communications device implementations are provided to clarify the range of utility of this device type.

17.1.1.1. Host to Host Communications

In this example of host to host communications only the SEND MESSAGE command is used.

A communications device (Comm A) is attached to a network (Net A) and another communications device (Comm B) is attached to a network (Net B). Comm A takes the initiator role and selects a communications device (Comm B), transmitting a packet to Comm B using the SEND MESSAGE command. After transmitting the packet on the attached network medium (Net B), the action required by the SEND MESSAGE command, Comm B assumes the initiator mode and selects Comm A as a communications device and uses a SEND MESSAGE command to transmit a packet on the network (Net A) attached to Comm A. Note that the SEND MESSAGE command is sufficient to perform complete transactions between communications devices if Comm A and Comm B act as initiators. This provides the function of a network bridge or gateway for high bandwidth intercommunication among nearby host processors.

17.1.1.2. Host to Device Communications

In this example of host to device communications the SEND MESSAGE and GET MESSAGE commands are used.

A host system, Host A, takes the initiator role and selects a communications device (Comm A) attached to a network (Net A), transmitting a packet on network Net A to some other network node(s) using the SEND MESSAGE command to Comm A. Host A then assumes that a result will be obtained by the other network node(s) consistent with rules understood by all involved network devices. Host A generates a GET MESSAGE command to obtain packets from other network nodes. If there are no packets that need handling pending at Comm A, Comm A may disconnect until a packet arrives. Comm A can then complete the transaction, reconnect to Host A and the requested data packet can be returned to Host A. Note that Host A need not support target mode and Comm A need not support initiator mode to successfully complete an exchange between the two devices. The host system (Host A) can be replaced by communications device that is capable of acting as an initiator.

17.1.1.3. Multiple Role Communications

In this example of multiple role communications the SEND MESSAGE and GET MESSAGE commands are used. The device acts as both a host and a communications device depending on its needs and the requests made in the network.

A communications device (Comm A) is attached to a network (Net A) and another communications device (Comm B) is attached to a network (Net B). Comm A takes the initiator role and selects a communications device (Comm B). Comm A transmits a packet for network Net B using the SEND MESSAGE command to Comm B. Comm A then requests data packets from Net B by issuing a GET MESSAGE command to the Comm B device. Following these actions, Comm B assumes the

initiator role and transmits a packet using the SEND MESSAGE command to Comm A for transmission to network node(s) located on Net A. Comm B then requests data packets from Net A by issuing a GET MESSAGE command to Comm A.

17.2. Command Descriptions for Communications Devices

The commands for communications devices shall be as shown in Table 17-1.

Table 17-1: Commands for Communications Devices

Command Name	Operation Code	Type	Section	Page
CHANGE DEFINITION	40h	O	7.2.1	
GET MESSAGE(6)	08h	M	17.2.1	
GET MESSAGE(10)	28h	O	17.2.2	
GET MESSAGE(12)	A8h	O	17.2.3	
INQUIRY	12h	M	7.2.5	
LOG SELECT	4Ch	O	7.2.6	
LOG SENSE	4Dh	O	7.2.7	
MODE SELECT(6)	15h	O	7.2.8	
MODE SELECT(10)	55h	O	7.2.9	
MODE SENSE(6)	1Ah	O	7.2.10	
MODE SENSE(10)	5Ah	O	7.2.11	
READ BUFFER	3Ch	O	7.2.12	
RECEIVE DIAGNOSTIC RESULTS	1Ch	O	7.2.13	
REQUEST SENSE	03h	M	7.2.14	
SEND DIAGNOSTIC	1Dh	M	7.2.15	
SEND MESSAGE(6)	0Ah	M	17.2.4	
SEND MESSAGE(10)	2Ah	O	17.2.5	
SEND MESSAGE(12)	AAh	O	17.2.6	
TEST UNIT READY	00h	M	7.2.16	
WRITE BUFFER	3Bh	O	7.2.17	

Key M = Command implementation is mandatory.
O = Command implementation is optional.

All other operation codes for communications devices are reserved for future standardization.

17.2.1. GET MESSAGE(6) Command

Table 17-2: GET MESSAGE(6) Command

Bit	7	6	5	4	3	2	1	0
Byte								
0	Operation Code (08h)							
1	Logical Unit Number				Reserved			
2	(MSB)							
3	Allocation Length							
4	(LSB)							
5	Control							

The GET MESSAGE(6) command (Table 17-2) transfers data from the target to the initiator.

17.2.2. GET MESSAGE(10) Command

Table 17-3: GET MESSAGE(10) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (28h)							
1	Logical Unit Number			Reserved				
2	Reserved							
3	Reserved							
4	(MSB)							
5	Stream Selection				(LSB)			
6	Reserved							
7	(MSB)							
8	Allocation Length				(LSB)			
9	Control							

The GET MESSAGE(10) command (Table 17-3) transfers data from the target to the initiator.

The stream selection field specifies a further level of addressing for the data, so that it can be accessed by the target from the appropriate data stream.

Table 17-4: GET MESSAGE(12) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (A8h)							
1	Logical Unit Number				Reserved			
2	Reserved							
3	Reserved							
4	(MSB)							
5	Stream Selection							
6	(MSB)							
7								
8	Allocation Length							
9								
10	Reserved							
11	Control							

The GET MESSAGE(12) command (Table 17-4) transfers data from the target to the initiator.

The stream selection field specifies a further level of addressing for the data, so that it can be accessed by the target from the appropriate data stream.

Table 17-5: SEND MESSAGE(6) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (0Ah)							
1	Logical Unit Number				Reserved			
2	(MSB)							
3	Transfer Length							
4	(LSB)							
5	Control							

The SEND MESSAGE(6) command (Table 17-5) transfers data from the initiator to the target.

The transfer length specifies the length in bytes of data that shall be sent during the DATA OUT phase. A transfer length of zero indicates that no data shall be sent. This condition shall not be considered as an error.

17.2.5. SEND MESSAGE(10) Command

Table 17-6: SEND MESSAGE(10) Command

Bit Byte	7	6	5	4	3	2	1	0
0	Operation Code (2Ah)							
1	Logical Unit Number			Reserved				
2	Reserved							
3	Reserved							
4	(MSB)							
5	Stream Selection (LSB)							
6	Reserved							
7	(MSB)							
8	Transfer Length (LSB)							
9	Control							

The SEND MESSAGE(10) command (Table 17-6) transfers data from the initiator to the target.

The stream selection field specifies a further level of addressing for the data, so that it can be directed by the target to the appropriate data stream.

Please refer to the SEND MESSAGE(6) command for a description of the fields in this command.

17.2.6. SEND MESSAGE(12) Command

Table 17-7: SEND MESSAGE(12) Command

Bit Byte	7	6	5	4	3	2	1	0						
0	Operation Code (AAh)													
1	Logical Unit Number			Reserved										
2	Reserved													
3	Reserved													
4	(MSB)	Stream Selection						(LSB)						
5														
6	(MSB)	Transfer Length						(LSB)						
7														
8														
9														
10	Reserved													
11	Control													

The SEND MESSAGE(12) command (Table 17-7) transfers data from the initiator to the target.

The stream selection field specifies a further level of addressing for the data, so that it can be directed by the target to the appropriate data stream.

Please refer to the SEND MESSAGE(6) command for a description of the fields in this command.

17.3. Parameters for Communication Devices

17.3.1. Diagnostic Parameters

This section defines the descriptors and pages for diagnostic parameters used with communication devices.

The diagnostic page codes for communication devices are defined in Table 17-8.

Table 17-8: Diagnostic Page Codes

Page Code	Description	Section
00h	Supported Log Pages	7.3.1.1
01h-3Fh	Reserved (for all device type pages)	
40h-7Fh	Reserved	
80h-FFh	Vendor-specific pages	

17.3.2. Log Parameters

This section defines the descriptors and pages for log parameters used with communication devices.

The log page codes for communication devices are defined in Table 17-9.

Table 17-9: Log Page Codes

Page Code	Description	Section
01h	Buffer Over-Run/Under-Run Page	7.3.2.1
02h	Error Counter Page (Write) Page	7.3.2.2
03h	Error Counter Page (Read) Page	7.3.2.2
04h	Error Counter Page (Read Reverse) Page	7.3.2.2
05h	Error Counter Page (Verify) Page	7.3.2.2
07h	Last n Error Events Page	7.3.2.3
06h	Non-Medium Error Page	7.3.2.4
00h	Supported Log Pages	7.3.2.5
08h - 2Fh	Reserved	
30h - 3Eh	Vendor-specific pages	
3Fh	Reserved	

17.3.3. Mode Parameters

This section defines the descriptors and pages for mode parameters used with communication devices.

The mode parameter list, including the mode parameter header and mode block descriptor are defined in 7.3.3.

The medium-type code field is contained in the mode parameter header (see Table 7-61 and 7-62). This field is reserved for communications devices.

The device specific parameter field is contained in the mode parameter header (see Table 7-61 and 7-62). This field is reserved for communications devices.

The density code field is contained in the mode parameter block descriptor (see Table 7-63). This field is reserved for communications devices.

The mode page codes for communication devices are shown in Table 17-10.

Table 17-10: Mode Page Codes

Page Code	Description	Section
0Ah	Control Mode Page	7.3.3.1
02h	Disconnect-Reconnect Page	7.3.3.2
09h	Peripheral Device Page	7.3.3.3
01h	Reserved	
03h - 08h	Reserved	
0Bh - 1Fh	Reserved	
Fh	Returns all pages (valid only for the MODE SENSE command)	
00h	Vendor specific (does not require page format)	
20h - 3Eh	Vendor specific	

17.4. Glossary for Communications Devices

communications device. An SCSI device whose principal function is to communicate with one or more other systems, usually over distances that exceed the maximum cable length defined for SCSI.

external medium. The medium used by the communications device to send or receive information to or from one or more communications devices. The other communications devices may or may not use an SCSI interface.

external protocol. The protocol(s) used by the communications device to transfer information over the external medium. The external protocol(s) are not defined by this standard.

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Appendixes (These Appendixes are not part of the requirements of American National Standard X3.131-198x, but are included for information only.)

A. SCSI Signal Sequence Example

This Appendix is included to provide an example of the signal sequencing of an I/O process that includes most of the SCSI bus phases (Figure A-1). In this example, the target does not disconnect from the SCSI bus prior to completing the I/O process.

The following notes apply to Figure A-1:

NOTE: In a typical system, the computer's host adapter acts as the "initiator" and the peripheral device's controller acts as the "target." In general, this standard does not attempt to distinguish between a computer and its host adapter. These functions may be separate or merged. The term "initiator" encompasses both. The term "target" refers to the controller portion of the peripheral device, which may be separate (bridge controller) from the peripheral device or merged with it (embedded controller). The term "SCSI device" refers to a device that may be connected to the SCSI bus. An SCSI device may act in the initiator role, the target role, or both roles.

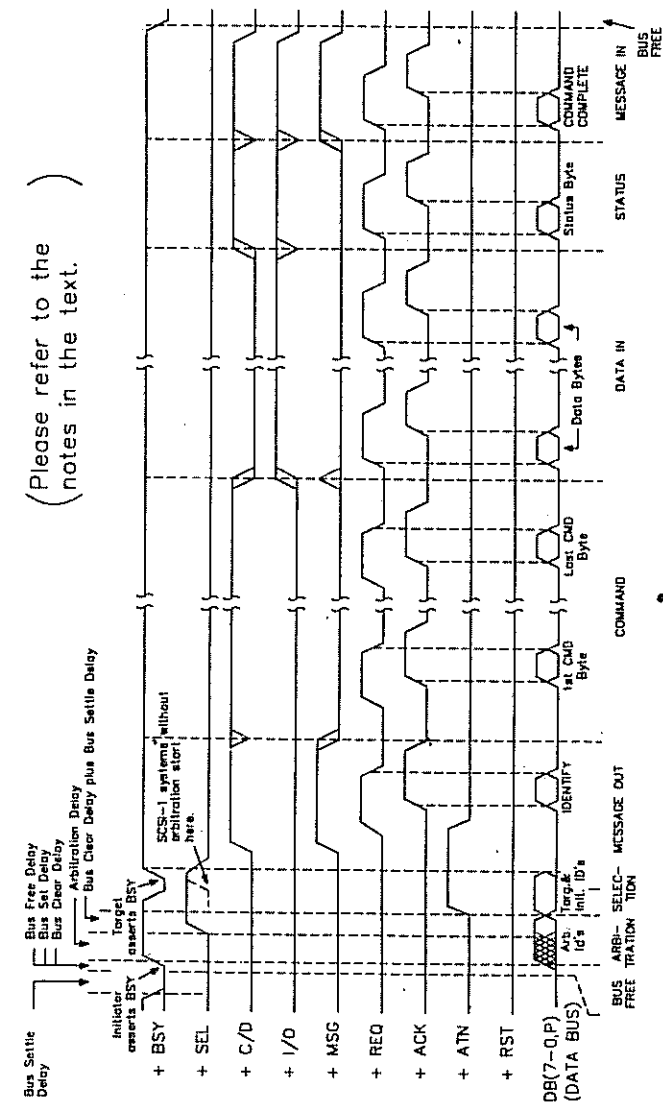


Figure A-1: SCSI Signal Sequence Example

DATA BUS NOTES:

- (1) DB(7) is the most significant bit.
- (2) DB(7) is the highest priority arbitration bit.
- (3) DB(P) is the data parity bit (odd). Parity is not valid during the ARBITRATION phase.

BUS PHASE NOTES:

BUS FREE phase. BUS FREE phase begins when the SEL and BSY signals are both continuously false for a bus settle delay. It ends when the BSY signal becomes true. (In the SCSI-1 single-initiator option, BUS FREE phase could also end when the SEL signal became true.)

ARBITRATION phase. This phase is documented as mandatory in SCSI-2. In SCSI-1, this phase was optional.

At least one bus free delay after first detecting BUS FREE phase, but no more than a bus settle delay after the bus was last free, the initiator asserts the BSY signal and its assigned SCSI device ID bit on the DATA BUS. The initiator waits an arbitration delay, then examines the DATA BUS. If a higher priority SCSI device ID bit is true, the initiator loses arbitration and may release the BSY signal and its SCSI ID bit. Otherwise, the initiator wins arbitration and asserts the SEL signal. All SCSI devices must release the BSY signal and their SCSI ID bit within a bus clear delay after the SEL signal becomes true (even if they have not yet examined the DATA BUS). The winning SCSI device waits at least a bus clear delay plus a bus settle delay after asserting the SEL signal before changing any signals on the bus.

SELECTION phase. The I/O signal is false during this phase to distinguish it from the RESELECTION phase.

NONARBITRATING SYSTEMS (only permitted in SCSI-1): In such systems, the initiator waits at least a bus clear delay after detecting BUS FREE phase, then it asserts the target's SCSI ID bit and, optionally, the initiator's SCSI ID bit on the DATA BUS. After at least two deskew delays, the initiator asserts the SEL signal.

ARBITRATING SYSTEMS: In such systems, the SCSI device that won arbitration has both the BSY and SEL signals asserted. After at least a bus clear delay plus a bus settle delay, it places both the target's and the initiator's SCSI ID bits on the DATA BUS. At least two deskew delays later, it releases the BSY signal.

ALL SYSTEMS: The target determines that it is selected when the SEL signal and its SCSI ID bit are true and the BSY and I/O signals are false for at least a bus settle delay. The target then asserts the BSY signal within a selection abort time after it last determined that it was still being selected. (The target is not required to respond to a selection within a selection abort time; but it must ensure that it will not assert the BSY signal more than a selection abort time after the initiator aborts a selection attempt.)

At least two deskew delays after the initiator detects the BSY signal is true, it releases the SEL signal.

MESSAGE OUT phase. During this phase the initiator sends an IDENTIFY message to the target. The target asserts the C/D and MSG signals and negates the I/O signal for the message transfer. After detecting the assertion of the REQ signal, the initiator negates the ATN signal before asserting the ACK signal. (Refer to the handshake procedure for the COMMAND phase.)

COMMAND phase. The target asserts the C/D signal and negates the I/O and MSG signals for all of the bytes transferred during this phase. The direction of transfer is from the initiator to the target.

HANDSHAKE PROCEDURE: The target asserts the REQ signal. Upon detecting the REQ signal is true, the initiator drives the DATA BUS to the desired value, waits at least one deskew delay plus a cable skew delay and then asserts the ACK signal. The initiator continues to drive the DATA BUS until the REQ signal is false.

When the ACK signal is true at the target, the target reads the DATA BUS and then negates the REQ signal.

When the REQ signal becomes false at the initiator, the initiator may change or release the DATA BUS and negate the ACK signal.

The target may continue to request command bytes by asserting the REQ signal again. The number of command bytes is determined by the group code (most significant 3 bits) that is contained in the first command byte.

DATA IN phase. The target asserts the I/O signal and negates the C/D and MSG signal for all of the bytes transferred during this phase. The direction of transfer is from the target to the initiator.

HANDSHAKE PROCEDURE: The target first drives the DATA BUS to their desired values, waits at least one deskew delay plus a cable skew delay, and then asserts the REQ signal. The target continues to drive the DATA BUS until the ACK signal is true.

When the REQ signal is true at the initiator, the initiator reads the DATA BUS and then asserts the ACK signal.

When the ACK signal is true at the target, the target may change or release the DATA BUS and negate the REQ signal.

When the REQ signal is false at the initiator, the initiator negates the ACK signal. After the ACK signal is false, the target may continue the transfer by driving the DATA BUS and asserting the REQ signal as described above.

DATA OUT phase (not shown in the figure). The target negates the C/D, I/O, and MSG signals for all of the bytes transferred during this phase. The direction of transfer is from the initiator to the target. (Refer to the handshake procedure for the COMMAND phase.)

STATUS phase. The target asserts the C/D and I/O signals and negates the MSG signal for the byte transferred during this phase. The direction of

transfer is from the target to the initiator. (Refer to the handshake procedure for the DATA IN phase.)

MESSAGE IN phase. The target asserts the C/D, I/O, and MSG signals during the byte transferred during this phase. Typically, a COMMAND COMPLETE message would be sent at this point. The direction of transfer is from the target to the initiator. (Refer to the handshake procedure for the DATA IN phase.)

BUS FREE phase. The target returns to BUS FREE phase by releasing the BSY signal. Both the target and the initiator release all bus signals within a bus clear delay after the BSY signal is continuously false for a bus settle delay.

B. Typical Bus Phase Sequence

This Appendix is included to provide an example of the SCSI bus phase sequence for a typical READ command (Tables B-1 and B-2). In this example, the target does not disconnect from the SCSI bus prior to completing the command.

Table B-1: Typical READ Command Phase Sequence

Bus Phase	Signals											Comment
	B S Y	S E L	A T N	M S G	C /	I /	R E Q	A C K	R E S	D B (7-0)	D B (P)	
BUS FREE	-	-	-	-	-	-	-	-	-	-	-	SCSI bus is available.
ARBITRATION	1	-	-	-	-	-	-	-	-	ID	X	Initiator tries to get the SCSI bus.
SELECTION	1	1	1	-	-	-	-	0	-	ID	I, T, V	Initiator has SCSI bus and selects a target.
"	-	1								ID	I, T, V	
"	1	1								ID	I, T, V	ATN is on.
"	1	-								X	X	
MESSAGE OUT	1	-	1	1	1	0	0	0	-	X	X	Target has control of the bus and gets the IDENTIFY message from the initiator.
"			1				1	0		X	X	
"			0				1	1		Message	V	
"			0				0	1		X	X	
"			0				0	0		X	X	
COMMAND	1	-	0	0	1	0	0	0	-	X	X	Target gets a command byte from the initiator. (This handshake is repeated for each byte.)
"							1	0		X	X	
"							1	1		Command	V	
"							0	1		X	X	
"							0	0		X	X	

Table B-2: Typical READ Command Phase Sequence (Continued)

Bus Phase	Signals											Comment
	B Y	S E L	A T N	M S G	C /	I /	R E Q	A C K	R S T	D B (7-0)	D B (P)	
DATA IN	1	-	0	0	0	1	0	0	-	X	X	Target sends data to the initiator. (This handshake is repeated for each byte.)
"							1	0	Read	Data	V	
"							1	1		X	X	
"							0	1		X	X	
"							0	0		X	X	
STATUS	1	-	0	0	1	1	0	0	-	X	X	Target sends status to the initiator.
"							1	0	Status		V	
"							1	1		X	X	
"							0	1		X	X	
"							0	0		X	X	
MESSAGE IN	1	-	0	1	1	1	0	0	-	X	X	Target sends a COMMAND COMPLETE message to the initiator.
"							1	0	Message		V	
"							1	1		X	X	
"							0	1		X	X	
"							0	0		X	X	
BUS FREE	-	-	-	-	-	-	-	-	-	-	-	SCSI bus is available.

Key: - = Signal driver is passive.
 0 = Signal is false.
 1 = Signal is true.
 "Blank" = Signal state is the same as the previous line.
 ID = SCSI ID for arbitration.
 ID I,T = SCSI ID of initiator and target.
 V = Parity is valid.
 X = The signal is not guaranteed to be in a known state.

C. SCSI System Operation

This Appendix is included to provide an explanation of the relationship of the various pieces of an SCSI system. This Appendix also provides additional information about the use of SCSI in a multi-tasking system. Such systems typically use a host adapter circuit to interface from the host memory to the SCSI bus. Although other architectures are possible (including native or embedded SCSI), the host adapter logic still exists as part of the system. The term "initiator" is used throughout this standard to encompass all such architectures. The term "host adapter" is used within this Appendix to refer to the logic that interfaces from the host memory to the SCSI bus.

C.1. Host Memory, Host Adapter, SCSI Target Relationship

The SCSI architecture utilizes the concept of host memory blocks for command, data, and status interchange between the host system and the SCSI target. In the middle of this exchange is the SCSI host adapter, which acts as the gateway between host memory and the SCSI bus. The host adapter is an important portion of the overall intelligence of SCSI. Along with providing an information path from the SCSI bus to the host bus, the host adapter is intimately involved in assuring data integrity and proper performance of the I/O subsystem.

In order to fully understand SCSI operation, the concepts of I/O memory blocks and nexus are detailed. Figure C1 presents an example block diagram of a single host/single peripheral SCSI I/O subsystem. The host memory contains three I/O blocks: command, data, and status. The SCSI disk controller target needs to read the command block and write to the status block in order to perform the task specified by the host in the command block. Likewise, the SCSI controller needs to either read or write the data block in host memory, depending on the task specified. The SCSI controller "reaches into host memory" via the SCSI host adapter. The host adapter must know the addresses of the command, data, and status blocks in order for it to "reach" into the right spot in memory. In other words, the host adapter must be given a pointer to the start of each block by the host. As the SCSI controller takes information from the command block, the memory pointer for the command block advances to the next byte. The same is true for the data and status pointers.

SCSI architecture provides for two sets of three pointers within the host adapter. The first set is known as the current (or active) pointer values. These are the pointers to the next command, data, or status byte to be transferred between the host memory and the SCSI controller. There is only one set of current pointers in the host adapter. The current pointers are shared among all devices and are used by the current device connected to the host adapter.

The second set is known as the saved pointer values. There is one set of saved pointers for each active I/O process. For command and status, these pointers always point to the start of the memory command block and memory status block. The saved data pointer points to the start of the data block at the beginning of each command. It remains at this value until the controller sends a SAVE DATA POINTER message to the host adapter which copies the value of the current data pointer into the saved data pointer. The controller may retrieve the saved value by sending a RESTORE POINTERS message. This moves

the value of each of the three saved pointers into the current pointers. Whenever an SCSI device disconnects from the bus, only the saved pointer values are retained. The current pointer values are set from the saved values upon the next reconnection. The current and saved pointers provide an efficient method to break up large transfers into smaller bursts, and to facilitate error retry and recovery.

C.2. SCSI READ Command Example

One method of understanding the "host, host adapter, SCSI controller" relationship is via an example. Consider the case of a multiple sector READ command that will cross a cylinder boundary on a direct-access device such as a disk.

The first activity in the I/O operation is for the system to create a command descriptor block in memory and determine where the data and status are to be written in host memory. The host then sends a command to the host adapter that includes the starting address (pointer) for each of the command, data, and status blocks and the SCSI address of the peripheral to perform the operation. In this example, there is only one SCSI controller and physical disk, but its address is required in order for the host adapter to select it.

Upon receiving the command, the host adapter arbitrates for the SCSI bus and wins (due to the lack of competing devices) and proceeds to select the target SCSI device with the ATN signal asserted. The ATTENTION condition indicates to the SCSI target that the initiator (the host adapter) has a message to send to the target. When the target responds to the SELECTION phase, an I_T nexus is established between the two devices.

After the SELECTION phase is completed, the target responds to the initiator's ATTENTION condition by receiving an IDENTIFY message from the initiator. This message, generated by the host adapter, indicates the desired logical unit number in the target and whether the initiator can support bus disconnect. In this example, the initiator supports disconnect. When the controller receives the IDENTIFY message, an I_T L nexus is established. The nexus uniquely identifies the relationship between the initiator and the specified logical unit of the target disk controller.

An additional message following the IDENTIFY may be sent for purposes of command queuing. If a QUEUE TAG message is sent, the I_T L nexus is replaced by an I_T L Q nexus. This I_T L Q nexus behaves in a similar manner as the I_T L nexus for purposes of pointer management; it merely permits more sets of pointers to be identified. In this example, however, command queuing is not used.

Input/output activity from this point are principally controlled by the target. The host adapter is simply an "arm" of the target used to reach into host memory. Utilizing this arm, the target reads in the command descriptor block (CDB). The host adapter is expected to ensure that the target does not reach outside its allocated blocks.

After decoding the instruction, the controller determines that a disk seek is required to get the starting data block. Since the SCSI bus will not be utilized until data has been read from the disk, the target controller disconnects from the bus. The disconnect process includes the transmission of

a SAVE DATA POINTER message and DISCONNECT message from the target to the host adapter. The host adapter responds to the SAVE DATA POINTER message by saving the current data pointer, which is still set to the start of the data block. (Strictly speaking, the target need not send the SAVE DATA POINTER message following the command phase since at that time the saved and current pointers are equal.) After transmission of the DISCONNECT message the target releases the BSY signal, freeing the bus.

Although the initiator host adapter and target disk controller are disconnected, the I/O process has not completed and the I_T L nexus still exists. Both devices know they have a command to finish and will return to that job at a later point in time. The ability to disconnect allows multiple I/O processes to occur simultaneously, utilizing a single physical bus. The logical connection is actually not just between the host adapter and the disk controller, but runs all the way from the host memory I/O block to the peripheral device (disk) performing the operation. (See Figure C-2 for a pictorial presentation of this concept.)

Once the target has started filling its data buffers, it can transmit data to the initiator, but first it must revive the connection. The reconnection process involves the target arbitrating for the bus and reselecting the host adapter. After the reselection is made, the target sends an IDENTIFY message to the host adapter to indicate which target logical unit is reconnecting. This information provides the correct logical connection via the I_T L nexus into host memory. After reconnection, the roles of the initiator and target are just as they were prior to disconnection. The target transfers data into host memory via the host adapter. The data transfer continues until the disk reaches the end of its cylinder and the disk controller determines that a second physical seek is required to complete the READ command. The target again performs a SAVE DATA POINTER message and a DISCONNECT message. However, this time the current data pointer is not at the beginning of the memory data block, and is required to ensure that the I/O process continues at the correct data block location. The saved value at disconnect reflects the change.

After seek completion and transfer of data into its buffer, the controller reconnects to the host adapter and completes the data transfer as requested by the READ command. At this point, the controller sends ending status into host memory via the host adapter. The final action of the target is to send the host adapter a COMMAND COMPLETE message and go to BUS FREE. The target has completed its operation and considers the I/O process ended.

Upon receipt of the COMMAND COMPLETE message, the host adapter signals the host that the I/O process is complete. This signal can be an interrupt or the setting of a flag read by the host in a polled I/O environment. This action by the host adapter breaks the logical connection between the host adapter and the I/O memory blocks of the host. The host reviews the status of the operation in the status block and proceeds to utilize the data transferred into the data block.

C.3. I/O Channel Concept

The I/O channel concept fully utilizes the high performance capability of the SCSI. The I/O channel is basically an intelligent SCSI host adapter that can maintain multiple simultaneous I/O processes between host memory I/O blocks and different SCSI devices.

The I/O channel utilizes a single direct memory access (DMA) path into host memory supporting the DMA operations of numerous SCSI peripherals. Since the SCSI bus is a single physical bus and most host computers have a single physical backplane bus, multiple DMA channels into memory are not necessary. In many implementations of a multiple DMA channel architecture, when a channel is accessing memory, all other channels are idle. In such implementations, a single channel supporting multiple I/O processes can supply the same performance as separate DMA peripherals. An obvious advantage to the host is lower system cost as well as the saving in backplane card slots.

In the READ command example discussed in C.2, the I/O channel is the SCSI host adapter. The host gives the I/O channel a command by providing it with pointers to the I/O memory blocks and the SCSI peripheral address. This establishes a logical connection between the host adapter and the host I/O memory blocks. The I/O channel then opens a sub-channel that is assigned the task of managing the physical link and nexus between the host adapter and the target controller. All physical connections and reconnections to the host adapter are managed by this sub-channel. The number of active or open sub-channels an I/O channel can support is totally dependent upon its design. The SCSI definition could, in theory, support an I/O channel with up to 56 sub-channels for simple I T L nexus, and many more if target routines and command queuing are implemented.

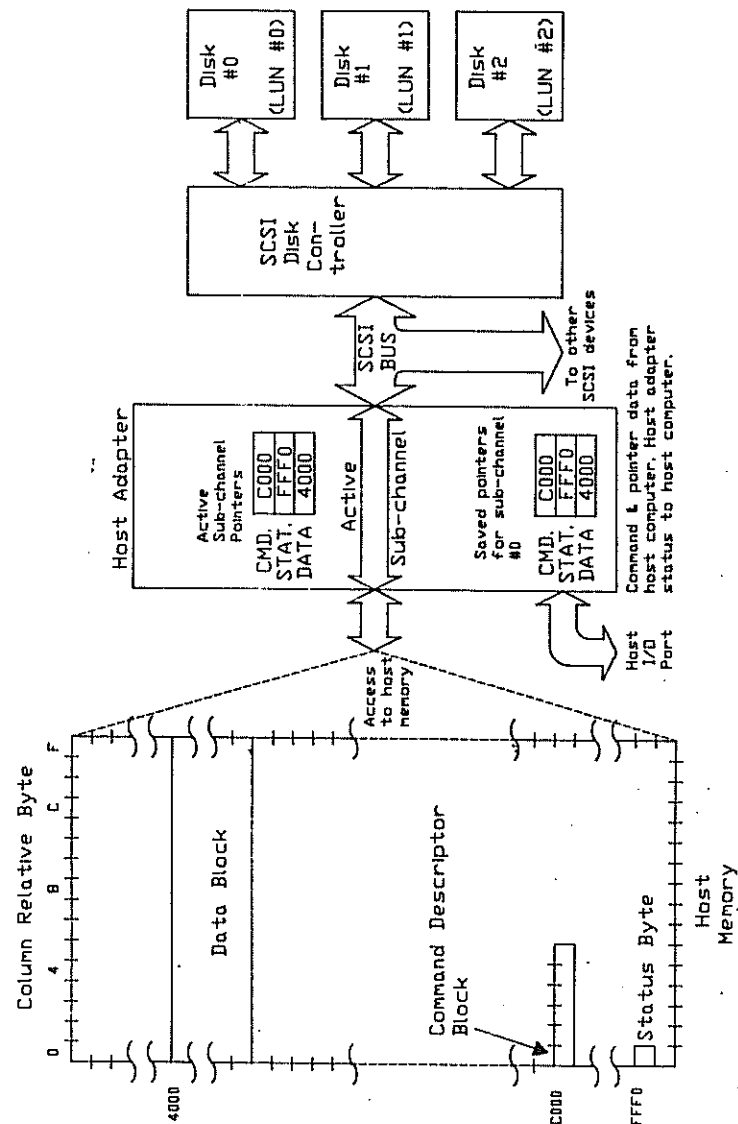


Figure C-1: Snapshot Prior to Initial Selection

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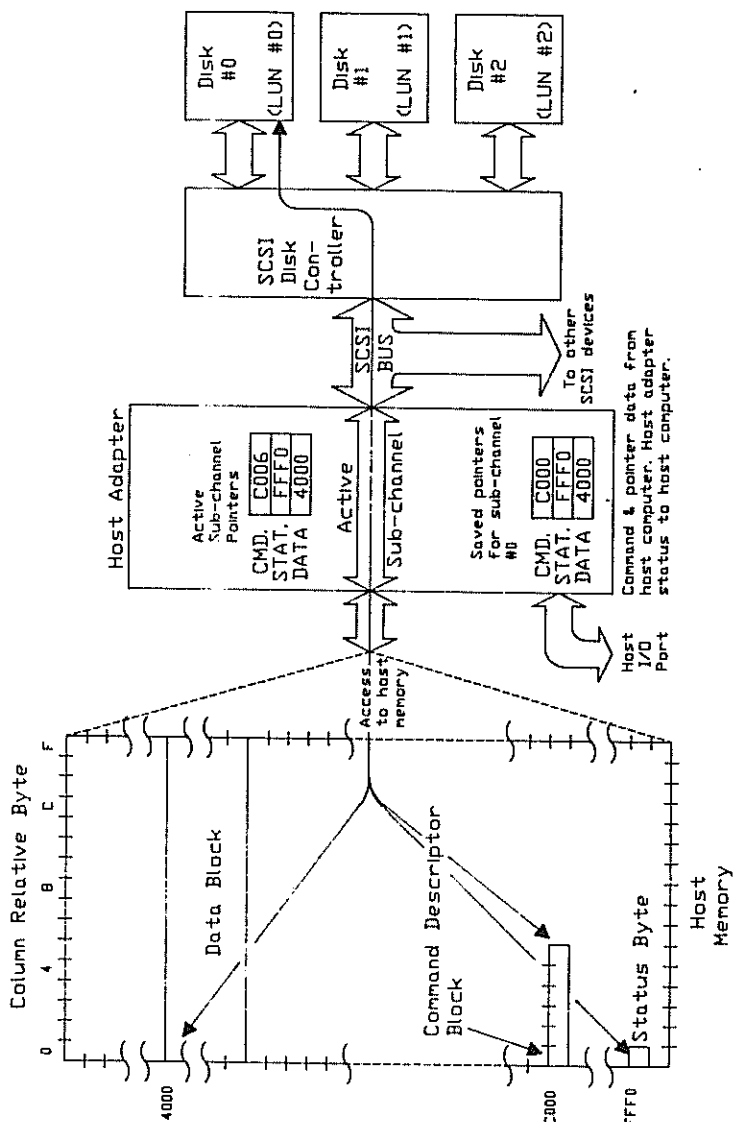


Figure C-2: Snapshot Prior to Data Transfer

Note: The following Appendix is not part of this American National Standard, but is included for information only.

D. Additional Medium Type and Density Code Standards

In Sections 8 and 9 of this standard, the MODE SELECT and MODE SENSE data define medium type codes and density codes for certain flexible disks and magnetic tapes. American National Standards are referenced for code values if a standard exists for that code value. In many cases, other standards or X3 draft documents also exist for a code value. Tables D-1 and D-2 in this Appendix provide additional references to those standards or draft documents.

DISCLAIMER: It is not the purpose of this Appendix to indicate that these standards are exactly equivalent to each other. However, these standards may be useful. Please refer to Sections 8 and 9 for additional information.

Table D-1: Direct-Access Medium-Type Codes

Code	Medium Type		
00h	See Section 8.		
01h	See Section 8.		
02h	See Section 8.		
Flexible-Disk Reference Standard(s)			
05h	X3.73-1980	ECMA-54	ISO 5654-1 : 1984 ISO 5654-2 : 1985
06h		ECMA-59	
09h	(None -- X388 has abandoned this project.)		
0Ah	X3.121-1984	ECMA-69	ISO 7065-1 : 1985 ISO 7065-2 : 1985
0Dh	X3.82-1980	ECMA-66	ISO 6596-1 : 1985 ISO 6596-2 : 1985
12h	X3.125-1985	ECMA-70	ISO 7487-1 : 1985 ISO 7487-2 : 1985 ISO 7487-3 : 1984
16h	X3.126-1986	ECMA-78	DIS 8378/1-1984 DIS 8378/2 DIS 8378/3
1Ah	X3B8/86-32 (Note 1)	ECMA-99	DIS 8630/1-1985 DIS 8630/2-1985
1Eh	X3.137 (Note 1)	ECMA-100	DIS 8860/1-1985 DIS 8860/2-1985
Direct-Access Magnetic Tape Standard(s)			
	ANSI	ECMA	ISO
40h	X3B5/85-151 (Note 2)	ECMA TC19/83/39	
44h	X3B5/85-151 (Note 2)	ECMA TC19/83/39	
80h-FFh	Vendor unique		
All others	Reserved		

NOTES:

(1) These listings are currently under development. Please check with the X3 Secretariat for information concerning status and availability.

(2) This draft document is for unrecorded miniature cartridge media. The usage referred to here is for serial GCR recording using a format known as QIC-100. Since Subcommittee X3B5 issues a new document number for each revision of their working draft document, please contact the chairman of X3B5 for the latest document number.

Table D-2: Sequential-Access Density Codes

Code Value	Density		
00h	See Section 9.		
	Magnetic Tape Reference Standard(s)		
	ANSI	ECMA	ISO
01h	X3.22-1983,	ECMA-62,	ISO 1863-1976
02h	X3.39-1986,	ECMA-62,	ISO 3788-1976
03h	X3.54-1986,	ECMA-62,	ISO 5652-1984
04h	Old format known as QIC-11		
05h	X3.136-1986,	ECMA-98	
06h	X3B5/85-194A (See Note)		
07h	X3.116-1986,	ECMA-79,	ISO 8063/1-1984
08h	X3B5/86-099 (See Note)		
09h	X3B5/86-055 (See Note)		
0Ah	X3B5/85-88 (See Note)		
0Bh	X3.55-1982, X3.56-1986,	ECMA-46,	ISO 4057-1979
80h -- FFh	Vendor unique		
All others	Reserved		

NOTE: Draft document. Subcommittee X3B5 assigns a new document number to each revision of their documents. Please contact the chairman of X3B5 for the latest document number.

E. Data Integrity and I/O Process Queuing

This Appendix demonstrates the practicality of having the target reorder I/O processes which have been queued for a specific logical unit under its control with a minimum of explicit direction by the initiator. A clear and precise written explanation was deemed appropriate. While this appendix is only directly applicable to direct-access devices, the same concepts can be applied to any SCSI device.

This appendix is not intended to indicate how command queuing must be implemented by the target in order to insure correct execution. Rather, it simply illustrates one possible implementation that does insure correctness at a reasonable cost (in overhead and performance) and is easy to analyze.

E.1. Glossary

Unless otherwise stated, all terms used in this Appendix are as defined in the body of this standard. The following terms are new:

correct execution sequence. Any sequence of execution from the I/O process queue(s) for a logical unit that both obeys the rules for I/O process queuing and which results in the state of the media, and the data returned to the initiator concerning the contents of the media, to be identical to those of a first-in first-out (FIFO) execution of the primary queue.

NOTE: The state of other components of the target, such as the buffer, are not guaranteed to be the same under different re-orderings that result in correct execution.

explicit ordered I/O process. This is an I/O process that includes an ORDERED QUEUE TAG message.

implicit ordered I/O process. This is an I/O process that includes a SIMPLE QUEUE TAG message, but the target has determined it will treat as an ordered I/O process for the purposes of queuing.

head of queue queue. This is the queue for a specific logical unit containing head of queue I/O processes for that logical unit.

LBA. An abbreviation for "logical block address".

ordered I/O process. This is an explicit or implicit ordered I/O process.

primary queue. This is the queue for a specific logical unit containing the ordered and unordered I/O processes for that logical unit.

Each primary queue can be divided into a series of one or more segments. Each segment normally consists of a sequence of I/O processes containing zero or more unordered I/O processes and one ordered I/O process such that the ordered I/O process is the last in the sequence and the unordered I/O processes are those which arrived after the ordered I/O process of the previous segment in the queue and before the ordered I/O process in this segment. The last segment in the queue is a special case which may not include an ordered I/O process.

For example, a queue containing commands in the following order:

U U O O U O U O O O U U U U U

can be divided into segments as follows:

(U U O) (O) (U O) (U O) (O) (O) (U U U U U)

where, "U" represents an unordered command and "O" represents an ordered command.

regeneration point. The point in time when no command is under execution and the first I/O process of a new segment in the primary queue is the next I/O process to be executed.

reordering rule. The algorithm used by a target to reorder commands in the primary queue of a logical unit.

state of the media. At any particular moment, the state is defined to be the complete mapping of logical block addresses to the data stored in those logical block addresses. Thus the state is a measure of the contents of the device.

TL. An abbreviation for "transfer length".

E.2. Thesis

The point of this Appendix is that the target can implement reordering rules which result in a correct execution sequence at:

- (1) low cost in command overhead,
- (2) high improvement in performance, and
- (3) without requiring the initiator to explicitly order commands (although such ordering is allowed).

Under any reordering rule, only the reordering done within a queue segment can make the execution sequence incorrect.

This follows directly from the definitions in the above glossary and the entire philosophy of I/O process queuing, under which the explicit ordering of a I/O process or the use of a head of queue I/O process indicates that the initiator is removing any control of order of execution from the target. Doing so shifts any risk that the resulting execution sequence may be "incorrect" from the target to the initiator. A sequence of execution is correct if for each queue segment the execution of I/O processes in that segment, if considered to be the total queue for the logical unit, would be considered to be correct.

Since the order of execution of head of queue I/O processes and the order of execution of queue segments is restricted to a single ordering by the rules of I/O process queuing, only reordering within a segment can create a deviation from the FIFO primary queue execution sequence which is always correct.

Assume all unordered I/O processes other than those containing READ(6), READ(10), WRITE(6), and WRITE(10) commands to be implicitly ordered by the target. (For simplicity, "read" is used for "READ(6) or READ(10)" and "write" is used for "WRITE(6) or WRITE(10)" in the following section.)

Note that this assumption does not significantly decrease the performance gains to be realized by reordering (since the remaining unordered I/O processes still make up over 99.9% of the I/O processes actually encountered during normal execution), nor increase the overhead (since a simple operation code check is all that is required), but will significantly simplify the analysis of reordering rules. Targets might be able to insure correct execution sequence without this restriction, but allowing such commands as MODE SELECT, RESERVE/RELEASE, and FORMAT to be reordered obviously leads to potential difficulties and much complexity for little gain.

The test for correct execution is made at regeneration points. Note that I/O processes cannot be reordered across regeneration points. This implies that halting execution (e.g., for an error) in the middle of a queue segment may leave the state of the media in an incorrect state. As always, it is up to the initiator to successfully perform recovery operations.

All segments (except for the last, which is treated as a special case) are finite, and any reordering algorithm eventually results in reaching a regeneration point. For the last segment, the target insures that all commands are executed in a finite period of time (i.e., starvation does not occur). Many popular reordering algorithms prevent starvation, and the assumption is that one such algorithm is implemented.

Thus the problem has been reduced to requiring that the reordering of I/O processes within a segment does not result in the return of data which differs from that of a FIFO execution nor leaves the media in a different state. Note that under any reordering of a segment, the ordered command is always constrained to be executed last. Thus as long as the data returned and the state of the media for the sequence of unordered I/O processes meets the correctness criteria, then the I/O processes in the segment as a whole are correctly executed.

All unordered I/O processes in a segment contain a variety of either the read or write commands. Consider the N unordered I/O processes in a segment to be numbered 1 to N. Then any reordering is uniquely defined by the $N!$ (N factorial) ordered pairs of I/O processes (x,y) , where each pair implies that I/O process x comes before I/O process y in the reordering.

If all the pairs were (read,read) pairs (i.e., all unordered commands were reads), then any reordering could not affect the state of the media (since it is never changed) nor the returned data. Similarly, if a pair was a (read,write), a (write,read), or a (write,write) then the reordering of these two commands could not affect correctness as long as the range of the specified logical block addresses for each command did not intersect.

Thus the above is both a necessary and sufficient condition for generating a correct execution sequence. However, the target need not generate the $N!$ pairs and perform the check required by theory. A more practical implementation of the above test would be the following:

First, any reordering of I/O processes implies that a sorting operation (usually with respect to the LBA of the command) be performed. The sort may result in an explicit data structure (i.e., a binary tree of pointers) or an implicit structure (i.e., the command descriptor blocks are reordered in an array, or an array of pointers to command descriptor blocks are reordered). In any event, denote T as the time required to perform such sorting, and denote A as the resulting sequence of execution.

This list is now sorted so that the LBA+TL of the immediately preceding command is $\leq$ the LBA of the next command. Note that LBA+TL is one more than the last LBA in the command, and this sort can be performed at a cost no greater than T (note that LBA+TL must be computed for each command anyway in order to perform a range check against the logical unit's maximum LBA, and that a more sophisticated data structure can reduce the incremental effort to perform this second sort considerably). This ordering is denoted as B .

For each segment, a I/O process has a position in both queues denoted by the pair (a,b) . The execution sequence is then determined as follows:

- 1) Attempt to execute I/O process in the ordering determined by A .
- 2) If $a = b$, then execute the I/O process.
- 3) If $a < b$, then scan A until a command equaling b is found. For all commands in A between a and this b , search B and keep track of the I/O process that appears last in B (denote this c). Now scan A again, but use c as the search target instead of b . Continue the search process, alternating between A and B , until all I/O processes to search for are exhausted. The result is a subsequence of I/O processes in A and B such that each I/O process in the subsequence in A appears in the subsequence in B and vice-versa, but the orders differ between the subsequences. These I/O processes should be executed in the original FIFO order (i.e., both re-orderings should be ignored).
- 4) when done, go to step 1) again until the queue is empty.

As an example, considering the following pairs of ordered LBA ranges:
 (0,3) (6,8) (7,12) (8,15) (20,23) (28,32) (31,35) (36,39) (37,38)
 (0,3) (7,12) (6,8) (8,15) (20,23) (31,35) (28,32) (37,38) (36,39)

Thus the execution order is:

(0,3)
 (6,8) (7,12) in FIFO order
 (8,15)
 (20,23)
 (28,32) (31,35) in FIFO order
 (36,39) (37,38) in FIFO order

Note that other execution sequences may be defined that provide greater performance (i.e., (read,read) sequences can be freely reordered) at a cost of greater overhead. But in the normal case of few intersections, the total overhead is $2 \times I$ plus a check per I/O process (this can grow to $N \times N$ checks in the worse case).

Finally, overhead should not be an issue in I/O process queuing. Overhead grows as the queue lengthens, but the opportunity to overlap queuing tasks with seek time and rotational latency also grows with the queue length. Thus most, if not all, of the queuing overhead can be effectively hidden.

Explicit ordering of I/O processes by the initiator can shift the the implementation burden from target to initiator, and this may have many practical benefits. Error recovery might prove easier to implement, and target resources might be more profitably used.

F. Power On Protocols - Recommended Initialization Procedure

This appendix describes the normal mechanisms for obtaining the information required for system initialization from SCSI-2 devices as well as all SCSI devices meeting conformance level 2 as defined in Appendix E of SCSI-1. This procedure documents the steps required to obtain this information and to achieve the desired initial states in the attached devices.

F.1. System Initialization

The following list of information is assumed necessary and sufficient for normal system initialization:

- 1) A list of each installed and powered on SCSI device for each SCSI address. SCSI devices that are not powered on are treated as not installed, assuming that the terminators are powered from a source other than the unpowered SCSI devices.
- 2) A list of the installed logical units for each SCSI device. Unpowered or failing logical units may not be completely identifiable.
- 3) The device type for each available logical unit.
- 4) The manufacturer and model for each available logical unit. (This information may not be available for SCSI-1 devices)
- 5) The critical device type information for each available logical unit. This information varies depending on the device type.
- 6) Extended functionality of SCSI devices such as target role capability in devices that are principally initiators, AEN capability, etc.

The following states are established for each attached logical unit that has power available and is not failing:

- 1) The ready state for each available logical unit, including any required medium initialization, but not formatting.
- 2) All error conditions associated with the starting process are cleared.
- 3) All UNIT ATTENTION conditions are cleared.
- 4) All data transfer parameters are established.
- 5) All pertinent system tuning parameters are established where known. Note that these may be modified dynamically to improve the performance characteristics of the system.

The following procedures show the sequences necessary to implement a system that initializes itself with a minimum of information available at power on time. Many real systems are not so generalized, having considerable information available about the proper configuration at power-on time. In those cases, the sequence steps that would have been necessary to obtain that predefined information may be skipped or ignored.

F.2. General Procedure for Initializing Devices

The system should execute the following steps to perform initialization. Some of the steps are detailed in subsequent paragraphs. Note that the text represents a primitive pseudo-code that can be converted to the appropriate software object code by those who implement device drivers.

F.2.1. General Procedure Executed by Initiators

Initiator Activities:

Power On: It is assumed that each SCSI device, as it is powered on, performs appropriate internal reset operations and internal test operations. Once powered on, initiators that have target capability should be prepared to respond to a selection within a system-specific time.

Reset: At power-on time, it is likely that an SCSI device has caused errors to the ongoing activities on the SCSI bus. A bus reset should be generated to notify attached devices that any activities that may have been occurring should be restarted.

Find Devices: Each SCSI address other than the initiator's SCSI address should be tested to determine if an SCSI device responds. If an SCSI device responds, an INQUIRY command to logical unit 0 should be executed. The information obtained indicates the device type, manufacturer, and model of the attached logical unit 0 if the response data format field is one or two. If the response data format field is zero, only the device type field is valid. In addition, the version of the command set supported by the device is indicated by the ANSI-approved version field.

Find logical units: Each possible logical unit number on the attached targets should be tested for existence using an INQUIRY command. Those found with a non-zero peripheral qualifier in the INQUIRY data should not be included in the list of available logical units. Each available logical unit should be added to the host configuration information, identifying the associated logical unit number, device type, manufacturer, and model.

Verify State: The verify state test (see F.2.3) should be made to clear any outstanding errors, capture and clear UNIT ATTENTION conditions, and determine the state of readiness of the available logical units. The logical units should be identified as ready, not ready, or failing by this test.

Device Initialization: The device undergoes a device-dependent initialization process. This process is described for direct-access devices, sequential-access devices, and processor devices. Other device initialization procedures are not described since they tend to be similar to one of these initialization procedures. The initialization process takes into account the state of the device as identified during the verify state test.

Device On-line: The successful completion of the device initialization process allows the device-table entry to be fully enabled. The device joins the system with all key parameters identified and initialized. The device state is known and may be presented to the system operator.

F.2.2. Procedure Executed by Temporary Initiators

A temporary initiator typically performs initiator operations only under the direction of a host processor. Such operations may include read and write commands associated with management of a COPY command. Other possible operations include issuing a SEND command associated with asynchronous event notification. As such, temporary initiators need not completely perform the last two steps defined above. Since all commands are managed by a host processor, temporary initiators normally need not recover information about the media density, the sparing algorithms, or other detailed information that may be required only by a host processor.

F.2.3. Verify State Test

The verify state test uses the following steps to identify any outstanding errors, clear any UNIT ATTENTION conditions, and determine the readiness of the devices. The verify state test should be executed against each available logical unit.

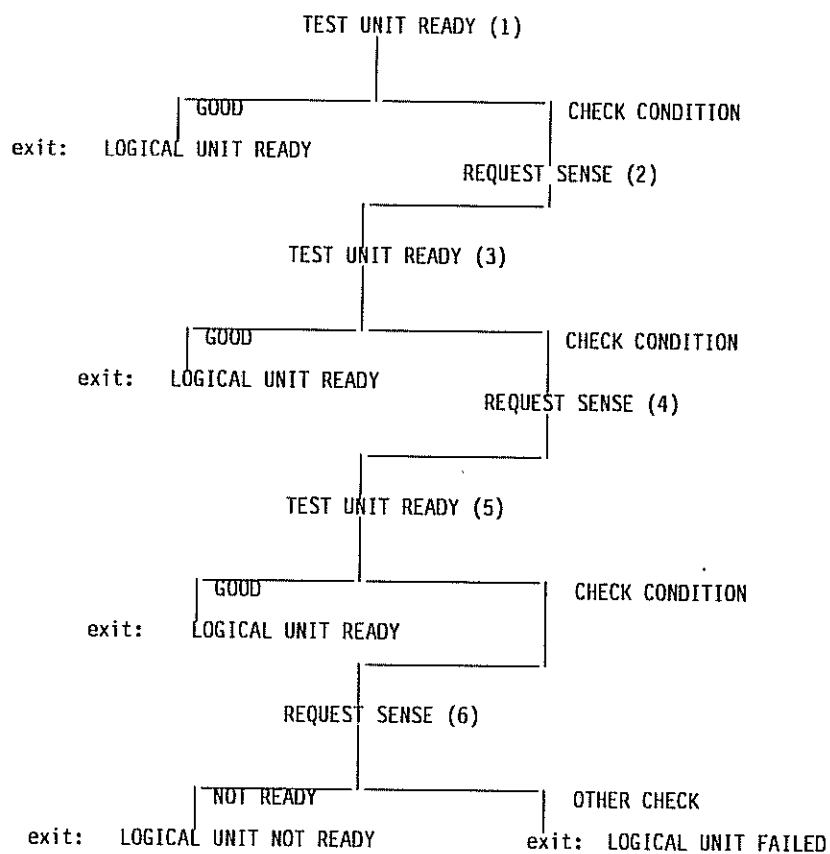


Figure F-1: Verify State Test

TEST UNIT READY (1):

This TEST UNIT READY command is used to determine if any outstanding CHECK CONDITION or UNIT ATTENTION condition exists. If not, the device is indicated to be ready.

REQUEST SENSE (2):

This REQUEST SENSE command is used to clear the outstanding CHECK CONDITION. Most SCSI-2 logical units return UNIT ATTENTION sense key in this sense information.

TEST UNIT READY (3):

This TEST UNIT READY command is used to see if the UNIT ATTENTION condition or other error was successfully cleared. In some special cases, another error may have been nested with the UNIT ATTENTION and this TEST UNIT READY command may also return CHECK CONDITION status.

REQUEST SENSE (4):

This REQUEST SENSE command is used to determine which error or exception was associated with the CHECK CONDITION status returned by the TEST UNIT READY (3) command. In addition, this REQUEST SENSE command is used to clear the outstanding CHECK CONDITION. This may be a NOT READY sense key or another unexpected error.

TEST UNIT READY (5):

This TEST UNIT READY command is used to see if all outstanding CHECK CONDITION statuses have finally been cleared. If so, the logical unit is identified as ready.

REQUEST SENSE (6):

This REQUEST SENSE command is used to determine why there is a persistent CHECK CONDITION status. If the sense key is NOT READY, the logical unit is identified as not ready. If the sense key indicates some other failure, the logical unit is identified as failing and the sense key is logged in the appropriate area.

IMPLEMENTORS NOTE: Commands that receive BUSY or RESERVATION CONFLICT status should be re-issued until some other status is received.

F.3. Direct-Access Device Initialization Procedure

The device-dependent initialization process for a direct-access device may be divided into three independent activities. The first activity enables the minimum logical function required for execution of READ commands on the boot device. The second activity is performed on all direct-access devices, including the boot device. It establishes all required initial parameters and operating conditions. The third activity is performed on direct-access devices that have never been formatted or initialized. This activity is normally performed by an initialization utility program.

F.3.1. Boot Device Initialization Procedure

It is assumed that the boot program and boot device have been prepared in such a manner that proper block lengths, data file contents, and logical addresses have been implemented by both the boot device and the boot program. The boot program prepares the boot device for operation in the following manner:

Verify Ready: The state of the device as determined by the verify state test (see F.2.3) is examined. If the test indicates that the required drive has failed, the boot device initialization is not performed and appropriate error indications are presented.

Start Device: A START STOP UNIT command should be issued with the start bit set to one. The Immed bit should be set to zero in order to guarantee that the returned status reflects the completion of the device start operation. A disconnect operation is likely to occur since the start process may take a considerable period of time. If system-controlled power sequencing of the peripheral devices is required, it is done by managing the timing relationship of the START STOP UNIT commands to different logical units.

If GOOD status is returned, the next step should be started. If CHECK CONDITION status is returned, a REQUEST SENSE command is issued to determine what error condition was detected. If an ILLEGAL REQUEST sense key is found, the START STOP UNIT command was not supported by the target or peripheral device and the next step should be started. If any other error is detected (BUSY status is not an error), the boot device initialization should be terminated and appropriate error indications should be presented.

Verify Ready / Spinning: A verify state test should be performed. If the device is ready, the next step should be started. If the device is not ready or failing, the boot device initialization should be terminated and appropriate error indications should be presented.

Boot: The boot READ commands can now be started on the boot device. It is assumed that the information read includes the programs that are required to continue the system initialization and bring-up process, including the necessary programs and device drivers to perform the other system initialization procedures.

F.3.2. General Direct-Access Device Initialization Procedure

A general direct-access device initialization procedure is defined below. The initialization procedure should be executed for each attached logical unit that has been identified as a direct-access device. Execution of this procedure may be overlapped from one logical unit to another. If the initiator supports only a limited range of devices, parts of this procedure may be skipped or simplified.

Verify Ready: The state of the device as determined by the verify state test (see F.2.3) should be examined. If the device has failed, the general direct-access device initialization should not be performed and appropriate error indications should be presented.

Start Device: A START STOP UNIT command should be issued with the start bit set to one. If GOOD status is returned, the next step should be started. If CHECK CONDITION status is returned, a REQUEST SENSE command should be issued to determine what error condition was detected. If an ILLEGAL REQUEST sense key is found, the START STOP UNIT command was not supported by the device and the next step should be started. If any other error is detected (BUSY status is not an error), the general direct-access device initialization procedure should be terminated on this logical unit and appropriate error indications should be presented.

Verify Ready / Spinning: A verify state test (see F.2.3) should be performed. If the device is ready, the next step should be started. If the device is not ready or failing, the general direct-access device initialization should be terminated for this logical unit and appropriate error indications should be presented.

Determine Parameters: If the ANSI-approved version field of the previously executed INQUIRY command was 0 or 1, the MODE SENSE information, if any, may be vendor specific and this function should be skipped unless required by the vendor specific initialization protocols.

If the ANSI-approved version field is 2, optional MODE SENSE information is by this standard. In this case, a MODE SENSE command should be executed with the page control field set to request current values and the page code field set to request all pages. A record should be made of the current values.

If a CHECK CONDITION status is returned to the MODE SENSE command, then a REQUEST SENSE command should be issued. If the sense key is ILLEGAL REQUEST, then the target does not support the MODE commands. In this case, the initiator should skip to the determine capacity step.

A second MODE SENSE command should be executed with the page control field set to request changeable values and the page code field set to request all pages. A record should be made of the changeable values.

Any errors that occur during the two MODE SENSE commands should be recorded and the initialization for the failing logical unit should be terminated.

Set Parameters: If the ANSI-approved version field is 0 or 1, the initialization operation may be vendor specific and may be executed according to the vendor's rules for the peripheral device. The system is assumed to have some other source of information concerning these requirements or it may skip this step, accepting the target's default parameters.

If the ANSI-approved version field is 2, the optional MODE SELECT command is defined by this standard. The actual requirements for the parameters are characteristic of the particular system and should be known to the system. The current values and the changeable values obtained from the previous MODE SENSE commands should be examined to see if the system's requirements are satisfied and if the parameters can be modified. If all values are correct, the remainder of this step may be skipped. If modifications need to be made to the changeable values, a MODE SELECT command should be issued to modify the appropriate pages. This may include modifying error recovery parameters or performance tuning parameters. Most geometry parameters should not be modified during general direct-access device initialization.

Determine Capacity: The capacity and block size of the logical unit are determined by issuing a READ CAPACITY command. The information is stored for access by the system device drivers.

The direct-access device is now fully initialized and all required information has been made available to the system. When all available non-failing devices have been initialized, the system initialization is considered complete.

F.3.3. Direct-Access Device Medium Initialization Procedure

The following initialization procedure is not part of normal power-up system initialization. It is assumed to be performed after completion of the general system initialization process but uses only the INQUIRY data information obtained during that process. It is performed to initialize the device medium and is normally performed only by an initialization utility program.

Determine Format Requirement: The requirement to perform a format operation is normally generated by an operator who has just installed a new device known to require formatting. It may also be generated by recognition that the device has information that is no longer valid and should be totally erased. It may also be generated by changes in system requirements, including different block sizes. Finally, reformatting may also be required to restructure the defect management.

The general direct-access device initialization procedure may have identified the device as failing because of the inability of the device to recover the READ CAPACITY parameters. The device is assumed to have been started during the general direct-access device initialization procedure. The verify state test should be executed again. The device should be ready according to that test. If the logical unit is not ready or failing, the direct-access device medium initialization procedure should be terminated and appropriate error indications should be presented.

If it was determined in the general direct-access device initialization procedure that the target does not support the MODE commands, then the initiator should either proceed to the perform format operation step or it should perform the determine format parameters and set format parameters steps in the vendor-specified manner.

Determine Format Parameters: If the ANSI-approved version field is 0 or 1, the direct-access device medium initialization procedure may be vendor specific and should be executed according to the vendor's rules for the peripheral device. The system is assumed to have some other source of information concerning these requirements or to be willing to accept the target's default format.

If the ANSI-approved version field is 2, a MODE SENSE command should be issued with the page control field set to current values and the page code field set to return all pages. A MODE SENSE command should be issued again with the page control field set to changeable values and the page code field set to return all pages. The information returned by the two MODE SENSE commands indicates what values should be provided by the system to complete the format parameters. If either of these MODE SENSE operations does not complete normally, the media initialization operation should be terminated and appropriate error indications should be presented.

Set Format Parameters: If the ANSI-approved version field is 0 or 1, the format requirements may be vendor specific and the appropriate commands should be known to the initialization utility or it should be willing to accept the target's default format. Those format preparation commands, if any, should be executed at this time.

If the ANSI-approved version field is 2 and the target supports the MODE commands, the logical unit should be prepared for medium formatting by executing a MODE SELECT command. The necessary formatting parameters are selected to meet the system requirements and are placed into the changeable value locations. The MODE SELECT command is then issued. If the command fails, the media initialization procedure should be terminated and appropriate error indications should be presented. If the command succeeds, the next step should be performed.

Perform Format Operation: After the appropriate format parameters are established, the FORMAT command should be executed. The FORMAT parameters depend on the system requirements and the device capabilities. These parameters should be made easily variable in the operating system architecture so that modifications can be performed when system or device requirements change. An error may be returned if improper format parameters are selected. If the FORMAT command fails, the media initialization procedure should be terminated and appropriate error indications should be presented. If the command succeeds, the device is fully operational and the next step should be performed.

Set Parameters: If the ANSI-approved version field is 0 or 1 or if the target does not support the MODE commands, the initialization operation may be vendor specific and may be executed according to the vendor's rules for the peripheral device. The system is assumed to have some other source of information concerning these requirements or it may skip this step, accepting the target's default parameters.

If the ANSI-approved version field is 2, the optional MODE SELECT command is defined by this standard. The actual requirements for the parameters are characteristic of the particular system and should be known to the system. The current values and the parameters established by the MODE SELECT and FORMAT commands should be examined to determine if the system requirements are satisfied and if the parameters should be modified. If all values are correct, the remainder of this step may be skipped. If modifications need to be made to the changeable values, a MODE SELECT command should be issued to modify the appropriate pages. This may include modifying error recovery parameters or performance tuning parameters. Most geometry parameters were established by the storing of parameters during the MODE SELECT and FORMAT commands and should not be modified.

Determine Capacity: The capacity and block size of the logical unit should be determined by issuing a READ CAPACITY command. The information should be stored for access by the system device drivers.

Upon completion of this procedure the device should be initialized and prepared to partake in system-oriented activities. Other system initialization operations may also be required, including the establishment of system volume labels, tables of contents, and other structures.

F.4. Sequential Access Device Initialization Procedure

The initialization process for a sequential-access device may be divided into two independent activities. The first activity establishes all required initial parameters and operating conditions for the identified devices. The second activity performs any required medium initialization for the available logical units.

F.4.1. General Sequential-Access Device Initialization

A general sequential-access device initialization procedure is defined below. The initialization procedure should be executed for each attached logical unit that has been identified as a sequential-access device. Execution of this procedure may be overlapped from one logical unit to another. If initiator supports only a limited range of devices, parts of this procedure may be skipped or simplified.

Verify Ready: The state of the device as determined by the verify state test (see F.2.3) should be examined. If the device has failed, the general direct-access device initialization should not be performed and appropriate error indications should be presented.

Start Device: A LOAD UNLOAD command should be issued with the load bit set one. If GOOD status is returned, the next step should be started. If CHECK CONDITION status is returned, a REQUEST SENSE command should be issued to determine what error condition was detected. If an ILLEGAL REQUEST sense key is found, the LOAD UNLOAD command was not supported by the device and the next step should be started. If any other error is detected, the device initialization procedure should be terminated on this logical unit and appropriate error indications should be presented.

Verify Ready / Loaded: If necessary, a verify state test (see F.2.3) should be performed. If the device and medium are ready, the next step should be started. If a NOT READY sense key is reported, manually loading the medium or activating a switch mechanism may be required to establish the ready state for the device. If any other error is detected, the device initialization procedure should be terminated and the appropriate error indications should be presented.

Determine Parameters: A READ BLOCK LIMITS command should be issued to determine the range of block sizes supported by the device. Following this command, a MODE SENSE command should be issued to determine additional operating parameters of the device. If the ANSI-approved version field of the previously executed INQUIRY command is 0 or 1, any MODE SENSE data following the header and block descriptor is vendor specific. If the ANSI-approved version field is 2, additional pages of MODE SENSE data may be available as defined in this standard. In this case, a MODE SENSE command should be issued with the page code field set to return all pages. If any unrecovered errors are detected during execution of the READ BLOCK LIMITS or MODE SENSE commands, the device initialization process should be terminated and the appropriate error indications should be presented.

Set Parameters: Specific system requirements may require that certain operating parameters be changed from the values reported in the previously executed MODE SENSE command. If changes are required, a MODE SELECT command should be issued to modify the appropriate parameters. This may include error recovery parameters, performance tuning parameters, or other basic operating parameters. If any unrecovered error occurs during this step, the device initialization process should be terminated and the appropriate error indications should be presented. If no change is required or no unrecovered error occurs, the general sequential-access device initialization procedure is complete.

F.5. Asynchronous Event Notification Initialization Procedure

A target using asynchronous event notification, must first execute an initialization procedure. This initialization procedure allows the target device to determine which SCSI devices are capable and willing to receive an asynchronous event notification. Parameters that affect asynchronous event notification within the target device is specified in the control mode page.

The initialization procedure is performed at power-up (after waiting the recommended 10 seconds for all devices to be able to respond and waiting the time specified in the control mode page). It may also be performed following a reset condition, or when a target becomes aware of another SCSI device, or following the issuance of the control mode page or prior to a device issuing an asynchronous event notification.

The target device that uses asynchronous event notification must determine which devices on the bus are capable of receiving an asynchronous event notification. This is done by the target device becoming a temporary initiator and selecting each SCSI device. If the SCSI device responds to selection, the verify state test (see F.2.3) is performed. If the verify state test fails, then the SCSI device does not support asynchronous event notification. If the verify state test succeeds then an INQUIRY command is issued to logical unit zero. The peripheral qualifier field in the INQUIRY data is examined to determine if the SCSI device is a processor device type and then the AENC bit is examined. An AENC bit of zero indicates that asynchronous event notification is not supported by the SCSI device. An AENC bit of one, indicates that asynchronous event notification is supported by the SCSI device.

Disabling of asynchronous event notification can be done by using a vendor-specific hardware mechanism (e.g., switch or jumper), or by issuing control mode pages to devices that support saved parameters.

G. Fast SCSI Skew Time

This Appendix is included to explain the skew budget for the fast SCSI option which is defined in Section 4.

Synchronous transfer rates using a transfer period between 100 ns and 200 ns are known as the "fast SCSI" option. Fast data transfer times have been tested using the following skew budget (Figure G-1) with the differential alternative using transceivers with 25 meters of 0.08042 square mm (28 AWG) twisted pair cable as specified in 4.2.3. The transceivers were subjected to a maximum temperature difference of 25 degrees celsius and a maximum of 200 mV of V_{CC} difference.

CC

FAST SCSI JITTER BUDGET			
#	parameter	+/-budget	
a	clock offset	5	TRANSMITTER
b	transmitting logic skew	3	
c	foil delay	1	
d	transmitter prop. delay skew	6	
e	foil delay	1	
f	drop cable prop. delay	1	
CONNECTOR			
g	external cable - skew between pairs	5	CABLE
h	distortion due to cable imbalance	1	
i	distortion due to intersymbol interference	2	
j	bias distortion	2	
CONNECTOR			
k	drop cable prop. delay	1	RECEIVER
l	foil delay	1	
m	receiver skew	9	
n	foil delay	1	
o	logic setup/hold	5	
TOTAL		44nS	

Figure G-1: Fast SCSI Jitter Budget

Table #	parameter in 4.7-8	value
g	Fast Cable Skew Delay	5
h - n	Fast Deskew Delay	~20
o	Fast Hold Time	~10
*	Fast Assertion Period	30
*	Fast Negation Period	30

Figure G-2: Mapping of Jitter to SCSI

NOTES:

- (1) Values preceded with "~" are rounded up from the numbers shown in the previous table.
- (2) The assertion and negation pulse widths are derived from isolated pulse measurements and represent a minimum pulse width with a satisfactory margin.

The maximum driver skew allowed was 6 ns ($t_{min.} - t_{max.}$) and the maximum receiver skew tested was 9 ns ($t_{min.}^{PLH} - t_{max.}^{PHL}$). Values greater than these could be used if other numbers could be reduced -- the sum is what is important.

Fast data transfer timing parameters were not tested for the single-ended transceiver option prior to publication of this standard.

Mapping the above jitter or skew budget to the SCSI format in 4.7 and 4.8 is done in Figure G-2.

H. Other SCSI Standardization Activities

This appendix provides information on other formal standardization activities related to SCSI.

H.1. SCSI-3 Standards Project

Accredited Standards Committee X3 has approved a project proposal to maintain and enhance the SCSI-2 standard. This project is assigned to the X3T9.2 Task Group which developed this standard and the SCSI-1 standard. Please contact the Chairman of X3T9.2 for further information concerning this project.

H.2. Digital Data Exchange for Color Electronic Prepress Systems

Accredited Standards Committee IT8 is developing a standard for the exchange of digital data between color electronic prepress systems and direct digital color proofers. These are devices that prepare color pictures for high quality color printing. Please contact the Secretary of IT8 for further information concerning this project.

H.3. Fiber Channel

Accredited Standards Committee X3 has approved a project proposal to develop a fiber optic channel physical layer for the Intelligent Peripheral Interface (IPI), SCSI, and the High Speed Channel (HSC). This project is assigned to the X3T9.3 Task Group. Please contact the Chairman of X3T9.3 for further information concerning this project.

I. Numeric Order Codes

This Appendix contains SCSI-2 additional sense codes and operation codes in numeric order as a reference. In the event of a conflict with the definitions of these codes in the body of this standard, those definitions should be regarded as correct.

Table I-1: ASC and ASCQ Assignments

ASC AND ASCQ ASSIGNMENTS					
D				= DIRECT ACCESS DEVICE	
T				= SEQUENTIAL ACCESS DEVICE	
L				= PRINTER DEVICE	
P				= PROCESSOR DEVICE	
W				= WRITE ONCE READ MULTIPLE DEVICE	
R				= READ ONLY (CD-ROM) DEVICE	
S				= SCANNER DEVICE	
O				= OPTICAL MEMORY DEVICE	
M				= MEDIA CHANGER DEVICE	
C				= COMMUNICATION DEVICE	
BYTE					
12	13	DTLPWRSOMC	DESCRIPTION	COMMENTS	

00	00	DTLPWRSOMC	NO ADDITIONAL SENSE INFORMATION		
00	01	T	FILEMARK DETECTED		
00	02	T S	END-OF-PARTITION/MEDIUM DETECTED		
00	03	T	SETMARK DETECTED		
00	04	T S	BEGINNING-OF-PARTITION/MEDIUM DETECTED		
00	05	T S	END-OF-DATA DETECTED	00 06 DTLPWRS	
00	11	R	AUDIO PLAY OPERATION IN PROGRESS	I/O PROC	
00	12	R	AUDIO PLAY OPERATION PAUSED	TERM.	
00	13	R	AUDIO PLAY OPERATION SUCCESSFULLY COMPLETED		
00	14	R	AUDIO PLAY OPERATION STOPPED DUE TO ERROR		
00	15	R	NO CURRENT AUDIO STATUS TO RETURN		
01	00	D W O	NO INDEX/SECTOR SIGNAL		
02	00	D WR OM	NO SEEK COMPLETE		
03	00	DTL W SO	PERIPHERAL DEVICE WRITE FAULT		
03	01	T	NO WRITE CURRENT		
03	02	T	EXCESSIVE WRITE ERRORS		
04	00	DTLPWRSOMC	LOGICAL UNIT NOT READY, CAUSE NOT REPORTABLE		
04	01	DTLPWRSOMC	LOGICAL UNIT IS IN PROCESS OF BECOMING READY		
04	02	DTLPWRSOMC	LOGICAL UNIT NOT READY, INITIALIZING COMMAND REQUIRED		
04	03	DTLPWRSOMC	LOGICAL UNIT NOT READY, MANUAL INTERVENTION REQUIRED		
04	04	DTL O	LOGICAL UNIT NOT READY, FORMAT IN PROGRESS		
05	00	DTL WRSOMC	LOGICAL UNIT DOES NOT RESPOND TO SELECTION		
06	00	D WR OM	NO REFERENCE POSITION FOUND (TRACK ZERO OR EQUIVALENT)		
07	00	DTL WRSOM	MULTIPLE PERIPHERAL DEVICES SELECTED		
08	00	DTL WRSOMC	LOGICAL UNIT COMMUNICATION FAILURE		
08	01	DTL WRSOMC	LOGICAL UNIT COMMUNICATION TIME-OUT		
08	02	DTL WRSOMC	LOGICAL UNIT COMMUNICATION PARITY ERROR		

00 06 DTLWRSOMC
I/O PROCESS
TERMINATED

Table I-1: ASC and ASCQ Assignments (Continued)

```

=====
09 00 DT WR 0 TRACK FOLLOWING ERROR
09 01 WR 0 TRACKING SERVO FAILURE
09 02 WR 0 FOCUS SERVO FAILURE
09 03 WR 0 SPINDLE SERVO FAILURE
0A 00 DTLPWRSMC ERROR LOG OVERFLOW
0B 00
0C 00 T S WRITE ERROR SENSE KEY SAYS WHETHER RECOVERED
0C 01 D W 0 WRITE ERROR RECOVERED WITH AUTO REALLOCATION
0C 02 D W 0 WRITE ERROR - AUTO REALLOCATION FAILED
0D 00
0E 00
0F 00
10 00 D W 0 ID CRC OR ECC ERROR
11 00 DT WRSO UNRECOVERED READ ERROR
11 01 DT W SO READ RETRIES EXHAUSTED
11 02 DT W SO ERROR TOO LONG TO CORRECT
11 03 DT W SO MULTIPLE READ ERRORS
11 04 D W 0 UNRECOVERED READ ERROR - AUTO REALLOCATE FAILED
11 05 WR 0 L-EC UNCORRECTABLE ERROR (L-EC CODES PRESENT AND L-EC ON)
11 06 WR 0 CIRC UNRECOVERED ERROR (L-EC UNAVAILABLE OR OFF)
11 07 W 0 DATA RESYNCHRONIZATION ERROR
11 08 T INCOMPLETE BLOCK READ (POSTAMBLE NOT FOUND)
11 09 T NO GAP FOUND
11 0A DT 0 MISCORRECTED ERROR
12 00 D W 0 ADDRESS MARK NOT FOUND FOR ID FIELD
13 00 D W 0 ADDRESS MARK NOT FOUND FOR DATA FIELD
14 00 DTL WRSO RECORDED ENTITY NOT FOUND
14 01 DT WR 0 RECORD NOT FOUND
14 02 T FILEMARK OR SETMARK NOT FOUND
14 03 T END-OF-DATA NOT FOUND
14 04 T BLOCK SEQUENCE ERROR
15 00 DTL WRSOM RANDOM POSITIONING ERROR
15 01 DTL WRSOM MECHANICAL POSITIONING ERROR
15 02 DT WR 0 POSITIONING ERROR DETECTED BY READ OF MEDIUM
16 00 D W 0 DATA SYNCHRONIZATION MARK ERROR
17 00 DT WRSO RECOVERED DATA WITH NO ERROR CORRECTION APPLIED
17 01 DT WRSO RECOVERED DATA WITH RETRIES
17 02 DT WR 0 RECOVERED DATA WITH POSITIVE HEAD OFFSET
17 03 DT WR 0 RECOVERED DATA WITH NEGATIVE HEAD OFFSET
17 04 WR 0 RECOVERED DATA WITH RETRIES AND/OR CIRC APPLIED
17 05 D WR 0 RECOVERED DATA USING PREVIOUS SECTOR ID
17 06 D W 0 RECOVERED DATA WITHOUT ECC. DATA AUTO-REALLOCATED.
18 00 DT WR 0 RECOVERED DATA WITH ERROR CORRECTION APPLIED
18 01 D WR 0 RECOVERED DATA WITH ERROR CORRECTION AND RETRIES APPLIED
18 02 D WR 0 RECOVERED DATA WITH ECC AND/OR RETRIES. DATA AUTO-REALLOCATED
18 03 R RECOVERED DATA WITH CIRC
18 04 R RECOVERED DATA WITH LEC
=====

```

Table I-1: ASC and ASCQ Assignments (Continued)

```

=====
19 00 D 0 DEFECT LIST ERROR
19 01 D 0 DEFECT LIST NOT AVAILABLE
19 02 D 0 DEFECT LIST ERROR IN PRIMARY LIST
19 03 D 0 DEFECT LIST ERROR IN GROWN LIST
1A 00 DTLPWRSMC PARAMETER LIST LENGTH ERROR
1B 00 DTLPWRSMC SYNCHRONOUS DATA TRANSFER ERROR
1C 00 D 0 DEFECT LIST NOT FOUND
1C 01 D 0 PRIMARY DEFECT LIST NOT FOUND
1C 02 D 0 GROWN DEFECT LIST NOT FOUND
1D 00 D W 0 MISCOMPARE DURING VERIFY OPERATION
1E 00 D W 0 RECOVERED ID WITH ECC CORRECTION
1F 00
20 00 DTLPWRSMC INVALID COMMAND OPERATION CODE
21 00 DT WR OM LOGICAL BLOCK ADDRESS OUT OF RANGE
21 01 M INVALID ELEMENT ADDRESS
22 00 D ILLEGAL FUNCTION FOR DEVICE TYPE
                                SHOULD USE 20 00, 24 00, OR 26 00
23 00
24 00 DTLPWRSMC INVALID FIELD IN CDB CHECK FIELD POINTER IN SENSE DATA
25 00 DTLPWRSMC LOGICAL UNIT NOT SUPPORTED
26 00 DTLPWRSMC INVALID FIELD IN PARAMETER LIST CHECK FIELD POINTER IN SENSE DATA
                                CHECK FIELD POINTER IN SENSE DATA
26 01 DTLPWRSMC PARAMETER NOT SUPPORTED CHECK FIELD POINTER IN SENSE DATA
26 02 DTLPWRSMC PARAMETER VALUE INVALID CHECK FIELD POINTER IN SENSE DATA
26 03 DTLPWRSMC THRESHOLD PARAMETERS NOT SUPPORTED CHECK FIELD POINTER IN SENSE DATA
27 00 DT W 0 WRITE PROTECTED
28 00 DTLPWRSMC NOT READY TO READY TRANSITION MEDIUM MAY HAVE CHANGED
28 01 M IMPORT / EXPORT ELEMENT ACCESSED
29 00 DTLPWRSMC POWER ON, RESET, OR BUS DEVICE RESET OCCURRED
2A 00 DTL WRSOMC PARAMETERS CHANGED
2A 01 DTL WRSOMC MODE PARAMETERS CHANGED
2A 02 DTL WRSOMC LOG PARAMETERS CHANGED
2B 00 DTLPWRSMC COPY CANNOT EXECUTE SINCE HOST CANNOT DISCONNECT
2C 00 DTLPWRSMC COMMAND SEQUENCE ERROR
2C 01 S TOO MANY WINDOWS SPECIFIED
2C 02 S INVALID COMBINATION OF WINDOWS SPECIFIED
2D 00 T OVERWRITE ERROR ON UPDATE IN PLACE
2E 00
2F 00 DTLPWRSMC COMMANDS CLEARED BY ANOTHER INITIATOR
30 00 DT WR OM INCOMPATIBLE MEDIUM INSTALLED
30 01 DT WR 0 CANNOT READ MEDIUM - UNKNOWN FORMAT
30 02 DT WR 0 CANNOT READ MEDIUM - INCOMPATIBLE FORMAT
30 03 DT CLEANING CARTRIDGE INSTALLED
31 00 DT W 0 MEDIUM FORMAT CORRUPTED
31 01 D L 0 FORMAT COMMAND FAILED
32 00 D W 0 NO DEFECT SPARE LOCATION AVAILABLE
32 01 D W 0 DEFECT LIST UPDATE FAILURE
33 00 T TAPE LENGTH ERROR
34 00
35 00
=====

```

Table I-1: ASC and ASCQ Assignments (Continued)

```

=====
36 00 L RIBBON/INK/TONER FAILURE
37 00 DTL WRSOMC ROUNDED PARAMETER
38 00
39 00 DTL WRSOMC SAVING PARAMETERS NOT SUPPORTED
3A 00 DTL WRSOMC MEDIUM NOT PRESENT
3B 00 TL SEQUENTIAL POSITIONING ERROR
3B 01 T TAPE POSITION ERROR AT BEGINNING-OF-MEDIUM
3B 02 T TAPE POSITION ERROR AT END-OF-MEDIUM
3B 03 L TAPE OR ELECTRONIC VERTICAL FORMS UNIT NOT READY
3B 04 L SLEW FAILURE
3B 05 L PAPER JAM
3B 06 L FAILED TO SENSE TOP-OF-FORM
3B 07 L FAILED TO SENSE BOTTOM-OF-FORM
3B 08 T REPOSITION ERROR
3B 09 S READ PAST END OF MEDIUM
3B 0A S READ PAST BEGINNING OF MEDIUM
3B 0B S POSITION PAST END OF MEDIUM
3B 0C S POSITION PAST BEGINNING OF MEDIUM
3B 0D M MEDIUM DESTINATION ELEMENT FULL
3B 0E M MEDIUM SOURCE ELEMENT EMPTY
3C 00
3D 00 DTLPWSOMC INVALID BITS IN IDENTIFY MESSAGE
3E 00 DTLPWSOMC LOGICAL UNIT HAS NOT SELF-CONFIGURED YET
3F 00 DTLPWSOMC TARGET OPERATING CONDITIONS HAVE CHANGED
3F 01 DTLPWSOMC MICROCODE HAS BEEN CHANGED
3F 02 DTLPWSOMC CHANGED OPERATING DEFINITION
3F 03 DTLPWSOMC INQUIRY DATA HAS CHANGED
40 00 D RAM FAILURE
40 NN DTLPWSOMC DIAGNOSTIC FAILURE ON COMPONENT NN (80H-FFH)
                                SHOULD USE 40 NN
                                USE ASCQ CODES 80H-FFH
41 00 D DATA PATH FAILURE
42 00 D POWER-ON OR SELF-TEST FAILURE
                                SHOULD USE 40 NN
                                SHOULD USE 40 NN
43 00 DTLPWSOMC MESSAGE ERROR
44 00 DTLPWSOMC INTERNAL TARGET FAILURE
45 00 DTLPWSOMC SELECT/RESELECT FAILURE
46 00 DTLPWSOMC UNSUCCESSFUL SOFT RESET
47 00 DTLPWSOMC SCSI PARITY ERROR
48 00 DTLPWSOMC INITIATOR DETECTED ERROR MESSAGE RECEIVED
49 00 DTLPWSOMC INVALID MESSAGE ERROR
4A 00 DTLPWSOMC COMMAND PHASE ERROR
4B 00 DTLPWSOMC DATA PHASE ERROR
4C 00 DTLPWSOMC LOGICAL UNIT FAILED SELF-CONFIGURATION
4D 00
4E 00 DTLPWSOMC OVERLAPPED COMMANDS ATTEMPTED
4F 00
50 00 T WRITE APPEND ERROR
50 01 T WRITE APPEND POSITION ERROR
50 02 T POSITION ERROR RELATED TO TIMING
51 00 T 0 ERASE FAILURE
52 00 T CARTRIDGE FAULT
=====

```

Table I-1: ASC and ASCQ Assignments (Continued)

```

=====
53 00 DTL WRSOMC MEDIA LOAD/EJECT FAILED
53 01 T UNLOAD TAPE FAILURE
53 02 DT WR OM MEDIUM REMOVAL PREVENTED
54 00 P SCSI TO HOST SYSTEM INTERFACE FAILURE
55 00 P SYSTEM RESOURCE FAILURE
56 00
57 00 R UNABLE TO RECOVER TABLE-OF-CONTENTS
58 00 0 GENERATION DOES NOT EXIST
59 00 0 UPDATED BLOCK READ
5A 00 DTLPWSOMC OPERATOR REQUEST OR STATE CHANGE INPUT (UNSPECIFIED)
5A 01 DT WR OM OPERATOR MEDIUM REMOVAL REQUEST
5A 02 DT W 0 OPERATOR SELECTED WRITE PROTECT
5A 03 DT W 0 OPERATOR SELECTED WRITE PERMIT
5B 00 DTLPWSOMC LOG EXCEPTION
5B 01 DTLPWSOMC THRESHOLD CONDITION MET
5B 02 DTLPWSOMC LOG COUNTER AT MAXIMUM
5B 03 DTLPWSOMC LOG LIST CODES EXHAUSTED
5C 00 D 0 RPL STATUS CHANGE
5C 01 D 0 SPINDLES SYNCHRONIZED
5C 02 D 0 SPINDLES NOT SYNCHRONIZED
5D 00
5E 00
5F 00
60 00 S LAMP FAILURE
61 00 S VIDEO ACQUISITION ERROR
61 01 S UNABLE TO ACQUIRE VIDEO
61 02 S OUT OF FOCUS
62 00 S SCAN HEAD POSITIONING ERROR
63 00 R END OF USER AREA ENCOUNTERED ON THIS TRACK
64 00 R ILLEGAL MODE FOR THIS TRACK
65 00
66 00
67 00
68 00
69 00
6A 00
6B 00
6C 00
6D 00
6E 00
6F 00
=====

```

Table I-1: ASC and ASCQ Assignments (Continued)

70 00		
71 00		
72 00		
73 00		
74 00		
75 00		
76 00		
77 00		
78 00		
79 00		
7A 00		
7B 00		
7C 00		
7D 00		
7E 00		
7F 00		

80 XX	\	
THROUGH	>	VENDOR UNIQUE.
FF XX	/	
XX 80	\	
THROUGH	>	VENDOR UNIQUE QUALIFICATION OF STANDARD ASC.
XX FF	/	
ALL CODES NOT SHOWN OR BLANK ARE RESERVED.		
=====		

Table I-2: SCSI-2 Operation Codes

=====

F = Flags: C = conflict in opcode usage for different devices

Device Columns M = Mandatory
 Key: O = Optional
 V = Vendor unique
 = Reserved (Blank)

D = Direct-Access Device
 T = Sequential-Access Device
 L = Printer Device
 P = Processor Device
 W = Write-Once Device
 R = CD-ROM Device
 S = Scanner Device
 O = Optical Memory Device
 M = Medium Changer Device
 C = Communication Device

F	OP	DTLPWRSOMC	Description	Comments

	00	MMMMMMMM	TEST UNIT READY	
C	01	M	REWIND	
C	01	O V 00 00	REZERO UNIT	
	02	VVVVVV V		
	03	MMMMMMMM	REQUEST SENSE	
C	04	O	FORMAT	
C	04	M O	FORMAT UNIT	
	05	VMVVVV V	READ BLOCK LIMITS	
	06	VVVVVV V		
	07	O	INITIALIZE ELEMENT STATUS	
	07	OVV O OV	REASSIGN BLOCKS	
C	08	M	GET MESSAGE (6)	
C	08	OMV 00 OV	READ (6)	
C	08	O	RECEIVE	
	09	VVVVVV V		
C	0A	M	PRINT	
C	0A	M	SEND	
C	0A	M	SEND MESSAGE (6)	
C	0A	OM O OV	WRITE (6)	
C	0B	O 00 OV	SEEK (6)	
C	0B	O	SLEW AND PRINT	
	0C	VVVVVV V		
	0D	VVVVVV V		
	0E	VVVVVV V		
	0F	VOVVVV V	READ REVERSE	
=====				

Table I-2: SCSI-2 Operation Codes (Continued)

```

=====
C 10 0 0 SYNCHRONIZE BUFFER
C 10 VM VVV WRITE FILEMARKS
11 VMVVVV SPACE
12 MMMMMMMMM INQUIRY
13 VOVVVV VERIFY
14 VOOVVV RECOVER BUFFERED DATA
15 OMO 000000 MODE SELECT (6)
C 16 M MM MO RESERVE
C 16 MM M RESERVE UNIT
C 17 M MM MO RELEASE
C 17 MM M RELEASE UNIT
18 00000000 COPY
19 VMVVVV ERASE
1A OMO 000000 MODE SENSE (6)
C 1B 0 LOAD/UNLOAD
C 1B 0 SCAN
C 1B 0 STOP PRINT
C 1B 0 00 0 STOP START UNIT
1C 000000000 RECEIVE DIAGNOSTIC RESULTS
1D MMMMMMMMM SEND DIAGNOSTIC
1E 00 00 00 PREVENT/ALLOW MEDIUM REMOVAL
1F
20 V VV V
21 V VV V
22 V VV V
23 V VV V
24 V VVM DEFINE WINDOW PARAMETERS
C 25 0 GET WINDOW PARAMETERS
C 25 M M M READ CAPACITY
C 25 M M READ CD-ROM CAPACITY
26 V VV
27 V VV
C 28 0 GET MESSAGE (10)
C 28 M MMM READ (10)
29 V VV 0 READ GENERATION
C 2A 0 SEND
C 2A 0 0 SEND MESSAGE (10)
C 2A M M M WRITE (10)
C 2B 0 LOCATE
2B 0 POSITION TO ELEMENT
C 2B 0 00 0 SEEK (10)
2C V 0 ERASE (10)
2D V 0 0 READ UPDATED BLOCK
2E 0 0 0 WRITE AND VERIFY (10)
2F 0 00 0 VERIFY (10)
=====

```

Table I-2: SCSI-2 Operation Codes (Continued)

```

=====
30 0 00 0 SEARCH DATA HIGH (10)
C 31 0 MEDIUM POSITION
C 31 0 00 0 SEARCH DATA EQUAL (10)
32 0 00 0 SEARCH DATA LOW (10)
33 0 00 0 SET LIMITS (10)
C 34 0 GET DATA STATUS
C 34 0 00 0 PRE-FETCH
C 34 0 READ POSITION
35 0 00 0 SYNCHRONIZE CACHE
36 0 00 0 LOCK/UNLOCK CACHE
37 0 0 READ DEFECT DATA
C 38 0 0 MEDIUM SCAN
39 00000000 COMPARE
3A 00000000 COPY AND VERIFY
3B 0000000000 WRITE BUFFER
3C 0000000000 READ BUFFER
3D 0 0 UPDATE BLOCK
3E 0 00 0 READ LONG
3F 0 0 0 WRITE LONG
40 0000000000 CHANGE DEFINITION
41 0 WRITE SAME
42 0 READ SUB-CHANNEL
43 0 READ TOC
44 0 READ HEADER
45 0 PLAY AUDIO (10)
46
47 0 PLAY AUDIO MSF
48 0 PLAY AUDIO TRACK/INDEX
49 0 PLAY TRACK RELATIVE (10)
4A
4B 0 PAUSE/RESUME
4C 0000000000 LOG SELECT
4D 0000000000 LOG SENSE
4E
4F
50
51
52
53
54
55 000 000000 MODE SELECT (10)
56
57
58
59
5A 000 000000 MODE SENSE (10)
5B
5C
5D
5E
5F
=====

```

Table I-2: SCSI-2 Operation Codes (Continued)

A0			
A1			
A2			
A3			
C A5	M	MOVE MEDIUM	
C A5	0	PLAY AUDIO (12)	
A6	0	EXCHANGE MEDIUM	
A7			
C A8	0	GET MESSAGE (12)	
A8	00 0	READ (12)	
A9	0	PLAY TRACK RELATIVE (12)	
AA	0	SEND MESSAGE (12)	
C AA	0 0	WRITE (12)	
AB			
AC	0	ERASE (12)	
AE	0 0	WRITE AND VERIFY (12)	
AF	00 0	VERIFY (12)	
B0	00 0	SEARCH DATA HIGH (12)	
B1	00 0	SEARCH DATA EQUAL (12)	
B2	00 0	SEARCH DATA LOW (12)	
B3	00 0	SET LIMITS (12)	
B4			
B5			
B5	0	REQUEST VOLUME ELEMENT ADDRESS	
B6			
B6	0	SEND VOLUME TAG	
B7	0	READ DEFECT DATA (12)	
B8			
B8	0	READ ELEMENT STATUS	
B9			
BA			
BB			
BC			
BE			
BF			

J. Vendor Identification

This Appendix contains the list of SCSI-2 vendor identifications as of the date of this document. The purpose of this list is to help avoid redundant usage of vendor identifications. Task Group X3T9.2 of Accredited Standards Committee X3 maintains an informal list of vendor identifications currently in use. Please contact the chairman of X3T9.2 prior to using a new vendor identification to avoid conflicts.

Table J-1: Vendor Identification List

ID	Organization
3M	3M Company
ADAPTEC	Adaptec
ADSI	Adaptive Data Systems, Inc. (a Western Digital subsidiary)
AMCODYNE	Amcodyne
ANAMATIC	Anamartic Limited (England)
ANCOT	ANCOT Corp.
ANRITSU	Anritsu Corporation
APPLE	Apple Computer, Inc.
ARCHIVE	Archive
ASPEN	Aspen Peripherals
AST	AST Research
AT&T	AT&T
BALLARD	Ballard Synergy Corp.
BULL	Bull Peripherals Corp.
CALIPER	Caliper (California Peripheral Corp.)
CAST	Advanced Storage Tech
CDC	Control Data or MPI
CHEROKEE	Cherokee Data Systems
CIE&YED	YE Data, C.Itoh Electric Corp.
CIPHER	Cipher Data Products
Ciprico	Ciprico, Inc.
CNGR SFW	Congruent Software, Inc.
COGITO	Cogito
COMPORT	Comport Corp.
CONNER	Conner Peripherals
CROSFLD	Crosfield Electronics
CSM, INC	Computer SM, Inc.
CYGNET	Cygnat Systems, Inc.
DATACOPY	Datacopy Corp.
DATAPT	Datapoint Corp.
DEC	Digital Equipment
DENON	Denon/Nippon Columbia
DEST	DEST Corp.
DGC	Data General Corp.

Table J-1: Vendor Identification List (Continued)

ID	Organization
DIGIDATA	Digi-Data Corporation
DILOG	Distributed Logic Corp.
DTC	Data Technology Corp.
DPT	Distributed Processing Technology
DXIMAGIN	DX Imaging
EMULEX	Emulex
EPSON	Epson
EXABYTE	Exabyte Corp.
FILENET	FileNet Corp.
FUJI	Fuji Electric Co., Ltd. (Japan)
FUJITSU	Fujitsu
FUTURE	Future Domain Corp.
GIGATAPE	GIGATAPE GmbH
GIGATEND	GigaTrend Incorporated
Goidelic	Goidelic Precision, Inc.
GOULD	Gould
HITACHI	Hitachi America Ltd or Nissei Sangyo America Ltd
HP	Hewlett Packard
IBM	International Business Machines
ICL	ICL
IGR	Intergraph Corp.
IMPRIMIS	Imprimis Technology Inc.
IOMEGA	Iomega
ISI	Information Storage inc.
JVC	JVC Information Products Co.
KODAK	Eastman Kodak
KONAN	Konan
KONICA	Konica Japan
LAPINE	Lapine Technology
LASERDRV	LaserDrive Limited
LMS	Laser Magnetic Storage International Company
MATSHITA	Matsushita
MAXTOR	Maxtor Corp.
MELA	Mitsubishi Electronics America
MELCO	Mitsubishi Electric (Japan)
MICROBTX	Microbotics Inc.
MICROP	Micropolis
MICROTEK	Microtek Storage Corp
MINSCRIB	Miniscribe
MOTOROLA	Motorola
NAI	North Atlantic Industries
NatSemi	National Semiconductor Corp.
NCL	NCL America

Table J-1: Vendor Identification List (Continued)

ID	Organization
NCR	NCR Corporation
NEC	NEC
NKK	NKK Corp.
NT	Northern Telecom
OSI	Optical Storage International
OPTIMEM	Cipher/Optimem
OPTOTECH	Optotech
OTL	OTL Engineering
PERTEC	Pertec Peripherals Corporation
PFTI	Performance Technology Inc.
PRAIRIE	PrairieTek
PTI	Peripheral Technology Inc.
PRIAM	Priam
QUALSTAR	Qualstar
QUANTUM	Quantum Corp.
RADSTONE	Radstone Technology
RICOH	Ricoh
RODIME	Rodime
RTI	Reference Technology
SEAGATE	Seagate
SIEMENS	Siemens
SMS	Scientific Micro Systems/OMTI
SONY	Sony Corporation Japan
SPERRY	Sperry (now Unisys Corp.)
STK	Storage Technology Corporation
SUN	Sun Microsystems, Inc.
SYSGEN	Sysgen
T-MITTON	Transmitton England
TANDBERG	Tandberg Data A/S
TANDON	Tandon
TEAC	TEAC Japan
TI-DSG	Texas Instruments
TOSHIBA	Toshiba Japan
UNISYS	Unisys
USDC	US Design Corp.
VERBATIM	Verbatim Corporation
VRC	Vermont Research Corp.
WangDAT	WangDAT
WANGTEK	Wangtek
WDIGTL	Western Digital
XEBEC	Xebec Corporation